

Epistemic Semantics in Guarded String Models

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Abstract

Constructive and computable multi-agent epistemic possible worlds models are as sets of guarded string models in an epistemic extension of Kleene Algebra with Tests (KAT). The account is framed as a formal language Epik (Epistemic KAT) for defining such models. The language is implemented by translation into the finite state calculus, and alternatively by modeling propositions as lazy lists in Haskell. The syntax-semantics interface for a fragment of English is defined by a categorial grammar.

1 Introduction and Related Work

Linguistic semantics in the Montague tradition proceeds by assigning propositional *semantic values* to disambiguated sentences of a natural language. A proposition is a set or class of *possible worlds*, which are often assumed to have the same nature and complexity as the world we occupy (Lewis, 1986). But alternatively, one can work with small idealized models, to illustrate and test ideas. The point of this paper is to extend standard idealized models to countable sets of worlds, and to constructively and computably model alternatives for epistemic agents.

Enter Epik, which is a systematic way of defining such models, and can be applied to natural language semantics, specifically, epistemic semantics and clausal embedding. The fundamental insight is to identify possible worlds with strings of primitive events, so that propositions are sets of (guarded) strings, whose regular sets have a rich

algebraic characterization (Kozen, 1997) and computational model (Kozen and Smith, 1997; Kozen, 2001). Epik leverages these algebras to define a mathematical model (Sections 2 & 3) and computationally interpret using finite state machines (Section 4) and lazy lists of strings (Section 5)).

Related Work. Epik synthesizes five antecedents. John McCarthy’s *Situation Calculus* is the source of the idea of constructing possible worlds as event sequences (McCarthy, 1963; Reiter, 2001). The algebraic theory of *Kleene Algebra with Tests* characterizes algebras with elements corresponding to propositions and event types in our application (Kozen, 2001). *Action models* in dynamic epistemic semantics introduce the technique of constructing epistemic models from primitive alternative relations on events, in order to capture the epistemic consequences of perceptual and communicative events (Baltag et al., 1999). Literature on *finite state methods in linguistic semantics* has used event strings and sets of event strings to theorize about tense and aspect in natural language semantics (Fernando, 2004, 2007; Carlson, 2009) and to express intensions (Fernando, 2017). Work on *finite state intensional semantics* has investigated how to do the semantics of intensional complementation in a setting where compositional semantics is expressed in a finite state calculus (Rooth, 2017; Collard, 2018).

Example. Consider an example event-sequence model called *The Concealed Coin*. Amy and Bob are seated at a table. There is a coin on the table under a cup, heads up. The coin could be heads-up H or tails-up T , and neither agent knows which it is. Call this initial possible world w_1 . Additional worlds can be constructed via the events in (1).

- (1) a_1 Amy peeks at H , by tipping the cup. Bob sees she’s peeking, but not what she sees.
- b_1 Bob peeks at H .

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- a_0 Amy peeks at T .
- b_0 Bob peeks at T .
- a_{01} Amy secretly turns the coin from T to H . She knows she turned the coin over, but not which side was face up. Bob thinks nothing happened.
- a_{10} Amy secretly turns the coin from H to T .
- b_{01} Bob secretly turns the coin from T to H .
- b_{10} Bob secretly turns the coin from H to T .

The worlds in (2) are examples constructed from the events in (1). Juxtaposition indicates the order of events; for example, w_3 indicates that after the events of w_2 , event b_1 occurs, which corresponds to, in the initial world, Amy surreptitiously peeking at a heads-up coin, and then Bob doing the same.

$$(2) \quad w_2 \triangleq w_1 a_1 \quad w_3 \triangleq w_2 b_1 \quad w_4 \triangleq w_3 a_{10} b_{01} b_1$$

Now, the truth (1) or falsity (0) of English sentences can be evaluated in each world w_i , i.e. *after* the events have happened. See (3) for examples, such as in w_4 , where Bob knows it's heads (because he peeked after turning the coin), but Amy does not (because Bob secretly turned the coin over).

(3)	w_1	w_2	w_3	w_4	Sentence
	0	1	1	0	Amy knows it's heads.
	0	0	1	1	Bob knows it's heads.
	0	0	1	0	Bob knows Amy knows it's heads.
	0	1	1	0	Bob knows Amy knows whether it's heads or tails.

To create realistic models, contradictory event sequences must be prohibited. For example, $a_1 a_{01}$ would require the coin to simultaneously heads-up and tails-up. To mechanize this, events in Epik come with Boolean pre-conditions and post-conditions. In the running example, the Boolean variables h (the coin is heads-up) and t (the coin is tails-up) represent the coin's state. Then, since Amy can only observe heads when the coin is heads-up, event a_1 has the invariant condition h . Similarly, Amy can only turn the coin from tails to heads if it shows tails, so a_{01} has the tails-up precondition t and the heads-up postcondition h .

Pre- and post- conditions are expressed using an operator “:” (read “and next”) that pairs Boolean formulas into an *effect formula*. Specifically, the effects for a_1 are written $h : h$, and for a_{01} are written $t : h$. Effect formulae are interpreted relations

on boolean valuations, as defined in Figure 1.

However, a coin cannot simultaneously show both heads and tails! Currently the precondition h of a_{10} only says that heads must be showing, and says nothing about the fact that tails must be face-down, indicated by the formula \bar{t} . The effect valuations are restricted to only the feasible ones via a *state formula*, shown in 4, and demonstrated for a_1 in 4. The state formula, also shown in 4, says that the coin is either heads-up or tails-up, but not both¹. Here juxtaposition represents conjunction, $+$ is disjunction, and overbar is negation.

With the state formula in hand, $a_1 a_{01}$ is evidently contradictory, the post-condition h of a_1 is incompatible with the pre-condition t of a_{01} . See Figure 1 for more details.

Given a set \mathbf{B} of state primitives and φ a state formula over \mathbf{B} , define $\mathcal{A}_{\mathbf{B}}^{\varphi}$ to be the set of valuations of \mathbf{B} that make formula φ true. The valuations for the coin example are shown in 4, using the sequence notation for valuations, e.g. $\bar{h}t$, wherein every state primitive is listed in fixed order, and left unmarked (indicating true) or marked with the overbar (indicating false). Valuations are called atoms, because they correspond to the atoms of a Boolean algebra of tests (Kozen, 2001).

(4)	\mathbf{B}	$\{h, t\}$
	state formula φ	$h\bar{t} + \bar{h}t$
	effect formula ζ_{a_1} for a_1	$h : h$
	$\mathcal{A}_{\mathbf{B}}^{\varphi}$	$\{\bar{h}\bar{t}, \bar{h}t\}$
	$\llbracket \zeta_{a_1} \rrbracket^{\varphi}$	$\{\langle \bar{h}\bar{t}, \bar{h}\bar{t} \rangle\}$

Thus far, the presentation is closely related to Kleene Algebra with Hypotheses (KA+H) (Cohen, 1994), which permits user-specification of further equations beyond the axioms of KA. Specifically, the state formula φ can be thought of as a hypothesis $\varphi \equiv 1$, and an event formula ζ_a for event a , is simply the hypothesis $a \equiv \sum_{\varphi, \varphi' \in \llbracket \zeta_a \rrbracket} \varphi; a; \varphi'$. Rather than rely on general techniques (Doumane et al., 2019; Hardin, 2002; Kozen and Smith, 1997), we realize our specific hypotheses directly.

The key facet that differentiates Epik from KA+H is the epistemic modality, which cannot be naively represented using universal hypotheses². Sections 2 and 3 discuss the model-theoretic and logical development of the epistemic operators.

¹Two Booleans constrained by a state formula, is a more illustrative example than using a single boolean and modelling tails as \bar{h} .

²The existence of this encoding is an interesting future topic

state formulas	$(a \in B)$
$\rho, \sigma, \varphi ::=$	$a \mid 0 \mid 1 \mid \rho + \sigma \mid \rho \sigma \mid \bar{\rho}$
effect formulas	
$\zeta, \eta ::=$	$\rho : \sigma \mid \zeta + \eta \mid \zeta \& \eta \mid \bar{\zeta}$
$\llbracket \rho : \sigma \rrbracket^\varphi \triangleq$	$\mathcal{A}_B^{\rho\varphi} \times \mathcal{A}_B^{\sigma\varphi}$
$\llbracket \zeta + \eta \rrbracket^\varphi \triangleq$	$\llbracket \zeta \rrbracket^\varphi \cup \llbracket \eta \rrbracket^\varphi$
$\llbracket \zeta \& \eta \rrbracket^\varphi \triangleq$	$\llbracket \zeta \rrbracket^\varphi \cap \llbracket \eta \rrbracket^\varphi$
$\llbracket \bar{\zeta} \rrbracket^\varphi \triangleq$	$\mathcal{A}_B^\varphi \times \mathcal{A}_B^\varphi \setminus \llbracket \zeta \rrbracket^\varphi$

Figure 1: Syntax of state formulas and syntax and semantics of effect formulas. Effect formulas denote relations between atoms. In a state formula, juxtaposition $\rho\sigma$ is conjunction.

2 Epistemic guarded string models

Epik is a specification language for possible worlds models that includes declarations of events and states, state formulas, effect formulas, and additional information. Figure 2 shows an Epik program that describes a possible worlds model for two agents with information about one coin, events of the agents semi-privately looking at the coin, and events of secretly turning the coin. The line beginning with `state` enumerates B . The line beginning with `restrict` gives the state formula. The lines beginning with `event` declare events and their effect formulas.

Finally, the lines beginning with `agent` define *event alternative* relations for the epistemic agents in the model. Each clause with an arrow has a single event symbol on the left, and a disjunction of alternative events on the right of the arrow. The interpretation of Amy’s alternatives for b_1 (Bob peeks at heads), is that when b_1 happens, for Amy either b_1 or b_0 (Bob peeks at tails) could be happening, indicating that she doesn’t know whether Bob saw heads or tails, only that he peeked at the coin.

Similarly, Her alternatives for a_{01} (she turns the coin over from heads to tails) are a_{01} and a_{10} , indicating that she doesn’t know, *a priori*, whether she’s turning the coin from H to T or from T to H . She has the same alternatives for the event a_{10} . Conversely, When Bob secretly turns the coin over, in event b_{01} or b_{10} , she doesn’t know anything has happened, so her alternative is the “no-operation” or “no-information” event `event o`. Bob’s event relation is symmetric.

This paper focuses on defining a concrete possible worlds model from an Epik specification. The models are an extension of guarded-string models for Kleene Algebra with Tests (KAT). This algebraic theory has model classes including guarded

state h t	agent amy
restrict h!t	o -> o
+ t!h	a1 -> a1
event o h:h + t:t	a0 -> a0
event a1 h:h	b1 -> b1 + b0
event a0 t:t	b0 -> b1 + b0
event b1 h:h	a10 -> a10 + a01
event b0 t:t	a01 -> a10 + a01
event a10 h:t	b10 -> o0 + o1
event a01 t:h	b01 -> o0 + o1
event b10 h:t	agent bob
event b01 t:h	<sim. swap a and b>

Figure 2: Epik program describing a possible-worlds event sequence model for two agents with information about one coin, and events of the agents semi-privately looking at the coin, and privately turning the coin.

string models, relational models, finite models, and matrix models. Our definitions and notation follow (Kozen, 2001). We add syntax and semantics to cover multi-agent epistemic semantics.

Guarded strings over a finite alphabet P are like ordinary strings, but with atoms over a set B alternating with the symbols from P . In the algebra described by Figure 2, P is the set of events $\{a_1, a_0, b_1, b_0, a_{01}, a_{10}, b_{01}, b_{10}, o\}$, and B is $\{h, t\}$. As we already saw in (4), \mathcal{A}_B^φ is $\{h\bar{t}, \bar{h}t\}$, for which we use the shorthand $\{H, T\}$. A guarded string over P and B is a string of events from P , alternating with atoms over B , and beginning and ending with atoms. In this construction, $w_1 = H$, $w_2 = Ha_1H$, $w_3 = Ha_1Hb_1H$, and $w_4 = Ha_1Hb_1Ha_{10}Tb_{01}Hb_1H$.

The discussion of (2) mentioned building worlds by incrementing worlds with events. This is accomplished in guarded string models with fusion product \diamond , a partial operation that combines two guarded strings, subject to the condition that the atom at the end of the the first argument is identical to the atom at the start of the second one. (5) gives some examples.

$$(5) \quad \begin{aligned} H b_1 H \diamond H a_1 H &= H b_1 H a_1 H \\ T b_{01} H \diamond T a_1 T &= \text{undefined} \end{aligned}$$

Rather than individual guarded strings, elements of a guarded string model for KAT are sets of guarded strings. In our application, these elements have the interpretation of propositions, which are sets of possible worlds. In a free guarded string model for KAT, any event can be adjacent to any atom in a guarded string that is an element of the underlying set for the algebra. We instead impose the constraints coming from the state and effect for-

mulas. (6) defines the well-formed guarded strings determined by an Epik specification. Condition (i) says that each atom is consistent with the state constraint, and condition (ii) says that each constituent token event $\alpha_i e_i \alpha_{i+1}$ is consistent with the effect constraint on e_i .³

- (6) Given P , B , a state formula φ , and an effect formula ζ_e for each event e in P , $\alpha_0 e_0 \dots e_n \alpha_{n+1}$ is well-formed iff
- (i) $\alpha_i \in \mathcal{A}_B^\varphi$ ($0 \leq i \leq n$), and
 - (ii) $\langle \alpha_i, \alpha_{i+1} \rangle \in \llbracket \zeta_{e_i} \rrbracket^\varphi$, ($0 \leq i \leq n$).

Well-formed guarded strings have the interpretation of worlds in the application to natural-language semantics. The set of possible worlds in the Kripke frame determined by an Epik specification is the set of well-formed guarded strings. At this point, we could say that any set of worlds is a proposition, so that the set of propositions is the power set of the set of worlds (Montague, 1975; Gallin, 1975). We will instead define a more restrictive set of propositions corresponding to the regular sets of strings. This is deferred to the next section. Certain sets of well-formed guarded strings have the additional interpretation of event types. An event-type is something that can “happen” in different worlds. For example, a_1 corresponds to the event type $\{H a_1 H\}$, and a_0 corresponds to the event type $\{T a_0 T\}$. Event types need not be singletons.

The construction so far defines a set of worlds from an Epik specification. Normally the set is countably infinite, though some choices of effect formulas can result in a finite set of worlds. The next step is to define an alternative relation R_a on worlds for each agent a . This will result in a general modal frame $\langle W, R_1, \dots, R_n, K \rangle$ consisting of a set of worlds, a world-alternative relation for each agent, and a set K of propositions, where each proposition is a subset of W (Chagrov, 1997).⁴

An Epik specification defines an alternative relation on bare events for each agent a , which is notated R_a , and lifted to a relation \hat{R}_a on worlds.

³An alternative is to define equations such as $\bar{\phi} = 0$ (from the state formula ϕ) and $a_1 = h a_1 h$ (from the effect formula $h : h$ for event a_1), and construct a quotient algebra from the equivalence relation generated by these equations. This results in equating sets of guarded strings in the free algebra that differ by guarded strings that are ill-formed according to the state and effect formulas. In the development in the text, we instead use a set of guarded strings that are well-formed according to the state and effect formulas as the representative of the equivalence class.

⁴As explained in the next section, K will not be the power set of W , rather it will consist of the regular subsets of W .

The basic idea is that when a world w is incremented with an event e , in the resulting world $w \diamond e$, epistemic alternatives for agent a are of the form $w' \diamond e'$, where w' is an alternative to w for a in w , and e' is an event-alternative to e for a .⁵ This needs to be implemented in a way that takes account of pre- and post-conditions for events. For this, our approach is to refer the definition of well-formed guarded strings. (7) defines an epistemic alternative relation on worlds from an alternative relation on bare events.

- (7) Let W be a set of guarded strings over events P and primitive tests B , and R be a relation on P . The corresponding relation \hat{R} on W holds between a guarded string $\alpha_0 e_0 \dots e_n \alpha_{n+1}$ in W and a guarded string q in W iff q is of the form $\alpha'_0 e'_0 \dots e'_n \alpha'_{n+1}$, where for $0 \leq n$, $e_i R e'_i$.

This requires that in an alternative world, each constituent event e'_i is an alternative to the corresponding event e_i in the base world. Compatibilities between events in the alternative world are enforced by the requirement that the alternative world is an element of W , so that state and effect formulas are enforced.

Consider a scenario like the one from Figure 1, but with an additional agent Cal. The base world Tb_0Tc_0T is one where the coin is tails, and first Bob looks at tails, and then Cal looks at tails. The first event b_0 has the alternatives b_0 and b_1 for Amy, and the second event c_0 has the alternatives c_0 and c_1 for Amy. This results in four combinations b_0c_0 , b_0c_1 , b_1c_0 , and b_1c_1 . But some of these worlds are bogus, such as b_0c_1 , which says that Bob peeked at the coin and saw tails, and Cal peeked at the same coin and impossibly saw heads. These contradictions are filtered by post- and pre-conditions of events in the alternative world, so that the set of alternatives for Amy in Tb_0Tc_0T is $\{Tb_0Tc_0T, Hb_1Hc_1H\}$, with two world-alternatives instead of four.

⁵In this it is important that the event-alternative relation for an agent is constant across worlds. We anticipate that the definition given here produces results equivalent to what is found in literature on event alternatives in dynamic epistemic semantics, though we have not verified this. That literature primarily focuses on mapping an epistemic model for a single time and situation to another, and uses general first-order models, rather than guarded string models. See Baltag et al. (1999), Van Ditmarsch et al. (2007), and articles in Van Ditmarsch et al. (2015). This literature is motivated by epistemic logic and AI planning, rather than computable possible worlds models in natural language semantics.

events	$e \in P$
p, q	$::= e \mid \sigma \mid p + q \mid pq \mid p^* \mid \neg p \mid \Diamond_a p$
$\Box_a p$	$\triangleq \neg \Diamond_a \neg p$
$p \wedge q$	$\triangleq \neg(\neg p + \neg q)$
$p \rightarrow q$	$\triangleq \neg p + q$
\bullet	$\triangleq \sum_{e \in P} e$

Figure 3: The language of Epik terms and key derived operators.

3 The logical language of Epistemic KAT

The standard language for Kleene algebra with tests has the signature $\langle K, +, \cdot, *, \neg, 0, 1 \rangle$ (Kozen, 2001). In a guarded string model for KAT, K is a set of sets of guarded strings, $+$ is set union, the operation \cdot is fusion product raised to sets, $*$ is Kleene star, the operation \neg is complement for tests, 0 is the empty set, and 1 is the set of atoms.⁶ To this we add a unary modal operation \Diamond_a for each agent, and a unary complement operation \neg on elements of K . Intuitively, $\Diamond_a p$ is the set of worlds where proposition p is epistemically possible for agent a . Propositional complement is included because natural languages have sentence negation. In addition, universal box modalities are defined as duals of existential diamond modalities.

With modalities and propositional negation added, the signature of n -agent epistemic KAT is $\langle K, +, \cdot, *, \neg, 0, 1, \neg, \Diamond_1, \dots, \Diamond_n \rangle$. Figure 3 defines the syntax of the language. Juxtaposition is used for product. Terms in this language are used to represent the propositional semantic values of English sentences. (8) gives some examples. To explain the first one, \bullet as defined in Figure 3 is the disjunction of the primitive events. Since a world is a well-formed sequence of events, \bullet^* is the set of worlds. Multiplying by the state symbol h in the term \bullet^*h has the effect of conjoining h with the atom at the end of the world. So \bullet^*h is the set of worlds where the coin ends heads-up.

- (8) \bullet^*t It's tails.
 \bullet^*h It's heads.
 $\bullet^*h \wedge \Box_a \bullet^*h$ Amy knows it's heads.
 $\Box_b(\bullet^*t \wedge \Box_a \bullet^*t + \neg \bullet^*t \wedge \Box_a \neg \bullet^*t)$
Bob believes Amy knows whether it's tails.

Standard Epistemic Modalities. Using our existential modal primitive \Diamond_a , and the dual encodings of \Box_a and \wedge , we can encode the standard modal operators expressing knowledge (\mathcal{K}_a) and belief

⁶0 has the dual role the identity for $+$ (union), and as False for operations on tests. 1 has the dual role of the identity for product (fusion product raised to sets), and True for tests.

$\llbracket 0 \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \emptyset$
$\llbracket 1 \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \mathcal{A}_B^\varphi$
$\llbracket b \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \mathcal{A}_B^{b,\varphi}$
$\llbracket \sigma \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \mathcal{A}_B^\varphi \setminus \llbracket \sigma \rrbracket^{B,P,\varphi,\zeta}$
$\llbracket e \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \{\alpha e \beta \mid \alpha \zeta_e \beta\}$
$\llbracket p + q \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \llbracket p \rrbracket^{B,P,\varphi,\zeta} \cup \llbracket q \rrbracket^{B,P,\varphi,\zeta}$
$\llbracket pq \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \left\{ x \diamond y \mid \begin{array}{l} x \in \llbracket p \rrbracket^{B,P,\varphi,\zeta} \\ y \in \llbracket q \rrbracket^{B,P,\varphi,\zeta} \\ x \diamond y \text{ is defined} \end{array} \right\}$
$\llbracket p^* \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \bigcup_{n \geq 0} \llbracket p^n \rrbracket^{B,P,\varphi,\zeta}$
$\llbracket \neg p \rrbracket^{B,P,\varphi,\zeta}$	$\triangleq \llbracket \bullet^* \rrbracket^{B,P,\varphi,\zeta} \setminus \llbracket p \rrbracket^{B,P,\varphi,\zeta}$
$\llbracket \Diamond_a p \rrbracket$	$\triangleq \{x \mid \exists y. x \hat{R}_a y \wedge y \in \llbracket p \rrbracket^{B,P,\varphi,\zeta}\}$

Figure 4: Interpretation of Epik terms as sets of guarded strings

(\mathcal{B}_a) as in (9).⁷

- (9) BELIEF $\mathcal{B}_a p \triangleq \Box_a p$
 KNOWLEDGE $\mathcal{K}_a p \triangleq p \wedge \mathcal{B}_a p$

Different types of reasoners (e.g. accurate, inaccurate, etc) are modeled using the event-alternatives in an Epik specification.⁸ The agents in Figure 1 do not always have reliable beliefs, because of the possibility of secret turning.

Guarded String Interpretation. A term p of the logical language is interpreted as a set of guarded strings $\llbracket p \rrbracket^{B,P,\varphi,\zeta}$, where superscript captures dependence on an Epik specification. Figure 4 defines the interpretation. The interpretation $\llbracket 1 \rrbracket^{B,P,\varphi,\zeta}$ of the multiplicative identity 1 is the set of atoms that satisfy the state constraint φ . Where b is a primitive Boolean, $\llbracket b \rrbracket^{B,P,\varphi,\zeta}$ is the set of atoms that satisfy the state constraint and where b is true. Where e is a primitive event, $\llbracket e \rrbracket^{B,P,\varphi,\zeta}$ is the set of guarded strings that have the form of e flanked by compatible atoms, as determined by the event formula ζ_e . The product pq is interpreted with fusion product raised to sets of guarded strings. Kleene star is interpreted as the union of exponents (p^n is the n -times product of p with itself, with $p^0 = 1$). Propositional complement is complement relative to the set of worlds. The epistemic formula $\Diamond_a p$ is interpreted with Kripke semantics for epistemic modality, as the pre-image of p under the world-alternative relation \hat{R}_a .

Summing up, given an Epik specification

⁷Deeper analysis of the lexical semantics of *know* requires adding modeling of presupposition (Collard, 2018). The grammar fragment in Section 6 does not model the presupposition of *know*, except as an entailment.

⁸See the discussion of modal axioms **T** and **D** below.

B, P, φ, ζ , term p (as defined syntactically in Figure 3) is interpreted as a set of guarded strings $\llbracket p \rrbracket^{B, P, \varphi, \zeta}$. Let $K^{B, P, \varphi, \zeta}$ be the sets that are interpretations of terms. Then $\langle K^{B, P, \varphi, \zeta}, +, \cdot, *, \neg, 0, 1, \neg, \diamond_{a_1}, \dots, \diamond_{a_n} \rangle$ is a concrete guarded string interpretation for the signature of epistemic KAT, with operations as in Figure 4 (e.g. the binary operation $+$ is union, and the unary operation \diamond_a is pre-image relative to \hat{R}_a). This provides a concrete n -agent general modal frame $\langle \llbracket \bullet \rrbracket^{B, P, \varphi, \zeta}, \hat{R}_1, \dots, \hat{R}_n, K^{B, P, \varphi, \zeta} \rangle$. The frame consists of a set of worlds, an epistemic-alternative relation for each agent, and a set of propositions, each of which is a set of worlds. It is used as a target for natural-language interpretation in Section 6.

Axiomatic Classification. To situate our logic as a modal logic, consider the soundness of the standard modal axioms given our semantics (Hughes et al., 1996). The axioms in (10)

- (10) **N** If p is valid, then $\Box_a p$ is valid
K $\Box_a(p \rightarrow q) \rightarrow \Box_a p \rightarrow \Box_a q$ is valid.

The axioms in (11) are not valid, but do hold when when the relation R_a has a certain shape.

- (11) **T** $\Box_a p \rightarrow p$ if $g \hat{R}_a g, \forall g$
D $\Box_a p \rightarrow \diamond_a p$ if \hat{R}_a is a function
4 $\Box_a p \rightarrow \Box_a \Box_a p$ if \hat{R}_a is idempotent

The condition on **4** is essentially just a restatement of the theorem in relational terms. The conditions on the remaining standard modal axioms, **B** and **5**, are similar.

4 Translation into the finite state calculus

The finite state calculus is an algebra of regular sets of strings and regular relations between strings that was designed for use in computational phonology and morphographemics (Kaplan and Kay, 1994; Beesley and Karttunen, 2003). Current implementations allow for the definition of functions on regular sets and relations (Hulden, 2009; Lindén et al., 2009; Karttunen, 2010). Such definitions are used here to construct of a model for epistemic KAT inside the finite state calculus (Figures 5&6).

The space of worlds is a set of ordinary (as opposed to guarded) strings. Bit sequences (sequences of 0's and 1's) encode atoms, and as before, these alternate with event symbols to encode a world. In this construction, w_3 of the example is the string $10 a1 10 b1 10$.

```
St
  Atoms such as 0110.
UnequalStPair
  Sequence of two unequal atoms such as 0110 0111.
define Wf0 ~[$ UnequalStPair];
  String that doesn't contain a non-matching pair of atoms.
define Squash St -> 0 || St _;
  Rewrite relation deleting the second of two atoms.
define Cn(X, Y)
  [[X Y] & Wf0] .o. Squash].1;
  KAT product.
define Kpl(X)
  [[X+] & Wf0] .o. Squash].1;
define Kst(X) St | Kpl(X);
  KAT Kleene plus and Kleene star. The Fst operation | is
  union.
```

Figure 5: Definition in Fst of KAT product and star.

Terms in the finite state calculus are interpreted as sets of strings, or for relational terms, as relations between strings. Computationally, the sets and relations are represented by finite state acceptors. As used here, a program in the Fst language of the finite state calculus is a straight-line program that defines a sequence of constants naming sets, constants naming relations, and functions (defined as macros) mapping one or more regular sets or relations to a regular set or relation.

Translating the Epik terms 0 , 1 , b , and e are straightforward: we simply convert the atoms as previously described, and decorate the events e with their compatible atoms. For example a_1 becomes an Fst term denoting $\{01a_101\}$. Fst has built-in operations of union ($|$), and intersection ($\&$), which are the sum and intersection operations in the guarded string algebra as represented in Fst. Fst set difference ($-$) is used to define propositional complement as the difference between the set of all worlds and the argument.

Defining KAT product using Fst's set-lifted string concatenation (denoted by juxtaposition $X Y$) requires more care. Naively concatenating strings with atoms (Boolean vectors) at both ends doubles atoms at the juncture, and does not enforce the requisite atom equality. To implement KAT product, we define the binary operation Cn , which concatenates strings in the string algebra, removes strings with non-matching atoms, and then deletes the second of two atoms to create a set of well-formed guarded strings (Figure 5). $Wf0$ is the set of ordinary strings with only equal pairs of atoms, as defined using Fst's containment operator $\$$. The *Squash* relation uses Fst's rewrite notation to delete atoms (elements of St) that are preceded

```

define RelKpl(R)
  Squash.i.o.Wf0.o.[R+].o.Wf0.o.Squash
   $\underbrace{\hspace{1cm}}_c \underbrace{\hspace{1cm}}_b \underbrace{\hspace{1cm}}_a \underbrace{\hspace{1cm}}_b \underbrace{\hspace{1cm}}_c$ 
  a Relational Kleene plus in the string algebra
  b Constrain domain and co-domain to contain
    no unmatched atoms.
  c Reduce doubled atoms to a single
    atom in the domain and co-domain.
define RelKst(R) [St.x.St] | Kpl(X);
  The Fst operation .x. is Cartesian product. R.i is the
  inverse of relation R.

```

Figure 6: Definition in Fst of the Kleene concatenation closure of a relation between guarded strings.

by another atom⁹. This relation is applied via the composition $(. \circ .)$ and codomain $(. \downarrow)$ operators.

Kleene plus is defined in a similar way, using Kleene plus in the string algebra, with checks for equality of atoms and deletion of atoms. Kleene star is defined from Kleene plus the well-formed atoms St , which implements 1.

Finally, the epistemic alternative relation for each agent, i.e. the `agent` relation from Figure 2 is implemented in Fst. Combining the agent relation with the effect formula gives a relation on events decorated with compatible atoms. Then the relation on worlds is constructed using the closure of the concatenation product operation. The concatenation product $R \circ S$ of two relations R and S is the set of pairs of the form $\langle x_1 x_2, y_1 y_2 \rangle$, where $x_1 R y_1$, and $x_2 S y_2$. In Fst, $R+$ is the closure of relation R with respect to this operation. Figure 6 defines the corresponding operation on sets of guarded strings as encoded in Fst.¹⁰ The epistemic alternative relation on worlds for an agent is then defined as the KAT relation concatenation closure RelKst of the decorated-event alternative relation for the agent.

5 Bounded Lazy Interpretation

We also implement the semantics of Epik terms using lazy lists in Haskell, rather than the direct interpretation as sets. Using lists sidesteps checking the set invariant (elements are unique) for large sets, such as \bullet^* , and laziness allows us to delay computing these large sets until they are actually

⁹This is a non-equal length regular relation. The finite state calculus includes such relations, and they can be used with relation composition and relation domain and co-domain. They are restricted in that the complement and set difference for non-equal length relations is not defined. Epistemic alternative relations are equal-length relations.

¹⁰Relation concatenation in Fst differs from relation composition $(. \circ .)$, and the closure under discussion here is the closure of the former rather than the latter.

$$\begin{aligned}
\langle p \rangle_0^{B,P,\phi,\zeta} &\triangleq [] \\
\langle 0 \rangle_n^{B,P,\phi,\zeta} &\triangleq [] \\
\langle 1 \rangle_n^{B,P,\phi,\zeta} &\triangleq [\mathcal{A}_B^\varphi] \\
\langle e \rangle_n^{B,P,\phi,\zeta} &\triangleq [\alpha e \beta \mid \alpha \zeta_e \beta] \\
\langle b \rangle_n^{B,P,\phi,\zeta} &\triangleq [\mathcal{A}_B^{b\psi}] \\
\langle p+q \rangle_n^{B,P,\phi,\zeta} &\triangleq \langle p \rangle_n^{B,P,\phi,\zeta} ++ \langle q \rangle_n^{B,P,\phi,\zeta} \\
\langle p;q \rangle_n^{B,P,\phi,\zeta} &\triangleq (\langle p \rangle_n^{B,P,\phi,\zeta} \diamond \langle q \rangle_n^{B,P,\phi,\zeta})|_n \\
\langle p^* \rangle_n^{B,P,\phi,\zeta} &\triangleq [] + (\langle p \rangle_n^{B,P,\phi,\zeta} \diamond \langle p^* \rangle_{n-i}^{B,P,\phi,\zeta})|_n \\
&\text{where } i = \max\{1, \min\{|g| \mid g \in \langle p \rangle_n^{B,P,\phi,\zeta}\}\} \\
\langle \neg p \rangle_n^{B,P,\phi,\zeta} &\triangleq \langle \bullet^* \rangle_n^{B,P,\phi,\zeta} \setminus \langle p \rangle_n^{B,P,\phi,\zeta} \\
\langle \diamond_a p \rangle_n^{B,P,\phi,\zeta} &\triangleq [g' \mid g' \hat{R}_a g, \text{for } g \text{ in } \langle p \rangle_n^{B,P,\phi,\zeta}]
\end{aligned}$$

Figure 7: Bounded interpretation using lazy lists

needed. To sidestep the infiniteness of models, we parameterize the interpretation function on a positive integer n and only produce guarded strings of length n or less.

The bounded interpretation into lists of strings is very similar to the unbounded interpretation into sets of strings, except for the bounds checking. The full details are shown in Figure 7. First note that when $n = 0$, the denotation is empty, denoted $[]$. Terms of the form 0 , 1 , e , and ψ have the same denotation as before, translated into a list (denoted $[S]$, for a set S). We compute atoms using BDDs, which concisely represent boolean functions (Lee, 1959).

We lift the remaining operators (except Kleene star) to their list equivalents: union becomes list append (written $++$); fusion product is lifted to lists instead of sets, negation is implemented using list difference \setminus , and the modal operator lifts the alternative relation over lists of strings¹¹. The only caveat to these direct interpretations is that we restrict the operators to have size $\leq n$, denoted as $l|_n$ for a list of guarded strings l .

The denotation of p^* uses the fact that p^* and $1 + p; p^*$ are equivalent, and decrements the size threshold on the recursive denotation of p^* by i , where i is the length of the longest (nonzero) string in the denotation of p , making sure to filter out guarded strings that are too long.

6 Syntax-semantics interface

English sentences are mapped to terms in the logical language via a semantically interpreted multimodal categorical grammar, consisting of a lexicon

¹¹Figure 7 depicts this using the list comprehension notation, which is analogous to set builder notation, except that it is written using square brackets. Element order is evoked by the keyword `for`, rather than using the unordered \forall .

ITEM	TYPE	SEMANTICS
Statives with expletive subject		
heads	$d \backslash_D t$	$\lambda x. \bullet^* h$
tails	$d \backslash_D t$	$\lambda x. \bullet^* t$
Agents		
Amy	e	\hat{R}_a
Bob	e	\hat{R}_b
Auxiliary verbs and expletive subjects		
it	d	<i>dummy</i>
is	$(d \backslash t) / (d \backslash_D t)$	$\lambda P x. P x$
it's	$t / (d \backslash_D t)$	$\lambda P x. P x$
isn't	$t / (d \backslash_D t)$	$\lambda P x. \neg P x$
doesn't	$(e \backslash t) / (d \backslash_V t)$	$\lambda P x. \neg P x$
Tensed attitude verbs		
knows	$(e \backslash t) /_M t$	$\lambda p R. p + \Diamond_{RP}$
believes	$(e \backslash t) /_M t$	$\lambda p R. \Diamond_{RP}$
Base form attitude verbs		
know	$(e \backslash_V t) /_M t$	$\lambda p R. p + \Diamond_{RP}$
believe	$(e \backslash_V t) /_M t$	$\lambda p R. \Diamond_{RP}$
Complementizers		
that	$((e \backslash t) /_M t) \backslash (e \backslash t) / t$	$\left(\begin{array}{c} \lambda p m R. \\ \neg(m(\neg p) R) \end{array} \right)$
whether	$((e \backslash t) /_M t) \backslash (e \backslash t) / t$	$\left(\begin{array}{c} \lambda p m R. \\ \neg(m(\neg p) R) \\ + \neg(m p R) \end{array} \right)$
Complementizers for base form verbs		
that	$((e \backslash_V t) /_M t) \backslash (e \backslash_V t) / t$	$\left(\begin{array}{c} \lambda p m R. \\ \neg(m(\neg p) R) \end{array} \right)$
whether	$((e \backslash_V t) /_M t) \backslash (e \backslash_V t) / t$	$\left(\begin{array}{c} \lambda p m R. \\ \neg(m(\neg p) R) \\ + \neg(m p R) \end{array} \right)$
Conjunction		
and	$(t \backslash t) / t$	$\lambda p q. p \wedge q$
or	$(t \backslash t) / t$	$\lambda p q. p + q$
and	$((e \backslash t) \backslash (e \backslash t)) / (e \backslash t)$	$\lambda p q x. p(x) \wedge q(x)$
or	$((e \backslash t) \backslash (e \backslash t)) / (e \backslash t)$	$\lambda p q x. p(x) + q(x)$

Figure 8: Categorical grammar lexicon, showing word form (column 1), a categorial type (column 2), and a semantic translation in Epik extended with lambda abstraction (column 3).

of words, their categorial types, and interpretations in a logical lambda language. The grammar covers basic statives (*it's heads*), *that*- and *whether*-complements, negation, and predicate and sentence conjunction. Figure 8 gives the lexicon.¹² The on-line supplement includes a version of this grammar that can be run using Barker and Shan's chart parser for categorial grammar (Barker and Shan, 2005).

¹²Category symbols use Lambek/Bar-Hillel notation for slashes, so that $(d \backslash t) / (d \backslash_D t)$ combines with $d \backslash_D t$ on the right to give a value that combines with d on the left to give t . In the semantics, lambda abstractions with multiple parameters are written $\lambda x y. e$ rather than $\lambda x. \lambda y. e$. d is the category of expletive *it*.

The grammar and semantics are optimized for a simple fragment of English concerned with clausal complementation. The agent names *Amy* and *Bob* contribute the epistemic alternative relations for those agents, rather than individuals. The root verb *believe* contributes existential modal force. The complementizers *that* and *whether* are the heads of their dominating clauses, and assemble an alternative relation, modal force, and proposition contributed by the complement. These complementizers introduce the dual via two negations, in order to express universal modal force.

Multimodal categories such as \backslash_D and \backslash_M are used to control the derivation—phrases with these top-level slashes can only combine syntactically as arguments. The semantic translations in the third column of Figure 8 use the logical language, incremented with lambda. The body of $\lambda x. \bullet^* h$, which is the semantic lexical entry for *heads*, is a term denoting the set of all worlds where the coin is heads, expressed as the set of all guarded strings that end with a Boolean valuation where the primitive proposition h (it's heads) is true. The body of $\lambda p. \lambda R. \Diamond_{RP}$, which is the semantic lexical entry of *believes*, is an term denoting the pre-image of the world-alternative relation contributed by the subject. This is not the right semantics for *Amy believes that it's heads*, because it has an existential modality \Diamond_{RP} , rather than an universal modality \Box_{RP} . This is corrected by the complementizer *that*, which introduces the dual.

Sentences are parsed with a chart parser for categorial grammar. The semantics for complex phrases are obtained by application of semantic translations, accompanied by beta reductions that eliminate all lambdas in logical forms for clauses. In consequence, the semantic term translating a sentence is an Epik logical term. Such a term designates a set of possible words (guarded strings). By way of example, (12a) is an English sentence with conjunction and several levels of clausal embedding. Using the grammar and parser, the sentence is mapped to the term in (12b). (12c) shows a simplified logical form constructed from (12b) using logical equivalences. Either term is compiled in an implementation of the finite state calculus to a finite state machine with 10 states and 23 edges, which accepts a countably infinite set of worlds.¹³

¹³Machine sizes need not be small, especially as the cardinality of B increases. A certain Epik model with fourteen primitive Booleans has the set of worlds represented by a finite state machine with 184,794 states and 257,881 edges.

Conjunction missing!

In this way the methodology “directly” represents the set of worlds denoted by (12a).

- (12) a. It’s tails and Amy knows that Bob knows that Amy knows whether it’s heads.
- b. $\bullet^*t \wedge \neg(\neg(\neg(\neg(\neg(\neg \bullet^*h \vee \Diamond_a \neg \bullet^*h) + \neg(\bullet^*h \wedge \Diamond_a \bullet^*h)) + \Diamond_b \neg(\neg(\neg \bullet^*h \vee \Diamond_a \neg \bullet^*h) + \neg(\bullet^*h \wedge \Diamond_a \bullet^*h)))) + \Diamond_a \neg(\neg(\neg(\neg \bullet^*h \vee \Diamond_a \neg \bullet^*h) + \neg(\bullet^*h \wedge \Diamond_a \bullet^*h)) + \Diamond_b \neg(\neg(\neg \bullet^*h \vee \Diamond_a \neg \bullet^*h) + \neg(\bullet^*h \wedge \Diamond_a \bullet^*h))))))$
- c. $\bullet^*t \wedge \mathcal{K}_a(\mathcal{K}_b(\mathcal{K}_a \bullet^*h + \mathcal{K}_a \neg \bullet^*h))$
- d. Amy knows that it’s tails.
- e. $\bullet^*t \wedge \neg \Diamond_a \neg \bullet^*t \quad (\equiv \mathcal{K}_a \bullet^*t)$

Sentence (12d) is assigned the logical form (12e) by the grammar. Logical relations between propositions are checked in the finite state calculus by checking set-theoretic relations between sets of worlds. For instance entailment $p \rightarrow q$ is decided by checking in an interpreter for the finite state calculus whether $p - q$ is non-empty. In the model defined by Figure 2, the propositions (12b) and (12e) are independent. They are equivalent in a version without secret flipping.

7 Discussion

Epistemic KAT is designed for use in linguistic semantics research, for computational linguistics research on model-theoretically grounded semantics, and for students in graduate course modules on intensional formal semantics.

There are straightforward extensions to additional linguistic phenomena, such as tense and perfective aspect (13a), and the combination of metaphysical modality and prospective aspect (13b).

- (13) a. Amy has learned that Bob had learned that it’s heads.
- b. Amy might learn that it’s heads.

The model framework is a constructive branching-time framework with metaphysical and epistemic modalities, which will be applicable in linguistic semantic research on combinations of tense, metaphysical modality, and epistemic complementation (Thomason, 1984; Abusch, 1998; Condoravdi, 2002). Connections with research on temporal constitution of events remain to be explored (Fernando, 2004, 2007; Carlson, 2009).

The syntactic part of the grammar formalism uses a standard approach to categorial grammar (Steedman, 2000; Bozsahin, 2012) to map between

strings (or trees or derivations) and logical terms, and is rigorous enough to admit computational implementations (Barker and Shan, 2005; Bozsahin, 2021).¹⁴ While this approach is simple and attractive for our applications, it would be possible to use Epik with other frameworks that support lambda extensions of logical languages.

The semantics presented here is exclusively concerned with events and worlds. Semantic type systems for natural language usually include a type for individuals (Gallin, 1975). In the coin example, the individuals are hidden in the primitive event symbols such as a_1 and b_1 . This situation would get worse in a more elaborate model with more agents and multiple coins. The solution to this should be to base the world construction on grounded event terms such as $\text{look}(a,k,1)$, for “agent a see that coin k is heads”, rather than atomic event symbols. This way of proceeding is found in research on situation calculus (Reiter, 2001). How to incorporate it in the scheme for Epik specifications and into the computational parts of the proposal is a topic for future research. In this it is not clear whether it is possible to introduce quantification over individuals, together with individual constants. Rooth (2017) develops an approach to a positive answer, based on introducing markers for witnesses for discourse referents in the construction of worlds.

The development here is concerned with defining concrete computable possible worlds models, and applying them in natural language semantics. Issues for further investigation are mathematical characterizations of epistemic KATs, e.g. sound and complete axioms and coalgebras.

Source code, examples, and instructions for Epik are available at <https://github.com/erichthewry/epik>.

¹⁴Our grammar uses basic categorial grammar, not additional features such as combinators, type raising, continuations, or multimodal slashes that interact with the grammar in non-trivial ways. These features become relevant in more sophisticated grammars. Multimodal categories are used for phrases such as a tenseless verb phrase that semantically are functional, but combine in the grammar only as arguments.

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