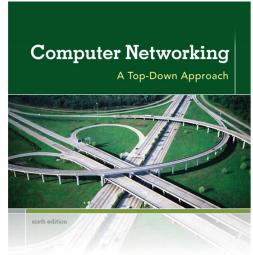
# CS335 Computer Networks

Application Layer 3



KUROSE ROSS

Computer
Networking: A Top
Down Approach
6th edition
Jim Kurose, Keith Ross
Addison-Wesley
March 2012

# Chapter 2: outline

- 2.1 principles of network applications
  - app architectures
  - app requirements
- 2.2 Web and HTTP
- 2.3 FTP
- 2.4 electronic mail
  - SMTP, POP3, IMAP
- **2.5 DNS**

- 2.6 P2P applications
- 2.7 socket programming with UDP and TCP

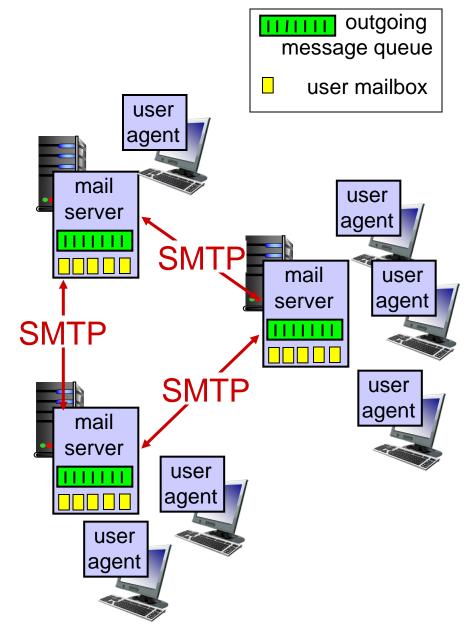
## Electronic mail

#### Three major components:

- user agents
- mail servers
- simple mail transfer protocol: SMTP

## **User Agent**

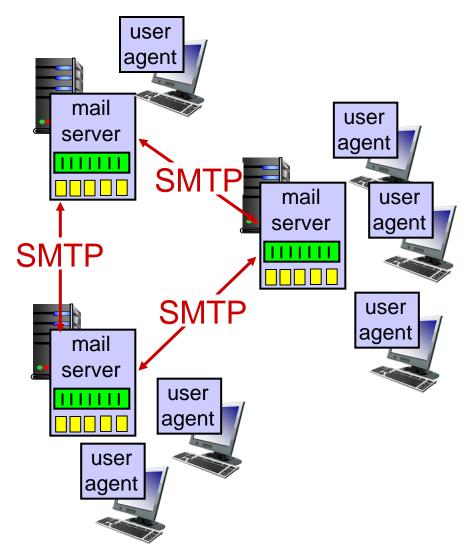
- \* a.k.a. "mail reader"
- composing, editing, reading mail messages
- e.g., Outlook, Thunderbird, iPhone mail client
- outgoing, incoming messages stored on server



## Electronic mail: mail servers

#### mail servers:

- mailbox contains incoming messages for user
- message queue of outgoing (to be sent) mail messages
- SMTP protocol between mail servers to send email messages
  - client: sending mail server
  - "server": receiving mail server



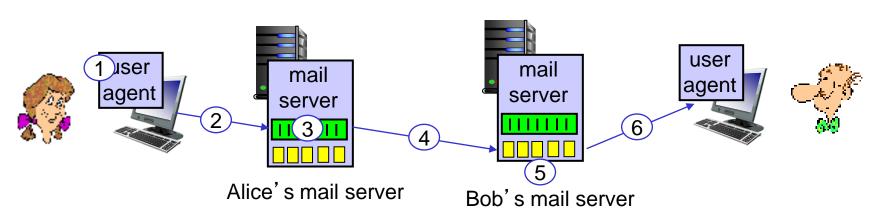
# Electronic Mail: SMTP [RFC 2821]

- uses TCP to reliably transfer email message from client to server, port 25
- direct transfer: sending server to receiving server
- three phases of transfer
  - handshaking (greeting)
  - transfer of messages
  - closure
- command/response interaction (like HTTP, FTP)
  - commands: ASCII text
  - response: status code and phrase
- messages must be in 7-bit ASCI

# Scenario: Alice sends message to Bob

- I) Alice uses UA to compose message "to" bob@someschool.edu
- 2) Alice's UA sends message to her mail server; message placed in message queue
- 3) client side of SMTP opens TCP connection with Bob's mail server

- 4) SMTP client sends Alice's message over the TCP connection
- 5) Bob's mail server places the message in Bob's mailbox
- 6) Bob invokes his user agent to read message



## Sample SMTP interaction

```
S: 220 hamburger.edu
                                   C: SMTP client
C: HELO crepes.fr
S: 250 Hello crepes.fr, pleased to meet you
C: MAIL FROM: <alice@crepes.fr>
S: 250 alice@crepes.fr... Sender ok
C: RCPT TO: <bob@hamburger.edu>
S: 250 bob@hamburger.edu ... Recipient ok
C: DATA
S: 354 Enter mail, end with "." on a line by itself
C: Do you like ketchup?
C: How about pickles?
C: .
S: 250 Message accepted for delivery
C: QUIT
S: 221 hamburger.edu closing connection
```

S: SMTP server

## Try SMTP interaction for yourself:

- \* telnet servername 25
- see 220 reply from server
- enter HELO, MAIL FROM, RCPT TO, DATA, QUIT commands

above lets you send email without using email client (reader)

## SMTP: final words

- SMTP uses persistent connections
- SMTP requires message (header & body) to be in 7-bit ASCII
- SMTP server uses
   CRLF.CRLF to
   determine end of message

#### comparison with HTTP:

- HTTP: pull
- SMTP: push
- both have ASCII command/response interaction, status codes
- HTTP: each object encapsulated in its own response msg
- SMTP: multiple objects sent in multipart msg

## Mail message format

SMTP: protocol for exchanging email msgs RFC 822: standard for text

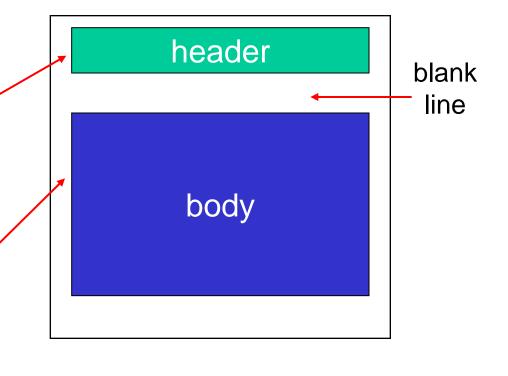
header lines, e.g.,

message format:

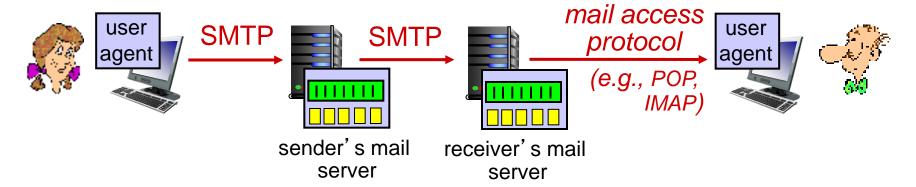
- To:
- From:
- Subject:

different from SMTP MAIL FROM, RCPT TO: commands!

- Body: the "message"
  - ASCII characters only



# Mail access protocols



- SMTP: delivery/storage to receiver's server
- mail access protocol: retrieval from server
  - POP: Post Office Protocol [RFC 1939]: authorization, download
  - IMAP: Internet Mail Access Protocol [RFC 1730]: more features, including manipulation of stored msgs on server
  - HTTP: gmail, Hotmail, Yahoo! Mail, etc.

## POP3 protocol

## authorization phase

- client commands:
  - user: declare username
  - pass: password
- server responses
  - +OK
  - -ERR

#### transaction phase, client:

- list: list message numbers
- retr: retrieve message by number
- dele: delete
- quit

```
S: +OK POP3 server ready
```

C: user bob

S: +OK

C: pass hungry

S: +OK user successfully logged on

C: list

S: 1 498

S: 2 912

S: .

C: retr 1

S: <message 1 contents>

S:

C: dele 1

C: retr 2

S: <message 1 contents>

S: .

C: dele 2

C: quit

S: +OK POP3 server signing off

# POP3 (more) and IMAP

#### more about POP3

- previous example uses POP3 "download and delete" mode
  - Bob cannot re-read email if he changes client
- POP3 "download-andkeep": copies of messages on different clients
- POP3 is stateless across sessions

#### **IMAP**

- keeps all messages in one place: at server
- allows user to organize messages in folders
- keeps user state across sessions:
  - names of folders and mappings between message IDs and folder name

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# DNS: domain name system

#### people: many identifiers:

SSN, name, passport #

#### Internet hosts, routers:

- IP address (32 bit) used for addressing datagrams
- "name", e.g.,www.yahoo.com -used by humans
- Q: how to map between IP address and name, and vice versa?

## Domain Name System:

- distributed database implemented in hierarchy of many name servers
- application-layer protocol: hosts, name servers communicate to resolve names (address/name translation)
  - note: core Internet function, implemented as applicationlayer protocol
  - complexity at network's "edge"

nslookup

## DNS as a Protocol

- Provide translation services to http, smtp, ftp, etc
- Port 53
- Supported by UDP
- ❖ Add additional delay that we did not mention<sup>©</sup>
  - You need to know who www.yahoo.com is.

## DNS: services, structure

#### **DNS** services

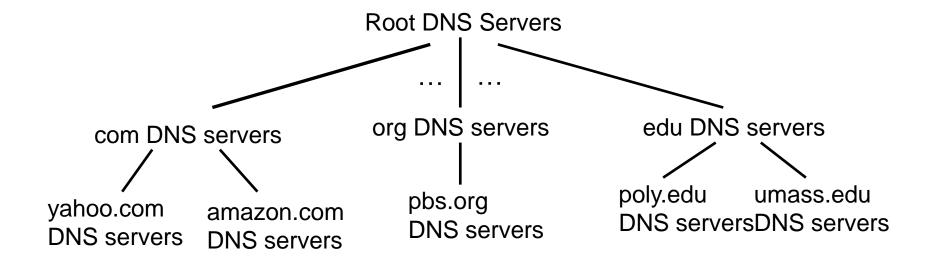
- hostname to IP address translation
- host aliasing
  - canonical, alias names
- mail server aliasing
- load distribution
  - replicated Web servers: many IP addresses correspond to one name

## why not centralize DNS?

- single point of failure
- traffic volume
- distant centralized database
- maintenance

A: doesn't scale!

## DNS: a distributed, hierarchical database

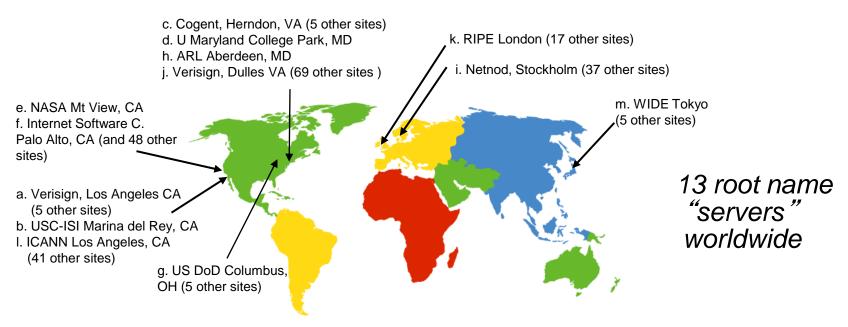


#### client wants IP for www.amazon.com; Ist approx:

- client queries root server to find com DNS server
- client queries .com DNS server to get amazon.com DNS server
- client queries amazon.com DNS server to get IP address for www.amazon.com

## DNS: root name servers

- contacted by local name server that can not resolve name
- root name server:
  - contacts authoritative name server if name mapping not known
  - gets mapping
  - returns mapping to local name server



# TLD, authoritative servers

#### top-level domain (TLD) servers:

- responsible for com, org, net, edu, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, ca, jp
- Network Solutions maintains servers for .com TLD
- Educause for .edu TLD

#### authoritative DNS servers:

- organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- can be maintained by organization or service provider

## Local DNS name server

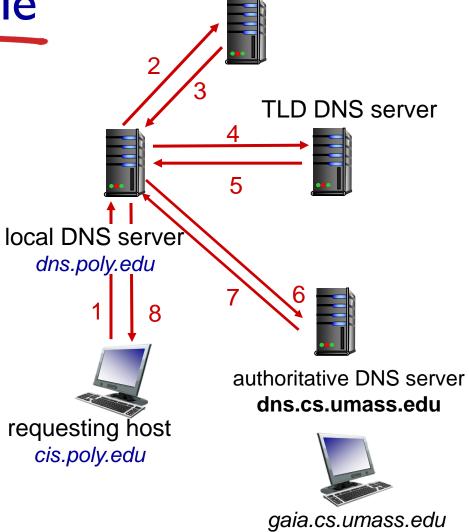
- does not strictly belong to hierarchy
- each ISP (residential ISP, company, university) has one
  - also called "default name server"
- when host makes DNS query, query is sent to its local DNS server
  - has local cache of recent name-to-address translation pairs (but may be out of date!)
  - acts as proxy, forwards query into hierarchy

DNS name resolution example

 host at cis.poly.edu
 wants IP address for gaia.cs.umass.edu

## iterated query:

- contacted server replies with name of server to contact
- "I don't know this name, but ask this server"

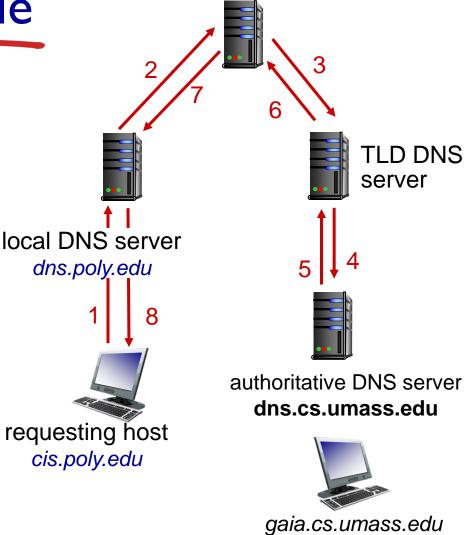


root DNS server

DNS name resolution example

## recursive query:

- puts burden of name resolution on contacted name server
- heavy load at upper levels of hierarchy?



root DNS server

# DNS: caching, updating records

- once (any) name server learns mapping, it caches mapping
  - cache entries timeout (disappear) after some time (TTL)
  - TLD servers typically cached in local name servers
    - thus root name servers not often visited
- cached entries may be out-of-date (best effort name-to-address translation!)
  - if name host changes IP address, may not be known Internet-wide until all TTLs expire
- update/notify mechanisms proposed IETF standard
  - RFC 2136

## **DNS** records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

## type=A

- name is hostname
- value is IP address

## type=NS

- name is domain (e.g., foo.com)
- value is hostname of authoritative name server for this domain

## type=CNAME

- name is alias name for some "canonical" (the real) name
- www.ibm.com is really servereast.backup2.ibm.com
- value is canonical name

## type=MX

 value is name of mailserver associated with name

# DNS protocol, messages

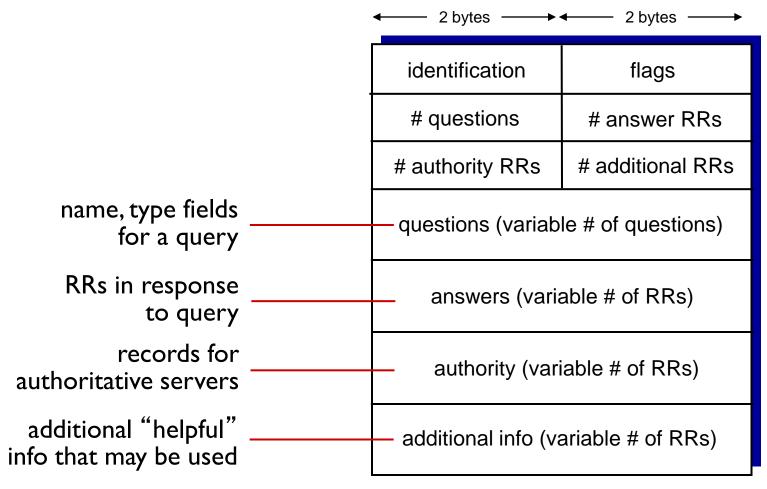
query and reply messages, both with same message format

#### msg header

- identification: I 6 bit # for query, reply to query uses same #
- flags:
  - query or reply
  - recursion desired
  - recursion available
  - reply is authoritative

2 bytes 2 bytes	
identification	flags
# questions	# answer RRs
# authority RRs	# additional RRs
questions (variable # of questions)	
answers (variable # of RRs)	
authority (variable # of RRs)	
additional info (variable # of RRs)	

# DNS protocol, messages



# Inserting records into DNS

- example: new startup "Network Utopia"
- register name networkuptopia.com at DNS registrar (e.g., Network Solutions)
  - provide names, IP addresses of authoritative name server (primary and secondary)
  - registrar inserts two RRs into .com TLD server: (networkutopia.com, dns1.networkutopia.com, NS) (dns1.networkutopia.com, 212.212.212.1, A)
- create authoritative server type A record for www.networkuptopia.com; type MX record for networkutopia.com

# Attacking DNS

#### DDoS attacks

- Bombard root servers
   with traffic
  - Not successful to date
  - Traffic Filtering
  - Local DNS servers cache IPs of TLD servers, allowing root server bypass
- Bombard TLD servers
  - Potentially more dangerous

#### Redirect attacks

- Man-in-middle
  - Intercept queries
- DNS poisoning
  - Send bogus relies to DNS server, which caches

## Exploit DNS for DDoS

- Send queries with spoofed source address: target IP
- Requires amplification

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