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# Assignment Four

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## 1 Entropy and Mutual Information

Let  $X$  be a discrete random variable on a finite space  $\mathcal{X}$  with  $|\mathcal{X}| = k$ .

1. (a) Prove that the entropy  $H(X) \geq 0$ , with equality only when  $X$  is a constant.

PROOF: WLOG we can assume that  $p(x) > 0 \forall x \in \mathcal{X}$ , (using that  $0 \cdot \log 0 = 0$ ). We have that  $H(X) = -\sum_x p(x) \log p(x) = \sum_x p(x) \log p(x)^{-1} \geq 0$ , since  $p(x) > 0$  and  $p(x)^{-1} \geq 1$ . If  $H(X) = 0$  then  $\exists \alpha$  such that  $p(\alpha)^{-1} = 1 \Rightarrow p(\alpha) = 1$ . Hence  $X$  must be a constant, as needed.

- (b) Let  $X \sim p$  and  $q$  be the Uniform distribution on  $\mathcal{X}$ . What is the relation between  $D(p||q)$  and  $H(X)$ .

CLAIM:  $D(p||q) = -H(X) + \log k$

PROOF:

$$\begin{aligned} D(p||q) &= \sum_x p(x) \log \frac{p(x)}{q(x)} \\ &= \sum_x p(x) \log p(x) - \sum_x p(x) \log q(x) \\ &= -H(X) + \sum_x p(x) \log k \\ &= -H(x) + \log k \end{aligned}$$

- (c) An upper bound for  $H(X)$  is  $\log k$  since  $H(X) = \log k - D(p||q) \Rightarrow H(X) \leq \log k$ .

We consider a pair of discrete random variables  $(X_1, X_2)$  defined over the finite set  $\mathcal{X}_1 \times \mathcal{X}_2$ . Let  $p_{1,2}$ ,  $p_1$  and  $p_2$  denote respectively the joint distribution, the marginal distribution of  $X_1$  and the marginal distribution of  $X_2$ . Define the mutual information as:

$$I(X_1, X_2) = \sum_{(x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2} p_{1,2}(x_1, x_2) \log \frac{p_{1,2}(x_1, x_2)}{p_1(x_1)p_2(x_2)}$$

We again assume WLOG that  $p(x_1, x_2) > 0 \forall (x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2$

2. (a) CLAIM:  $I(X_1, X_2) \geq 0$

PROOF: Notice that  $I(X_1, X_2) = D(p_{1,2}||p_1p_2) \geq 0$  by the positiveness of  $D(\cdot||\cdot)$ .

- (b) We want to express  $I(X_1, X_2)$  as a function of  $H(X_1)$ ,  $H(X_2)$  and  $H(X_1, X_2)$ .

$$\begin{aligned} I(X_1, X_2) &= \sum_{(x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2} p_{1,2}(x_1, x_2) \log \frac{p_{1,2}(x_1, x_2)}{p_1(x_1)p_2(x_2)} \\ &= \sum_{(x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2} p_{1,2}(x_1, x_2) \log p_{1,2}(x_1, x_2) - p_{1,2}(x_1, x_2) \log p_1(x_1)p_2(x_2) \\ &= -H(X_1, X_2) - \sum_{(x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2} \left( p_1(x_1)p_2(x_2) \log p_1(x_1) - p_1(x_1)p_2(x_2) \log p_2(x_2) \right) \\ &= -H(X_1, X_2) - \sum_{j=1}^2 \sum_{(x_1, x_2) \in \mathcal{X}_1 \times \mathcal{X}_2} p_1(x_1)p_2(x_2) \log p_j(x_j) \\ &= -H(X_1, X_2) - \sum_{j=1}^2 \sum_{x_j \in \mathcal{X}_j} p_j(x_j) \log p_j(x_j) \\ &= -H(X_1, X_2) + H(X_1) + H(X_2) \end{aligned}$$

And so we can represent  $I(X_1, X_2)$  using  $H(X_1)$ ,  $H(X_2)$  and  $H(X_1, X_2)$ , as needed.

- (c) From the previous result we have that  $I(X_1, X_2) \geq 0 \Rightarrow H(X_1) + H(X_2) \geq H(X_1, X_2)$ , and so the maximal entropy of  $(X_1, X_2)$  is  $H(X_1) + H(X_2)$ . By definition this only occurs when  $I(X_1, X_2) = 0$ , which only occurs if  $p_{1,2}(x_1, x_2) = p_1(x_1)p_2(x_2) \forall x_1, x_2 \in \mathcal{X}_1 \times \mathcal{X}_2$ . This can be seen directly from the definition of  $I$  and using the strict positivity of  $p(x_1, x_2)$ .

## 2 HMM – Implementation

We use an HMM model to account for the possible temporal structure of some data. We consider the following HMM model: the chain  $(z_t)_{t=1}^T$  has  $K = 4$  possible states, with an initial probability distribution  $\pi \in \Delta_4$  and a probability transition matrix  $A \in \mathbb{R}^{4 \times 4}$  where  $A_{ij} = p(z_t = i | z_{t-1} = j)$  and conditionally on the current state  $z_t$ , we have observations obtained from Gaussian emission probabilities  $x_t | (z_t = k) \sim N(x_t | \mu_k, \Sigma_k)$ .