

ragni-cas - A Pure *Ruby* Automatic Differentiation Library for Fast Prototyping of Interfaces

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Abstract

Ca. 100 words

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1. Motivation and significance

Ruby is a purely object-oriented scripting language that allows to express several programming paradigms. It was designed in the mid-1990s by Yukihiro Matsumoto (also known as *Matz*), and it is internationally standardized since 2012 as ISO/IEC 30170.

With the advent of the *Internet of Things*, a written from scratch version of the *Ruby* interpreter called *mRuby* (*eMbedded Ruby*) has been published on *GitHub* by Matsumoto in 2014. The new interpreter is a lightweight implementation aimed at both low power devices and personal computer that complies with the standard. *mRuby* has a completely new API, and it is designed to be embedded in a complex project as a front-end interface — e.g. a numerical optimization suite may use *mRuby* to get problem input definitions.

The *Ruby* code-base exposes a a large set of utilities in core and standard library. This set of tool can be furthermore expanded through libraries, also

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16 known as *gems*. Even the high number of gems deployed and available, there
17 is no library that implements a **symbolic automatic differentiation** (AD)
18 engine that also handles some basic computer algebra routines that is cross
19 compatible with all the different *Ruby* interpreter.

20 *Ruby* has matured its fame as a web oriented language with *Rails*, and
21 can efficiently generate code in other languages. An AD-capable gem is the
22 fundamental step to develop rapidly a specific code generator for well known
23 software — e.g. IPOPT.

24 The library that is described in this work, is a gem implemented in pure
25 *Ruby* code and thus compatible with all interpreter that complies with the
26 standard, that is able to perform symbolic AD and some simple computer
27 algebra operations. In particular the library aims at:

- 28 • be an instrument for rapid development of prototype interface for nu-
29 merical algorithms — including the *mRuby* engine or exporting gener-
30 ated code — that can be in different languages;
- 31 • rapidly generate descriptions of mathematical models, with easy to
32 implement workaround for numerical issues, changing on request how
33 the code is exported, and how functions are formulated in the target
34 language;
- 35 • creating a complete open-source CAS system for the *Ruby* language,
36 that is be compatible with all the interpreters that comply with the
37 standard, as a long-term ambitious impact.

38 This is not the first gem that tries to implement a CAS. The available
39 computer algebra library for *Ruby* are:

40 ***Rucas*** gem at early stage and with discontinued developing status; it im-
41 plements basic simplification routines. There is no AD method, but it

42 is one of the milestones. The development is discontinued since 2010.

43 ***Symengine*** is a wrapper for the C++ library *symengine*. The back-end
 44 library is very complete, but it is compatible only with the mainstream
 45 *Ruby* interpreter. At the moment, the *SciRuby* project reports the gem
 46 as broken, and removed it from its codebase. From a direct test, when
 47 performing AD of a function, the engine returns always `nil`.

48 2. Software description

49 2.1. Software Architecture

50 *ragni-cas* is an object oriented AD gem that supports some simple com-
 51 puter algebra routines such as *simplifications* and *substitutions*. When gem
 52 is required, automatically overloads some methods of the `Fixnum` and `Flo-`
 53 `at` classes, to make them compatible with the fundamental symbolic objects.

54 Each symbolic function is an object modeled by a class, that inherits from
 55 a common virtual ancestor: `CAS::Op(operation)`. An operation encapsulates
 56 sub-operations recursively, building a linked graph, that is the mathematical
 57 equivalent of function composition:

$$(f \circ g) \tag{1}$$

58 When a new operation is created, it is appended to the graph. The num-
 59 ber of branches are determined by the parent container class of the current
 60 symbolic function. There are three possible containers. Single argument
 61 functions — e.g. `sin(·)` — have as closest parent the `CAS::Op` class, that
 62 links to one sub-graph. Functions with two arguments — e.g. difference
 63 or exponential function — inherit from `CAS::BinaryOp`, that links to two
 64 subgraphs. Functions with arbitrary number of arguments — e.g. sum and

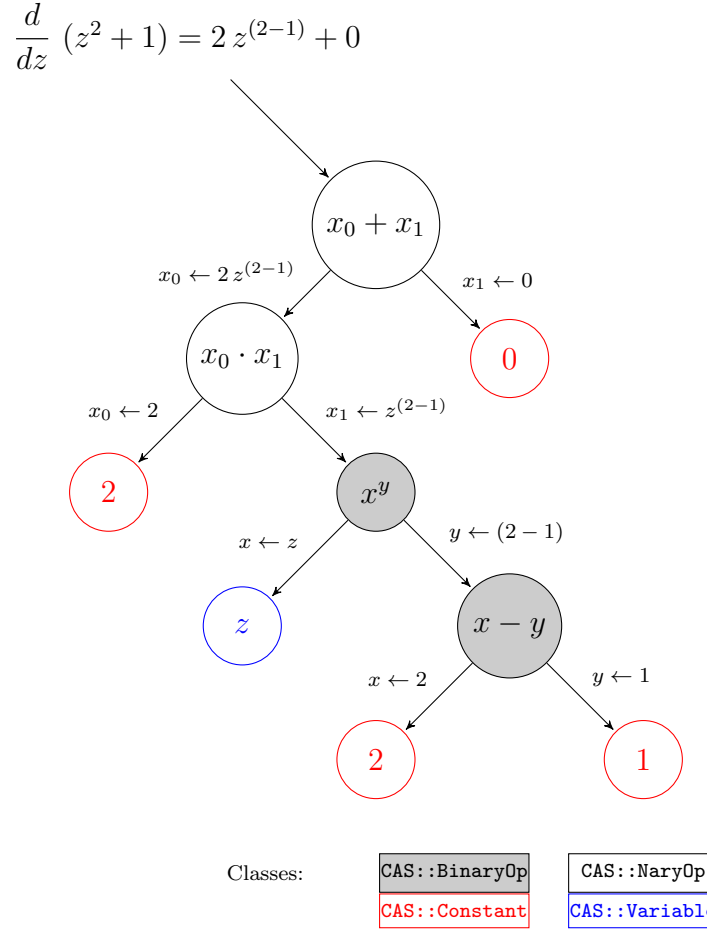


Figure 1: Example graph from the first function reported in listing 1

product — have as parent the `CAS::NaryOp`¹, that links to an arbitrary number of subgraph. Figure 2.1 contains an example of graph. The different kind of containers allows to introduce some properties like *associativity* and *commutativity*. Each container exposes the subgraphs as instance properties. Containers structure is shown in Figure 2.1.

Terminal leaf of the graph are the two classes `CAS::Constant` and `CAS::Variable`. The first is a node for a simple numerical value, while the latter

¹Please note that this container is still at experimental stage

72 represents an independent variable, that can be used to perform derivatives
 73 and evaluations. As for now, those nodes are only scalar numbers, with plans
 74 to define also the vectorial and matricial extensions in the next major release.

75 Automatic differentiation (`CAS::Op#diff`) crosses the graph until it reaches
 76 the ending node. The terminal node is the starting point for derivatives
 77 accumulation, the mathematical equivalent of the chain rule:

$$(f \circ g)' = (f' \circ g) g' \quad (2)$$

78 The recursiveness is used also for simplifications (`CAS::Op#simplify`), sub-
 79 stitutions (`CAS::Op#subs`) and evaluations (`CAS::Op#call`).

80 2.2. Software Functionalities

81 2.2.1. Basic Functionalities

82 The main functionality of the library is the **AD**, that can be performed
 83 with respect to an independent variable (`CAS::Variable`), even for implicit
 84 functions. The differentiation is done by a method of the `CAS::Op`, having a
 85 `CAS::Variable` as argument:

Listing 1: Differentiation example

```

86
87 x = CAS.vars 'x'           # creates a variable
88 f = x ** 2 + 1             # define a symbolic expression
89 f.diff(x)                  # derivative w.r.t. x
90 # => 2 * x ^ (2 - 1) + 0
91 g = CAS.declare :g, f      # creates implicit function
92 g.diff(x)                  # derivative w.r.t. x
93 # => (x ^ (2 - 1) * 2) * Dg[0](x ^ 2)
94

```

95 Resulting graph still contains a zero, since **simplifications** are not ex-
 96 ecuted automatically. Each node of the graph contains some rules for sim-
 97 plify itself. Simplification are called recursively inside the graph, exactly like

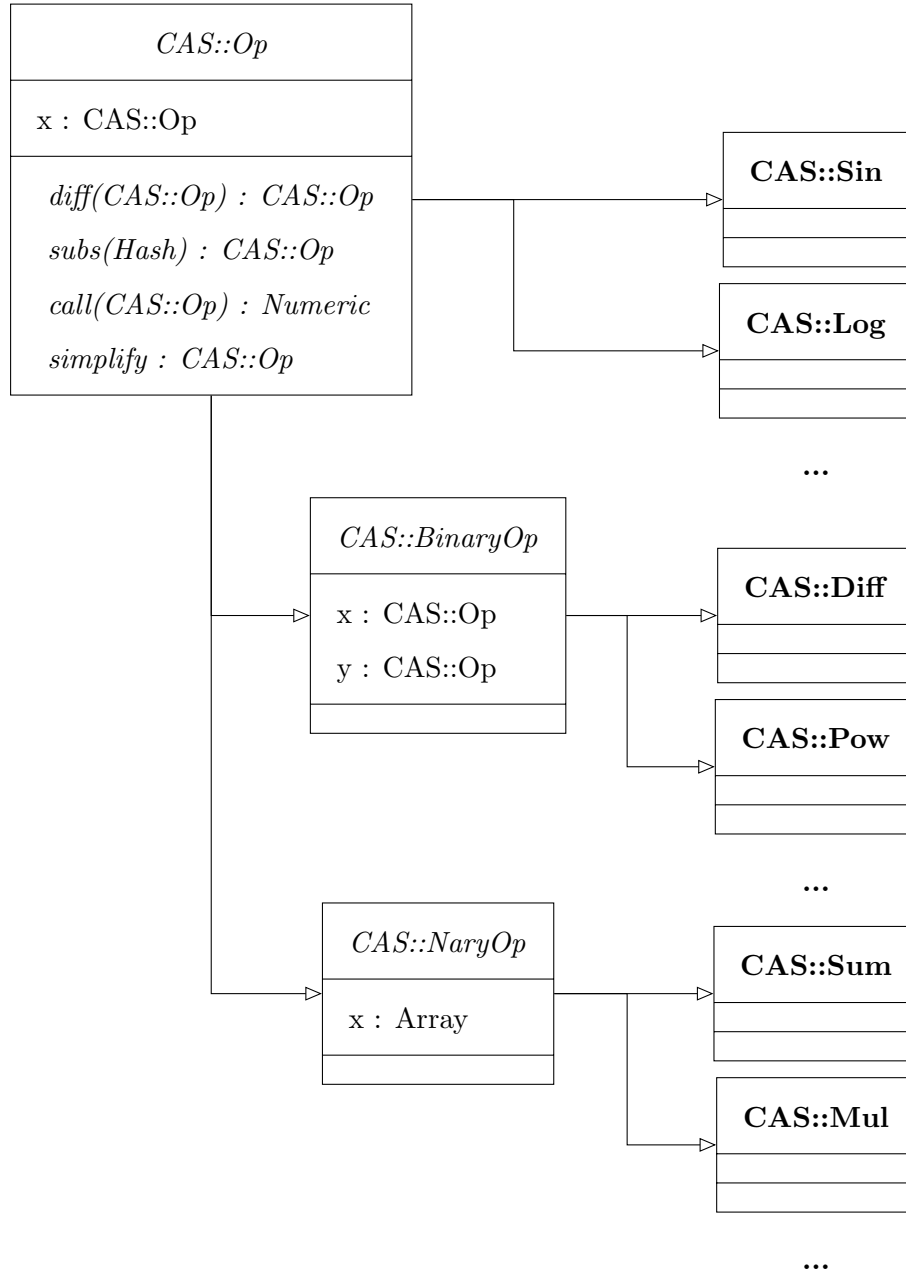


Figure 2: Simplified version of classes interface and inheritance

98 AD, bringing the strong limitation of not handling simplifications that come
 99 from *heuristic expansion* of sub-graphs — e.g. simplification through the use
 100 of trigonometric identities. Those simplification can be achieved manually

101 using **substitutions**.

Listing 2: Simplification example

```
102
103 x, y = CAS::vars 'x', 'y'          # creates two variables
104 f = CAS.log( CAS.sin( y ) )        # symbolic expression
105 f.subs y: CAS.asin(CAS.exp(x))    # perform substitution
106 f.simplify                         # simplify expression
107 # => x
108
```

109 The graph can be **evaluated** substituting defining some values for the
110 independent variable in a feed dictionary. The graph is recursively reduced
111 to a single numeric value:

Listing 3: Graph evaluation example

```
112
113 \label{code:example-call}
114 x = CAS::vars 'x'                  # creates a variable
115 f = x ** 2 + 1                     # define a symbolic expression
116 f.call x: 2                        # evaluate for x = 2
117 # => 5
118
```

119 Symbolic functions can be used to create expressions — e.g. $f(\cdot) \geq g(\cdot)$
120 — or piecewise functions — e.g. $\max(f(\cdot), g(\cdot))$:

Listing 4: Expressions and Piecewise functions

```
121
122 x = CAS::vars 'x'
123 f = x ** 2
124 g = 2 * x + 1
125 f.greater_equal g
126 # => ((x)^(3) >= ((2 * x) + 1))
127 CAS::max f, g
128 # => (((x)^(3) >= ((2 * x) + 1)) ? (x)^(3) : ((2 * x) + 1))
129
```

130 Expression are stored in a special container class, called **CAS::Condition**.

131 2.2.2. *Metaprogramming and Code-Generation*

132 The library is developed explicitly for **generation of code**, and in some
133 case also **metaprogramming**. Expressions, once manipulated, can be easily
134 exported as source code (in a defined language —i.e. the following example in
135 standard *Ruby* code) the function is used as a prototype for a callable *closure*
136 (`Proc` object):

Listing 5: Graph evaluation example

```
137  
138 x = CAS::vars 'x' # creates a variable  
139 f = CAS::log(CAS::sin(x)) # define a symbolic function  
140  
141 proc = f.as_proc # exports callable lambda  
142 proc.call 'x' => Math::PI/2  
143 # => 0.0  
144
```

145 Must be noted that making a closure of the graph is like making a snapshot,
146 and any further modifications to the graph will not update the callable object.
147 This drawback is balanced by faster execution time of the `Proc`: when a
148 graph needs only to be evaluated in a iterative algorithm, and not to be
149 manipulated, transforming it in a *closure* reduces the execution time per
150 loop.

151 Code generation should be flexible enough to export a graph in a user's
152 target language. Generation functions are usually included in specific plugins.
153 Users can expand exporting capabilities by writing specific exportation rules,
154 overriding method when required, or describing their own plugin:

Listing 6: Example of Ruby exportation plugin

```
155  
156 # Definition  
157 module CAS  
158 {  
159     # . . .
```



```

160     CAS::Variable => Proc.new { "#{name}" }
161     CAS::Sin      => Proc.new { "Math.sin(#{x.to_ruby})" },
162     # . . .
163   }.each do |cls, prc|
164     cls.send(:define_method, :to_ruby, &prc)
165   end
166 end
167
168 # Usage
169 x = CAS.vars 'x'
170 (CAS.sin(x)).to_ruby
171 # => Math.sin(x)
172

```

173 Included plugins may implement some advanced features such as code
174 optimization: this is an example with the *C* plugin:

Listing 7: Calling optimized-C exporter

```

175
176 require 'ragni-cas/c-opt'
177
178 x, y = CAS.vars :x, :y
179 f = x ** y + CAS.log(CAS.sin(x ** y))
180
181 CLib.create "example" do
182   implements_as "func", f
183 end
184

```

185 library created contains the following source (the header is omitted for brevity):

Listing 8: Calling optimized-C exporter

```

186
187 // Source file for library: example.c
188
189 #include "example.h"
190
191 double func(double x, double y) {

```

```

192     double __t_0 = pow(x, y);
193     double __t_1 = sin(__t_0);
194     double __t_2 = log(__t_1);
195     double __t_3 = (__t_0 + __t_2);
196
197     return __t_3;
198 }
199
200 // end of example.c
201

```

202 3. Illustrative Examples

203 As example, an implementation of a trapezoidal integration algorithm
204 will be provided, using the automatic differentiation as support library for
205 the function to be integrated. The result of the algorithm is the order of
206 convergence of the integration scheme.

207 Given a function $F(x)$, and its derivative $f(x)$, and the trapezoidal rule
208 for the integration between the interval $[a, b]$:

$$\hat{F}(a, b, n) = \frac{b-a}{n} \left(\frac{f(a) + f(b)}{2} + \sum_{k=1}^n f\left(a + k \frac{b-a}{n}\right) \right) \quad (3)$$

209 where n is the number of the point of the mesh. The error of the approxi-
210 mation is:

$$E[n] = F(b) - F(a) - \hat{F}(a, b, n) \quad (4)$$

211 This error shows a direct relation:

$$E[n] \propto C n^{-p} \quad (5)$$

212 where p is the convergence order. Using a different value for n , for example
213 $2n$:

$$\frac{E[n]}{E[2n]} = 2^p \quad \rightarrow \quad p = \log_2 \left(\frac{E[n]}{E[2n]} \right) \quad (6)$$

214 **4. Impact**

215 **5. Conclusions**

216 **Acknowledgements**

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 218 the public, commercial, or not-for-profit sectors.

219 **Current code version**

Nr.	Code metadata description	Please fill in this column
C1	Current code version	0.0.0
C2	Permanent link to code/repository used for this code version	github.com/MatteoRagni/cas-rb & rubygems.org/gems/ragni-cas
C3	Legal Code License	MIT
C4	Code versioning system used	<i>git</i> (GitHub)
C5	Software code languages, tools, and services used	<i>Ruby</i>
C6	Compilation requirements, operating environments	<i>Ruby</i> $\geq 2.x$, <i>pry</i> for testing console (optional)
C7	If available Link to developer documentation/manual	rubydoc.info/gems/ragni-cas
C8	Support email for questions	info@ragni.me

Table 1: Code metadata (mandatory)