CSE 4600 Homework 5: Semaphores

Highlights of this assignment:

- 1. Definition of Semaphore
- 2. Why Semaphores?
- 3. Functions for Posix Semaphores
 - sem init
 - sem wait
 - sem post
 - sem getvalue
 - sem destroy
- 4. Homework Assignment

1. Definition of Semaphore

The original meaning of a semaphore was:

"A visual signaling apparatus with flags, lights, or mechanically moving arms, as one used on a railroad." (from http://dictionary.reference.com)

What do railroad tracks and computers have in common? Well, railroad semaphores were the inspiration of E. W. Dijkstra solution to the mutual exclusion problem in operating systems.

"...consider a stretch of railroad in which there is a single track over which only one train at a time is allowed.

Guarding this track is a semaphore. A train must wait before entering the single track until the semaphore is in a state that permits travel. When the train enters the track, the semaphore changes state to prevent other trains from entering the track. A train that is leaving this section of track must again change the state of the semaphore to allow another train to enter."

(from Oracle Documentation: Programming with Synchronization Objects.)

Railroad semaphores serve as a signal to indicate whether or not a single track has a train on it. Similarly, semaphores in operating systems serve as a signal to indicate whether or not a shared resource is currently being used. The following is a definition of the operating systems semaphores (from The Free On-line Dictionary of Computing, http://www.foldoc.org/, Editor Denis Howe):

"[Semaphores are] the classic method for restricting access to shared resources (e.g. storage) in a multiprocessing environment. They were invented by Dijkstra and first used in T.H.E operating system.

A semaphore is a protected variable (or abstract data type) which can only be accessed using the following operations:

```
s = s + 1;
}
Init(s, v)
Semaphore s;
int v;
{
    s = v;
}
```

P and V stand for Dutch "Proberen", to test, and "Verhogen", to increment. The value of a semaphore is the number of units of the resource which are free (if there is only one resource a "binary semaphore" with values 0 or 1 is used). The P operation busy-waits (or maybe sleeps) until a resource is available whereupon it immediately claims one. V is the inverse, it simply makes a resource available again after the process has finished using it. Init is only used to initialise the semaphore before any requests are made. The P and V operations must be indivisible, i.e. no other process can access the semaphore during the their execution.

To avoid busy-waiting, a semaphore may have an associated queue of processes (usually a FIFO). If a process does a P on a semaphore which is zero the process is added to the semaphore's queue. When another process increments the semaphore by doing a V and there are tasks on the queue, one is taken off and resumed."

Note: These operations must be performed as an atomic hardware instruction to prevent two processes from gaining control of the same semaphore. The ideas above are concepts, but only safe when performed through one hardware instruction.

2. Why Semaphores

Read the following code badent.c,

```
#include <pthread.h>
#include <semaphore.h>
#include <stdio.h>
#include <stdlib.h>
#define NITER 1000000
int cnt = 0;
void * Count(void * a)
    int i, tmp;
    for (i = 0; i < NITER; i++)
        cnt++;
int main(int argc, char * argv[])
    pthread t tid1, tid2;
    if (pthread create (&tid1, NULL, Count, NULL))
      printf("\n ERROR creating thread 1");
      exit(1);
    if(pthread create(&tid2, NULL, Count, NULL))
      printf("\n ERROR creating thread 2");
      exit(1);
```

Compile it using

```
gcc badcnt.c -o xbadcnt -lpthread
```

Run the executable file and observe the ouput.

Quite unexpected! Since cnt starts at 0, and both threads increment it NITER times, we should see cnt equal to 2*NITER at the end of the program. What happens?

Threads can greatly simplify writing elegant and efficient programs. However, there are problems when multiple threads share a common address space, like the variable ont in our earlier example.

To understand what might happen, let us analyze this simple piece of code:

```
THREAD 1 THREAD 2 b = data; b++; data = a; data = b;
```

Now if this code is executed serially (for instance, THREAD 1 first and then THREAD 2), there are no problems. However threads execute in an arbitrary order, so consider the following situation:

Thread 1	Thread 2	data
a = data;		0
a = a + 1;		0
	b = data; // 0	0
	b = b + 1;	0
data = a; // 1		1
	data = b; // 1	1

So data could end up 1 or 2, and there is **NO WAY** to know which value! It is completely non-deterministic!

The solution to this is to provide functions that will block a thread if another thread is accessing data that it is using.

Pthreads may use semaphores to achieve this.

3. Functions for Posix Semaphores

All POSIX semaphore functions and types are prototyped or defined in semaphore.h. To define a semaphore object, use

```
#include <semaphore.h>
sem_t sem_name;
```

sem init()

sem init initializes a semaphore.

```
int sem_init(sem_t *sem, int pshared, unsigned int value);
```

Notes:

- sem points to a semaphore object to initialize
- pshared is a flag indicating whether or not the semaphore should be shared with fork()ed processes. Linux Threads does not currently support shared semaphores
- value is an initial value to set the semaphore to

Example of use:

```
sem_init(&sem_name, 0, 10);
```

sem_wait()

To wait on a semaphore, use sem wait:

```
int sem_wait(sem_t *sem);
```

• If the value of the semaphore is negative, the calling process blocks; one of the blocked processes wakes up when another process calls <code>sem_post</code>.

Example of use:

```
sem_wait(&sem_name);
```

sem_post()

To increment the value of a semaphore, use sem_post:

```
int sem_post(sem_t *sem);
```

• It increments the value of the semaphore and wakes up a blocked process waiting on the semaphore, if any.

Example of use:

```
sem_post(&sem_name);
```

sem_getvalue()

To find out the value of a semaphore, use sem getvalue:

```
int sem_getvalue(sem_t *sem, int *valp);
```

• gets the current value of sem and places it in the location pointed to by valp.

Example of use:

```
int value;
sem_getvalue(&sem_name, &value);
printf("The value of the semaphors is %d\n", value);
```

sem destroy()

To destroy a semaphore, use, use sem_destroy:

```
int sem_destroy(sem_t *sem);
```

• destroys the semaphore; no threads should be waiting on the semaphore if its destruction is to succeed.

Example of use:

```
sem_destroy(&sem_name);
```

Using semaphores - a short example

Consider the problem we had before and now let us use semaphores:

```
Declare the semaphore global (outside of any funcion):
sem_t mutex;
Initialize the semaphore in the main function:
sem_init(&mutex, 0, 1);
```

Thread 1	Thread 2	data
sem wait (&mutex);		0
	sem wait (&mutex);	0
a = data;	/* blocked */	0
a = a + 1;	/* blocked */	0
data = a;	/* blocked */	1
sem post (&mutex);	/* blocked */	1
/* blocked */	b = data;	1
/* blocked */	b = b + 1;	1
/* blocked */	data = b;	2
/* blocked */	sem post (&mutex);	2
[data is fine. The data race is gone.]		

4. Homework Assignment

Activity 1

Use the example above as a guide to fix the program badent.c, so that the program always produces the expected output (the value 2*NITER). Make a copy of badent.c into goodent.c before you modify the code. To compile a program that uses pthreads and posix semaphores, use

```
gcc -o goodcnt goodcnt.c -lpthread
```

Activity 2

Given the following code,

```
hw5.c
gcc hw5.c -o hw5 -lpthread
#include <unistd.h>
#include <stdlib.h>
#include <stdio.h>
#include <pthread.h>
#include <semaphore.h>
#define TRUE 1
#define BUFFER SIZE 5
#define LOOPS TO
/* the semaphores */
sem t full, empty, mutex;
/* the buffer */
int buffer[BUFFER SIZE];
/* buffer counter */
int in, out;
void *producer(void *param); /* the producer thread */
void *consumer(void *param); /* the consumer thread */
void initializeData() {
   /* Create the mutex lock */
   sem init(&mutex, 0, 1);
   /st Create the full semaphore and initialize to 0 st/
   sem init(&full, 0, 0);
   /* Create the empty semaphore and initialize to BUFFER_SIZE */
sem_init(&empty, 0, BUFFER_SIZE);
   /* Get the default attributes */
   pthread attr init(&attr);
   /* init buffer */
   in = 0;
   out = 0;
/* Producer Thread */
void *producer(void *param) {
   int item;
   int i = 0;
   while(i < LOOPS)
      /* generate a random number */
      item = rand()%1000;
      //ADD YOUR CODE
```

```
/* Consumer Thread */
void *consumer(void *param) {
  int item;
   int i = 0;
   while(i < LOOPS)
   // ADD YOUR CODE
}
int main(int argc, char *argv[]) {
   /* Loop counter */
   int i;
   /* Verify the correct number of arguments were passed in */
   if(argc != 3)
      fprintf(stderr, "USAGE:./pc \n");
   int numProd = atoi(argv[1]); /* Number of producer threads */
   int numCons = atoi(argv[2]); /* Number of consumer threads */
   /* Initialize the app */
   initializeData();
   /* Create the producer threads */
   for(i = 0; i < numProd; i++) {
     /* Create the thread */
      pthread create (&tid, &attr, producer, NULL);
   /* Create the consumer threads */
   for(i = 0; i < numCons; i++) {
     /* Create the thread */
      pthread create (&tid, &attr, consumer, NULL);
   /* Sleep for the specified amount of time */
   sleep(1);
   /* Exit the program */
   printf("Exit the program\n");
   exit(0);
```

You will be adding code in the producer and consumer functions where the comments indicate to "ADD YOUR CODE".

If you do not know how to start the producer function, you might want to think about the answers to these questions:

- If the "empty" semaphore is 0, what will happen when you perform a sem_wait on the "empty" semaphore? How does the "empty" semaphore change to a positive integer?
- What is the producer doing with the buffer array? Storing values or accessing values?
- What index will you be working with in the producer: in or out?
- When the producer is done, which semaphore will need to have one added to it (ie. "full" or "empty")?

If you do not know how to start the consumer function, you might want to think about the answers to these questions:

- If the "full" semaphore is 0, what will happen when you perform a sem_wait on the "full" semaphore? How does the "full" semaphore change to a positive integer?
- What is the consumer doing with the buffer array? Storing values or accessing values?
- What index will you be working with in the consumer: *in* or *out*?
- When the consumer is done, which semaphore will need to have one added to it (ie. "full" or "empty")?

Note: these questions do not need to be handed in.

A sample run is the following (yours might be slightly different):

```
Producer [140078164084480] produced 383
Producer [140078164084480] produced 915
Consumer [140078138906368] consumed 383
Producer [140078155691776] produced 886
Producer [140078155691776] produced 335
  Consumer [140078130513664] consumed 915
Producer [140078164084480] produced 793
   Consumer [140078138906368] consumed 886
Producer [140078155691776] produced 386
Producer [140078147299072] produced 777
Consumer [140078130513664] consumed 335
   Consumer [140078130513664] consumed 793
Producer [140078164084480] produced 492
Consumer [140078138906368] consumed 386
   Consumer [140078122120960] consumed 777
Producer [140078147299072] produced 421
Producer [140078147299072] produced 27
Consumer [140078122120960] consumed 492
   Consumer [140078138906368] consumed 421
Producer [140078155691776] produced 649
Producer [140078155691776] produced 59
   Consumer [140078138906368] consumed 27
Consumer [140078138906368] consumed 649
Producer [140078155691776] produced 763
Producer [140078164084480] produced 362
  Consumer [140078130513664] consumed 59
Consumer [140078130513664] consumed 763
Producer [140078147299072] produced 690
Producer [140078147299072] produced 426
Consumer [140078138906368] consumed 362
Consumer [140078122120960] consumed 690
   Consumer [140078130513664] consumed 426
Producer [140078164084480] produced 540
Producer [140078155691776] produced 926
Producer [140078155691776] produced 211
Producer [140078147299072] produced 172
Consumer [140078122120960] consumed 540
   Consumer [140078138906368] consumed 926
   Consumer [140078130513664] consumed 211
Producer [140078147299072] produced 567
Producer [140078155691776] produced 368
Consumer [140078122120960] consumed 172
   Consumer [140078122120960] consumed 567
Producer [140078147299072] produced 782
   Consumer [140078138906368] consumed 368
```

Deliverables:

Submit 3 files to Canvas:

- 1. goodcnt.c with the semaphore version of a mutex
- 2. producer-consumer code hw5.c with the 3 semaphores
- 3. A script log file *hw5log.txt* showing:
 - compiling badcnt.c
 - running of badent
 - compiling *goodcnt.c*
 - running of goodent
 - compiling producer-consumer hw5.c code
 - running of producer-consumer code

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