

Conflict-free Replicated Data Types (CRDTs)

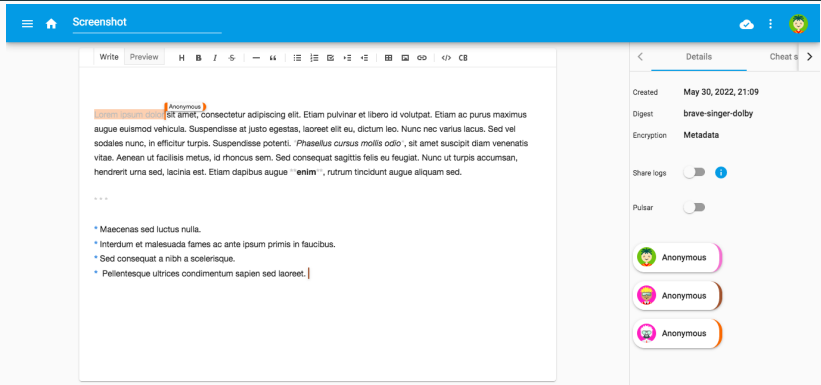
An Overview

Matthieu Nicolas (`matthieu.nicolas@inria.fr`)

02/05/2024



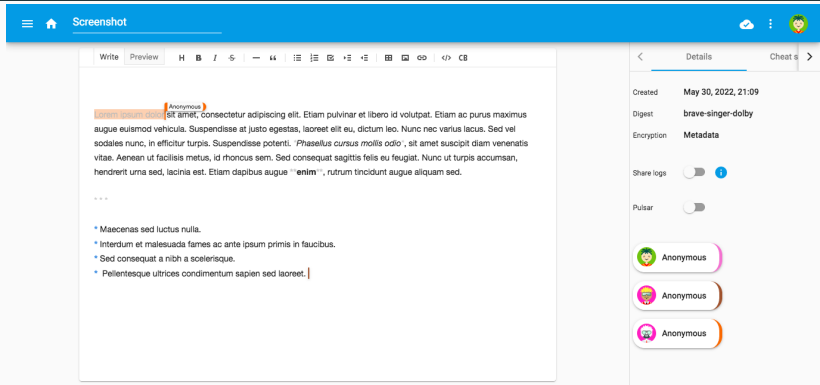
MUTE [Nic+17]



Local First Software [Kle+19]

- Peer-to-Peer (P2P) collaborative applications

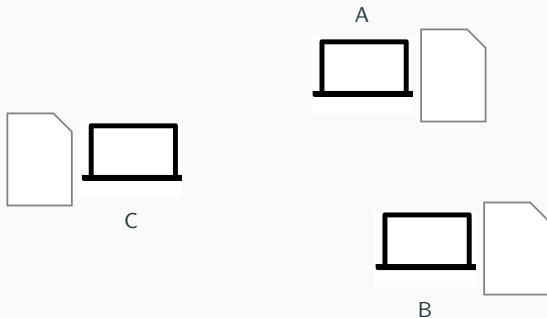
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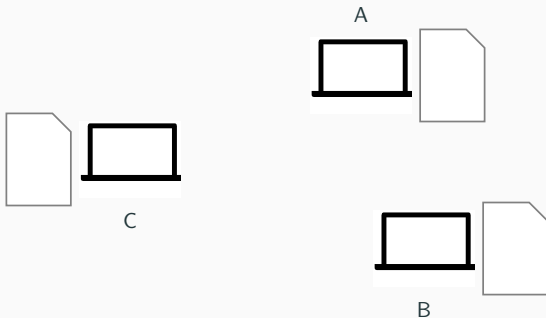
Local First Software [Kle+19]

- Peer-to-Peer (P2P) collaborative applications
- Always available
- Ensure ownership and privacy of data

Data replication in P2P systems

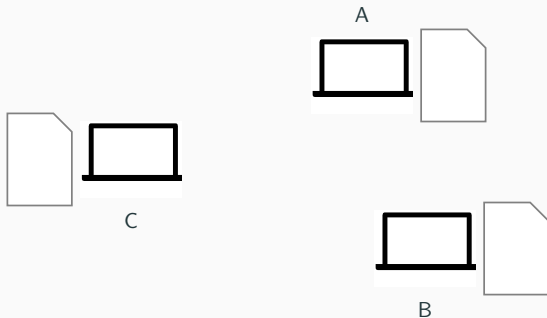


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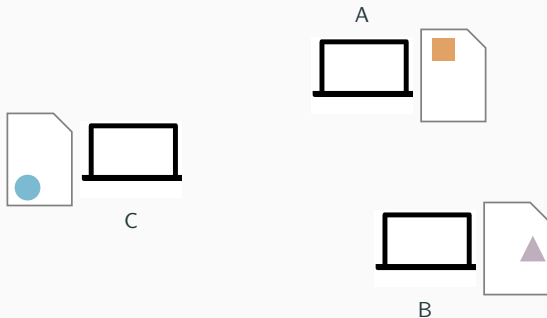
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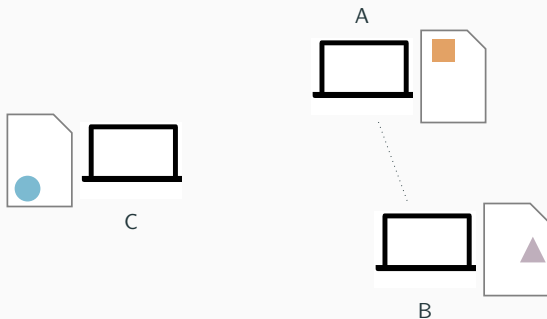
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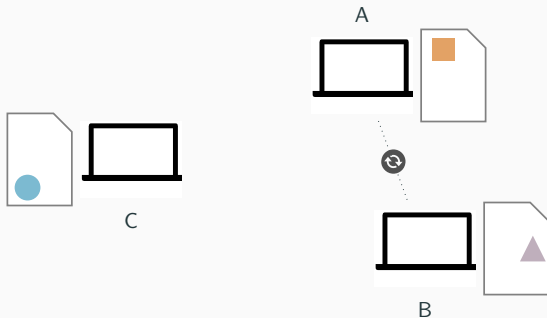
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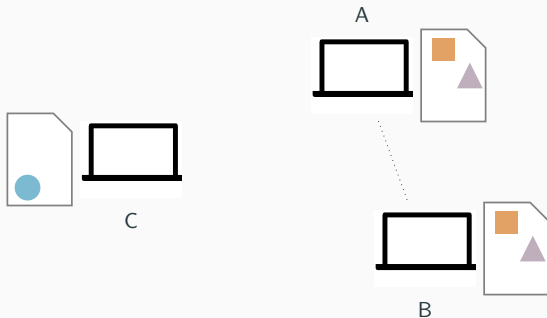
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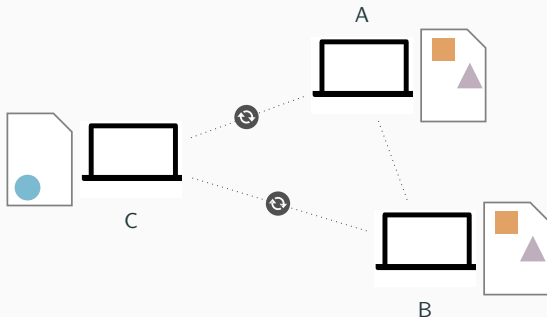
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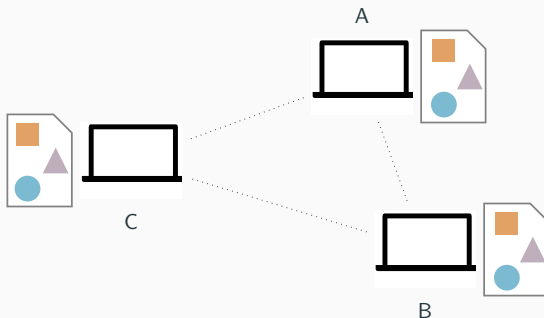
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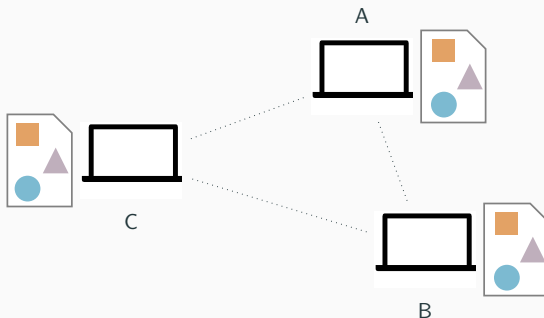
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- ... Despite different integration orders of updates

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Require conflict resolution mechanisms

Conflict-free Replicated Data Types (CRDTs)
are a **family of conflict resolution mechanisms**
for P2P systems

Conflict-free Replicated Data Types (CRDTs) [Sha+11]

- New specifications of existing Data Types, e.g. *Set* or *Sequence*
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- Rely on the lattice theory . . .
- . . . More specifically, **CRDTs are join-semilattices**

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- Exists a set of minimal states of the join-semilattice, the *irreducible elements*

Design of CRDTs

- Several CRDTs may be designed for a given data type ...
- ... Each offering different trade-offs

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What impact the design of a given CRDT [Pre18]

- Conflict Resolution Semantics
- Synchronisation Model

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What impact the design of a given CRDT [Pre18]

- Conflict Resolution Semantics
 - Synchronisation Model
-
- Impact their overhead in terms of computation, memory and bandwidth

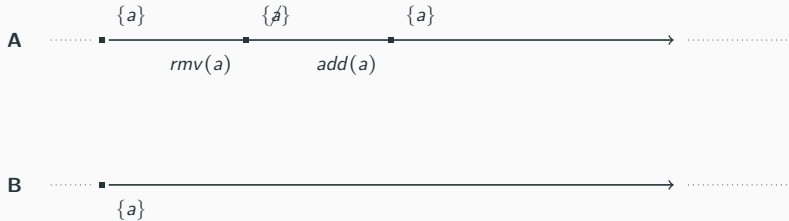
Conflict Resolution Semantics

- Distributed setting **allows new scenarios**



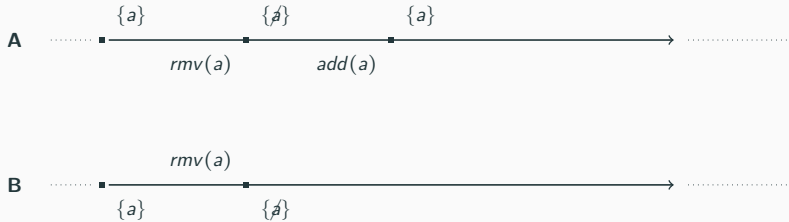
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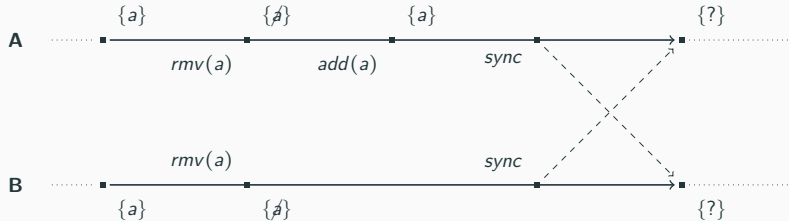
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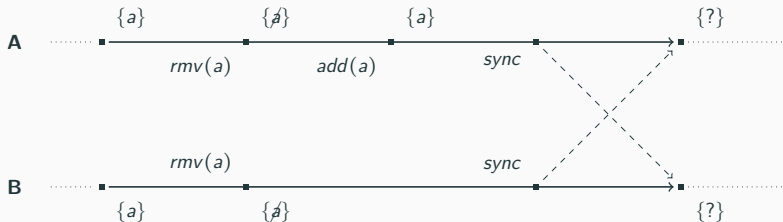
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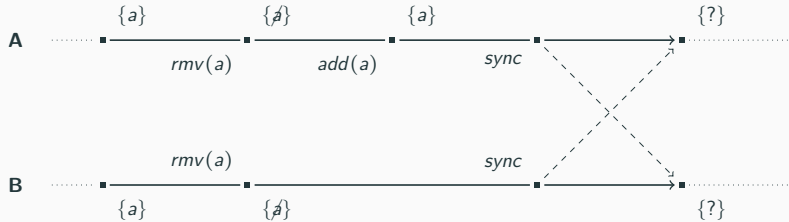
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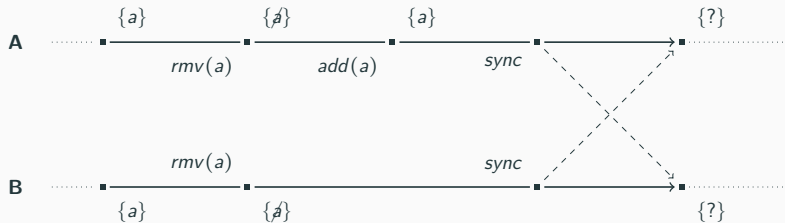
- **What should be the final state** in this scenario ?
- Designing a CRDT consists in **defining its behaviour** in such cases

Conflict Resolution Semantics - Case study of the Set



Several semantics proposed:

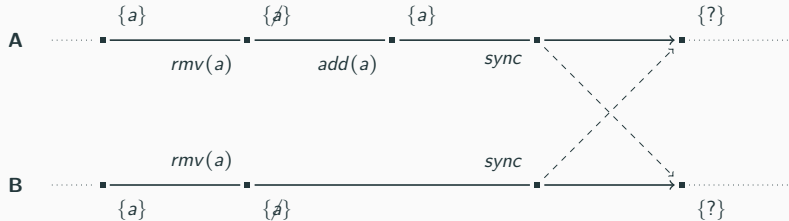
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- *Add-Wins*: $add(a)$ has priority over concurrent $rmv(a) \implies \{a\}$

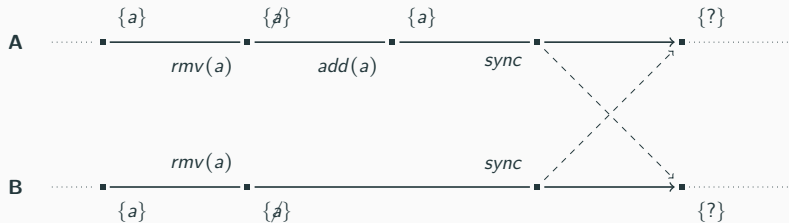
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- *Remove-Wins*: $rmv(a)$ has priority over concurrent $add(a) \implies \{\emptyset\}$
- *Causal-Length* [YR20]: The last action of the longest chain of updates determines the presence (or not) of the element $\implies \{a\}$

Synchronisation Models

To converge

- Nodes have to propagate changes ...
- ... And integrate those of others

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- Delta-based synchronisation [ASB18]
 - "Best of the two worlds" approach

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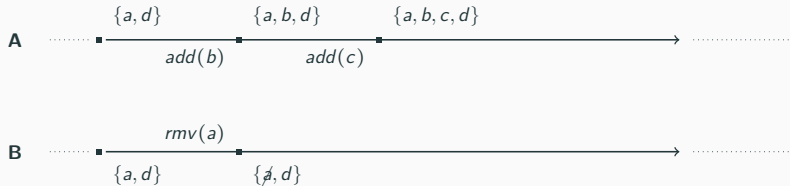
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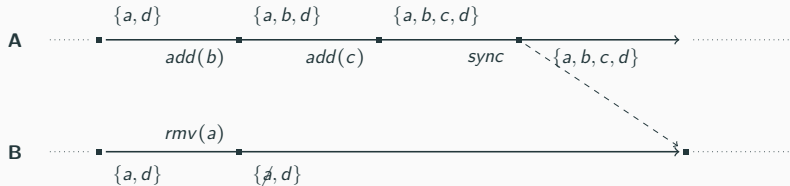
State-based synchronisation

- Send periodically current state to other nodes



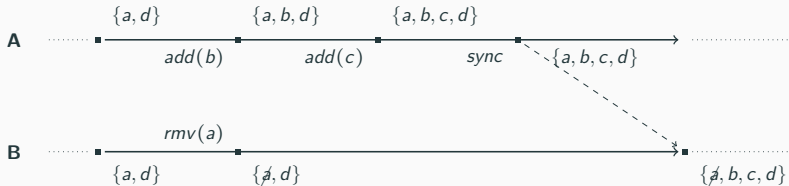
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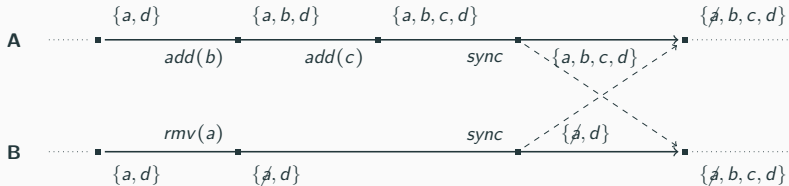
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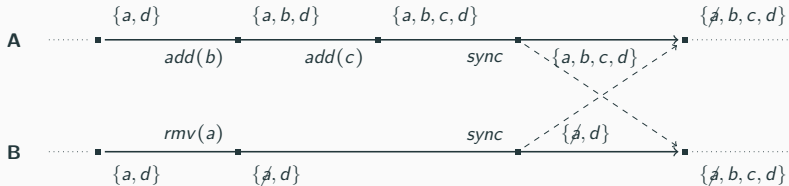
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- With merge, a commutative, associative and idempotent function

Strengths

- No assumptions on the network reliability
- i.e. messages may be lost, re-ordered or duplicated w/o impact

State-based synchronisation - Review

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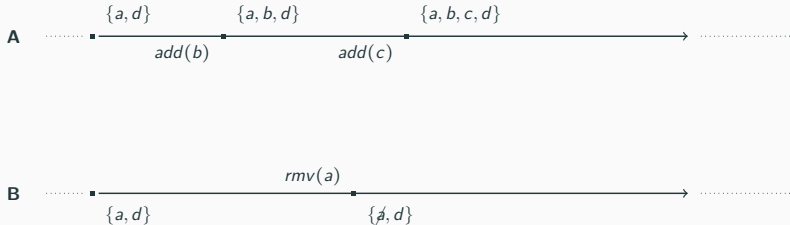
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- Depending on data type, states expensive to broadcast. . .
- . . . And merge expensive to execute

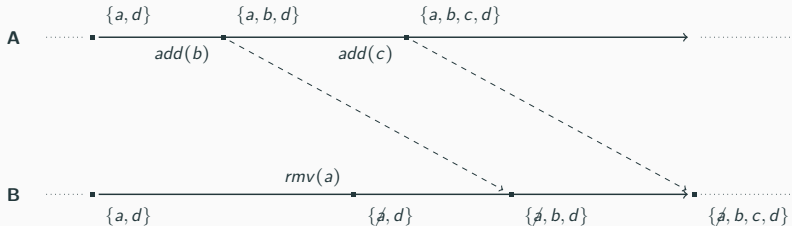
Operation-based synchronisation

- Encode updates as arbitrary messages, called *operations*
- An operation corresponds to one or several *irreducible elements*



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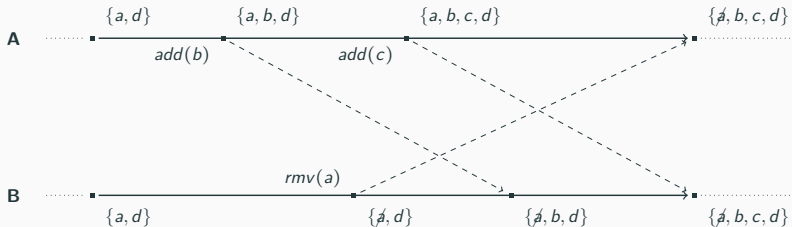
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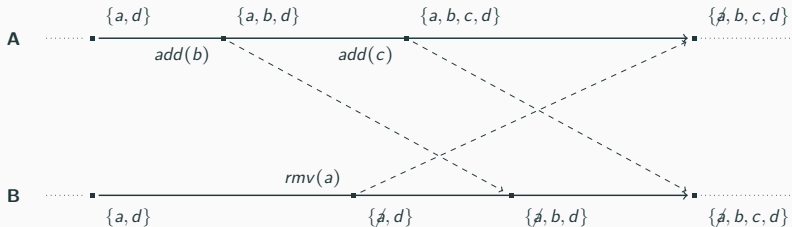
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- Upon reception, apply operations on current state
- Concurrent operations must be commutative

Strengths

- Designing operations is straightforward
- Operations usually cheap to broadcast and apply

Operation-based synchronisation - Review

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- Hides/delegates complexity to delivery of operations
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- Weak to network failures
- Have to pair Op-based CRDTs to a message delivery service
 - To re-order and/or de-duplicate operations
 - To retrieve lost operations using anti-entropy mechanisms

CRDTs are new specifications of Data Types

- Enable nodes to collaboratively edit data w/o coordination
- Ensure Strong Eventual Consistency

To recap

CRDTs are new specifications of Data Types

- Enable nodes to collaboratively edit data w/o coordination
- Ensure Strong Eventual Consistency

Several CRDTs designed per data types

- With different behaviours in case of conflict. . .
- . . . and different techniques to broadcast updates
- E.g. 10+ CRDTs for *Sequence*

- CRDTs for many data types
 - *Register*, *Set* [Pre18], *Sequence* [Roh+11; WUM09], *JSON* [KB17], *Tree* [Kle+22] ...
- Libraries providing CRDTs to build new applications
 - Yjs [Yjs], Automerge [Aut]...

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 - Yjs [Yjs], Automerge [Aut]...
- Used in collaborative applications
 - Teletype (Atom's pair-programming feature), Apple Notes...
- Used in multi-master replication for distributed databases
 - Redis [Lab], Microsoft Azure CosmosDB [DB]...

- Designing CRDTs for new use cases
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Ongoing researches

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- Designing CRDTs for systems with malicious nodes

Thanks for your attention, any questions?



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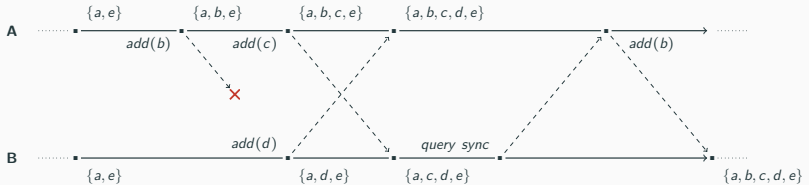
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Back-up slides

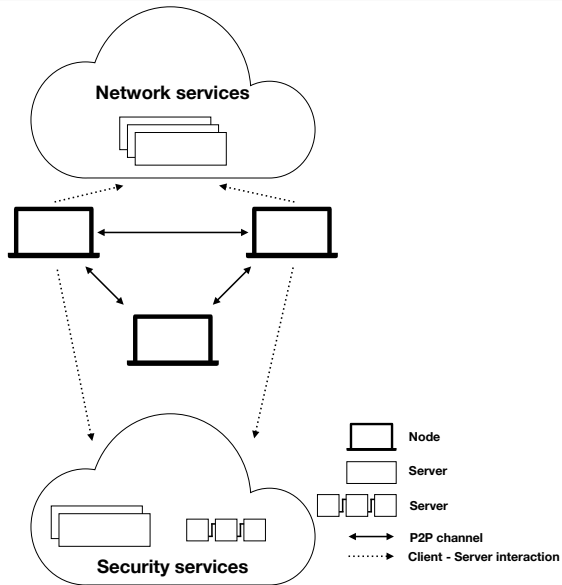
Operation-based synchronisation - network failure



Synchronisation Models - Summary

	State-based	Op-based	Delta-based
Integrate updates by merging states	✓	✗	✓
Integrate updates by irreducible elts	✗	✓	✓
Handle natively network failures	✓	✗	✓
Suited for real-time systems	✗	✓	✓

MUTE System Architecture



MUTE Software Architecture

