

Optimizing City-to-City Transportation with Ford-Fulkerson and Dijkstra's Algorithm

Northeastern University
College of Engineering

EECE 2560: Fundamentals of Engineering Algorithms
Kennedy Scheimreif, Anna Valades, Matthew Geisel

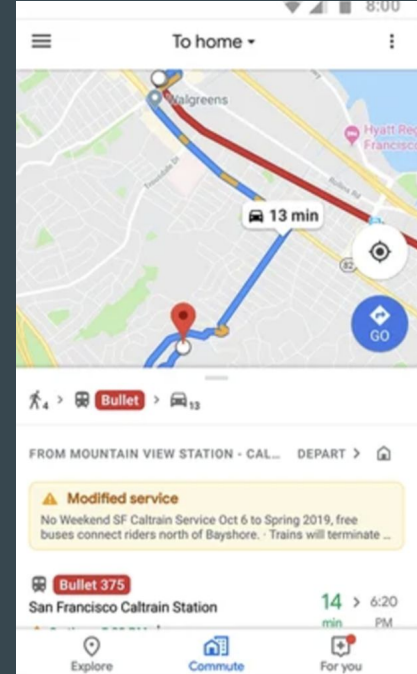


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Existing Technology

Travel Booking Industry worth \$520 billion

- Many mass-transit systems must be booked independently (i.e. Amtrak, the T)
- Multi-modal transportation systems (Google Maps)
- Most booking sites not optimized for multiple people
 - Booking groups of 10+ difficult & pricey



jetBlue

6:00 AM – 8:06 AM
JetBlue

3 hr 6 min
BOS–ORD

Nonstop

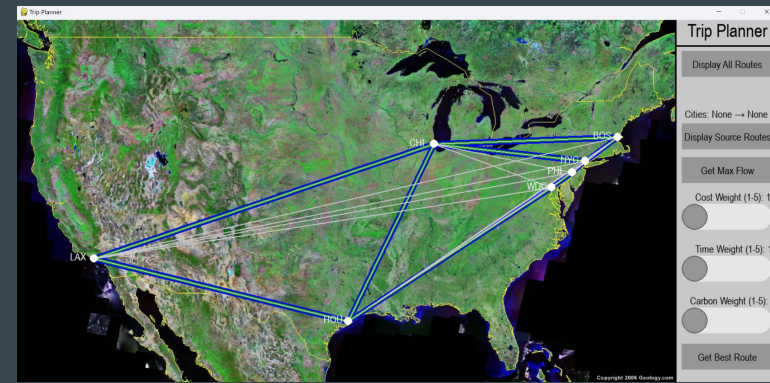
161 kg CO₂e
+27% emissions ⓘ

\$133
round trip



Our Solution

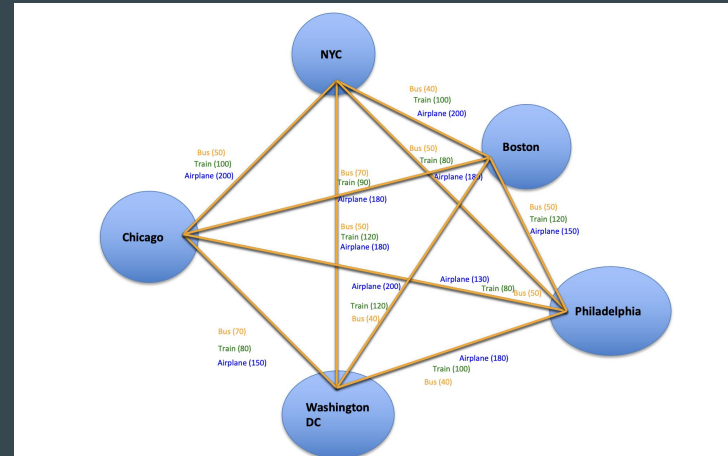
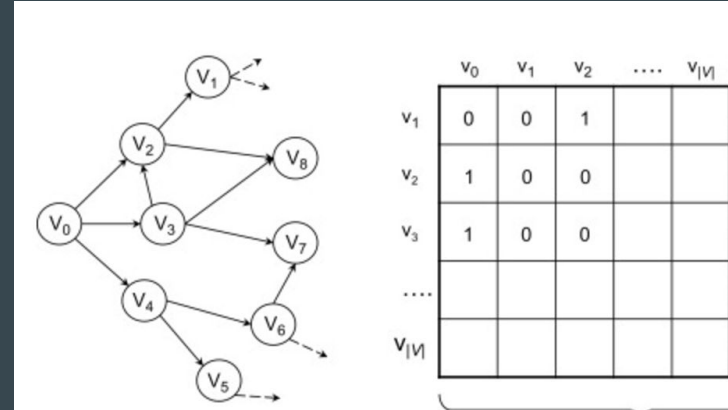
- **Ford-Fulkerson** and **Breadth-First Search** to find maximum flow
- **Dijkstra's Algorithm** to optimize people flow using different modes of transportation and user's traveling preferences
- Real-time data from APIs are expensive
 - **CSV** file as a stand-in
- Backend uses **C++** and the frontend uses **Python** and is integrated with a custom **socket communication protocol**



- Efficient
- Scalable
- User-centric

Methodology: Ford-Fulkerson Algorithm

- **Adjacency Matrix:** 2D matrix where the rows and columns are nodes in the graph.
- **Residual Graph:** a modified version of the original matrix that shows the residual capacity (capacity - current flow).



Pseudocode + Time Complexity Used Pt 1

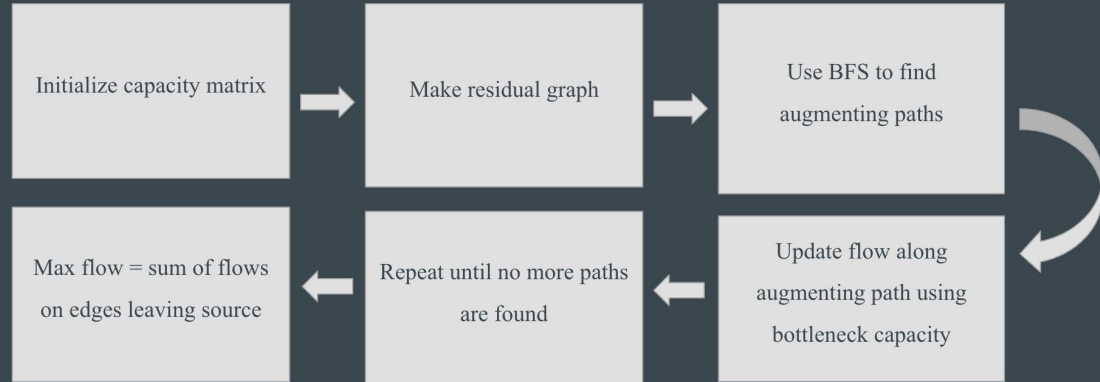
```
1 findAugmentingPath(source, sink, parent, residual):
2     initialize visited array to false
3     empty queue q
4     push source node into q
5     mark source as visited
6     set parent[source] = -1
7
8     while q not empty:
9         curr = q.pop()
10
11         for each node next from 0 to size of residual:
12             if next is not visited and residual[curr][next] > 0:
13                 set next as visited
14                 set parent[next] = curr
15                 if next == sink:
16                     return True
17             add next to the queue
18     return False
```

- **Augmenting Path:** a path from the source to the sink where there is a remaining capacity along all edges in the path (bottleneck determines amount of additional flow).
- **BFS:** uses a queue to find the augmenting paths by marking nodes as visited and tracking the path.
- **Time Complexity:** $O(n+m)$

Pseudocode + Time Complexity Used Pt 2

Ford-Fulkerson Algorithm

```
1 calculateMaxFlow(sourceCity, destinationCity):
2     if sourceCity or destinationCity is invalid:
3         print error message
4         return
5
6     initialize maxFlow = 0
7     convert sourceCity and destinationCity to their index values
8     copy capacity graph into residual graph
9     initialize parent array
10
11     while findAugmentingPath(source, sink, parent, residual):
12         set pathFlow = very large number
13         traverse the augmenting path from sink to source:
14             for each node v in the path:
15                 u = parent[v]
16                 pathFlow = MIN(pathFlow, residual[u][v])
17
18         traverse the augmenting path again to update residual
19         graph:
20             for each node v in the path:
21                 u = parent[v]
22                 residual[u][v] -= pathFlow
23                 residual[v][u] += pathFlow
24
25         increment maxFlow by pathFlow
26
27     print "Maximum flow from sourceCity to destinationCity is
28         maxFlow"
```



Time Complexity:

- $O(n^2 \cdot m)$

Pseudocode + Time Complexity Pt 3

Get City Connections:

```
Algorithm getCityConnections(city)
begin
  Print "Cities directly connected to city:"
  Initialize connectionFound as false
  for each route in routes do
    if route.sourceCity == city then
      Set connectionFound to true
      Print "Destination: route.destinationCity — Mode: route.mode — Travel Time:
        route.time hours — Cost: $route.cost"
    end
  end
  if connectionFound is false then
    Print "No direct connections found for city."
  end
end
```

Time Complexity:

- Iterate through the routes to find connections: $O(R)$ where R is number of routes
- Print results $O(1)$
- Total Time Complexity: $O(R)$

Get Busiest Routes:

```
Algorithm getBusiestRoutes()
begin
  if routes is empty then
    Print "No routes available in the network."
    return
  end
  Sort routes in descending order of capacity:
  Compare capacity of route a and route b
  if a.capacity > b.capacity then
    Return a.capacity > b.capacity
  end
  Print "Busiest Routes by Capacity:"
  Set limit to the smaller of 5 or routes.size()
  for i from 0 to limit-1 do
    Let route = routes[i]
    Print "route.sourceCity to route.destinationCity — Mode: route.mode — Capacity:
      route.capacity people — Travel Time: route.time hours"
  end
end
```

Time Complexity:

- Check if routes is empty: $O(1)$
- Sort routes by capacity using introsort which is a combination of quicksort and heap sort: $O(R \log R)$
- Print top 5 routes: $O(5)$
- Total Time Complexity: $O(R \log R)$

Pseudocode + Time Complexity Pt 4

Dijkstra's Algorithm

1. Create **vectors** of weights & routes (weights set to very large)
2. Go through each valid weight city with **2 nested for loops**
 - a. Go through every **route** that includes the city
 - i. Calculate weight of route (cost / scalar) * weight + curr_weight
 1. Time: 600 mins, Carbon: 0.5 tons, Cost: \$300
 - ii. If the other city's weight is **more**, set its route & weight
3. Start at destination city
 - a. Add its route to **return vector**
 - b. Go to the other city in route

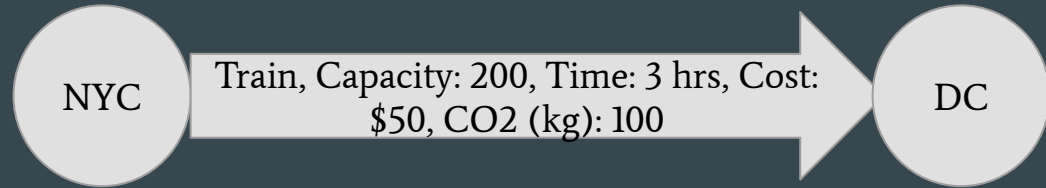
Input: src, dst(string), timeWeight, costWeight, impactWeight(int)

Returns: vector<Route>

Time Complexity: $O(n^2 * m)$ (n is # cities, m is # routes per city)

Data Structures Pt 1

- **Route Class:** represents individual connections between cities with attributes and is stored as a **vector** of **objects** with the following data types:
 - Source City: **string**
 - Destination City: **string**
 - Capacity: **int**
 - Time: **double**
 - Mode: **string**
 - Cost: **double**
 - Environmental Impact: **double**
- **CityNetwork Class:** manages the overall network of cities and routes.
- **Parent Array:** tracks the parent of each node during BFS for path construction

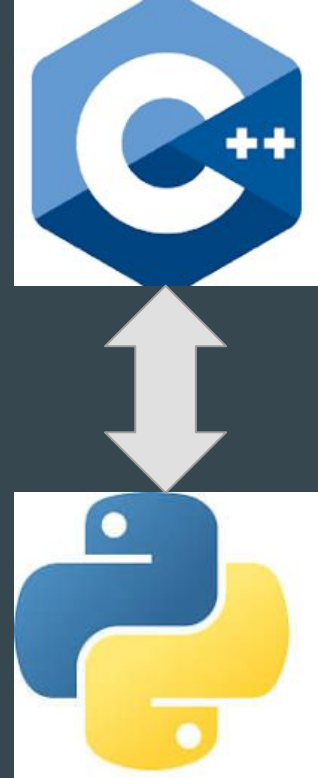


Data Structures Pt 2

- **Queue:** used in the BFS to find augmenting paths where there is a queue of integers corresponding to cities
- **Vectors:** used to store a dynamic list of data elements such as cities, routes, and graph edges.
- **FlowResult struct:** stores data that gives result of Ford-Fulkerson algorithm
 - Flow: **int**
 - Total max flow
 - Connections: **vector<vector<string>>**
 - Stores vectors of city names
 - Capacities: **vector<int>**
 - Stores the number of people sent through each connection

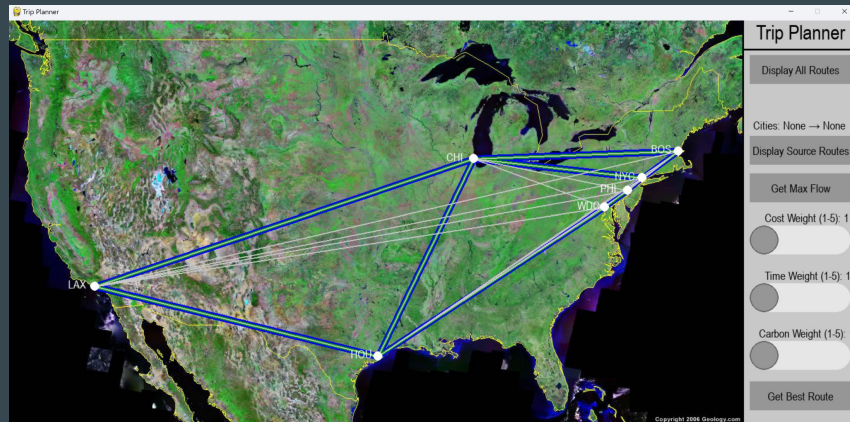
Backend - Frontend Integration

- **Socket:** Way that two processes can communicate (port 8001 used)
 - **C++: OS-Specific** (only Windows driver developed)
 - **Python: OS-Independent** (advantage of Python)
 - **Frontend-Independent** (easily integratable)
- Developed **custom communication protocol** of strings
 - Frontend: *Command Parameters*
 - “Display *src*” (set to All for all)
 - “Flow *src dst*”
 - “Route *src dst cost time impact*” (last 3 are 1-5)
 - Backend: Response
 - Display/Route: “*src,dst,mode;...* ”
 - Flow: “*capacity:city,city,...; ...* ”

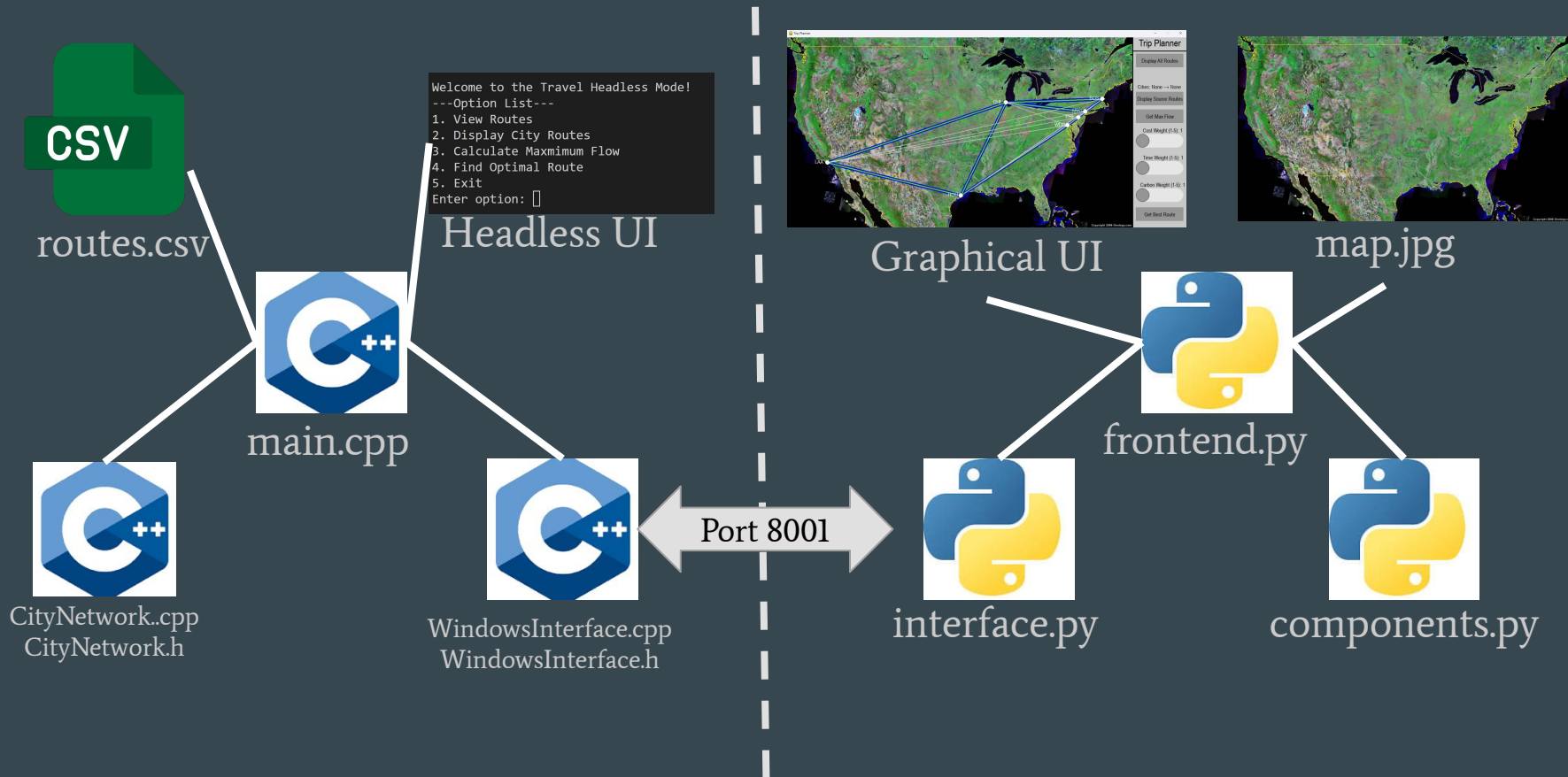


Frontend Implementation

- Uses **PyGame** library (common for GUIs)
 - Loops every ~ 0.01 seconds
 - Displays **airplane**, **bus**, **train**
 - Except in max flow
- On press: send, recv, process, display
 - **Vector** (Python list) stores each route [src, dst, mode] & selected cities
 - For **display & optimize**: just processes message ($O(n)$)
 - **Flow**: has to sum weights of sub-parts ($O(n^2 * m)$)
 - For each route, check if it already exists. If so, add the capacity.
- **Mercator Projection** used: easy to translate coordinates to position
 - $[x, y] = [21.433 * \text{long} + 2686.177, -28.185 * \text{lat} + 1420.191]$

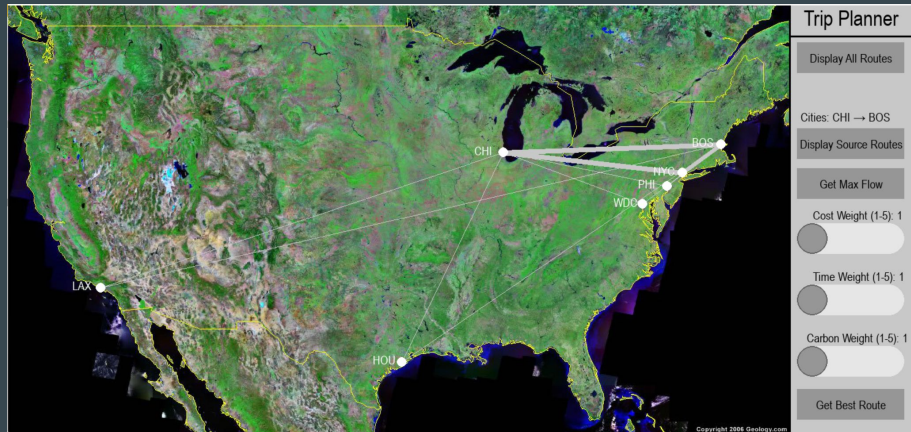


Overall Code Structure



Functionality Examples

BOS -> CHI Max Flow



```
Enter the source city: BOS
Enter the destination city: CHI
Maximum flow: 200
Paths used:
80 sent via CHI -> BOS
80 sent via CHI -> NYC -> BOS
10 sent via CHI -> WDC -> BOS
10 sent via CHI -> HOU -> BOS
10 sent via CHI -> LAX -> BOS
10 sent via CHI -> PHI -> BOS
```

BOS -> WDC

Impact-Optimized Route



```
Enter the source city: BOS
Enter the destination city: WDC
Enter the time weight: 1
Enter the cost weight: 1
Enter the environmental weight: 5
```

Result:

```
PHI -> WDC (Train): cost- 50, time- 120, tons of carbon- 0
NYC -> PHI (Train): cost- 25, time- 60, tons of carbon- 0
BOS -> NYC (Train): cost- 30, time- 240, tons of carbon- 0
```

Live Demonstration

Analysis and Results: Time Spent on Project Per Person

- **Backend:** 1.5 hr/week
- **Frontend:** 1 hr/week
- **Testing and Optimization:**
1 hr/week
- **Reports and Presentation:**
.5 hr/week
- **Total:** ~32 hours over the
past 2 months per person

Date	Week	Coding Progress	Technical Report Progress	PowerPoint Progress	Team member	
2024-10-03 00:00:00	Week 1	Not Started	Started	Not Started	Kennedy	
2024-10-04 00:00:00	Week 1	Not Started	Started	Not Started	Anna	
2024-10-05 00:00:00	Week 1	Not Started	Started	Not Started	Matt	
2024-10-06 00:00:00	Week 1	Not Started	Started	Not Started	Kennedy	
2024-10-07 00:00:00	Week 1	Not Started	Started	Not Started	Anna	
2024-10-08 00:00:00	Week 1	Not Started	Started	Not Started	Matt	
2024-10-09 00:00:00	Week 1	Not Started	Started	Not Started		
2024-10-10 00:00:00	Week 2	Not Started	Started	Not Started	Kennedy	
2024-10-11 00:00:00	Week 2	Not Started	Started	Not Started	Anna	
2024-10-12 00:00:00	Week 2	Not Started	Started	Not Started	Matt	
2024-10-13 00:00:00	Week 2	Not Started	Started	Not Started	Kennedy	
2024-10-14 00:00:00	Week 2	Not Started	Started	Not Started	Anna	
2024-10-15 00:00:00	Week 2	Not Started	Started	Not Started	Matt	
2024-10-16 00:00:00	Week 2	Not Started	Not Started	Not Started		
2024-10-17 00:00:00	Week 3	Started	Started	Not Started	Kennedy	
2024-10-18 00:00:00	Week 3	Started	Started	Not Started	Anna	
2024-10-19 00:00:00	Week 3	Started	Started	Not Started	Matt	
2024-10-20 00:00:00	Week 3	Started	Started	Not Started	Kennedy	
2024-10-21 00:00:00	Week 3	Started	Started	Not Started	Anna	

Analysis and Results: Key Findings and Interpretation of Results

- **Customer Use**
 - The system could be integrated in an app, helping users plan affordable, time-efficient, and environmentally sustainable trips
- **Industry Use**
 - City planners can use this tool to simulate and optimize transit networks, identify bottlenecks, and allocate resources more effectively
- **Combines Theory and Practical Applicability:**
 - Maps user preferences to real-world scenarios
 - Adaptable for both small-scale and large-scale use cases

Discussion

Implications of Findings:

- Time-effective process for determining the most optimal trip
- Scalable algorithms able to be used for more routes

Project Limitations:

- Number of cities
- Fixed capacities, travel times, and cost
- Doesn't include private modes of transportation (cars)

Conclusion

- **Recommendations for Future Work:**
- Real time data integration with API
 - Consider price fluctuations over weekends and holidays
- Add more cities
- Consider how business class and other travel options affect price

References

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