COMP15111 Notes

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1 Lecture 1: Introduction

1.1 A Computational Model

The simplest, earliest, commonest, most important computational model is the
Von-Neumann Imperative Procedural Computer Model

According to this model, a computer can:

1. Store information

2. Manipulate the stored information

3. Make decisions depending on the stored information

1.2 Simple View Of A Computer

$$Memory \Leftrightarrow Bus \Leftrightarrow Processor$$

1.2.1 Memory

Memory is a set of locations which can hold information, such as numbers(or
programs). Each memory location has a unique (numerical) address, and there
are typically thousands of millions of different locations. There are various ways
of depicting memory; a common one is a 'hex dump' that often looks something
like this:

Address	Values (8 bit numbers)	Characters
00000000	48 65 6c 6c 6f 0a	Hello.

Each item that is in the memory has a unique address.

Run the command *hexdump* to generate hexdumps.

1.2.2 Bus

A bus is a bidirectional communication path. It is able to transmit addresses
and numbers between components inside the computer.

1.2.3 Processor

The processor obeys a sequence of instructions, commonly referred to as a program. Historically the processor was often referred to as a CPU, however, this
is inappropriate nowadays since typical processors consist of several processing
cores.

1.3 Three-address instructions

Every kind of processor has a different set of instructions, real world examples include: Pentium, ARM and others

Each three-address instruction:

1. Copies the values from any two memory locations and sends them to the processor (source operands)

2. Copies some operation e.g. adds the copied numbers together

3. Copies the result back from the processor into a third memory location (destination operand)

For example, if we wanted to convert the Java code $sum = a + b$; into a three-address instruction we would:

1. Identify the two *source operands*: a holds 2, b holds 3

2. Perform the *operation*: $2 + 3 = 5$

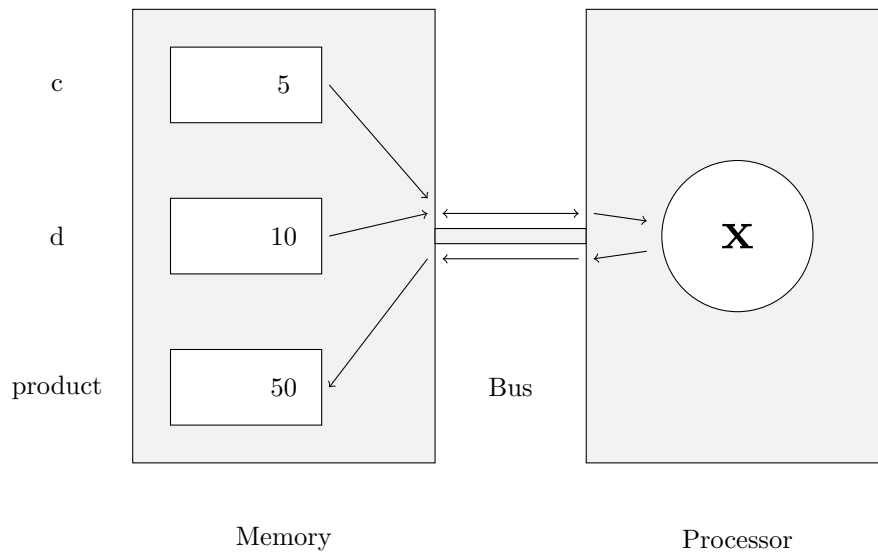3. Let the variable *sum* equal the answer 5. This is the *destination operand*

1.3.1 Three address example

Question: Convert the Java code *product = c * d;* into the three-address style and draw a two box view of it.

First we need to re-write the Java code in the three-address style:

$$product \leftarrow c * d$$

Now we can draw the box view of it:

c   5

d   10

X

product   50   Bus

Memory   Processor

1.4 Registers

Reduce accesses to main memory by using registers:

1. Very small, very fast memory within the processor

2. Each register can hold a single value

Processors contain up to a few dozen registers e.g. ARM: 16 registers - R0 to R15

1.5 Instruction Styles

Make instructions as simple as possible while allowing for the best use of registers. One-address style (e.g. Intel Pentium):

1. Only use (at most) 1 memory location in each instruction

2. Use registers for the other operands

$$register \leftarrow register + memory location$$

Load-store style (e.g. ARM): Use registers for all three operands

$$register \leftarrow register + register$$

So we need extra instructions to get at memory locations:

1. **Load** (from memory): register ← memory location

2. **Store** (to memory): memory location → register

Sum = a + b + c; Becomes: Register1 ← a (i.e. load from a)
Register2 ← b (i.e. load from b)
Register3 ← register 1 + register2 (i.e. a+b)
Register4 ← c (i.e. load from c)
Register5 ← register3 + register4 (i.e. (a+b)+c)
sum → register5 (i.e. store to sum)
One load or store instruction per variable (a, b, c, sum)
One arithmetic instruction per operation (+, +)
Lots of very simple, very fast instructions

2 Lecture 2

Computers obey programs which are sequences of instructions. Instructions are coded as values in memory. The sequences are held in memory adjacent memory locations. Values in memory can be interpreted as:

- Numbers (in several different ways)
- Instructions
- Text
- Colours
- Music
- Anything you want

Values are often represented as numbers for convenience.

2.1 Assembly Language

Assembly language is a means of representing machine instructions in a human readable form.
Each type of processor has its own assembly language but they typically have a lot in common:

- A mnemonic specifies the type of operation
- A destination a register on this case
- And one or more sources also registers
- Possibly with a comment too

2.2   ARM instructions

ARM has many instructions but we only need three categories:

- Memory operations
- Processing operations
- Control flow

Memory operation move data between the memory and the registers. Processing operations perform calculations using value already in registers. Control flow instructions are used to make decisions, repeat operations etc.

2.3   ARM memory instructions

Memory operations load a register from the memory or store a register value to the memory. e.g. LDR R1, a means: R1 ← a
e.g. STR R5, sum means: R5 → sum (i.e. sum ← R5)

a and *sum* are aliases for the addresses of memory locations.

2.4   ARM processing instructions

Processing operations such as addition, subtraction, multiplication.

e.g. ADD R3, R1, R2 means: R3 ← R1 + R2

2.5   ARM control instructions

Fundamentally, these are branches to other code sequences. Often, branches are made conditional to allow decisions to be made.

e.g. B somewhere means: branch to somewhere
e.g. BEQ elsewhere means: branch to elsewhere IF previous result was equals
e.g. BNE wherever means: branch to wherever IF previous result was not equal

2.6   Stored programs and the Program Counter

A computer can make decisions, and choose which instructions to obey next depending upon the results of those decisions. How? First we need to see how the sequence of instructions is controlled. Von-Neumann Model: memory holds both instructions and numbers - **stored program Program Counter**

(PC) register: holds the address of the memory location containing the next instruction to b obeyed (executed).

ARM uses register 15 as its PC

2.7   Fetch-Execute Cycle

Start with PC containing the address of (the memory location holding) the first instruction of a program.

Repeatedly:

1. **Fetch**: copy the instruction, pointed to by the PC, from memory and set PC to point to the next instruction

2. **Execute**: obey the instruction (exactly as before)

ARM:

1. 'Resets' to (starts at) address 00000000

2. Instructions each occupy 4 memory locations, so PC increases by 4 in each fetch

2.8   Decision Making

Linear sequences of instructions are limiting. To make a decision, the computer must change (or not) to a different sequence of instructions.

e.g. a 1 pound discount on items worth 20 pounds or more. Decision: compare the total and 20 pounds to see if it is larger, then depending on result, either perform action or not. Action: subtract 1 pound from the total.

Computers have no intelligence, so spell out details. **Formalise**: if total ≥ 20 pounds then subtract 1 pound from total
Rewrite: if total < 20 pounds then dont
subtract 1 pound from total
Encode: as ARM instructions

2.9   Compare And Branch