

# SuperScalar Processors

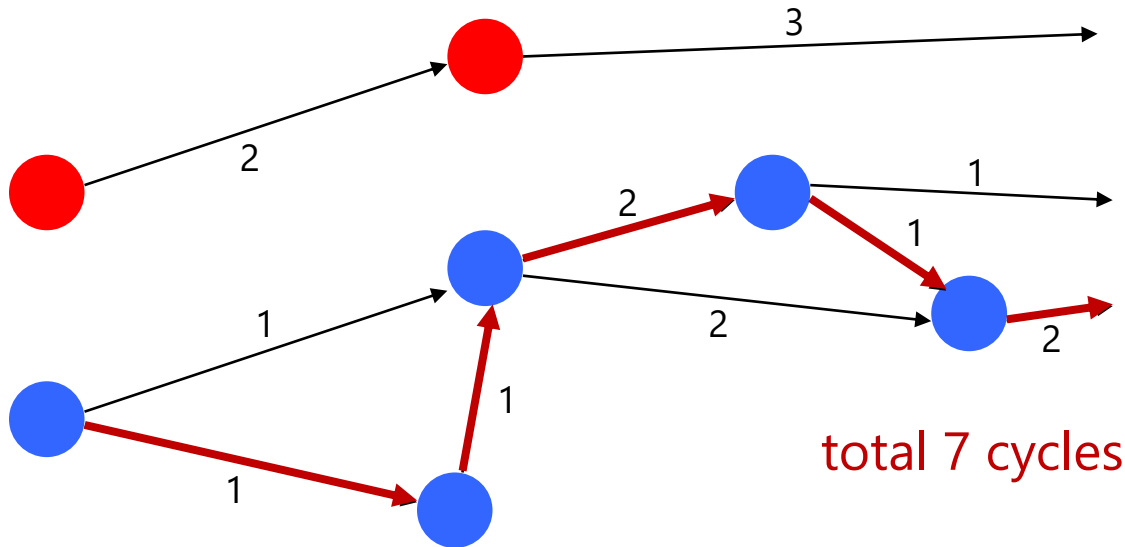
CS/COE 1541 (Fall 2020)  
Wonsun Ahn

# In-order vs. Out-of-order superscalars

- **Superscalar**: a wide-issue processor that does dynamic scheduling
  - Extracts instruction level parallelism (ILP) within the processor
- **In-order** superscalar: **does not reorder** instructions
  - Only detects hazards between instructions to insert bubbles
  - Only extracts ILP that arises from given ordering of instructions
  - The processor simulated in Project 1
- **Out-of-order** superscalar: **does reorder** instructions
  - Reorders instructions to remove hazards and increase utilization
  - Typically results in higher performance compared to in-order
  - But dynamic reordering consumes lots of power
- Out-of-order sounds more exciting so let's talk about that

# Name of the game is still ILP

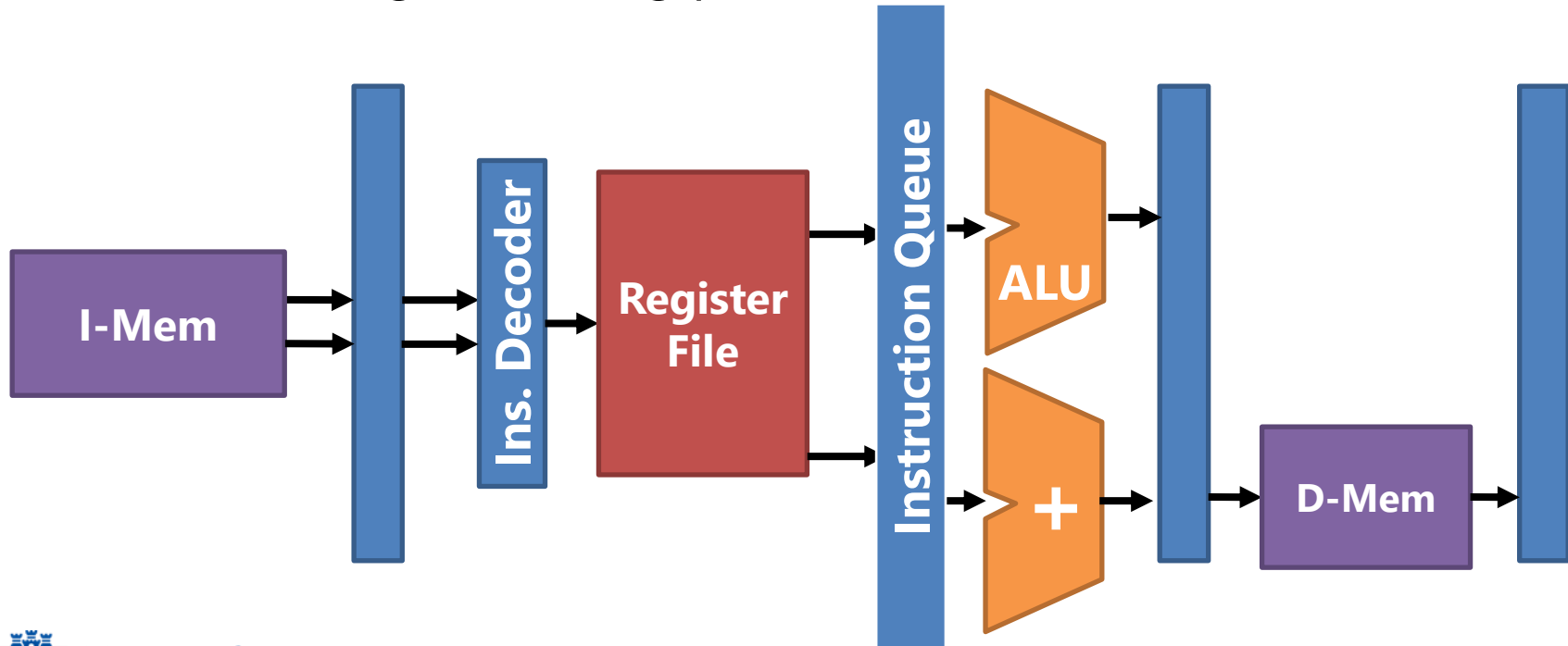
- The processor internally constructs the data dependency graph
- The processor tries to take advantage of ILP as much as possible
  - By executing the red nodes in parallel with the blue nodes



*illustration courtesy of Dr. Melhem*

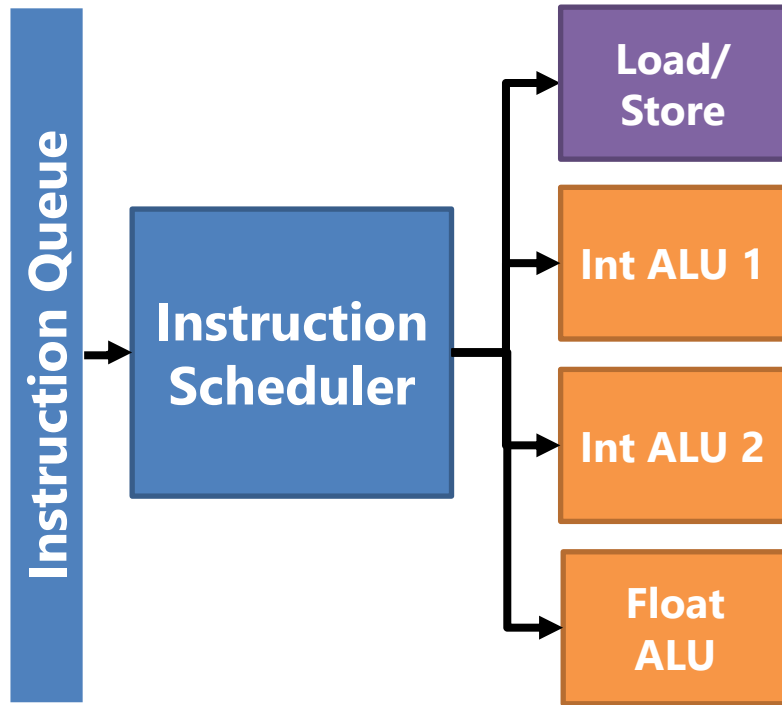
# Instruction Queue

- In order to expose ILP, superscalars need a big **instruction window**
  - Just like the compiler did for VLIWs
  - HW structure for storing instructions is called ***instruction queue***
  - Now EX stage has a big pool of instructions to choose from



# Instruction Queue

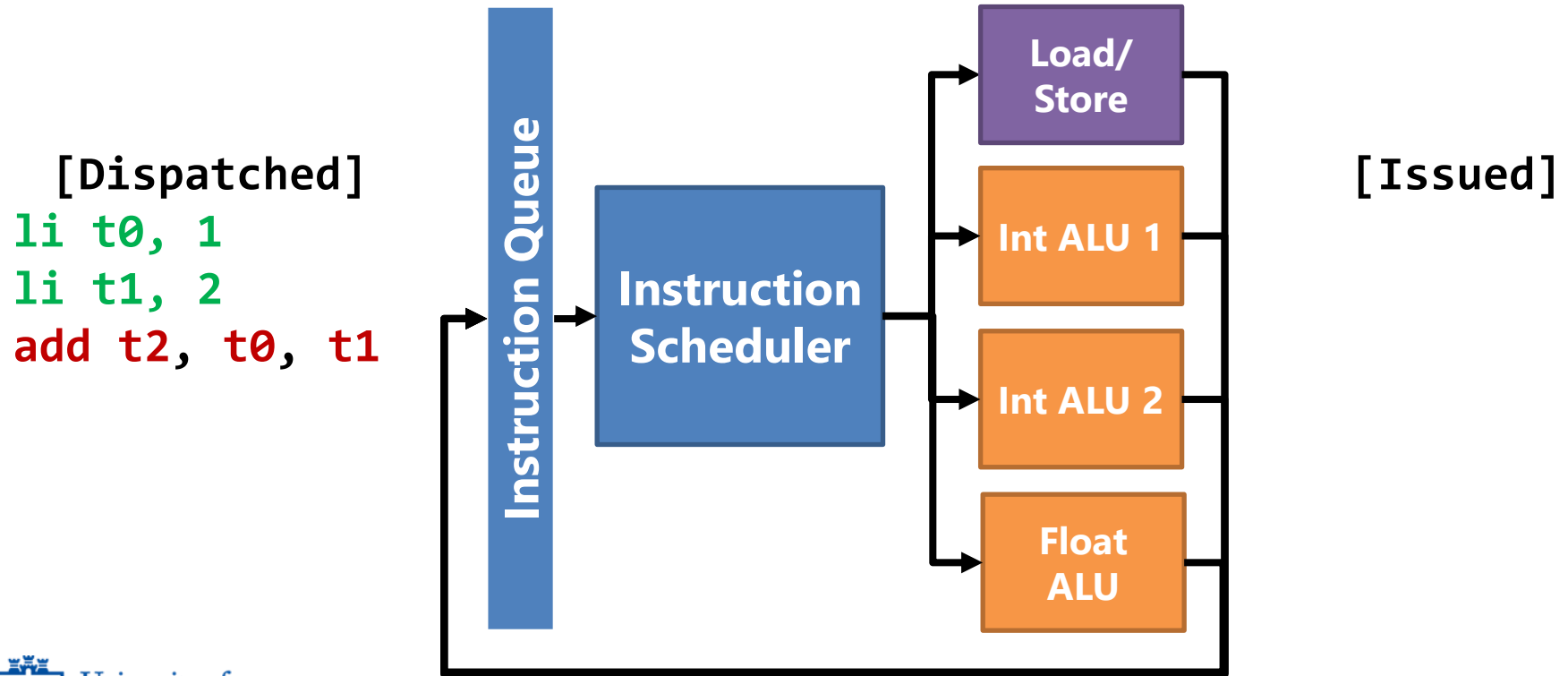
- In order to expose ILP, superscalars need a big **instruction window**
  - Just like the compiler did for VLIWs
  - HW structure for storing instructions is called ***instruction queue***



- At **ID**, instructions are decoded
  - And **dispatched** to the i-queue
- At **EX**, ready instructions are chosen from the instruction queue
  - Ready as in operands are available
  - And **issued** to an EX unit
- Typically queue is always full
  - Insts start queueing up when insts fail to issue due to hazards

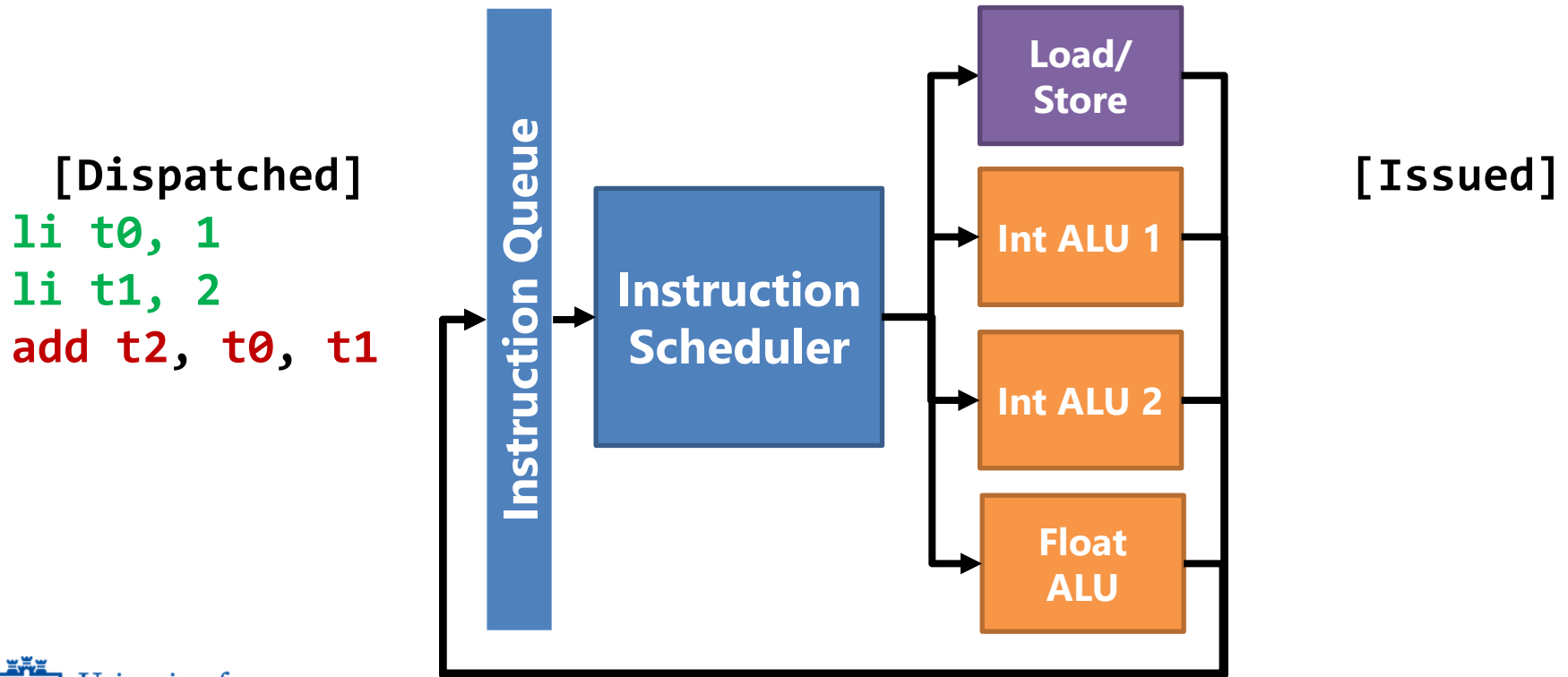
# Scheduling the Instruction Queue

- Now we have pool of instructions. When do they become ready?
  - Ready operands and instructions are in **green**
  - Not ready operands and instructions are in **red**



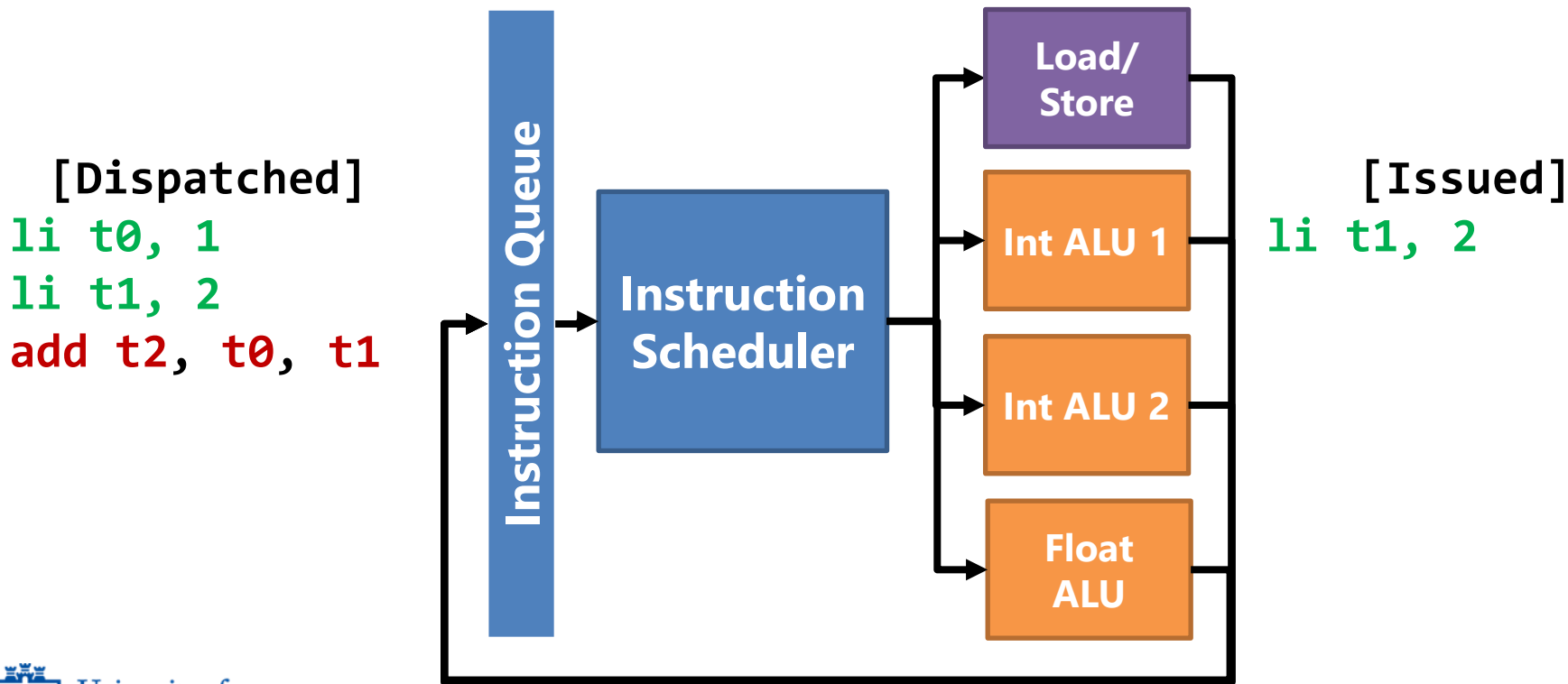
# Scheduling the Instruction Queue

- Initially both **li t0, 1** and **li t1, 2** are ready
  - The **li** instruction does not have any register operands
  - Instruction scheduler has a choice of what to issue



# Scheduling the Instruction Queue

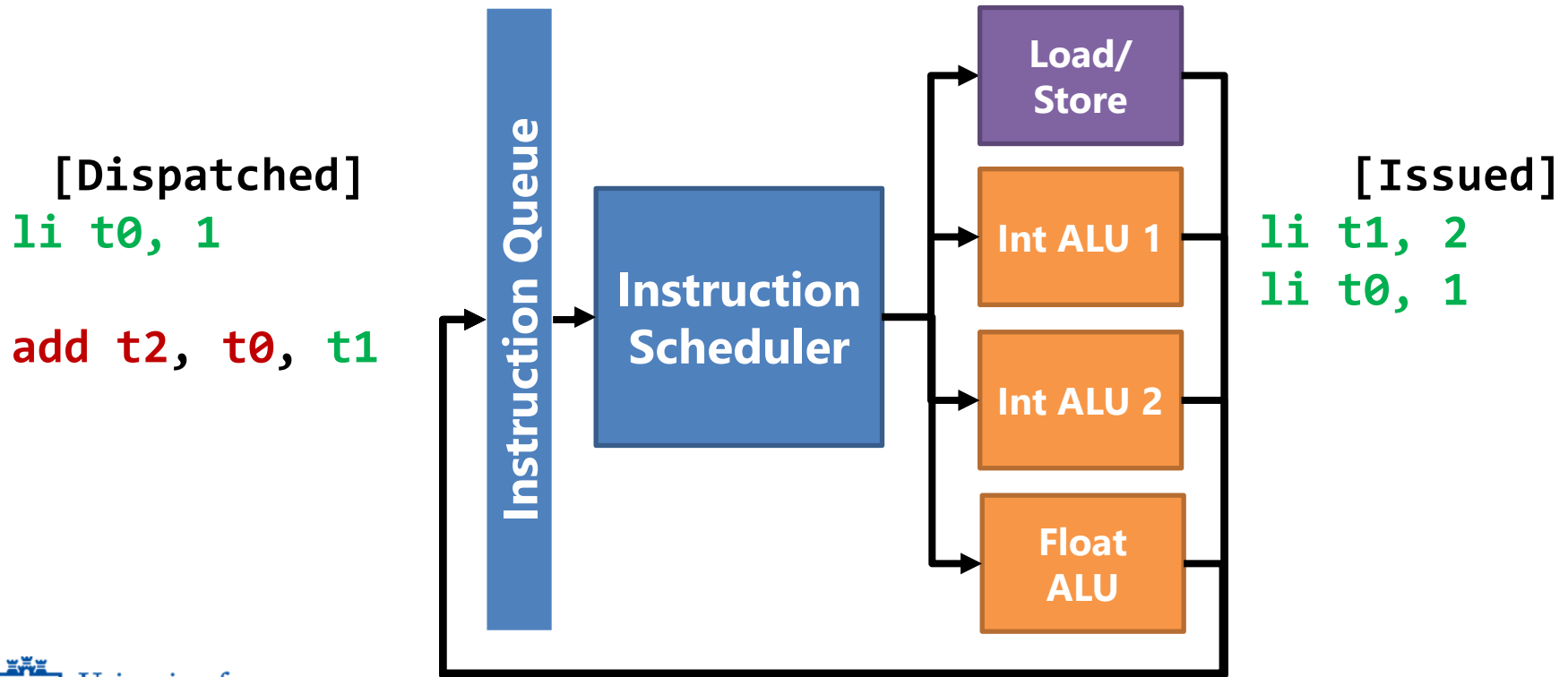
- Let's say the scheduler issues **li t1, 2** first
- Then the **t1** operand becomes ready after it completes





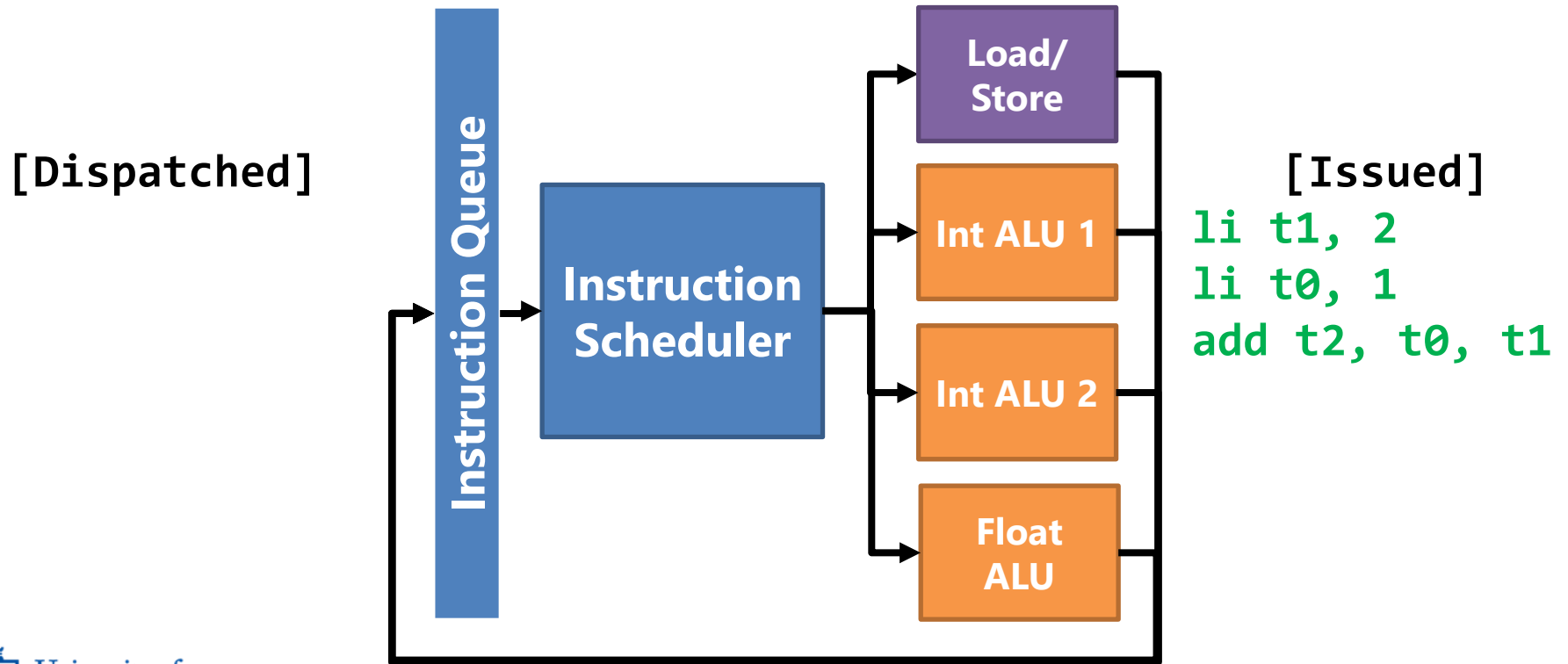
# Scheduling the Instruction Queue

- Now the only ready instruction **li t0, 1** issues
- Then the **t0** operand becomes ready after it completes
- Now **add t2, t0, t1** is finally ready to issue



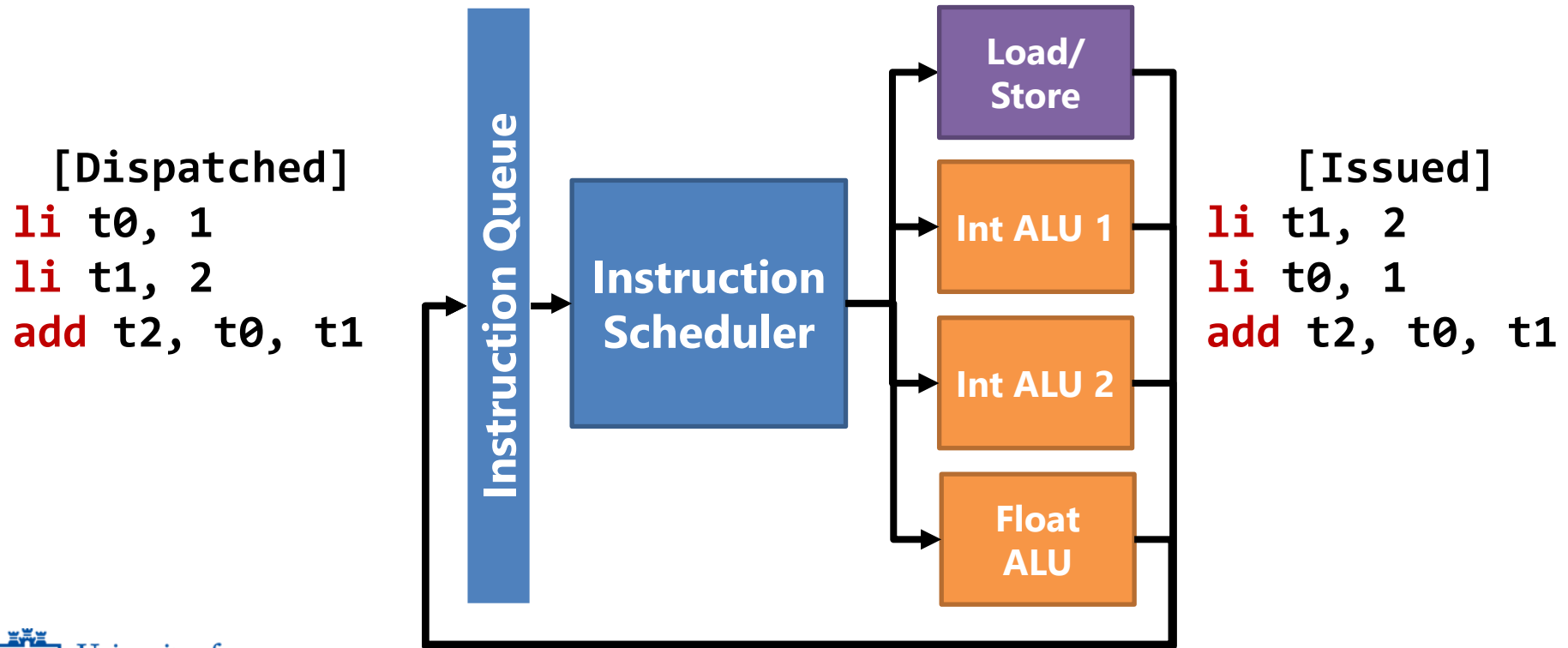
# Scheduling the Instruction Queue

- And we are done!



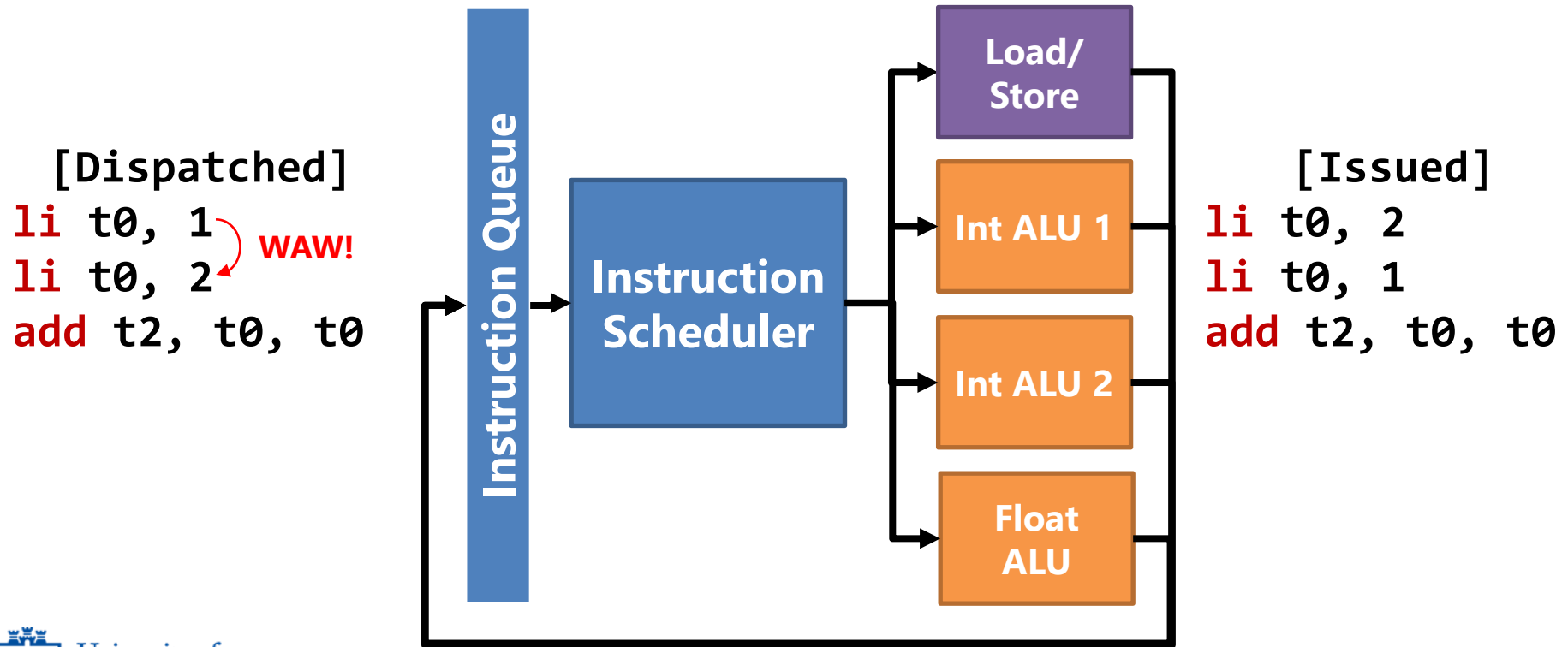
# Scheduling the Instruction Queue

- Note how we reordered **li** t0, 1 and **li** t1, 2
  - There are no dependencies between the two, so no issues
  - Also, RAW dependency with **add** t2, t0, t1 was enforced



# What if we had a WAW dependency?

- Reordering **li t0, 1** and **li t0, 2** still allowed (both are ready)
  - Now **t2 = 4** in original code, but **t2 = 2** during execution!
  - How do we disallow this from happening?

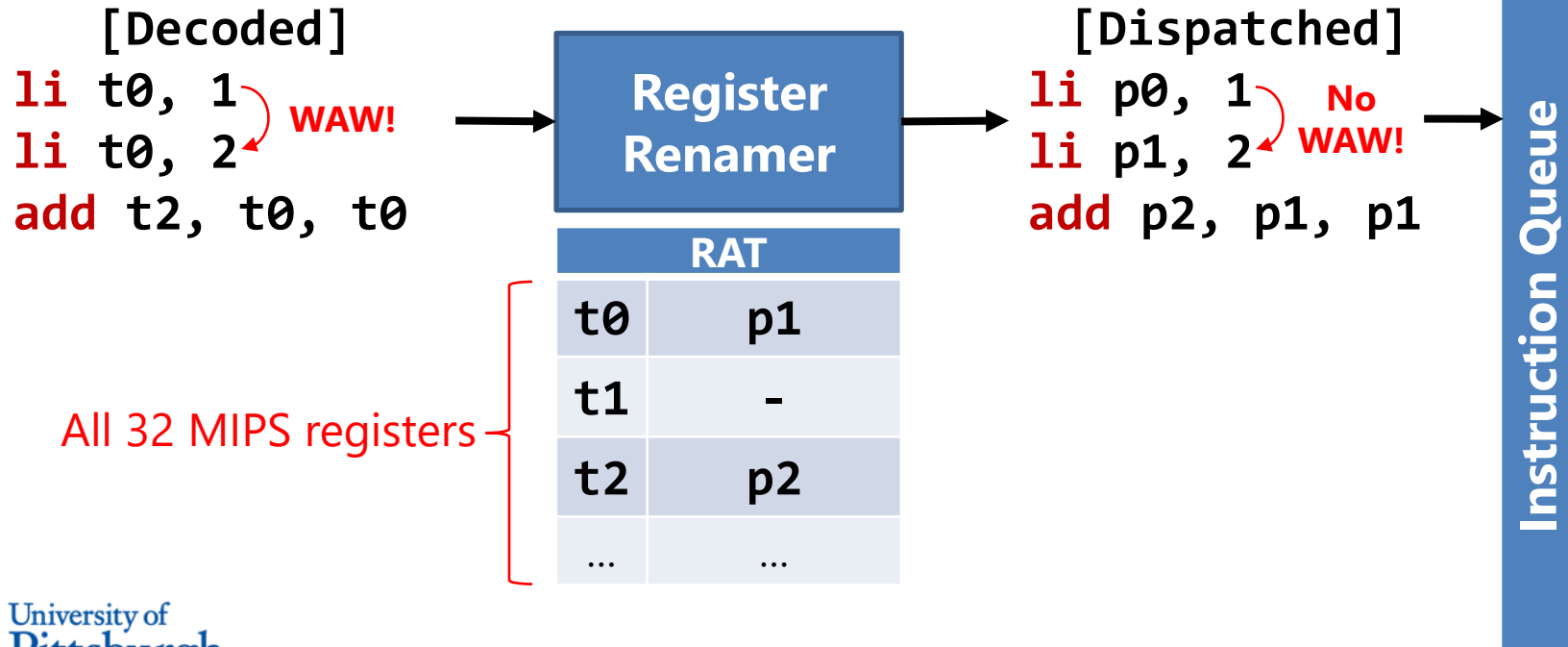


# WAW and WAR dependencies are tricky

- RAW (true) dependencies are automatically enforced
  - Instructions cannot issue until all operands are ready (written)
- WAW and WAR dependencies are not enforced
  - There is no data passing between the two instructions
  - The two instructions can become ready in any order
- We could somehow enforce WAW and WAR dependencies
  - But there is a better solution: **register renaming!**
  - Remember? That's what the compiler did to remove WAW/WAR.

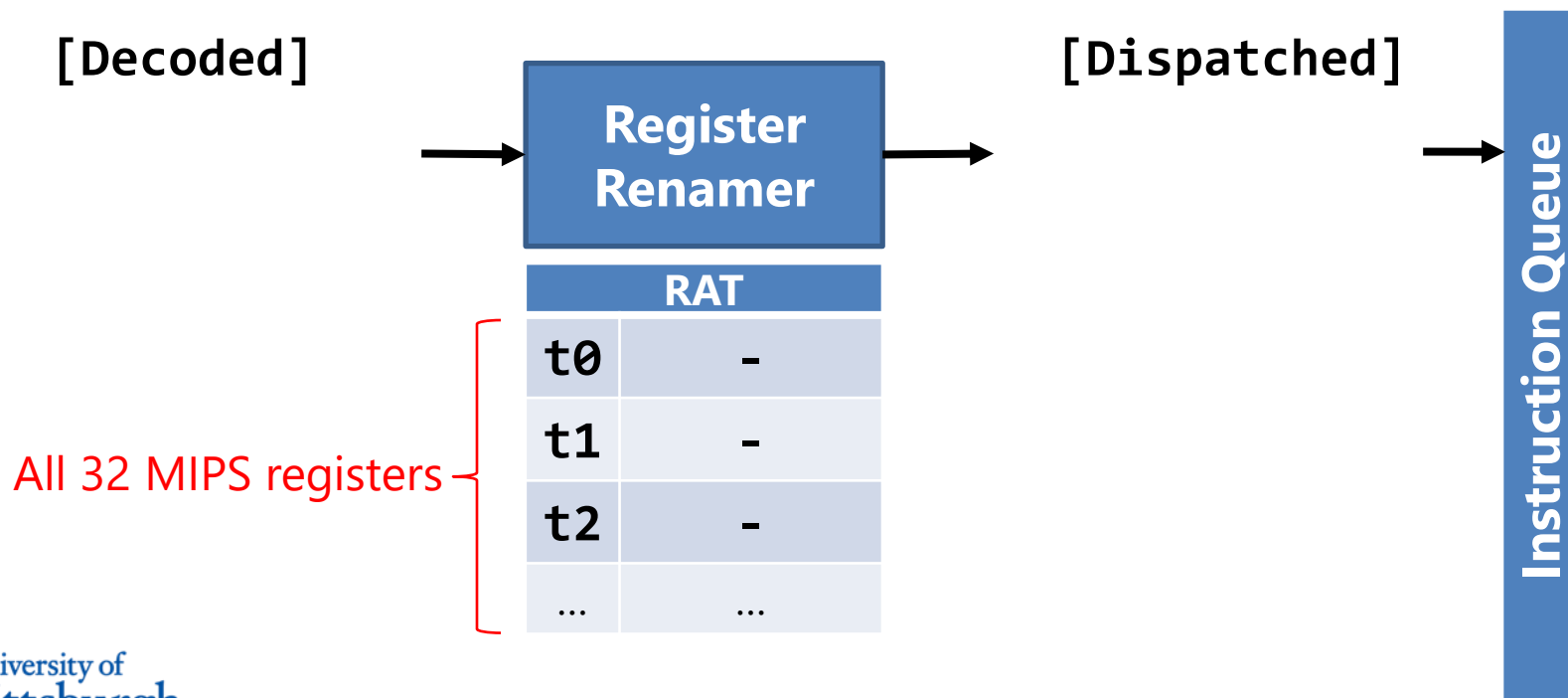
# Register Renamer and the RAT

- As soon as decode, **Register Renamer** renames all registers
  - Done with the help of the **Register Alias Table (RAT)**
  - RAT** is current mapping between **architectural** and **physical** registers
    - Architectural registers: Registers in ISA used in programs (t0, t1, t2, ...)
    - Physical registers: Renamed registers used in processor (p0, p1, p2, ...)



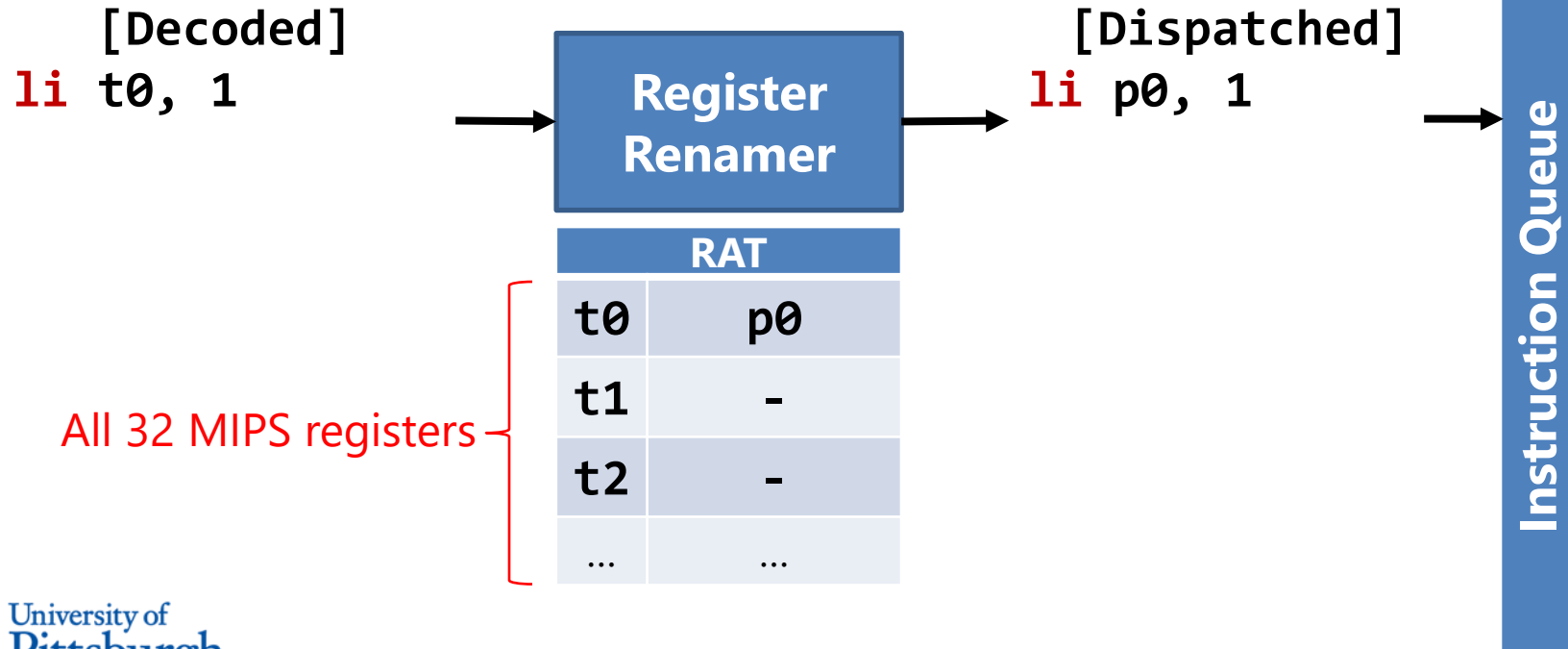
# Register Renamer and the RAT

- So how does the RAT work?
- Initially, no assignments have been done, so mapping is empty.



# Register Renamer and the RAT

1. **li** **t0**, 1 is decoded, destination **t0** is renamed to **p0**





# Register Renamer and the RAT

1. **li** **t0**, 1 is decoded, destination **t0** is renamed to **p0**
2. **li** **t0**, 2 is decoded, destination **t0** is renamed to **p1**
  - Remember the single assignment rule?
  - A new value always gets a new register

[Decoded]  
**li** **t0**, 1  
**li** **t0**, 2



[Dispatched]  
**li** **p0**, 1  
**li** **p1**, 2

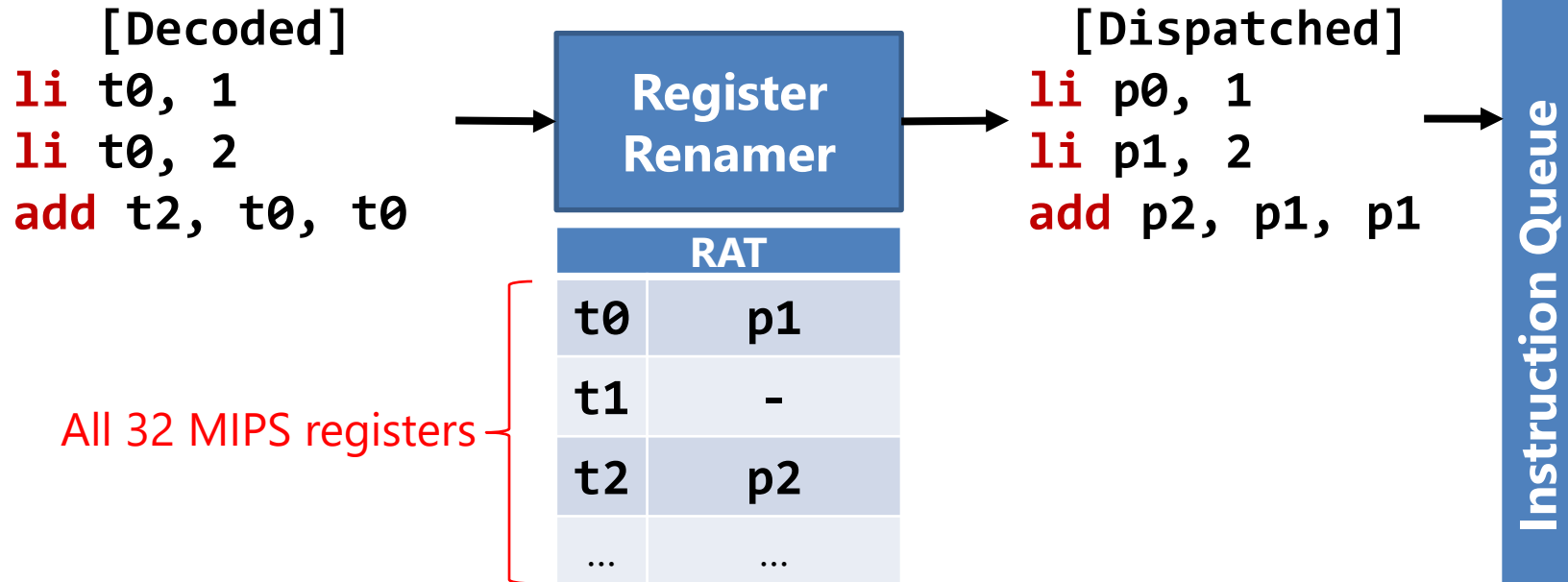


RAT	
t0	p1
t1	-
t2	-
...	...

All 32 MIPS registers

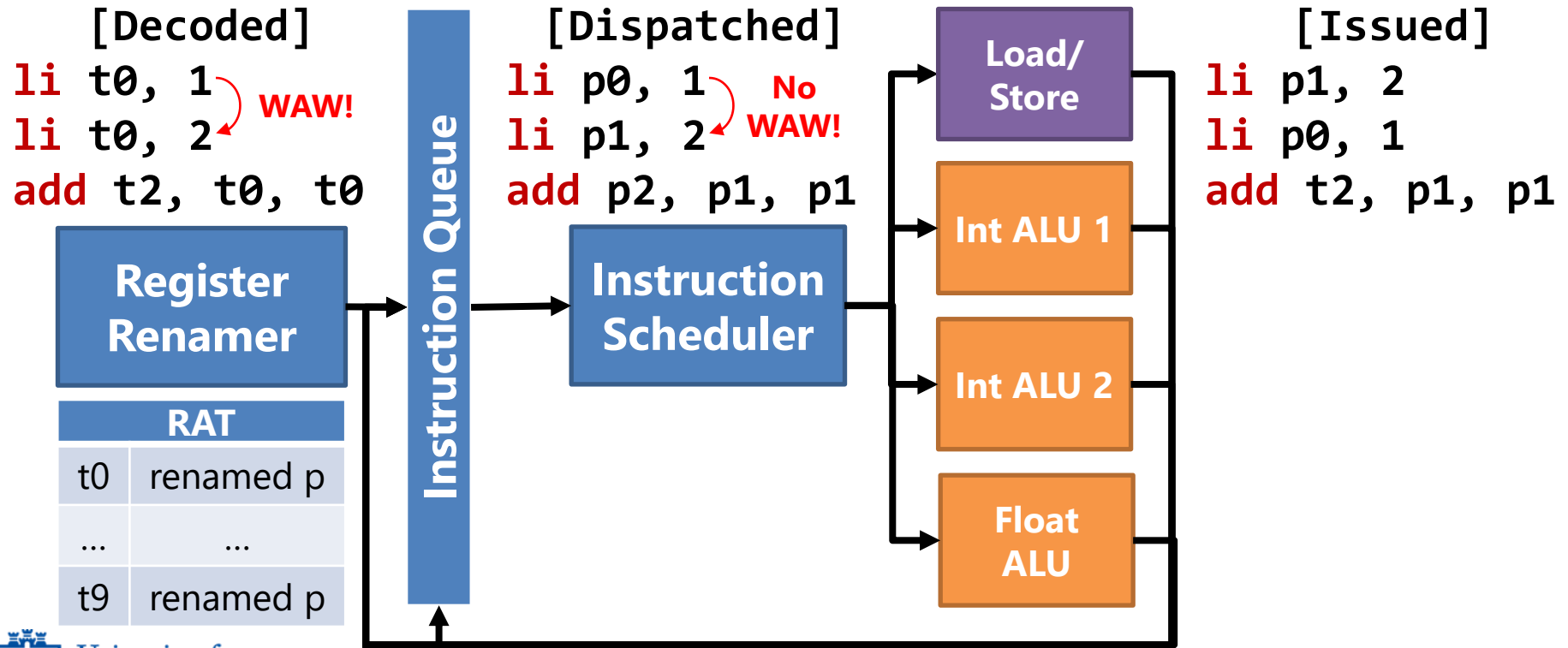
# Register Renamer and the RAT

1. **li** **t0**, 1 is decoded, destination **t0** is renamed to **p0**
2. **li** **t0**, 2 is decoded, destination **t0** is renamed to **p1**
3. **add** **t2**, **t0**, **t0** is decoded:
  - Two **t0** input registers use current mapping **p1**
  - Destination register **t2** is renamed to **p2**



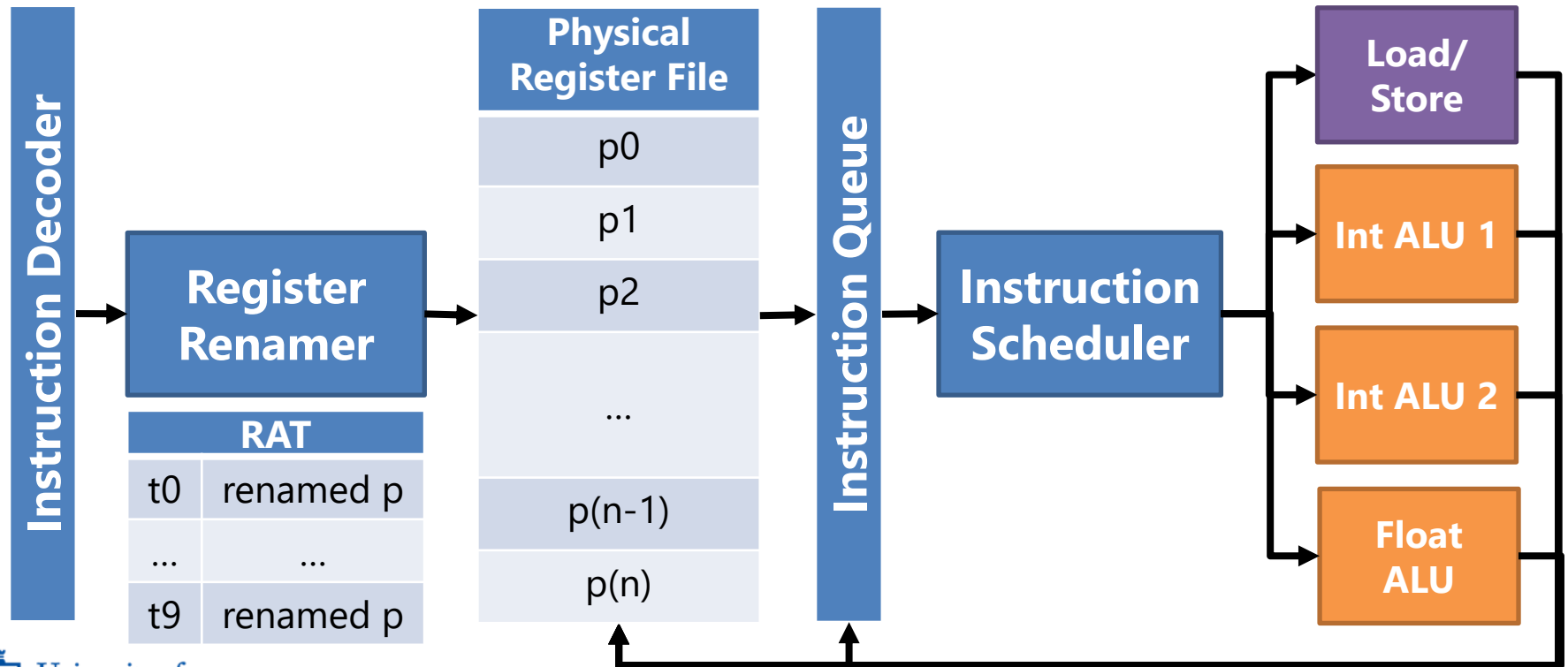
# Register Rename Removes all WAW/WAR Deps

- By the time instructions are dispatched to i-queue
  - All architectural registers have been renamed to physical registers
  - All WAR and WAW dependencies have been removed



# All Computation Done using Physical Registers

- Now ID stage (dispatch) reads registers for physical register file
- All data forwarding also done based on physical registers



# Limits on IPC

- We already discussed limits on pipelining.
- Time to discuss limits on IPC improvements through wide-issue!
- There is a **fundamental limit** to achievable IPC
  - Amount of **ILP (Instruction Level Parallelism)** in code
  - Remember the data dependence graph?
  - ILP is constrained only by true **RAW** dependencies
  - How about control dependencies?
    - Not a fundamental limit → can elide using branch prediction
- ILP is a **property of the program**, not the processor
  - This limit applies to both VLIW and superscalar processors

# ILP present in different programs

- After renaming, theoretical limit of IPC is 35 ~ 4003!

Benchmark	IPC No Renaming	IPC Register Renaming	IPC Memory Renaming	IPC r29 Removed	IPC 10K Window
compress95	3.12	26.25	73.88	226.33	18.89
cc1	3.61	39.79	41.63	239.96	86.45
go	2.50	49.15	53.77	141.46	70.71
ijpeg	2.41	55.47	93.60	94.11	52.94
li	3.56	19.60	19.61	81.45	27.70
m88ksim	2.76	19.93	62.06	363.26	20.50
perl	3.47	82.01	127.57	153.05	128.84
vortex	4.57	26.26	26.27	271.97	92.04
applu	2.82	106.65	2037.61	2076.06	78.67
apsi	3.6	54.89	183.44	1224.86	79.56
fpppp	3.33	103.62	774.13	1837.96	134.62
hydro2d	3.09	144.80	147.67	242.08	52.14
mgrid	3.34	1876.11	3933.03	4003.44	286.48
su2cor	3.22	38.21	34.81	55.56	47.60
swim	3.10	112.08	112.08	275.21	89.15
tomcatv	3.61	32.85	61.47	119.67	58.91
turb3d	3.42	370.98	482.24	3652.46	0
wave5	3.25	29.28	35.71	35.71	0

Matthew Postiff et al. "The Limits of Instruction Level Parallelism in SPEC95 Applications". ACM SIGARCH Computer Architecture News, 1999

# Limits of Wide-Issue Processors

- Achieving the theoretical limit would be awesome
  - But in real life, processors are typically no more than 10-wide
- Practical limits on Wide-Issue Processors
  - Practical limits on the **IPC** you can achieve
    - Instruction queue size (impacts scheduling window)
    - Physical register file size (also impacts scheduling window)
    - Number and types of execution units (i.e. width of processor)
  - Upsizing above structures negatively impacts **cycle time**
    - Time to search and schedule instruction queue
    - Time to access register file (increased size and number of ports)

# Exceptions

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# Exceptions Review

- **Exception**: an event which causes the CPU to stop the normal flow of execution and go somewhere else (the exception handler)
- There are mainly two causes of exceptions:
  - **Software exceptions** (or **traps**): Triggered by a program instruction
    - Trap instruction: typically used to call OS routines (system calls)
    - Page fault: instruction accessed a page not mapped to memory
    - Divide-by-0: instruction performed a divide-by-0 arithmetic
    - Arithmetic overflow: instruction overflowed MAX\_INT of register
  - **Hardware exceptions** (or **interrupts**): Triggered by hardware event
    - User has typed on the keyboard
    - A network packet has arrived
    - A file block read has completed
- In all cases, the OS **exception handler** is invoked

# Handling exceptions

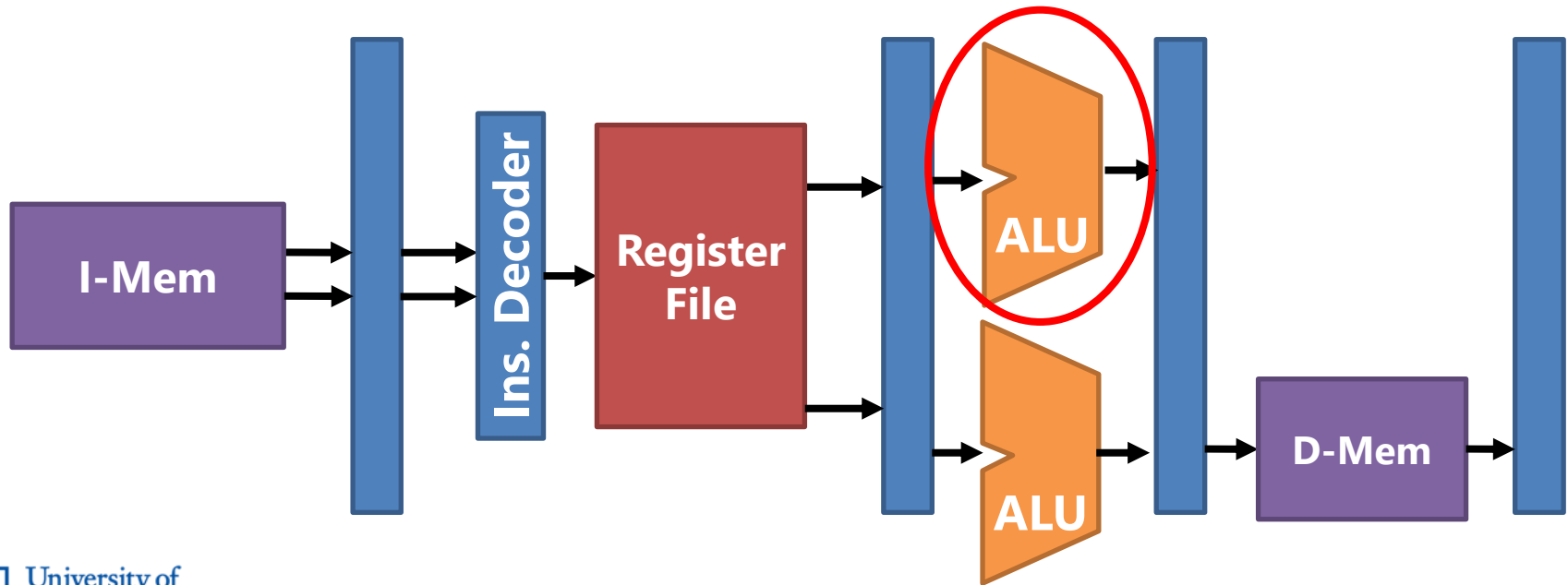
- What happens when an exception is triggered:
  1. Processor stops execution of user program.
  2. Processor stores information about exception (cause, PC).
  3. Processor jumps to the OS exception handler.
  4. Handler **creates backup** of program **register values** in memory.
  5. Handler inspects exception info and handles it accordingly.
    - While **overwriting** some of the registers that were backed up.
  6. Handler **restores** program **register values** from memory.
  7. Processor resumes execution of user program.
- Processor must provide precise register values at point of exception
  - Otherwise, when processor resumes, program will malfunction
  - Guaranteeing this is called a **precise exception**

# Rules for Precise Exceptions

1. **All instructions before** the exception must have executed
  2. **No instructions after** the exception must execute
- **Architectural state:** the state visible to the ISA (i.e. software)
    - State in architectural registers (For MIPS: t0, t1, t2, ...)
    - State in data memory
  - Architectural state at point of exception must reflect above rules

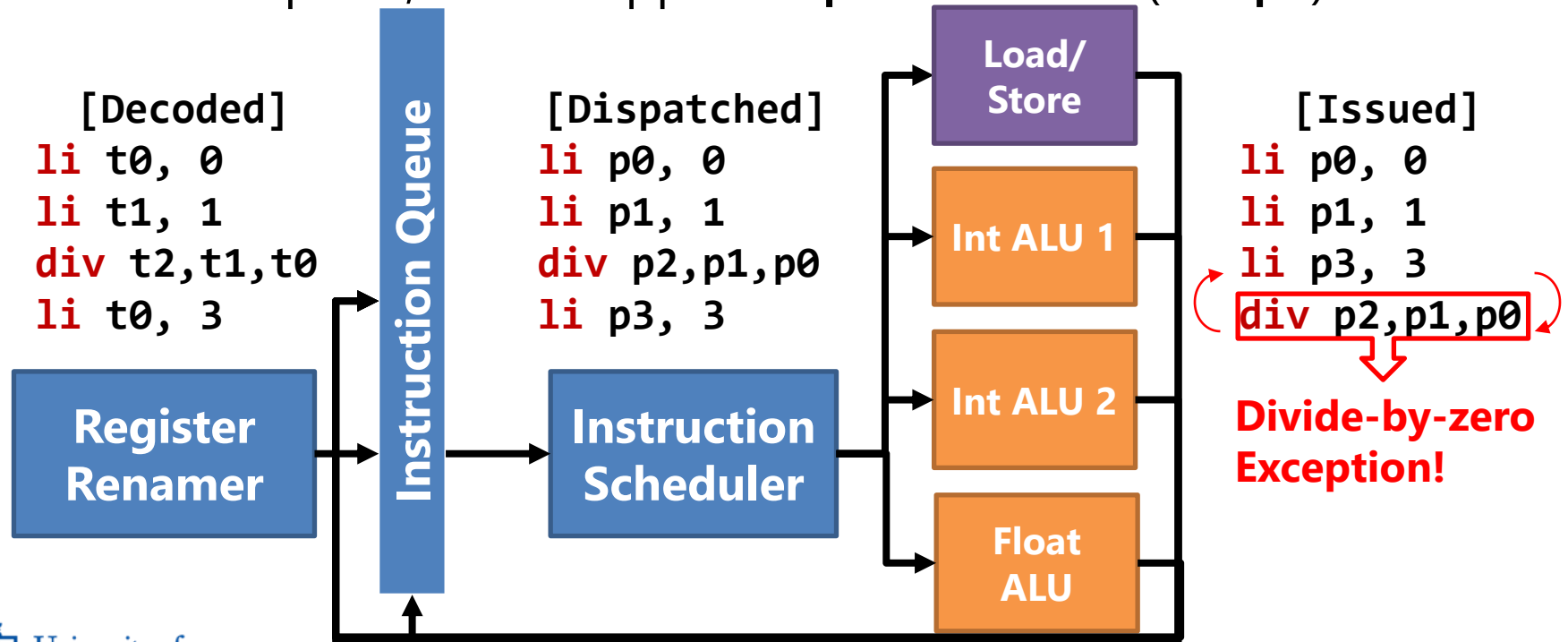
# Precise Exceptions in In-order Processors is Easy

- Exceptions are typically detected at the EX stage
  - Stage where all arithmetic happens as well as address calculations
- On exception, flush EX and all previous stages (ID and IF)
  - Since in-order, guarantees instructions following EX do not writeback
  - Only state leading up to instruction at EX will be written to reg / mem



# Precise Exceptions in Out-of-order Processors is Hard

- Suppose **div** t2,t1,t0 and **li** t0, 3 issue out-of-order as below
  - **div** p2,p1,p0 triggers a divide-by-zero exception (p0 = 0)
  - But at point of exception, t0 appears to be 3 due to **li** p3, 3!
    - At that point, t0 is mapped to p3 in the RAT (not p0)

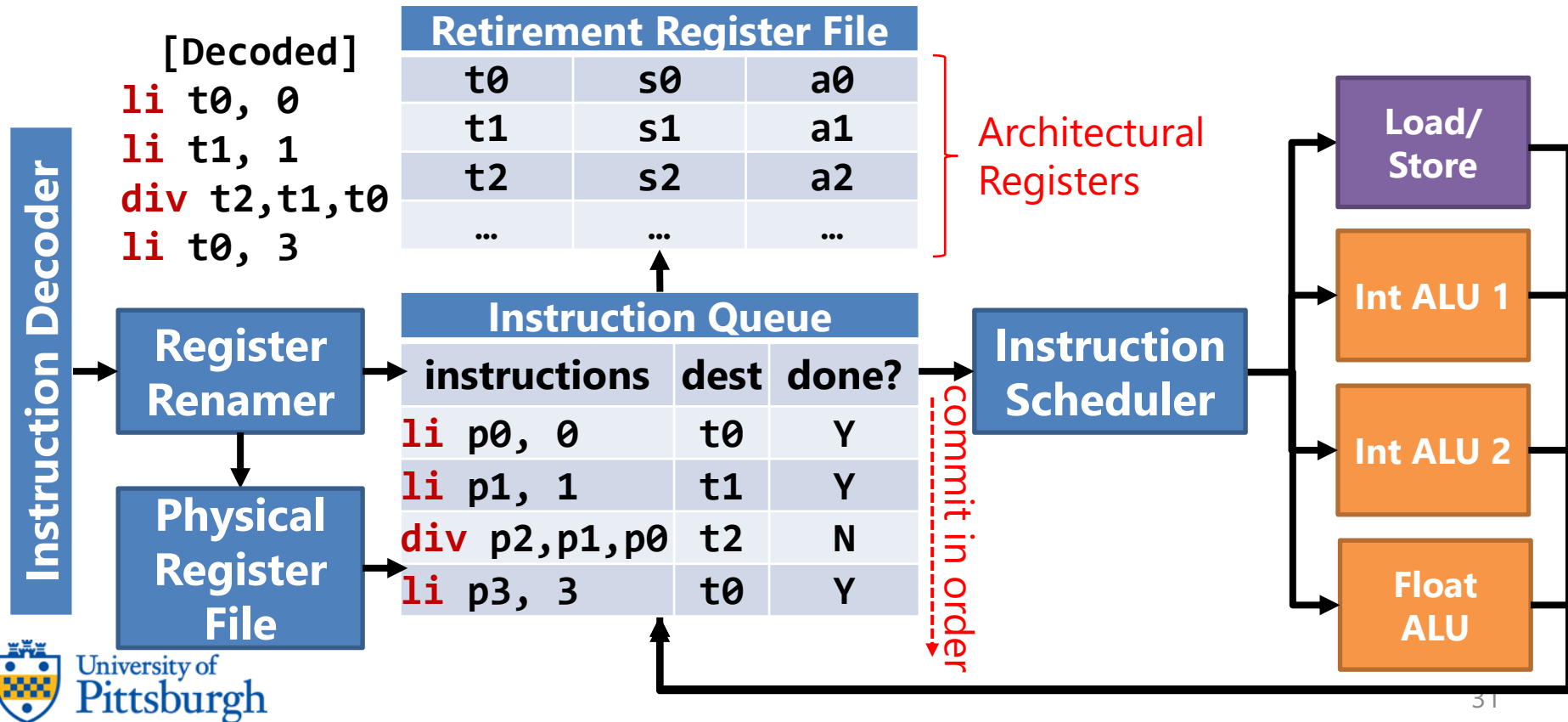


# Precise Exceptions in Out-of-order Processors is Hard

- This is the challenge with out-of-order processors
  - Instructions execute and complete out-of-order
  - For precise exceptions, instructions must appear to complete in-order
- Solution: update architectural state in-order
  - When instructions **execute**, have them only update “**internal**” state
    - Physical registers
    - Store queue (MEM queues up stores instead of performing them)
  - Internal state is transferred to **visible** state during in-order **commit**
    - Physical registers are copied to architectural registers
    - Store queue entries are written to memory

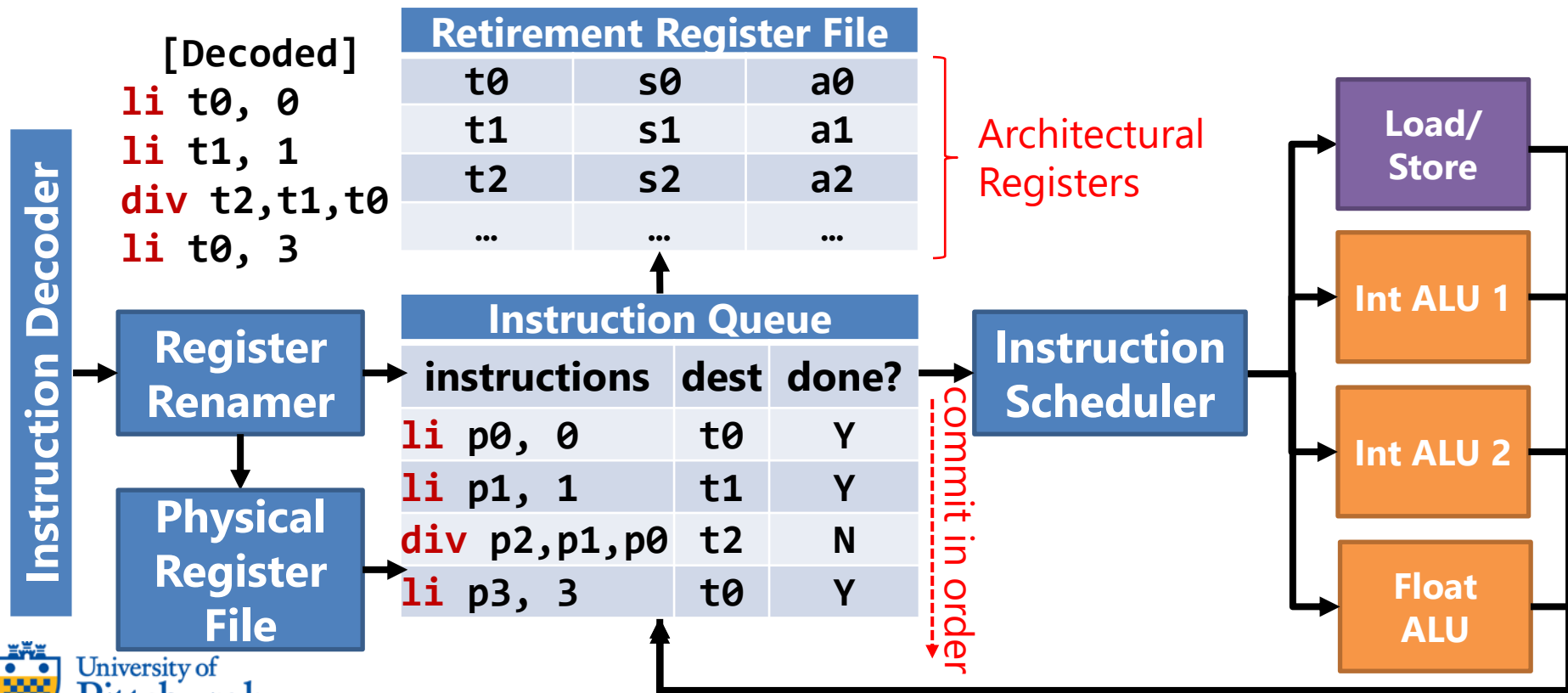
# In-order Commit

- Decoded instructions are **stored** to i-queue **in-order**
- Instructions **execute out-of-order** (updating done? field)
- Done instructions **commit** from i-queue **in-order** to Retirement Register File



# In-order Commit Example: Cycle 1

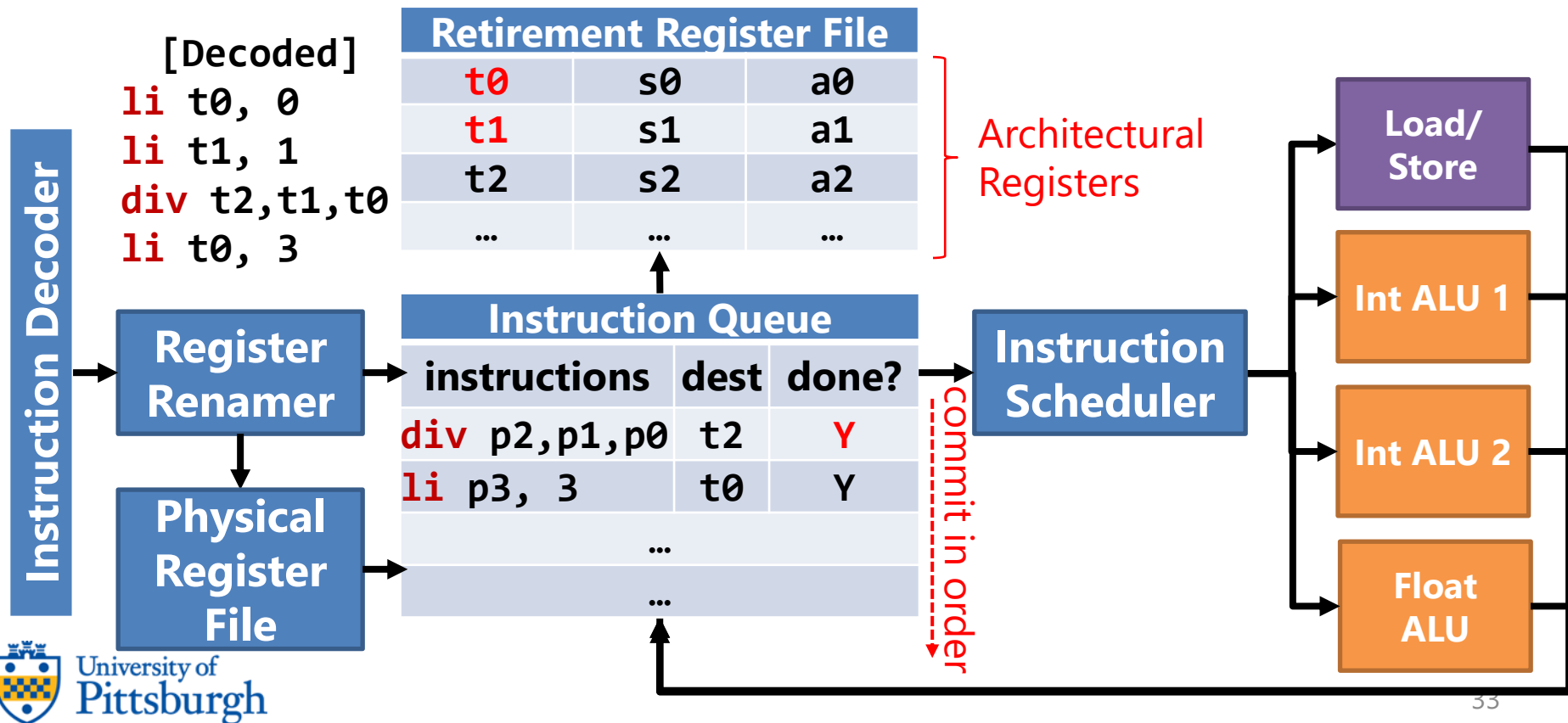
- At this point, all **li** instructions have completed but not the **div**
- li p0, 0** and **li p1, 1** can commit on the next cycle
  - But not **li p3, 3**, 3 since we have in-order commit!





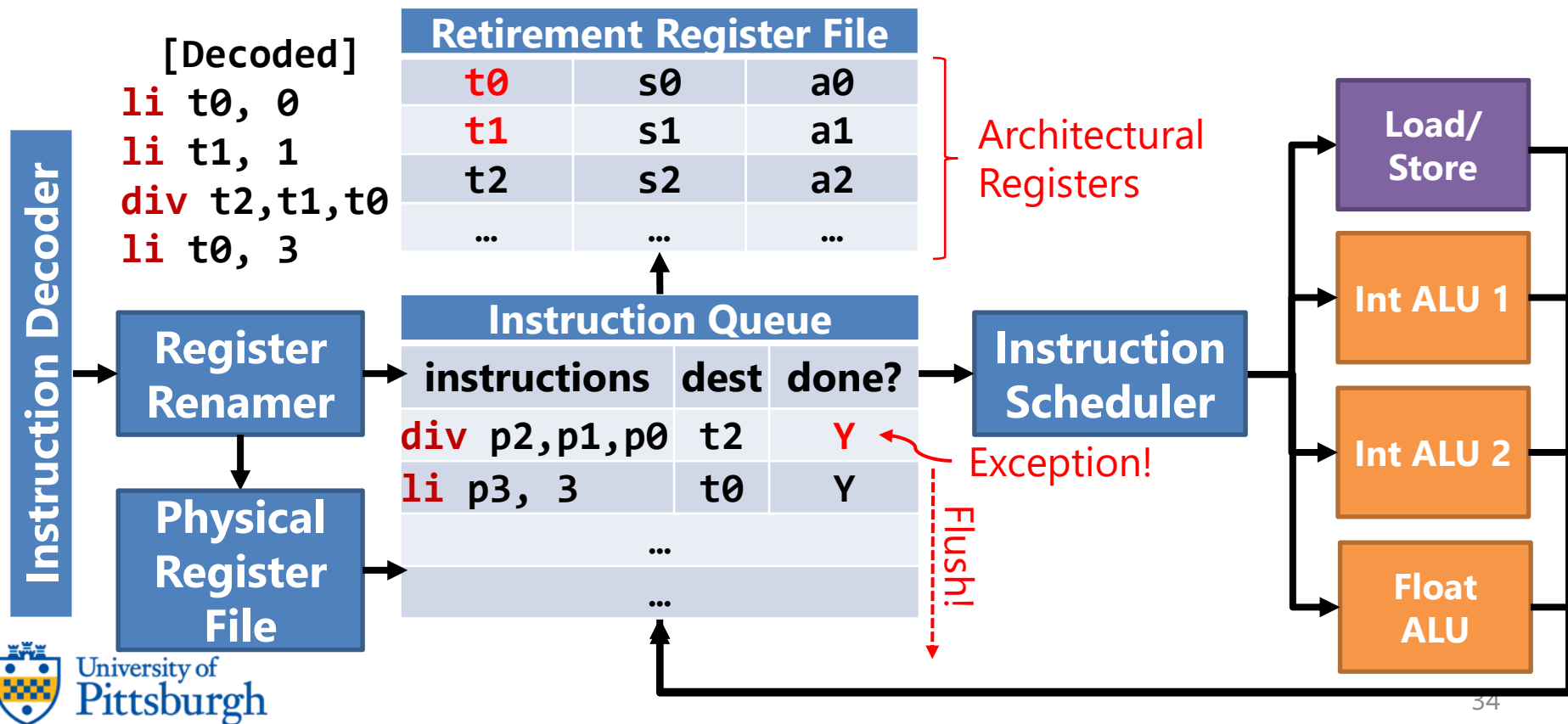
# In-order Commit Example : Cycle 2

- **li** p0, 0 and **li** p1, 1 have committed updating **t0** and **t1**
- **div** p2,p1,p0 has completed execution and is finally ready to commit
  - On completion, **div** sets an “exception bit” in i-queue (not shown here)



# In-order Commit Example : Cycle 3

- An exception is raised for **div** p2,p1,p0 when it tries to commit
  - Because “exception bit” has previously been set on completion
- **Retirement Register File** contains a **precise** architectural state



# Real Life Superscalars

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# The ARM Cortex-A8 architecture

- The ARM Cortex-A8 is an **in-order superscalar** processor
  - Notice the use of the **architectural register file**

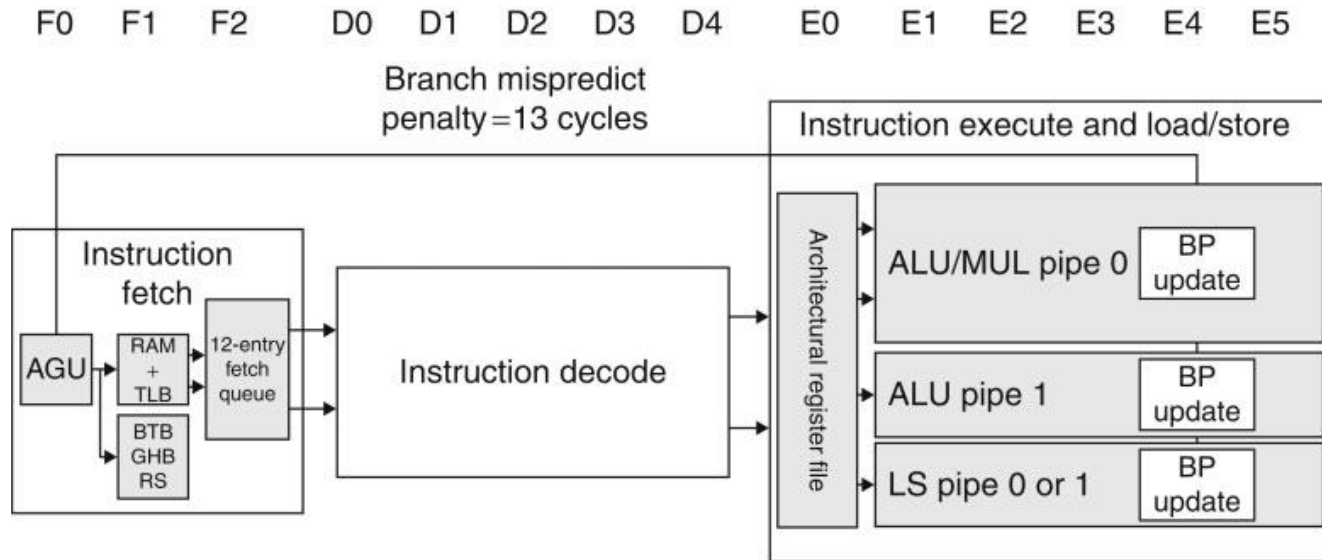
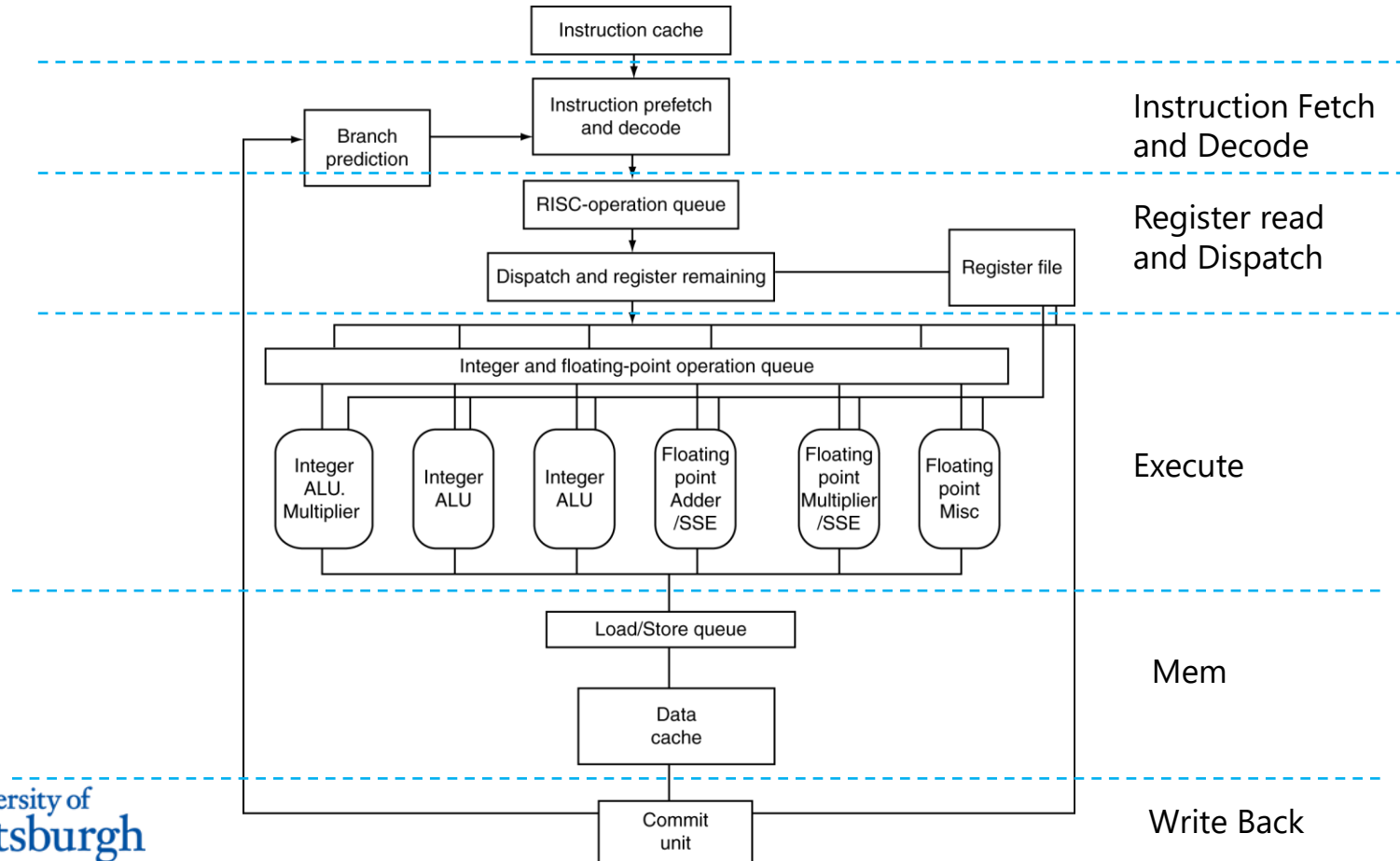


FIGURE 4.75 The A8 pipeline. The first three stages fetch instructions into a 12-entry instruction fetch buffer. The *Address Generation Unit* (AGU) uses a *Branch Target Buffer* (BTB), *Global History Buffer* (GHB), and a *Return Stack* (RS) to predict branches to try to keep the fetch queue full. Instruction decode is five stages and instruction execution is six stages.

# The AMD Opteron X4 Microarchitecture

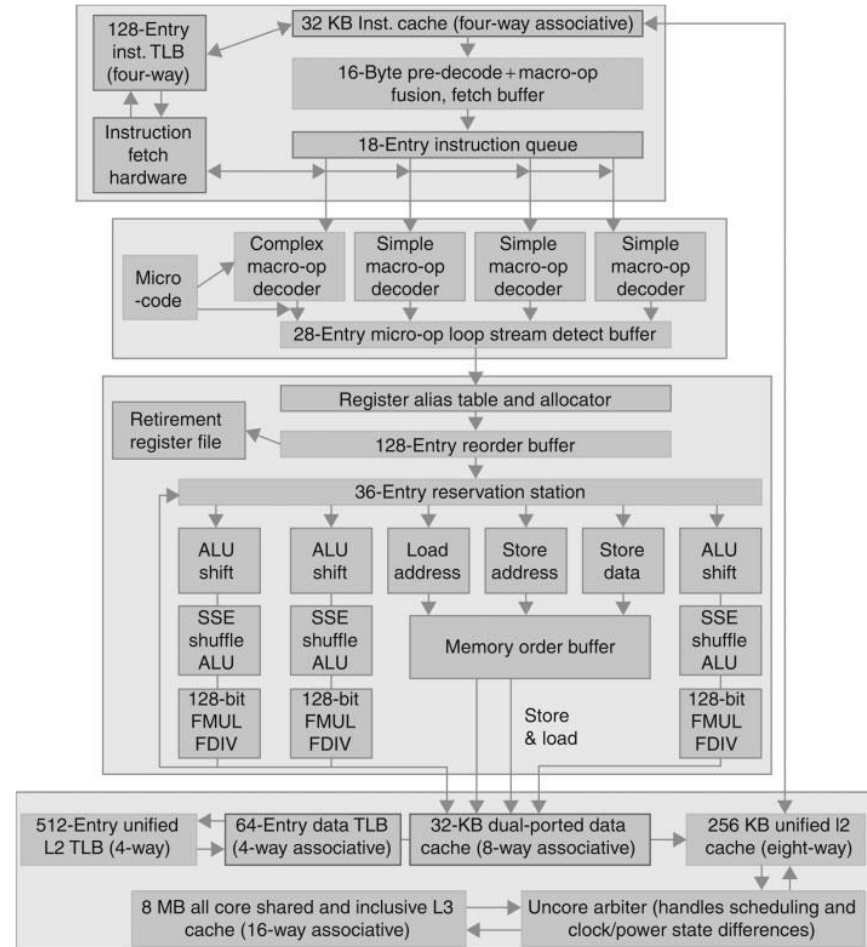
- The AMD Opteron is an **out-of-order superscalar** processor
  - **Commit unit** oversees retiring instructions from reorder buffer



# The Intel Core i7 architecture

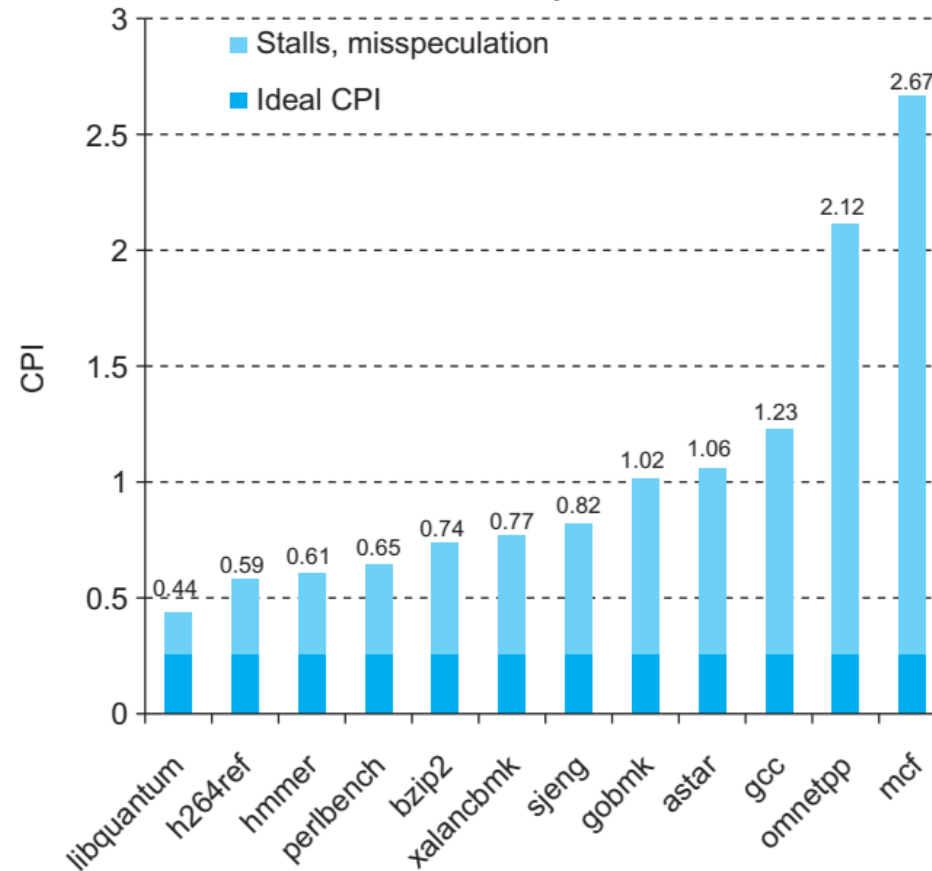
- The Intel Core i7 is another **out-of-order superscalar** processor

FIGURE 4.77 The Core i7 pipeline with memory components. The total pipeline depth is 14 stages, with branch mispredictions costing 17 clock cycles. This design can buffer 48 loads and 32 stores. The six independent units can begin execution of a ready RISC operation each clock cycle.



# Intel Core i7 Performance

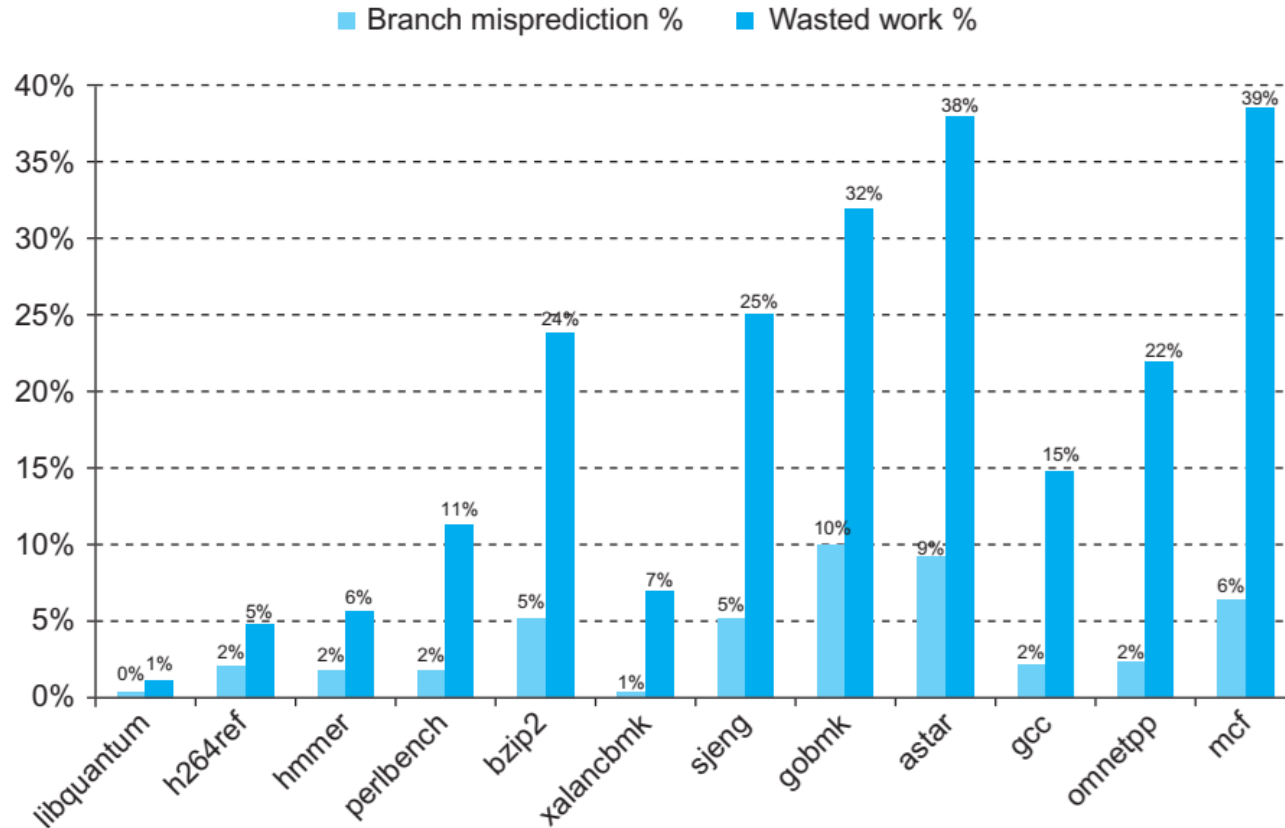
- Ideal CPI = 0.25 since this is a 4-wide processor



**FIGURE 4.78 CPI of Intel Core i7 920 running SPEC2006 integer benchmarks.**

# Intel Core i7 Impact of Branch Misprediction

- Due to deep pipeline, tiny misprediction can have outsized impact



**FIGURE 4.79** Percentage of branch mispredictions and wasted work due to unfruitful speculation of Intel Core i7 920 running SPEC2006 integer benchmarks.