

Vorlesung 4

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Kapitel 1

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1 Übung 1

$\text{ite } \text{true } s \ t \rightarrow_{\delta} (\lambda bxy.bxy)\text{true } s \ t \rightarrow_{\beta} (\lambda xy.(\text{true})xy) \ s \ t \rightarrow_{\beta} (\lambda y.(\text{true})sy) \ t \rightarrow_{\beta} \text{true } s \ t \rightarrow_{\delta} (\lambda xy.y)st \rightarrow_{\beta} (\lambda y.s)t \rightarrow_{\beta} s$

$\text{ite } \text{false } s \ t \rightarrow_{\delta} (\lambda bxy.bxy)\text{false } s \ t \rightarrow_{\beta} (\lambda xy.(\text{false})xy) \ s \ t \rightarrow_{\beta} (\lambda y.(\text{false})sy) \ t \rightarrow_{\beta} \text{false } s \ t \rightarrow_{\delta} (\lambda xy.x)st \rightarrow_{\beta} (\lambda y.y)t \rightarrow_{\beta} t$

Wir schreiben $\text{ite } x \ y \ z = \text{if } x \text{ then } y \text{ else } z$.

Direkt ohne beta-reduktion anwendbar.

$\text{not } b = b \ \text{false } \text{true}$.

$\text{xor } b1 \ b2 = \text{if } b1 \text{ then } (\text{if } b2 \text{ then } \text{false } \text{else } \text{true}) \text{ else } b2$

$\text{implikation } b1 \ b2 = \text{if } b1 \text{ then } b2 \text{ else } \text{true}$

2 Übung 2

$\text{fact } \lceil 2 \rceil \rightarrow_{\beta} \text{if } \lceil 2 \rceil \leq \lceil 1 \rceil \text{ else } \lceil 2 \rceil \text{fac}(\lceil 2 \rceil - \lceil 1 \rceil) \rightarrow_{\beta}^* \text{if } \text{false } \text{else } \lceil 2 \rceil \text{fac}(\lceil 2 \rceil - \lceil 1 \rceil) \rightarrow_{\beta}^* \text{if } \text{false } \text{else } \lceil 2 \rceil \text{fac}(\lceil 1 \rceil) \rightarrow_{\beta} \text{fac}(\lceil 1 \rceil) \rightarrow_{\beta} \lceil 2 \rceil \text{if } \lceil 1 \rceil \leq \lceil 1 \rceil \text{ else } \lceil 2 \rceil \text{fac}(\lceil 1 \rceil - \lceil 1 \rceil) \rightarrow_{\beta}^* \text{fac}(\lceil 1 \rceil) \rightarrow_{\beta} \lceil 2 \rceil \text{if } \text{true } \text{else } \lceil 2 \rceil \text{fac}(\lceil 1 \rceil - \lceil 1 \rceil) \rightarrow_{\beta} \lceil 2 \rceil \lceil * \rceil \lceil 1 \rceil = \lceil 2 \rceil$

$\text{odd } n = \text{if } n == \lceil 0 \rceil \text{ then } \text{false } \text{else } (\text{if } \lceil 2 \rceil == 1 \text{ then } \text{true } \text{else } \text{odd } (n - \lceil 2 \rceil))$

$\text{halve } n = \text{if } n \leq \lceil 1 \rceil \text{ then } \lceil 0 \rceil \text{ else } 1 + \text{halve } (n - \lceil 2 \rceil)$

3 Übung 3

Via normaler Reduktion, schauen, was nie reduziert wird, dass muss das Problem sein, weil normale sub immer NF erreicht, wenn sie existiert:

a)
$$\underbrace{(\lambda xy.y(\lambda z.x))(uu)(\lambda v.v((\lambda w.w)(\lambda w.w)))}_{\text{lambda leftmost}} \rightarrow_{\beta} \underbrace{(\lambda y.y(\lambda z.(uu)))(\lambda v.v((\lambda w.w)(\lambda w.w)))}_{\text{outermost}} \rightarrow_{\beta} \underbrace{(\lambda v.v((\lambda w.w)(\lambda w.w)))(\lambda z.(uu))}_{\text{outermost-leftmost, lambda aussen}} \rightarrow_{\beta} \underbrace{(\lambda z.(uu))((\lambda w.w)(\lambda w.w))}_{\text{outermost-leftmost}} \rightarrow_{\beta}^*$$

uu

Problem: $((\lambda w.w)(\lambda w.w))$ b) $(\lambda u.u(\lambda y.z))(\lambda x.x((\lambda v.v)(\lambda v.v)))$

$$\rightarrow_{\beta} \underbrace{(\lambda u.u(\lambda y.z))(\lambda x.x((\lambda v.v)(\lambda v.v)))}_{outermost} \rightarrow_{\beta} \underbrace{(\lambda x.x((\lambda v.v)(\lambda v.v)))(\lambda y.z)}_{outermost} \rightarrow_{\beta} (\lambda y.z)((\lambda v.v)(\lambda v.v)) \rightarrow_{\beta} z$$

Problem $((\lambda v.v)(\lambda v.v))$

2.

$$U = (\lambda f.fI(\Omega\Omega))(\lambda xy.xx)$$

applikativ:

$$\begin{aligned} & (\lambda f.fI(\underbrace{\Omega\Omega}_{innermost}))(\lambda xy.xx) \rightarrow_a (\lambda f.fI(\underbrace{((\lambda x.xx)\Omega)}_{innermost}))(\lambda xy.xx) \\ & \rightarrow_a (\lambda f.fI((\lambda x.xx)(\lambda x.xx)))(\lambda xy.xx) \\ & \rightarrow_a (\lambda f.fI((\lambda x.xx)(\lambda x.xx)))(\lambda xy.xx) \end{aligned}$$

unendliche Schleife

Normale:

$$\begin{aligned} & \rightarrow_n \underbrace{(\lambda f.fI(\Omega\Omega))(\lambda xy.xx)}_{outermost} \\ & \rightarrow_n \underbrace{(\lambda xy.xx)I(\Omega\Omega)}_{leftmost} \\ & \rightarrow_n \underbrace{(\lambda y.II)(\Omega\Omega)}_{outermost} \\ & \rightarrow_n \underbrace{II}_{outermost} \end{aligned}$$

4 Übung 4

twice fst (pair (pair true false) true) applikativ:

$$\begin{aligned} & \underbrace{(\lambda fx.f(fx))fst}_{leftmost-innermost} \text{ (pair (pair true false) true)} \\ & \underbrace{(\lambda fx.f(fx))(\lambda p.p(\lambda xy.x))}_{leftmost-innermost} (\text{pair}(\text{pair true false})\text{true}) \\ & (\lambda x.(\lambda p.p(\lambda xy.x))((\lambda p.p(\lambda xy.x)) x)) \underbrace{(\text{pair}(\text{pair true false})\text{true})}_{leftmost-innermost} \\ & (\lambda x.(\lambda p.p(\lambda xy.x))((\lambda p.p(\lambda xy.x)) x)) ((\lambda ab \text{ select.select } a \text{ } b) \text{ (pair true false) true}) \\ & (\lambda x.(\lambda p.p(\lambda xy.x))((\lambda p.p(\lambda xy.x)) x)) ((\lambda ab \text{ select.select } a \text{ } b) ((\lambda ab \text{ select.select } a \text{ } b) \text{ true false) true}) \end{aligned}$$

normal:

$$\begin{aligned} & \underbrace{twice\text{fst}}_{leftmost-outermost} \text{ (pair (pair true false) true)} \\ & \underbrace{(\lambda fx.f(fx))fst}_{leftmost-outermost} \text{ (pair (pair true false) true)} \\ & \underbrace{(\lambda x.fst(fst x))(\text{pair}(\text{pair true false})\text{true})}_{leftmost-outermost} \\ & (\underbrace{fst}_{leftmost-outermost} (fst (\text{pair}(\text{pair true false})\text{true}))) \\ & (\lambda p.(\lambda xy.x)) \underbrace{(fst (\text{pair}(\text{pair true false})\text{true}))}_{leftmost-outermost} \end{aligned}$$

$((\lambda p.(\lambda xy.x))((\lambda p.(\lambda xy.x)) (pair(pair\ true\ false)true)))$