Divide and Conquer Master Theorem Problem

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1) T(n) = 8T(n/3) + n^2
a = 8, b = 3, f(n) = n^2, log_3 8 < 2
n^{\log_{h} a} = n^{\log_{3} 8} < n^2 \Rightarrow n^2 = n^{\log_{3} 8 + \epsilon}, where \in is equal to (2 - \log_{3} 8), therefore > 0.
n^2 = \Omega(n^{\log_3 8 + \epsilon})
af(n/b) \le c*f(n), where c<1
8f(n/3)^2 \le c^*f(n)^2
8(n^2/9) \le c^*n^2, therefore c \ge 8/9
Thus, case 3 applies and T(n) = \Theta(n^2)
2) T(n) = 10T(n/3) + n^2
a = 10, b = 3, f(n) = n^2, log_3 10 > 2
n^{\log_{b} a} = n^{\log_{3} 10} > n^2 \Rightarrow n^2 = n^{\log_{3} 10 - \epsilon}, where \epsilon is equal to (\log_{3} 10 - 2), therefore \epsilon > 0.
Thus, case 1 applies and T(n) = \Theta(n^{\log_3 10})
3) T(n) = 16T(n/4) + n^2 log^3 n
a = 16, b = 4, f(n) = n^2 \log^3 n, \log_4 16 = 2
n^{\log_b a} = n^{\log_4 16} = n^2 = f(n) = n^2.
Therefore, T(n) = \Theta(n^{\log_b a} * \log^{k+1}(n)), where k = 3.
Thus, case 2 applies and T(n) = \Theta(n^2 \log^4(n))
4) T(n) = 9T(n/3) + n^3
a = 9, b = 3, f(n) = n^3, log_3 9 = 2
n^{\log_{h} a} = n^{\log_{3} 9} < n^{3} \Rightarrow n^{2+\epsilon}, where \in is equal to 1, therefore > 0.
af(n/b) \le c^*f(n), where c<1
9f(n/3)^3 \le c^*f(n)^3
9(n^3/27) \le c^*n^3, therefore c \ge 9/27 = 1/3
Thus, case 3 applies and T(n) = \Theta(n^3)
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