

# Mathematics of Machine Learning - Summer School

## Lecture 3

### Bernstein's Concentration Inequalities. Fast Rates

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# Sub-Gaussian and Bernstein Random Variables

## Sub-Gaussian (Definition 6.5)

A random variable  $X$  is *sub-Gaussian* with *variance proxy*  $\sigma^2 > 0$  if

$$\mathbf{E} e^{\lambda(X - \mathbf{E}X)} \leq \exp(\sigma^2 \lambda^2 / 2) \quad \text{for any } \lambda \in \mathbb{R}$$

- ▶  $\psi^*(\varepsilon) = \varepsilon^2 / (2\sigma^2)$
- ▶ **Bounded r.v.'s**: if  $a \leq X - \mathbf{E}X \leq b$  then  $\sigma^2 = \frac{(b-a)^2}{4}$  (Hoeffding's Lem. 2.1)

## One-sided Bernstein's condition (Definition 7.1)

A random variable  $X$  satisfies the *one-sided Bernstein's condition* with  $b > 0$  if

$$\mathbf{E} e^{\lambda(X - \mathbf{E}X)} \leq \exp\left(\frac{(\mathbf{Var} X) \lambda^2 / 2}{1 - b\lambda}\right) \quad \text{for any } \lambda \in [0, 1/b)$$

- ▶  $\psi^*(\varepsilon) = \frac{\mathbf{Var} X}{b^2} h\left(\frac{b\varepsilon}{\mathbf{Var} X}\right)$  with  $h(u) = 1 + u - \sqrt{1 + 2u}$  for  $u > 0$
- ▶ **Bounded above r.v.'s**: if  $X - \mathbf{E}X \leq c$  then  $b = c/3$  (Proposition 7.4)

# Hoeffding's Inequality vs Bernstein's Inequality

Consider  $X_1, \dots, X_n \sim X$  i.i.d. bounded in  $[-c, c]$

## ► Upper-tail bounds:

$$\mathbf{P}\left(\frac{1}{n} \sum_{i=1}^n X_i - \mathbf{E}X \geq \epsilon\right) \leq e^{-n\epsilon^2/(2c^2)} \quad (\text{Hoeffding's})$$

$$\mathbf{P}\left(\frac{1}{n} \sum_{i=1}^n X_i - \mathbf{E}X \geq \epsilon\right) \leq \exp\left(-\frac{n\epsilon^2/2}{\mathbf{Var}X + c\epsilon/3}\right) \quad (\text{Bernstein's})$$

## ► Upper-confidence bounds:

$$\mathbf{P}\left(\frac{1}{n} \sum_{i=1}^n X_i - \mathbf{E}X < \sqrt{\frac{2c^2 \log(1/\delta)}{n}}\right) \geq 1 - \delta \quad (\text{Hoeffding's})$$

$$\mathbf{P}\left(\frac{1}{n} \sum_{i=1}^n X_i - \mathbf{E}X < \frac{c}{3n} \log(1/\delta) + \sqrt{\frac{2(\mathbf{Var}X) \log(1/\delta)}{n}}\right) \geq 1 - \delta \quad (\text{Bernstein's})$$

If  $\mathbf{Var}X = 0$  then we get fast rate  $\Rightarrow$  need to understand **noise** in learning

# Back to Binary Classification

To understand main ideas to get fast rate, consider binary classification:

- ▶  $Z_i = (X_i, Y_i) \in \mathbb{R}^d \times \{-1, 1\}$
- ▶ Admissible action set  $\mathcal{A} \subseteq \mathcal{B} := \{a : \mathbb{R}^d \rightarrow \{-1, 1\}\}$
- ▶ **True** loss function  $\ell(a, (x, y)) = 1_{a(x) \neq y}$

$$r(a) = \mathbf{P}(a(X) \neq Y) \quad a^* \in \operatorname{argmin}_{a \in \mathcal{A}} r(a) \quad a^{**} \in \operatorname{argmin}_{a \in \mathcal{B}} r(a)$$

$$R(a) = \frac{1}{n} \sum_{i=1}^n 1_{a(X_i) \neq Y_i} \quad A^* \in \operatorname{argmin}_{a \in \mathcal{A}} R(a)$$

The **Bayes decision rule**  $a^{**}$  reads

$$a^{**}(x) \in \operatorname{argmax}_{\hat{y} \in \mathcal{Y}} \mathbf{P}(Y = \hat{y} | X = x) = \begin{cases} 1 & \text{if } \eta(x) > 1/2 \\ -1 & \text{if } \eta(x) \leq 1/2 \end{cases}$$

with the **unkown** regression function  $\eta(x) := \mathbf{P}(Y = 1 | X = x)$   
( $\eta$  captures noise of **unkown** generative model)

# Regression Function: Excess Risk and Bayes Risk

(Theorem 7.6)

For any  $a \in \mathcal{B}$   $r(a) - r(a^{**}) = \mathbf{E}[|2\eta(X) - 1|1_{a(X) \neq a^{**}(X)}]$

$$r(a^{**}) = \mathbf{E} \min\{\eta(X), 1 - \eta(X)\} \leq \frac{1}{2}$$

$r(a^{**}) = 1/2$  if and only if  $\eta(X) = 1/2$  ( $Y$  contains no information on  $X$ )

- ▶  $\eta$  **close to**  $1/2$ : large Bayes risk large; small excess risk
- ▶  $\eta$  **away from**  $1/2$ : small Bayes risk large; large excess risk

# Fast Rate: Massart's Condition

## Massart's Noise Condition (Definition 7.7)

There exists  $\gamma \in (0, 1/2]$  such that

$$\mathbf{P}\left(\left|\eta(X) - \frac{1}{2}\right| \geq \gamma\right) = 1$$

( $\gamma = 0$  would mean condition is void)

## Fast Rate in Binary Classification (Theorem 7.10)

Let  $a^{**} \in \mathcal{A}$  so that  $a^* = a^{**}$ . If Massart's condition holds with  $\gamma \in (0, 1/2]$ ,

$$\mathbf{P}\left(r(A^*) - r(a^*) \leq \frac{\log(|\mathcal{A}|/\delta)}{\gamma n}\right) \geq 1 - \delta$$

Fast rate if  $|\mathcal{A}| < \infty$

- ▶ Massart's condition is strong:  $\eta$  uniformly bounded away from  $1/2$
- ▶ Weaker conditions:  $\eta$  arbitrarily close to  $1/2$ , but with small probability

# Proof of Theorem 7.10 (Part I)

- Error decomposition:  $r(A^\star) - r(a^\star) \leq R(a^\star) - R(A^\star) - (r(a^\star) - r(A^\star))$

$$\begin{aligned} G(a) &:= R(a^\star) - R(a) - (r(a^\star) - r(a)) = R(a^\star) - R(a) - \mathbf{E}[R(a^\star) - R(a)] \\ &= \frac{1}{n} \sum_{i=1}^n (g(a, Z_i) - \mathbf{E}g(a, Z_i)) \end{aligned}$$

with  $g(a, z) = 1_{a^\star(x) \neq y} - 1_{a(x) \neq y}$

- The above yields  $r(A^\star) - r(a^\star) \leq G(A^\star)$
- Bernstein's inequality for bounded random variables yields, for any  $a \in \mathcal{A}$ ,

$$\mathbf{P}(G(a) \geq \varepsilon) \leq \exp \left( - \frac{n \mathbf{Var} g(a, Z)}{b^2} h \left( \frac{b\varepsilon}{\mathbf{Var} g(a, Z)} \right) \right)$$

- Setting the right-hand side to  $\delta/|\mathcal{A}|$ , using that  $h^{-1}(u) = u + \sqrt{2u}$  for  $u > 0$

$$\begin{aligned} \mathbf{P} \left( G(A^\star) < \frac{b}{n} \log(|\mathcal{A}|/\delta) + \sqrt{\frac{2(\mathbf{Var} g(A^\star, Z)) \log(|\mathcal{A}|/\delta)}{n}} \right) \\ \geq \mathbf{P} \left( \bigcap_{a \in \mathcal{A}} \left\{ G(a) < \frac{b}{n} \log(|\mathcal{A}|/\delta) + \sqrt{\frac{2(\mathbf{Var} g(a, Z)) \log(|\mathcal{A}|/\delta)}{n}} \right\} \right) \geq 1 - \delta \end{aligned}$$

## Proof of Theorem 7.10 (Part II)

- As for any  $a \in \mathcal{A}$  we have  $|g(a, Z)| = 1_{a(X) \neq a^*(X)}$ , then

$$\mathbf{Var} g(a, Z) \leq \mathbf{E}[g(a, Z)^2] = \mathbf{P}(a(X) \neq a^*(X))$$

and from Theorem 7.6 and Massart's noise condition we have

$$r(a) - r(a^*) = \mathbf{E}[2\eta(X) - 1 | 1_{a(X) \neq a^*(X)}] \geq 2\gamma \mathbf{P}[a(X) \neq a^*(X)],$$

which yields  $\mathbf{Var} g(a, Z) \leq \frac{1}{2\gamma}(r(a) - r(a^*))$

- Using that  $r(A^*) - r(a^*) \leq G(A^*)$ , we can conclude

$$\mathbf{P}\left(r(A^*) - r(a^*) < \frac{2}{3n} \log(|\mathcal{A}|/\delta) + \sqrt{\frac{(r(A^*) - r(a^*)) \log(|\mathcal{A}|/\delta)}{\gamma n}}\right) \geq 1 - \delta.$$

The proof follows by solving the expression in the event with respect to the excess risk  $r(A^*) - r(a^*)$ , using that  $x < 2\alpha/3 + \sqrt{x\alpha/\gamma}$  for  $x \in [0, 1]$ , with  $\alpha > 0$  and  $\gamma \in (0, 1/2]$ , implies  $x < \alpha/\gamma$ .



# Interpolation Slow and Fast Rate: Tsybakov's Condition

## Tsybakov's Noise Condition (Definition 7.11)

There exist  $\alpha \in (0, 1)$ ,  $\beta > 0$ , and  $\gamma \in (0, 1/2]$  such that, for all  $t \in [0, \gamma]$ ,

$$\mathbf{P}\left(\left|\eta(X) - \frac{1}{2}\right| \leq t\right) \leq \beta t^{\alpha/(1-\alpha)}$$

## Interpolation Slow and Fast Rate in Binary Classification (Theorem 7.13)

Let  $a^{**} \in \mathcal{A}$ . If Tsybakov's condition holds for  $\alpha \in (0, 1)$ ,  $\beta > 0$ ,  $\gamma \in (0, 1/2]$ ,

$$\mathbf{P}\left(r(A^*) - r(a^*) \leq c \left(\frac{\log(|\mathcal{A}|/\delta)}{n}\right)^{\frac{1}{2-\alpha}}\right) \geq 1 - \delta$$

for a given constant  $c$  that depends on  $\alpha, \beta, \gamma$ .

- ▶ if  $\alpha \rightarrow 0$  then we recover slow rate (condition becomes void)
- ▶ if  $\alpha \rightarrow 1$  then we recover fast rate (condition recovers Massart's)

Note:  $A^*$  does **not** depend on  $\alpha$ : it automatically adjusts to the noise level!