Distributed Systems COMP 212

Lecture 3-4

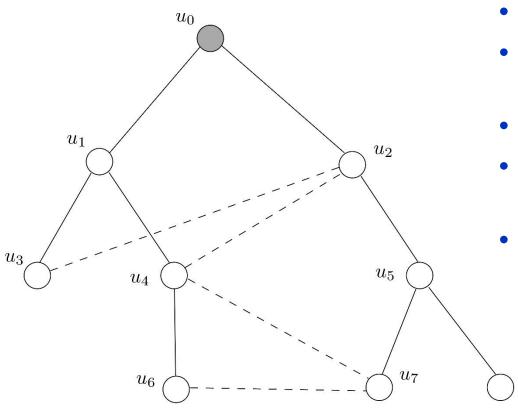
Othon Michail



Flooding/Broadcast

Broadcast given Spanning Tree

 We start from the case in which a spanning tree of the network is given



- Network G = (V, E)
- $E' \subseteq E$ specifies a spanning tree T = (V, E')
- Root: u_0 (leader)
- Processors know T in a distributed way
- Each u_i knows:
 - a parent_i

 u_8

a set children_i

Broadcast given Spanning Tree

Problem:

- u₀ has some information it wishes to send to all processors
 - e.g., a message (M)
 - additionally all nodes must have terminated in the end

Solution: Informal Description

- Root u_0 sends $\langle M \rangle$ to all channels leading to its children and terminates
- When a u_i receives $\langle M \rangle$ through the channel from its parent
 - it sends (M) to all channels leading to its children
 and
 - terminates

An Alternative Round

A round:

- 1. all nodes read incoming messages
- 2. all nodes update their state
- 3. all nodes generate new messages and put them in transit
- all messages are transmitted over the channels and the next round begins
- 1-3: Local computation by processors
- 4: Transmission of messages handled by the network (this step could even come first)

Equivalent to the previous type of round

Use the one that is more convenient to you

Solution: Pseudocode

Algorithm Spanning tree broadcast

State of processor u_i :

- parent_i: holds a processor index or nil; u_i 's parent
- $children_i$: holds a set of processor indices (possibly empty); u_i 's children
- Boolean $terminated_i$: indicates whether u_i has terminated (1) or not (0)

Solution: Pseudocode

Algorithm Spanning tree broadcast Initially u_0 knows $\langle M \rangle$

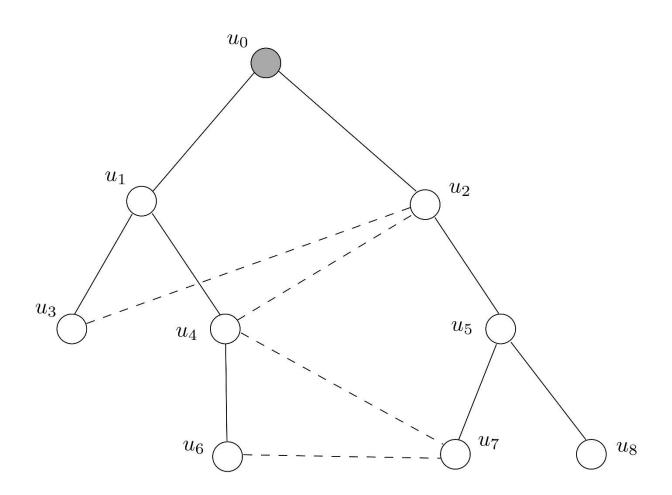
Code for leader (u_0) : send $\langle M \rangle$ to all children terminate

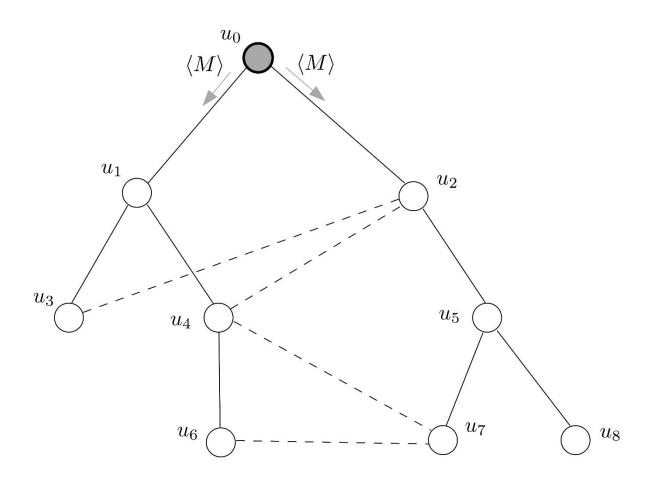
Code for non-leader:

upon receiving (M) from parent:

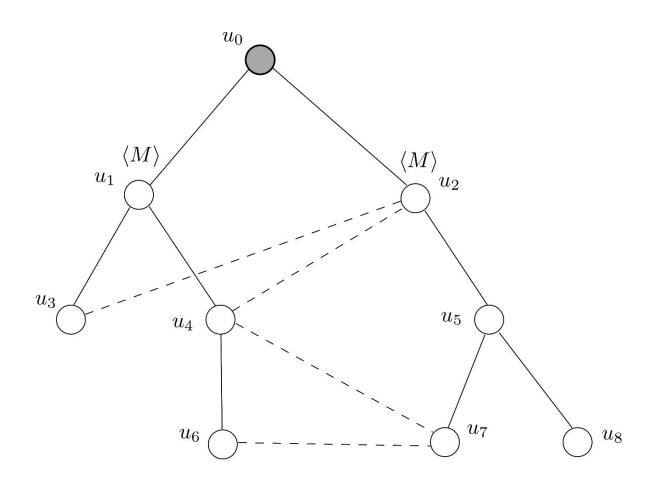
send (M) to all children

terminate

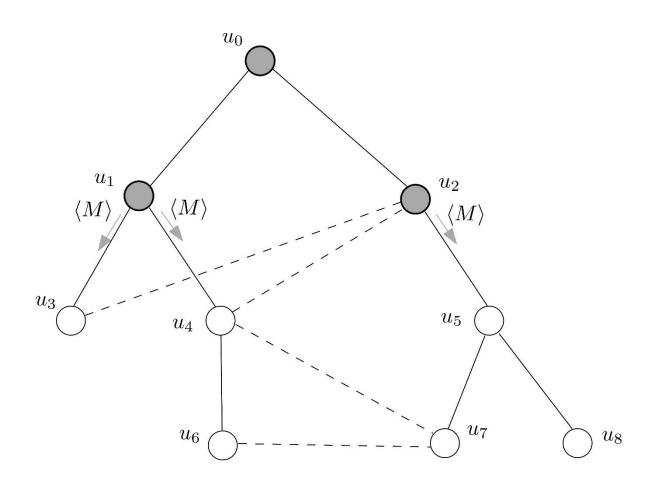




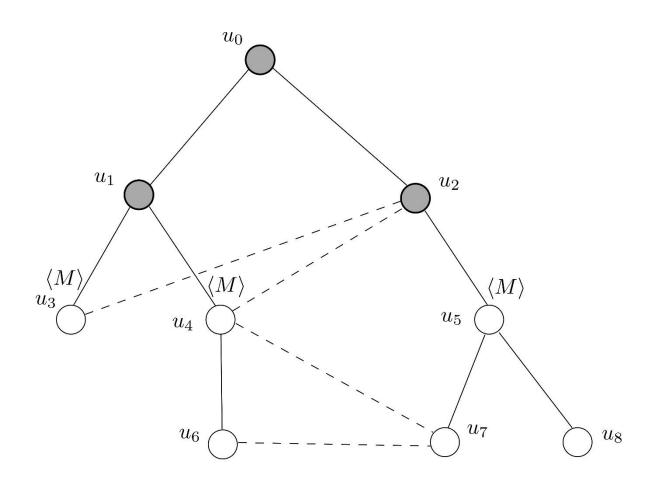
round = 1



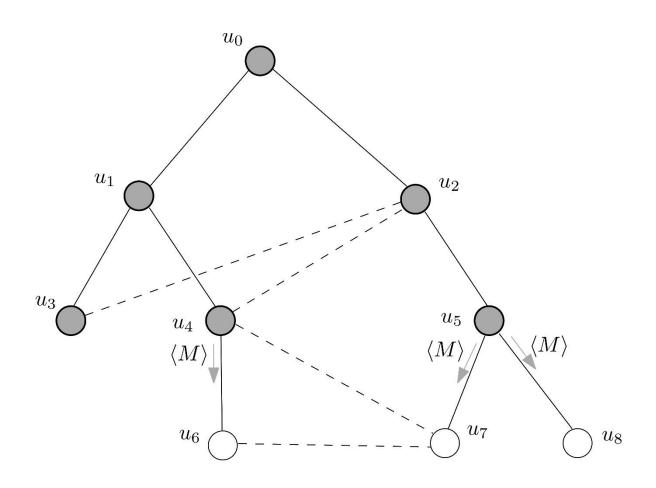
round = 1



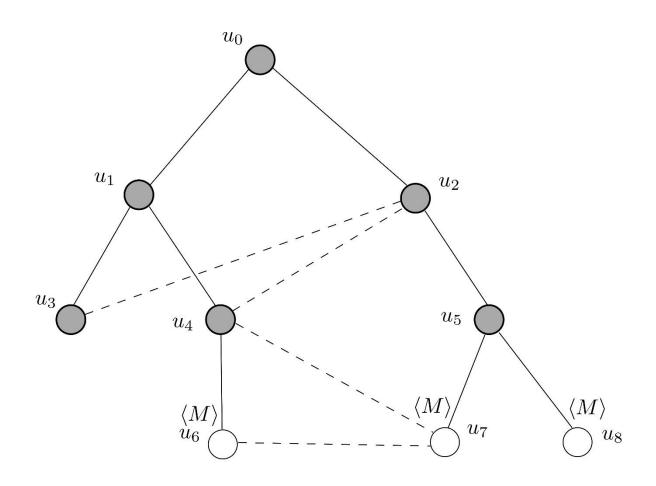
$$round = 2$$



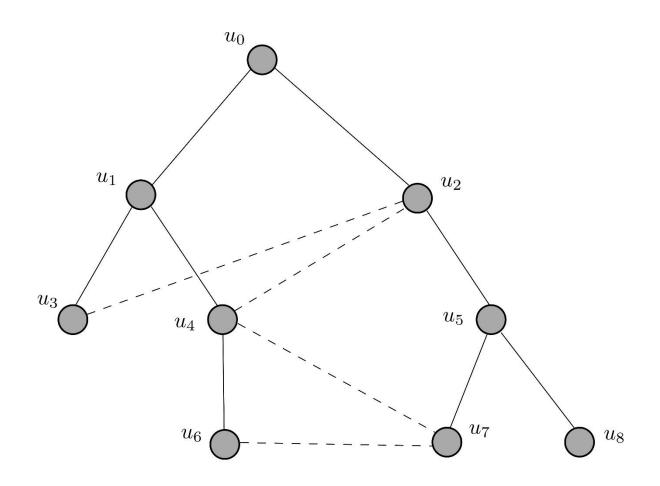
round = 2



$$round = 3$$



round = 3



round = 4

Correctness and Performance

When we devise an algorithm we typically should

- 1. Convince that it is correct
- 2. Analyse its performance

Correctness:

- Usually a proof that the algorithm does as expected
- Performance:
 - Time Complexity (e.g., #rounds required)
 - Space Complexity (e.g., memory used by processors)
 - Communication Complexity (e.g., total #messages transmitted, size of messages)