## CS 170 DIS 13

## Released on 2018-04-24

## 1 Follow the regularized leader

- (a) **Follow the leader.** You are playing T rounds of the following game: At round t you pick one of n strategies; your payoff for picking strategy i is  $A(t,i) \in [0,1]$ . You try the following algorithm: at each iteration pick the strategy which gave the highest average payoff so far (on the first iteration, you pick strategy 1).
  - Give an example of payoffs for T = 100 and n = 2, where your algorithm obtains a payoff of 0, but sticking to either i = 1 or i = 2 would have given you a payoff of almost 50.
- (b) Follow the randomized leader. The reason the algorithm above didn't do so well, is because when we deterministically jump from one strategy to another, an adversarially chosen set of strategies can be designed to thwart the algorithm.

To trick such adversaries, we want to use a randomized strategy; at time t we pick our strategy i at random from distribution  $D_t$ . Let  $p_t(i) \ge 0$  denote the probability that we assign to strategy i (i.e.  $\sum_{i=1}^{n} p_t(i) = 1$ ).

The previous algorithm ("Follow the leader") corresponds to setting  $D_t$  that maximizes

$$\sum_{i=1}^{n} \left( p_t(i) \cdot \sum_{\tau \in \{1, \dots, t-1\}} [A(\tau, i)] \right).$$

Why is this no better? (Hint: convexity).

(c) Follow the regularized leader. Instead, it is common to add a /regularized term that favors smoother distributions. A commonly used regularizer is the entropy function, i.e. we want to use pick i from the distribution that maximizes

$$\sum_{i=1}^{n} \left( p_t(i) \cdot \sum_{\tau \in \{1, \dots, t-1\}} [A(\tau, i)] - \eta p_t(i) \ln p_t(i) \right). \tag{1}$$

(Here,  $\eta > 0$  is a parameter that we can tweak to balance exploration and exploitation. Notice also that  $\ln p_t(i) \leq 0$ .)

In this exercise you will show that following the regularized leader with the entropy regularizer is the same as Multiplicative Weights Update!

Show that for any distribution  $p_t$ , (1) is at most

$$\eta \cdot \ln \left( \sum_{i=1}^{n} e^{\sum_{\tau \in \{1,\dots,t-1\}} [A(\tau,i)]/\eta} \right) \tag{2}$$

(Hint: you may use the inequality  $\sum_{i=1}^{n} p_t(i) \cdot \ln(y_i) \leq \ln(\sum_{i=1}^{n} p_t(i) \cdot y_i)$  for any vector  $\vec{y}$ .)

(d) Show that for some choice of  $\epsilon$  (which depends on  $\eta$ ), when computing  $p_t$  using Multiplicative Weight Update, (1) is equal to (2). What is the dependence of  $\epsilon$  on  $\eta$ ?