

Mathematics of Machine Learning - Summer School

Lecture 5 Convex Loss Surrogates. Gradient Descent

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Recall Results on Binary Classification

- ▶ $Z_i = (X_i, Y_i) \in \mathbb{R}^d \times \{-1, 1\}$
- ▶ Admissible action set $\mathcal{A} \subseteq \mathcal{B} := \{a : \mathbb{R}^d \rightarrow \{-1, 1\}\}$
- ▶ **True** loss function $\ell(a, (x, y)) = 1_{a(x) \neq y} = \varphi^*(a(x)y)$ with $\varphi^*(u) := 1_{u \leq 0}$

$$r(a) = \mathbf{P}(a(X) \neq Y) \quad a^* \in \operatorname{argmin}_{a \in \mathcal{A}} r(a) \quad a^{**} \in \operatorname{argmin}_{a \in \mathcal{B}} r(a)$$

$$R(a) = \frac{1}{n} \sum_{i=1}^n 1_{a(X_i) \neq Y_i} \quad A^* \in \operatorname{argmin}_{a \in \mathcal{A}} R(a)$$

So far we have proved:

$$\mathbf{P}\left(r(A^*) - r(a^*) \lesssim \sqrt{\frac{\mathbf{VC}(\mathcal{A})}{n}} + \sqrt{\frac{\log(1/\delta)}{n}}\right) \geq 1 - \delta$$

Problem: In general, computing A^* is NP hard!

Idea: Define convex relaxation of the original problem

Convexity

Convex function (Definition 8.1)

A function $f : \mathbb{R}^d \rightarrow \mathbb{R}$ is *convex* if for every $x, \tilde{x} \in \mathbb{R}^d, \lambda \in [0, 1]$ we have

$$f(\lambda x + (1 - \lambda)\tilde{x}) \leq \lambda f(x) + (1 - \lambda)f(\tilde{x})$$

Convex set (Definition 8.2)

A set \mathcal{A} is *convex* if for every $a, \tilde{a} \in \mathcal{A}, \lambda \in [0, 1]$ we have

$$\lambda a + (1 - \lambda)\tilde{a} \in \mathcal{A}$$

Convex Loss Surrogates

Convex loss surrogate (Definition 8.3)

A function $\varphi : \mathbb{R} \rightarrow \mathbb{R}_+$ is called a *convex loss surrogate* if:

- convex
- non-increasing
- $\varphi(0) = 1$

True loss:

$$\varphi^*(u) = 1_{u \leq 0}$$

Exponential loss:

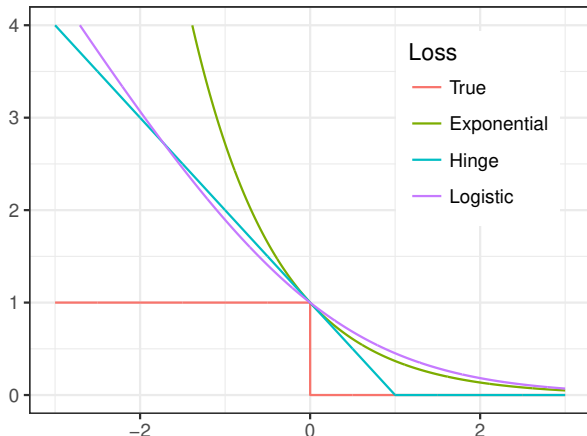
$$\varphi(u) = e^{-u}$$

Hinge loss:

$$\varphi(u) = \max\{1 - u, 0\}$$

Logistic loss:

$$\varphi(u) = \log_2(1 + e^{-u})$$



Convex Soft Classifiers

- ▶ **Soft** classifiers $\mathcal{A}_{\text{soft}} \subseteq \mathcal{B}_{\text{soft}} := \{a : \mathbb{R}^d \rightarrow \mathbb{R}\}$
- ▶ If $a \in \mathcal{B}_{\text{soft}}$, corresponding **hard** classifier is given by $\text{sign}(a)$

1. Linear functions with convex parameter space:

$$\mathcal{A}_{\text{soft}} = \{a(x) = w^\top x + b : w \in \mathcal{C}_1 \subseteq \mathbb{R}^d, b \in \mathcal{C}_2 \subseteq \mathbb{R}\}$$

$\mathcal{C}_1, \mathcal{C}_2$ are convex sets

2. Majority votes (Boosting):

$$\mathcal{A}_{\text{soft}} = \{a(x) = \sum_{i=1}^m w_i h_i(x) : w = (w_1, \dots, w_m) \in \Delta_m\}$$

Δ_m is the m -dim. simplex and $h_1, \dots, h_m : \mathbb{R}^d \rightarrow \mathbb{R}$ are *base classifiers*

Empirical φ -Risk Minimization

If φ and $\mathcal{A}_{\text{soft}}$ are convex, we are left with a convex problem

$$R_\varphi(a) = \frac{1}{n} \sum_{i=1}^n \varphi(a(X_i)Y_i)$$

$$A_\varphi^* \in \underset{a \in \mathcal{A}_{\text{soft}}}{\operatorname{argmin}} R_\varphi(a)$$

Zhang's Lemma

$$r_\varphi(a) = \mathbf{E} \varphi(a(X)Y)$$

$$a_\varphi^{**} \in \operatorname{argmin}_{a \in \mathcal{B}_{\text{soft}}} r_\varphi(a)$$

$$r(a) = \mathbf{E} \varphi^*(a(X)Y) = \mathbf{P}(a(X) \neq Y)$$

$$a^{**} \in \operatorname{argmin}_{a \in \mathcal{B}} r(a)$$

Zhang's Lemma (Lemma 8.5)

Let $\varphi : \mathbb{R} \rightarrow \mathbb{R}_+$ be a convex loss surrogate. For any $\tilde{\eta} \in [0, 1]$, $\tilde{a} \in \mathbb{R}$, let

$$H_{\tilde{\eta}}(\tilde{a}) := \varphi(\tilde{a})\tilde{\eta} + \varphi(-\tilde{a})(1 - \tilde{\eta}), \quad \tau(\tilde{\eta}) := \inf_{\tilde{a} \in \mathbb{R}} H_{\tilde{\eta}}(\tilde{a}).$$

Assume that there exist $c > 0$ and $\nu \in [0, 1]$ such that

$$\left| \tilde{\eta} - \frac{1}{2} \right| \leq c(1 - \tau(\tilde{\eta}))^\nu \quad \text{for any } \tilde{\eta} \in [0, 1]$$

Then, for any $a : \mathbb{R}^d \rightarrow \mathbb{R}$ we have

$$\underbrace{r(\operatorname{sign}(a)) - r(a^{**})}_{\text{excess risk hard classifier}} \leq 2c \underbrace{(r_\varphi(a) - r_\varphi(a_\varphi^{**}))^\nu}_{\text{excess } \varphi\text{-risk soft classifier}}$$

Zhang's Lemma: Examples

► **Exponential loss:**

$$\tau(\tilde{\eta}) = 2\sqrt{\tilde{\eta}(1 - \tilde{\eta})}$$

$$c = 1/\sqrt{2}$$

$$\nu = 1/2$$

► **Hinge loss:**

$$\tau(\tilde{\eta}) = 1 - |1 - 2\tilde{\eta}|$$

$$c = 1/2$$

$$\nu = 1$$

► **Logistic loss:**

$$\tau(\tilde{\eta}) = -\tilde{\eta} \log_2 \tilde{\eta} - (1 - \tilde{\eta}) \log_2 (1 - \tilde{\eta})$$

$$c = 1/\sqrt{2}$$

$$\nu = 1/2$$

Zhang's Lemma shows that we can reliably focus on convex problems

Elements of Convex Theory

Subgradients (Definition 8.8)

Let $f : \mathcal{C} \subset \mathbb{R}^d \rightarrow \mathbb{R}$. A vector $g \in \mathbb{R}^d$ is a *subgradient* of f at $x \in \mathcal{C}$ if

$$f(x) - f(y) \leq g^T(x - y) \quad \text{for any } y \in \mathcal{C}$$

The set of subgradients of f at x is denoted $\partial f(x)$.

Subgradients yield **global** information (**uniform** lower bounds)

Convexity and subgradients (Theorem 8.9)

Let $f : \mathcal{C} \subseteq \mathbb{R}^d \rightarrow \mathbb{R}$ with \mathcal{C} convex:

f is convex \implies for any $x \in \text{int}(\mathcal{C})$, $\partial f(x) \neq \emptyset$

f is convex \iff for any $x \in \mathcal{C}$, $\partial f(x) \neq \emptyset$

If f is convex and differentiable at x , then $\nabla f(x) \in \partial f(x)$

Convex functions that are differentiable allow to infer **global** information (i.e., subgradients) from **local** information (i.e., gradients)

This is why convex problems are “typically” amenable to computations...
To prove algorithms converge we need additional local-to-global properties

Are Convex Problems Easy to Solve?

- *Convex hull*: $\text{conv}(\mathcal{T}) := \left\{ \sum_{j=1}^m w_j t_j : w \in \Delta_m, t_1, \dots, t_m \in \mathcal{T}, m \in \mathbb{N} \right\}$
- *Epigraph*: $\text{epi}(f) := \{(x, t) \in \mathcal{D} \times \mathbb{R} : f(x) \leq t\}$.

Proposition 8.6

$$\min_{t \in \mathcal{T}} c^\top t = \min_{t \in \text{conv}(\mathcal{T})} c^\top t, \quad \max_{t \in \mathcal{T}} c^\top t = \max_{t \in \text{conv}(\mathcal{T})} c^\top t.$$

Proof: As $\mathcal{T} \subseteq \text{conv}(\mathcal{T})$, we have $\min_{t \in \mathcal{T}} c^\top t \geq \min_{t \in \text{conv}(\mathcal{T})} c^\top t$. Other direction:

$$\begin{aligned} \min_{t \in \text{conv}(\mathcal{T})} c^\top t &= \min_{m \in \mathbb{N}} \min_{t_1, \dots, t_m \in \mathcal{T}} \min_{(w_1, \dots, w_m) \in \Delta_m} c^\top \left(\sum_{j=1}^m w_j t_j \right) \\ &= \min_{m \in \mathbb{N}} \min_{t_1, \dots, t_m \in \mathcal{T}} \min_{(w_1, \dots, w_m) \in \Delta_m} \sum_{j=1}^m w_j c^\top t_j \geq \min_{t \in \mathcal{T}} c^\top t. \end{aligned}$$

Proposition 8.7

For any $f : \mathcal{D} \subseteq \mathbb{R}^d \rightarrow \mathbb{R}$, $\min_{x \in \mathcal{D}} f(x) = \min_{(x,t) \in \mathcal{C}} t$ with $\mathcal{C} = \text{conv}(\text{epi}(f))$.

Any minimization problem can be written in a convex form!

Local-to-Global Properties

► **Convex:** $f(y) \geq f(x) + \nabla f(x)^T(y - x) \quad \forall x, y \in \mathbb{R}^d$

► **α -Strongly Convex:**

$$\exists \alpha > 0 \text{ such that } f(y) \geq f(x) + \nabla f(x)^T(y - x) + \frac{\alpha}{2} \|y - x\|_2^2 \quad \forall x, y \in \mathbb{R}^d$$

► **β -Smooth:**

$$\exists \beta > 0 \text{ such that } f(y) \leq f(x) + \nabla f(x)^T(y - x) + \frac{\beta}{2} \|y - x\|_2^2 \quad \forall x, y \in \mathbb{R}^d$$

► **γ -Lipschitz:**

$$\exists \gamma > 0 \text{ such that } f(x) - \gamma \|y - x\|_2 \leq f(y) \leq f(x) + \gamma \|y - x\|_2 \quad \forall x, y \in \mathbb{R}^d$$

	Strongly convex?	Smooth?	Lipschitz?
Exponential loss (in \mathbb{R})	NO	NO	NO
Hinge loss (in \mathbb{R})	NO	NO	YES
Logistic loss (in \mathbb{R})	NO	YES	YES

However, we typically only need the domain to be a compact set of \mathbb{R}

Recap

- ▶ Training data: $(X_1, Y_1), \dots, (X_n, Y_n) \in \mathcal{X} \times \{-1, 1\}$, with $\mathcal{X} \subseteq \mathbb{R}^d$
- ▶ Loss function: $\varphi : \mathbb{R} \rightarrow \mathbb{R}_+$ (**convex**: reasonable by Zhang's lemma)
- ▶ Predictors $\mathcal{A} = \{x \in \mathbb{R}^d \rightarrow a_w(x) : w \in \mathcal{W}\}$ (\mathcal{W} **convex** in many cases)
NB. There are many settings where \mathcal{A} is **not** convex (e.g., neural networks)

Risk minimization:

$$\begin{array}{ll} \underset{w}{\text{minimize}} & r(w) = \mathbf{E}\varphi(a_w(X)Y) \\ \text{subject to} & w \in \mathcal{W} \end{array} \quad \Rightarrow \quad \text{Let } w^* \text{ be a minimizer}$$

Empirical risk minimization:

$$\begin{array}{ll} \underset{w}{\text{minimize}} & R(w) = \frac{1}{n} \sum_{i=1}^n \varphi(a_w(X_i)Y_i) \\ \text{subject to} & w \in \mathcal{W} \end{array} \quad \Rightarrow \quad \text{Let } W^* \text{ be a minimizer}$$

$$r(W) - r(w^*) \leq \underbrace{R(W) - R(W^*)}_{\text{Optimization}} + \underbrace{\sup_{w \in \mathcal{W}} \{r(w) - R(w)\} + \sup_{w \in \mathcal{W}} \{R(w) - r(w)\}}_{\text{Statistics}}$$

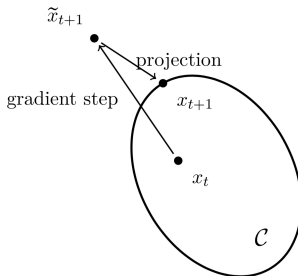
Projected Subgradient Method

Goal: $\min_{x \in \mathcal{C}} f(x)$ with f convex, \mathcal{C} convex and compact

Projected Subgradient Method

$$\begin{aligned}\tilde{x}_{t+1} &= x_t - \eta_t g_t, \text{ where } g_t \in \partial f(x_t) \\ x_{t+1} &= \Pi_{\mathcal{C}}(\tilde{x}_{t+1})\end{aligned}$$

with the projection operator $\Pi_{\mathcal{C}}(y) = \operatorname{argmin}_{x \in \mathcal{C}} \|x - y\|_2$.



Non-Expansivity of Projections

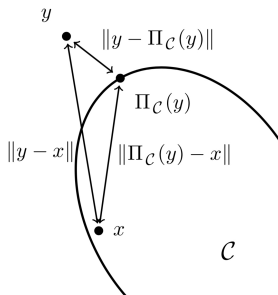
Non-expansivity (Proposition 9.2)

Let $x \in \mathcal{C}$ and $y \in \mathbb{R}^d$. Then,

$$(\Pi_{\mathcal{C}}(y) - x)^{\top} (\Pi_{\mathcal{C}}(y) - y) \leq 0$$

which implies $\|\Pi_{\mathcal{C}}(y) - x\|_2^2 + \|y - \Pi_{\mathcal{C}}(y)\|_2^2 \leq \|y - x\|_2^2$ and, in particular,

$$\|\Pi_{\mathcal{C}}(y) - x\|_2 \leq \|y - x\|_2$$



First Order Optimality Condition

First Order Optimality Condition (Proposition 8.10)

Let f be convex, and \mathcal{C} be a closed set on which f is differentiable. Then,

$$x^* \in \operatorname{argmin}_{x \in \mathcal{C}} f(x) \iff \nabla f(x^*)^\top (x^* - x) \leq 0 \quad \text{for any } x \in \mathcal{C}$$

Proof of Proposition 9.2. This is a direct consequence of Proposition 8.10 since $\Pi_{\mathcal{C}}(y)$ is a minimizer of the function $z \rightarrow f_y(z) = \|y - z\|_2$, and $\nabla f_y(z) = (z - y)/\|z - y\|_2$.

Results for Lipschitz Functions

A function f is γ -Lipschitz on \mathcal{C} if there exists $\gamma > 0$ such that (equivalent)

- ▶ For every $x, y \in \mathcal{C}$, $f(x) - \gamma\|x - y\|_2 \leq f(y) \leq f(x) + \gamma\|x - y\|_2$
- ▶ For every $x, y \in \mathcal{C}$, $|f(y) - f(x)| \leq \gamma\|x - y\|_2$
- ▶ For every $x \in \mathcal{C}$, any subgradient $g \in \partial f(x)$ satisfies $\|g\|_2 \leq \gamma$

Projected Subgradient Method—Lipschitz (Theorem 9.3)

- ▶ Function f is γ -Lipschitz
- ▶ Assume $\|x_1 - x^*\|_2 \leq b$

Then, the projected subgradient method with $\eta_s \equiv \eta = \frac{b}{\gamma\sqrt{t}}$ satisfies

$$f\left(\frac{1}{t} \sum_{s=1}^t x_s\right) - f(x^*) \leq \frac{\gamma b}{\sqrt{t}}$$

It is not a descent method: the value function can increase in one time step

Proof of Theorem 9.3)

- Convexity yields:

$$f\left(\frac{1}{t} \sum_{s=1}^t x_s\right) - f(x^*) \leq \frac{1}{t} \sum_{s=1}^t f(x_s) - f(x^*) \leq \frac{1}{t} \sum_{s=1}^t g_s^\top (x_s - x^*)$$

- Using $2a^\top b = \|a\|_2^2 + \|b\|_2^2 - \|a - b\|_2^2$ and $g_s = \frac{1}{\eta}(x_s - \tilde{x}_{s+1})$:

$$\begin{aligned} g_s^\top (x_s - x^*) &= \frac{1}{\eta} (x_s - \tilde{x}_{s+1})^\top (x_s - x^*) \\ &= \frac{1}{2\eta} (\|x_s - x^*\|_2^2 + \|x_s - \tilde{x}_{s+1}\|_2^2 - \|\tilde{x}_{s+1} - x^*\|_2^2) \\ &= \frac{1}{2\eta} (\|x_s - x^*\|_2^2 - \|\tilde{x}_{s+1} - x^*\|_2^2) + \frac{\eta}{2} \|g_s\|_2^2 \\ &\leq \frac{1}{2\eta} (\|x_s - x^*\|_2^2 - \|x_{s+1} - x^*\|_2^2) + \frac{\eta}{2} \|g_s\|_2^2 \end{aligned}$$

where we used that $\|\tilde{x}_{s+1} - x^*\|_2 \geq \|x_{s+1} - x^*\|_2$ by Proposition 9.2.

- Summing from $s = 1$ to t :

$$f\left(\frac{1}{t} \sum_{s=1}^t x_s\right) - f(x^*) \leq \frac{1}{2\eta t} (\|x_1 - x^*\|_2^2 - \|x_{t+1} - x^*\|_2^2) + \frac{\eta \gamma^2}{2} \leq \frac{b^2}{2\eta t} + \frac{\eta \gamma^2}{2}$$

Minimizing the right-hand side we have $\eta = \frac{b}{\gamma\sqrt{t}}$ which yields the result.

Results for Smooth Functions

A function f is β -smooth on \mathcal{C} if there exists $\beta > 0$ such that (equivalent)

- ▶ For every $x, y \in \mathcal{C}$, $f(y) \leq f(x) + \nabla f(x)^\top (y - x) + \frac{\beta}{2} \|y - x\|_2^2$
- ▶ For every $x, y \in \mathcal{C}$, $|\nabla f(y) - \nabla f(x)| \leq \beta \|x - y\|_2$ (gradient is β -Lipschitz)
- ▶ For every $x \in \mathcal{C}$, $\nabla^2 f(x) \preceq \beta I$ (if f is twice-differentiable)

Projected Gradient Descent—Smooth (Theorem 9.4)

- ▶ Function f is β -smooth
- ▶ Assume $\|x_1 - x^*\|_2 \leq b$

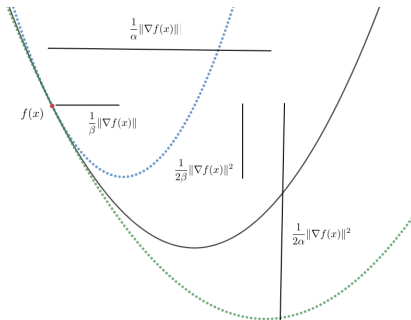
Then, projected gradient descent with $\eta_s \equiv \eta = 1/\beta$ satisfies

$$f(x_t) - f(x^*) \leq \frac{3\beta b^2 + f(x_1) - f(x^*)}{t}$$

In the case of smooth functions, gradient descent is a natural algorithm...

Interpretation for Smooth Functions

... it is the algorithm that at each time step moves to the point in \mathcal{C} that maximizes the guaranteed local decrease given by the quadratic function that uniformly upper-bounds the function f at the current location



$$\begin{aligned} \operatorname{argmin}_{y \in \mathcal{C}} \left\{ f(x) + \nabla f(x)^\top (y - x) + \frac{\beta}{2} \|y - x\|_2^2 \right\} &= \operatorname{argmin}_{y \in \mathcal{C}} \left\{ \left\| \left(x - \frac{1}{\beta} \nabla f(x) \right) - y \right\|_2^2 \right\} \\ &\equiv \Pi_{\mathcal{C}} \left(x - \frac{1}{\beta} \nabla f(x) \right) \end{aligned}$$

Results for Smooth and Strongly Convex Functions

A function f is α -strongly convex on \mathcal{C} if there is $\alpha > 0$ such that (equivalent)

- ▶ For every $x, y \in \mathcal{C}$, $f(y) \geq f(x) + \nabla f(x)^\top (y - x) + \frac{\alpha}{2} \|y - x\|_2^2$
- ▶ For every $x \in \mathcal{C}$, $\nabla^2 f(x) \succeq \alpha I$ (if f is twice-differentiable)

Gradient Descent—Smooth and Strongly Convex (Theorem 9.5)

- ▶ Assume $\mathcal{C} = \mathbb{R}^d$ (same type of result holds for projected gradient descent)
- ▶ Function f is α -strongly convex and β -smooth

Then, gradient descent with $\eta_s \equiv \eta = 1/\beta$ satisfies

$$f(x_t) - f(x^*) \leq \left(1 - \frac{\alpha}{\beta}\right)^{t-1} (f(x_1) - f(x^*))$$

Proof: (see illustration on the previous slide)

- ▶ Guaranteed progress in one step: $f(x_{s+1}) \leq f(x_s) - \frac{1}{2\beta} \|\nabla f(x_s)\|_2^2$
- ▶ Lower bound on objective function: $f(x^*) \geq f(x_s) - \frac{1}{2\alpha} \|\nabla f(x_s)\|_2^2$

Oracle Complexity, Lower Bounds, Accelerated Methods

► Convergence rates:

	L -Lipschitz	β -smooth
Convex	$O(\gamma b / \sqrt{t})$	$O((\beta b^2 + c)/t)$
α -strongly convex	$O(\gamma^2 / (\alpha t))$	$O(e^{-t\alpha/\beta} c)$

where $\|x_1 - x^*\|_2 \leq b$ and $f(x_1) - f(x^*) \leq c$

► Oracle complexities:

	L -Lipschitz	β -smooth
Convex	$O(\gamma^2 b^2 / \varepsilon^2)$	$O((\beta b^2 + c)/\varepsilon)$
α -strongly convex	$O(\gamma^2 / (\alpha \varepsilon))$	$O((\beta/\alpha) \log(c/\varepsilon))$

► Optimal rates (lower bounds)

	L -Lipschitz	β -smooth
Convex	$\Omega(\gamma a / (1 + \sqrt{t}))$	$\Omega(\tilde{b}^2 \beta / (t + 1)^2)$
α -strongly convex	$\Omega(\gamma^2 / (\alpha t))$	$\Omega(\alpha \tilde{b}^2 e^{-t\sqrt{\alpha/\beta}})$

where $a := \max_{x \in \mathcal{C}} \|x\|_2$ and $\tilde{b} := \max_{x, y \in \mathcal{C}} \|x - y\|_2$

Apart from Lipschitz, optimal rates are achieved only by **accelerated** algorithms

NB. Quantities α, β, γ and a, b, c, \tilde{b} depend implicitly on dimension d

Back to Learning: Linear Predictors with ℓ_2 Ball

Risk minimization:

$$\begin{array}{ll} \underset{w}{\text{minimize}} & r(w) = \mathbf{E}\varphi(w^\top XY) \\ \text{subject to} & \|w\|_2 \leq c_2^{\mathcal{W}} \end{array} \quad \Rightarrow \quad \text{Let } w^* \text{ be a minimizer}$$

Empirical risk minimization:

$$\begin{array}{ll} \underset{w}{\text{minimize}} & R(w) = \frac{1}{n} \sum_{i=1}^n \varphi(w^\top X_i Y_i) \\ \text{subject to} & \|w\|_2 \leq c_2^{\mathcal{W}} \end{array} \quad \Rightarrow \quad \text{Let } W^* \text{ be a minimizer}$$

$$r(\overline{W}_t) - r(w^*) \leq \underbrace{R(\overline{W}_t) - R(W^*)}_{\text{Optimization}} + \underbrace{\sup_{w \in \mathcal{W}} \{r(w) - R(w)\} + \sup_{w \in \mathcal{W}} \{R(w) - r(w)\}}_{\text{Statistics}}$$

$$\mathbf{E} \text{Statistics} \leq \frac{4c_2^{\mathcal{X}} c_2^{\mathcal{W}} \gamma_\varphi}{\sqrt{n}}$$

$$\text{Optimization} \leq \frac{2c_2^{\mathcal{X}} c_2^{\mathcal{W}} \gamma_\varphi}{\sqrt{t}}$$

Principled approach: Enough to run algorithm for $t \sim n$ time steps
(ONLY BASED ON UPPER BOUNDS!)