POLITECNICO DI TORINO TESTING AND FAULT TOLERANCE

REPORT: SW ASSIGNMENT MARCH ALGORITHM AND SBST

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1 Introduction

The software assignment was based on two portions of code that are used to:

- A March test algorithm on the RAM
- A set of SBST routines that are scheduled in portion of 1000 clock cycles, used to test three functional units: ALU, Multiplier and Register File.

1.1 March Algorithm

The March test has been put at the startup of the system, in order to perform a complete check of the data memory. The method, called <code>_run_march_algorithm</code> that is defined in the <code>march_algorithm.S</code> file, is called in the <code>crt0.S</code> just after the initialization of all registers to 0.

The base idea of the March Test is to perform a sequence of March Elements. Each March Element is composed of a sequence of write/read operations that must be performed in a cells and also for all the others.

The algorithm used in this assignment is the March LA, developed by Van de Goor, in 1996 has a complexity of 22n and it's able to detect Stuck-at-faults, Address-faults, Transition-faults and Coupling-faults, linked in any way. It composed of the following March elements:

- \(\psi \) (w0)
- \(\pha\) (r0, w1, w0, w1, r1)
- \(\pha\) (r1, w0, w1, w0, r0)
- \Downarrow (r0, w1, w0, w1, r1)
- \Downarrow (r1, w0, w1, w0, r0)
- ↓ (r0)

Moreover in order to make the test complete, the algorithm does not simply write 0 and 1 in the memory locations, but uses $\lfloor \log_2(m) \rfloor + 1$ patterns, where m is the word size. This allows to test also the intra-word coupling faults. So, the total complexity of the march test is:

$$6 \cdot 22 \cdot n = 132n$$

For this reason the _run_march_algorithm procedure acts only as a wrapper in order to call another procedure, called _march_LA.

This last procedure, is used to perform entirely the March LA algorithm, by passing one parameter on the a0 register that is used to define the 0 pattern that will be used when performing $w\theta$ and $r\theta$. The 1 pattern is simply computed by the produced, by negating a0.

This allows to not duplicate the code and execute the March LA for all the 6 patterns.

In order to speed up the end of the test, in case of error, the procedure terminates as soon as possible, returning 1 in case of an error and 0 otherwise.

In order to perform the test over the entire data ram, two sections have been defined:

- 1. ___RAM_START
- 2. RAM END

The first one, is used to define the starting point of the March LA while the second one its end. The end address is loaded into the t2 register but it can't be used as is. In fact, since it define the end point, the loop must avoid to write in that precise location (outside the dataram); for this reason, a subtraction of 4 is performed. The algorithms is 1:1 the implementation of the March LA, and by using t1 and t2 as starting and ending references, all the locations are read/written accordingly to the March algorithm.

1.1.1 Problems

One of the C library used in the linker script (newlib) creates in the .elf file a section in .data with a struct called <code>__rent</code> whose address is defined by the symbol <code>__global_impure_pointer</code>. The struct contains information about the errno, used by the ANSI C standard, stdout, stin and stderr and other callbacks. At the beginning, the .elf loaded during the simulation contains that information, and running the March LA destroys everything. In particular, the crt0 executes the function <code>memset</code> for the initialization of the bss segment, that reads data from the <code>__rent</code> structure, that has been erased by the memory test. The problem is that the RAM sections are (improperly) initialized by the testbench at power-on and this has been solved by copying at the startup all sections into a portion of a Flash memory. In this way, at the end of the march algorithm, data saved into the flash are moved into the RAM.

1.2 Results

As stated before, the March LA is a powerful algorithm to test memories that has the drawback of being complex from the point of view of the computational complexity. Moreover, the complexity increases even more, if the intra-coupling faults are tested. Having a RAM, which size is 0x40000, the time needed to test it with this implementation is: 212992660ns, that corresponds to more or less 0.2s.