COMP318 Ontologies and Semantic Web

RDF - Part 14



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Where were we

RDF entailment rules

RDFS entailment rules

Bijection between literals and surrogate bonds

- (gl inverse of lg) If G contains the triple
 - :u :a :n. we can infer
 - :u :a :1. However:
 - this rule can be applied only when _:n identifies a bnode that was introduced earlier by weakening the literal:1 via the rule |g|
 - This inference rule is necessary to bring back a literal that has been substituted by a blank node using rule lg, then some other inference rule produced a triple with this blank node (_:n) in the object position, and rule gl can now be used to bring this literal back!
 - E.g.:Murray atp:name "Andy Murray".

 :atp:name rdfs:range atp:PlayerName.
 - Would entail "Andy Murray" a atp:PlayerName. which is problematic. Why?

Bijection between literals and surrogate bonds

- The latter triple is not a valid RDF triple
 - a literal should not appear in the subject position!
- thus it will not be inferred (the domain of the ?v variable in the rdfs3 rule would prevent the inference). And so, to achieve the valid inference:
 - :AndyMurray a atp:PlayerName .

Requires the surrogate blank node (:AndyMurray) to be used through the rule lg.

The inverse rule gl then allows surrogates to "travel" back as literals into the object position, though
examples of such behaviour are not necessarily intuitive.

Reasoning engines

- Systems that perform inference are often called reasoning engines or reasoners.
 - **Reasoner engine**: a system that infers new information based on the contents of a knowledge base.
 - This can be accomplished using rules and a rule engine, triggers on a database or RDF store, decision trees, tableau algorithms, or even programmatically using hard-coded business logic
 - A reasoner must be compliant to the semantics of the ontology language it supports
 - Hence, an ontology language must state its semantics in a formal way
- The RDFS reasoner uses entailment rules that are supposed to capture the intended semantics

Soundness and Completeness

- Theorem. A graph G1 RDFS-entails a graph G2 if there is a graph G1' which has been derived from G1 via the rules lg, gl, rdfax, rdf1, rdf2, rdfsax and rdfs1 ... rdfs13 such that:
 - G1' simply entails G2 or
 - G1' contains an XML clash.
- The inference rules for RDFS-entailment we presented previously are sound but not complete (ter Horst, 2005).

Example

The following graph:

```
ex:isHappilyMarriedTo rdfs:subPropertyOf _:bnode.
_:bnode rdfs:domain ex:Person.
ex:john ex:isHappilyMarriedTo ex:mary .
```

• The triple ex:john rdf:type ex:Person . is a semantic consequence of the graph above, but this cannot be derived from the inference rules

Decidability and complexity

- RDFS entailment is decidable, even though one has to deal with the infinite number of axiomatic triples:
 - due to the fact that the RDF vocabulary for encoding lists includes property names rdf:_i for all i ≥ 1, with several RDFS axiomatic triples for each rdf: i
- The problem of deciding whether a graph G1 RDFSentails another graph G2 is NP-complete. The problem becomes polynomial if G2 contains no blank nodes

RDFS entailment in state of the art systems

 Existing RDF stores (Jena, Sesame, Virtuoso, Oracle, etc) offer implementations of RDFS entailment together with ways of querying the stored graphs through SPARQL

 Implementations may be based on applying the rules in a backward chaining or a forward chaining fashion

RDFS entailment cheatsheet

RDFS entailment patterns.

	If S contains:	then S RDFS entails recognizing D:
rdfs1	any IRI aaa in D	aaa rdf:type rdfs:Datatype .
rdfs2	aaa rdfs:domain XXX . yyy aaa zzz .	yyy rdf:type XXX.
rdfs3	aaa rdfs:range XXX . yyy aaa zzz .	ZZZ rdf:type XXX.
rdfs4a	xxx aaa yyy .	XXX rdf:type rdfs:Resource .
rdfs4b	xxx aaa yyy.	<pre>yyy rdf:type rdfs:Resource .</pre>
rdfs5	<pre>XXX rdfs:subPropertyOf yyy . yyy rdfs:subPropertyOf ZZZ .</pre>	XXX rdfs:subPropertyOf ZZZ.
rdfs6	XXX rdf:type rdf:Property .	XXX rdfs:subPropertyOf XXX.
rdfs7	aaa rdfs:subPropertyOf bbb . xxx aaa yyy .	xxx bbb yyy .
rdfs8	XXX rdf:type rdfs:Class .	XXX rdfs:subClassOf rdfs:Resource .
rdfs9	XXX rdfs:subClassOf yyy . ZZZ rdf:type XXX .	ZZZ rdf:type yyy .
rdfs10	XXX rdf:type rdfs:Class .	XXX rdfs:subClassOf XXX .
rdfs11	XXX rdfs:subClassOf yyy . yyy rdfs:subClassOf ZZZ .	XXX rdfs:subClassOf ZZZ.
rdfs12	XXX rdf:type rdfs:ContainerMembershipProperty .	XXX rdfs:subPropertyOf rdfs:member .
rdfs13	XXX rdf:type rdfs:Datatype .	XXX rdfs:subClassOf rdfs:Literal .

RDFS entailment

```
 Given the graph G below,
  (d:Poe, o:wrote, d:TheGoldBug .)
  (d:TheGoldBug, rdf:type, o:Novel .)
  (d:Baudelaire, o:translated, d:TheGoldBug .)

  (d:Poe, o:wrote, d:TheRaven .)
  (d:TheRaven,rdf:type,o:Poem .)
  (d:Mallarme´, o:translated, d:TheRaven .)

  (d:Mallarme´,o:wrote,_:b .)
  (_:b, rdf:type, o:Poem .)

  <o:Poem rdfs:subClassOf ex:Literature .>
  <o:Novel rdfs:subClassOf ex:Literature .>
```

• And the following graph S, determine if G entails (using simple and RDFS entailment) S, and explain why.

```
S= <d:Poe wrote _:c .> <_:c rdf:type ex:Literature .>
```

Comp 318 11

RDFS entailment

```
1. (d:Poe, o:wrote, d:TheGoldBug .)
2. (d:TheGoldBug, rdf:type, o:Novel .)
3. (d:Baudelaire, o:translated,
d:TheGoldBug .)
4. (d:Poe, o:wrote, d:TheRaven .)
5. (d:TheRaven, rdf:type, o:Poem .)
6. (d:Mallarme', o:translated,
d:TheRaven .)
7. (d:Mallarme',o:wrote,_:b .)
(_:b, rdf:type, o:Poem .)
8. <o:Poem rdfs:subClassOf ex:Literature .>
9. <o:Novel rdfs:subClassOf
ex:Literature .>
S= <d:Poe wrote :c .>
< :c rdf:type ex:Literature .>
```

```
From RDFS 9 applied to 9 and 2

10. (d:TheGoldBug, rdf:type,
o:Literature .)

Apply SE1 to 1 and we obtain
<d:Poe o:wrote _:c>
with -:c -> d;TheGoldBug

Apply SE2 to 10, and we obtain
11.(_:c, rdf:type, o:Literature .)
```

So, the graph S is RDFS entailed

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