

Capitolo 1

Implementazione di procedure di decisione per frammenti Binding in Vampire

L'algoritmo di decisione, la classificazione, Il preprocessing

1.1 Preprocessing

Preprocess	«typedef» BindingFormulaMap: DHMap<Literal*, Formula*>
+prb: Problem +fragment: Fragment - _bindingFormulas: BindingFormulaMap - _booleanToLiteral: BooleanToLiteralBindingMap - _literalToBoolean: LiteralToBooleanBindingMap - _bindingClauses: BindingClauseMap - _sat2Fo: SAT2FO - _clauses: SATClauseStack - _literals: LiteralList*	«typedef» BooleanToLiteralBindingMap: DHMap<Literal*, LiteralList*>
	«typedef» LiteralToBooleanBindingMap: DHMap<Literal*, Literal*>
+Preprocess(prb: Problem) +ennf() +topBooleanFormula() +naming() +nnf() +satClausify() - _newBooleanBinding(): Literal* - _newBindingLiteral(lit: Literal*): Literal* - _addBindingFormula(formula: Formula*): Formula* - _getSingleLiteralSatClause(literal: Literal*): SATClauseStack* - _topBooleanFormula(formula: Formula*): Formula* +isBooleanBinding(literal: Literal*): bool +isBindingLiteral(literal: Literal*): bool +getLiteralBindings(booleanBinding: Literal*): LiteralList* +getBooleanBinding(literalBinding: Literal*): Literal* +getSatClauses(literal: Literal*): SATClauseStack* +literals(): LiteralList* +satClauses(): SATClauseStack* +toSAT(literal: Literal*): SATLiteral +maxSatVar(): unsigned	«typedef» BindingClauseMap: DHMap<Literal*, SAT::SATClauseStack*>

Figura 1.1: Struttura del Preprocessing

1.1.1 Boolean Top Formula

1.1.2 Forall-And

1.1.3 SAT-Clausification

1.2 Procedura di Decisione

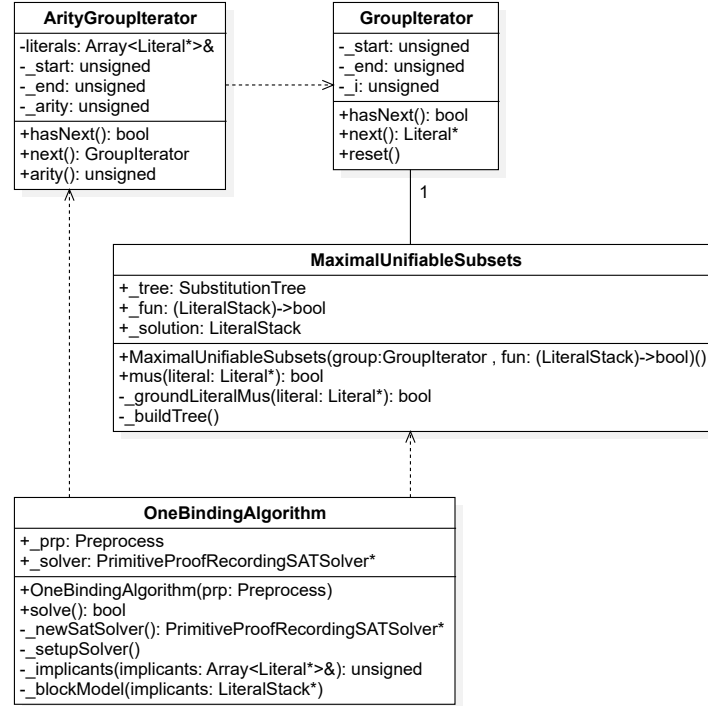


Figura 1.2: Struttura dell'algoritmo di decisione

1.2.1 Implicants Sorting

1.2.2 Maximal Unifiable Subsets

Algorithm 1: Maximal Unifiable Subsets

Firma: $\text{mus}(\text{literal})$

Input: literal un puntatore ad un letterale

Output: \top o \perp

GlobalData: S una mappa da letterali a bool

if $S[\text{literal}]$ **then**

 | **return** \top ;

end

if literal is ground **then**

 | **return** $\text{groundLiteralMus}(\text{literal})$;

end

$S[\text{literal}] = \top$;

$\text{res} := \text{mus}(\text{literal}, \emptyset)$;

$S[\text{literal}] = \perp$;

return res ;

Algorithm 2: Maximal Unifiable Subsets

Firma: $\text{mus}(\text{literal}, \text{FtoFree})$

Input: literal un puntatore ad un letterale, FtoFree un puntatore ad una lista di letterali

Output: \top o \perp

GlobalData: \mathbf{S} una mappa da letterali a interi, **fun** una funzione da lista di letterali a bool, **tree** un SubstitutionTree

$\text{isMax} := \top$;

$\text{uIt} = \text{tree.getUnifications}(\text{query} : \text{literal}, \text{retrieveSubstitutions} : \text{true})$;

$\text{toFree} := \emptyset$;

while $\text{uIt.hasNext}()$ **do**

$(u, \sigma) := \text{uIt.next}()$;

if $S[u] = 0$ **then**

$S[u] = 1$;

$l := \text{literal}^\sigma$;

if $l = \text{literal}$ **then**

$u' := u^\sigma$;

if $u' = u$ **then**

$\text{FtoFree} := \text{FtoFree} \cup \{u\}$;

end

else

$\text{toFree} := \text{toFree} \cup \{u\}$;

end

end

else

$\text{isMax} = \perp$;

if $\neg \text{mus}(l, \text{toFree})$ **then**

return \perp ;

end

$S[u] = -1$;

$\text{toFree} := \text{toFree} \cup \{u\}$;

end

end

end

if isMax **then**

if $\neg \text{fun}(\{x \mid S[x] = 1\})$ **then**

return \perp ;

end

end

while $\text{toFree} \neq \emptyset$ **do**

$S[\text{toFree.pop}()] = 0$;

end

return \top ;

Algorithm 3: Maximal Unifiable Subsets Ground

Firma: $\text{groundMus}(\text{literal})$

Input: *literal* un puntatore ad un letterale ground

Output: \top o \perp

GlobalData: **S** una mappa da letterali a interi, **fun** una funzione da lista di letterali a bool,
tree un SubstitutionTree

if $S[\text{literal}] \neq 0$ **then**

return \top ;

end

$uIt = \text{tree.getUnifications}(\text{query} : \text{literal}, \text{retrieveSubstitutions} : \text{true});$

$\text{solution} := \emptyset;$

while $uIt.hasNext()$ **do**

$(u, \sigma) := uIt.next();$

if $S[u] = 0$ **then**

if u is ground **then**

$S[u] = -1;$

end

$\text{solution} := \text{solution} \cup \{u\};$

end

end

return $\text{fun}(\text{solution});$

1.2.3 Algoritmo Finale

Algorithm 4: Algoritmo di decisione

Firma: solve(*prp*)

Input: *prp* il problema pre-processato

Output: \top o \perp

satSolver := newSatSolver();

satSolver.addClauses(*prp.clauses*);

while *satSolver.solve*() = SATISFIABLE **do**

res := \top ;

implicants := getImplicants(*satSolver*, *prp*);

implicants := sortImplicants(*implicants*);

if *implicants* contains only ground Literals **then**

return \top ;

end

agIt := ArityGroupIterator(*implicants*);

while *res* And *agIt.hasNext*() **do**

maximalUnifiableSubsets := SetupMus(*group*, *internalSat*);

foreach *lit* \in *group* **do**

if \neg *maximalUnifiableSubsets.mus*(*lit*) **then**

res := \perp ;

blockModel(*maximalUnifiableSubsets.getSolution*());

Break;

end

end

if *res* = \top **then**

return \top ;

end

end

end

return \perp ;

Algorithm 5: Sat interna

Firma: internalSat(*literals*)

Input: *literals* una lista di letterali

Output: \top o \perp

if *literals.length* = 1 And *getSatClauses*(*literals.top*()) *length* = 1 **then**

return \top ;

end

satSolver := newSatSolver();

foreach *l* \in *literals* **do**

satSolver.addClause(*getSatClauses*(*l*));

end

return *satSolver.solve*() = SATISFIABLE;

Algorithm 6: getImplicants

Firma: getImplicants(solver, prp)**Input:** *solver* un sat solver, *prp* il problema pre-processato**Output:** Una lista letterali*implicants* := \emptyset ;**foreach** $l \in prp.literals()$ **do** *satL* := *prp.toSat*(*l*); **if** *solver.trueInAssignment*(*satL*) **then** **if** *prp.isBooleanBinding*(*l*) **then** *implicants* := *implicants* \cup *prp.getLiteralBindings*(*l*); **end** **else** *implicants* := *implicants* \cup {*l*}; **end** **end****end****return** *implicants*;

1.3 Algoritmo di Classificazione

(Input formula rettificata senza true e false)

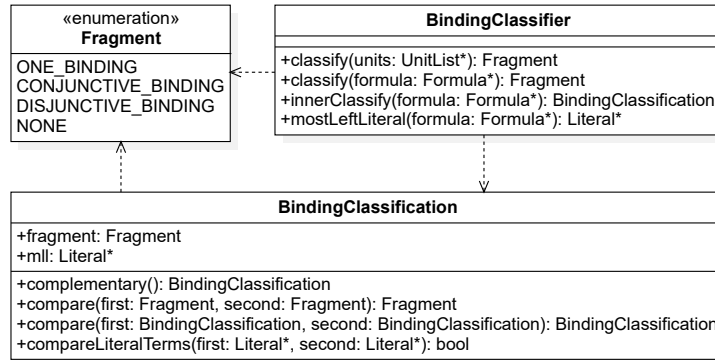


Figura 1.3: Classificatore

Algorithm 7: Classificatore esterno

Firma: $\text{classify}(\varphi)$ **Input:** φ Una formula rettificata

Output: Un elemento dell'enumerazione Fragment

```
switch  $\varphi$  do
  case Literal do
    | return ONE_BINDING;
  end
  case  $A[\wedge, \vee]B$  do
    | return  $\text{compare}(\text{classify}(A), \text{classify}(B))$ ;
  end
  case  $\neg A$  do
    | return  $\text{classify}(A).\text{complementary}()$ ;
  end
  case  $[\forall, \exists]A$  do
    |  $\text{sub} := \varphi$ ;
    |  $\text{connective} :=$  connective of  $\varphi$ ;
    | repeat
    |   |  $\text{sub} :=$  subformula of sub;
    |   |  $\text{connective} :=$  connective of sub;
    | until  $\text{connective} \notin \{\forall, \exists\}$ ;
    |  $(\text{fragment}, -) := \text{innerClassify}(\text{sub})$ ;
    | return  $\text{fragment}$ ;
  end
  case  $A \Leftrightarrow B$  do
    | return  $\text{compare}(\text{classify}(A \Rightarrow B), \text{classify}(B \Rightarrow A))$ ;
  end
  case  $A \oplus B$  do
    | return  $\text{classify}(A \Leftrightarrow B).\text{complementary}()$ ;
  end
  case  $A \Rightarrow B$  do
    | return  $\text{compare}(\text{classify}(\neg A), \text{classify}(B))$ ;
  end
end
```

Algorithm 8: Classificatore interno

Firma: $\text{innerClassify}(\varphi)$ **Input:** φ Una formula rettificata

Output: Una coppia (Fragment, Literal)

```
switch  $\varphi$  do
  case Literal  $l$  do
    | return (ONE_BINDING,  $l$ );
  end
  case  $A[\wedge, \vee]B$  do
    | return innerCompare(innerClassify( $A$ ), innerClassify( $B$ ), connective of  $\varphi$ );
  end
  case  $\neg A$  do
    | return innerClassify( $A$ ).complementary();
  end
  case  $A[\Rightarrow, \Leftrightarrow, \oplus]B$  do
    | return innerCompare(innerClassify( $A$ ), innerClassify( $B$ ), connective of  $\varphi$ );
  end
  else
    | return (None, null);
  end
end
```

Algorithm 9: Compare esterno

Firma: $\text{compare}(A, B)$ **Input:** A, B due elementi dell'enumerazione Fragment

Output: Un elemento dell'enumerazione Fragment

```
if  $A = B$  then
  | return  $A$ ;
end
if One_Binding  $\notin \{A, B\}$  then
  | return None;
end
return max( $A, B$ );
```

Algorithm 10: Compare interno

Firma: $\text{innerCompare}(A, B, \text{con})$ **Input:** A, B due coppie (Fragment, Literal), con un connettivo

Output: Una coppia (Fragment, Literal)

```
switch  $A.\text{first}, B.\text{first}, \text{con}$  do
  case  $\text{One\_Binding}, \text{One\_Binding}, \_$  do
    if  $A.\text{second}$  has same terms of  $B.\text{second}$  then
      | return  $A$ ;
    end
    else if  $\text{conn} = \wedge$  then
      | return  $(\text{Conjunctive\_Binding}, \text{null})$ ;
    end
    else if  $\text{conn} = \vee$  then
      | return  $(\text{Disjunctive\_Binding}, \text{null})$ ;
    end
  end
  case  $[\text{One\_Binding}, \text{Conjunctive\_Binding} \mid \text{Conjunctive\_Binding}, \text{One\_Binding}], \wedge$  do
    | return  $(\text{Conjunctive\_Binding}, \text{null})$ ;
  end
  case  $[\text{One\_Binding}, \text{Disjunctive\_Binding} \mid \text{Disjunctive\_Binding}, \text{One\_Binding}], \vee$  do
    | return  $(\text{Disjunctive\_Binding}, \text{null})$ ;
  end
  case  $\text{Conjunctive\_Binding}, \text{Conjunctive\_Binding}, \wedge$  do
    | return  $(\text{Conjunctive\_Binding}, \text{null})$ ;
  end
  case  $\text{Disjunctive\_Binding}, \text{Disjunctive\_Binding}, \vee$  do
    | return  $(\text{Disjunctive\_Binding}, \text{null})$ ;
  end
end
return  $(\text{None}, \text{null})$ ;
```
