

# Analysis of CPU-Scheduling Algorithms and Real time Scheduling

## Operating System

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8 December, 2016

# Outline

- 1 Project Overview
  - Introduction
- 2 Scheduling
  - Scheduling Criteria
- 3 CPU Scheduling Algorithms
  - Scheduling Algorithms
  - Simulation Considering Uniprocessor
  - Simulation Results
- 4 Real time Scheduling
  - Why Real time Scheduling and It's required parameters
  - Deadline Criteria and General Classes of Real time Scheduling algorithms
  - Implementation
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## Processor Aim

The aim of processor scheduling is to assign processes to be executed by the processor or processors over time, in a way that meets system objectives, such as **response time, throughput, and processor efficiency**. In many systems, this scheduling activity is broken down into three separate functions: **long-, medium-, and short-term** scheduling, and those functions are handled by schedulers.

# Introduction

## Processor Aim

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## Scheduling Mechanism in Support

As per the above claim scheduling is a crucial task to manage the efficiency of system. so here we have incorporated two kind of CPU scheduling schemes and analyze how they match up with certain scheduling criteria.

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# Scheduling Criteria

## User Oriented Criteria

- Turnaround time
- Waiting time
- Deadline

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- Throughput
- CPU Utilization
- Fairness



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## System Oriented Criteria

- Throughput
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  - Fairness
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- High quality CPU scheduling algorithms generally focus on criterias such as throughput, CPU utilization rate, response time, turnaround time, and waiting time.

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# Scheduling Algorithms

## Generalize Scheduling Schemes

- Preemptive Scheduling.
- Non Preemptive Scheduling.
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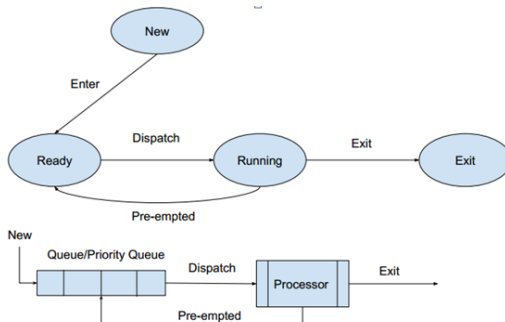


Figure : Scheduling State transactions diagram

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## Simulated Scheduling Algorithms

- FCFS - First Come First Serve
- SRN - Shortest Remaining Next
- RR - Round Robin(Quantum-6)
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## Analysis Criteria

- Average Waiting Time
- Average Turn Around Time
- Deadline Miss
- Average Throughput

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# Average Waiting Time

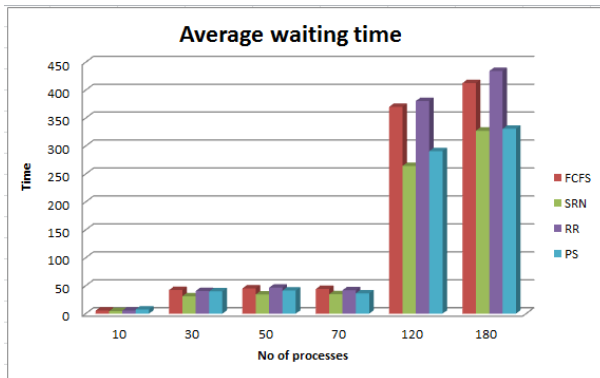
$$\text{Average waiting time} = \frac{\sum_{i=1}^n w_i}{n} \quad (1)$$

- $n$  = number of processes,  $w_i$  = waiting time of  $i^{\text{th}}$  process

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# Average Turn Around Time

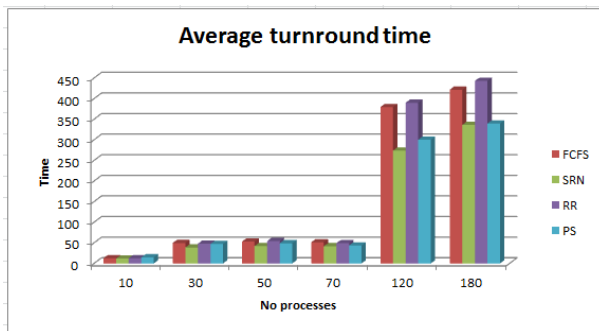
$$\text{Average turnaround time} = \frac{\sum_{i=1}^n (w_i + e_i)}{n} \quad (2)$$

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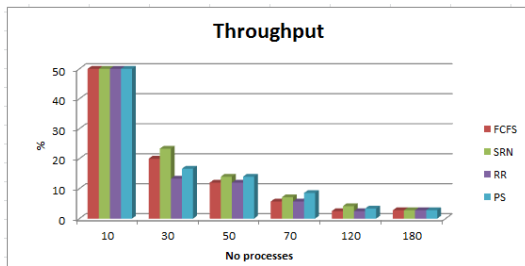
$$\text{Average Throughput} = \frac{n'}{T} \quad (3)$$

- The number of processes completed successfully per time unit.
- $n$  = number of processes ,  $T$  = Time taken for completion of  $n$  processes ,  $t$  = time value where  $t < T$   $n'$  = number of processes with finish time  $< t$
- Here, average through put means number of processes completed successfully before given time.

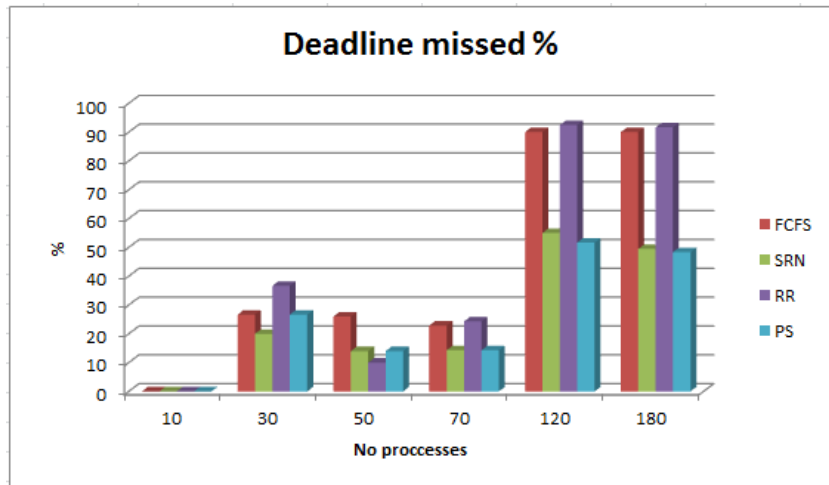
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# Deadline Missed



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- From the previous deadline miss result we can say that this type of scheduling algorithms, may break down performance of Real time systems very badly because those systems are depended on fixed deadlines.

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# Why Real time Scheduling and It's required parameters

- From the previous deadline miss result we can say that this type of scheduling algorithms, may break down performance of Real time systems very badly because those systems are depended on fixed deadlines.
- A common characteristic of real time system is that their requirements specification includes timing information in the form of **deadlines** which traditional scheduling algorithms not emphasize
- Based on deadlines we can classify the deadline consequences
  - **Hard Deadlines**
  - **Soft Deadlines**

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# Deadline Criteria and General Classes of Real time Scheduling algorithms

## Deadline Criteria

$$C \leq D$$

$$C \leq D \leq T \text{ For Periodic Process}$$

where C - Computation time , D - Deadline , T - Period

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## General Classes of Real time Scheduling algorithms

- Static table-driven approaches
- Static priority-driven preemptive approaches
- Dynamic planning-based approaches
- Dynamic best effort approaches:

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# Assumptions and Algorithms

## Assumptions

- $C_i \leq D_i = P_i$  (i.e. The processes have computation time less than their deadline, and the deadline is equal to their period).
- Computation times for a given process are constant.
- No precedence relations exist between processes.
- No inter-process communication or synchronization is permitted.
- Context switches have zero cost.
- All processes are allocated to a single processor.



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## Implemented algorithms

- RMS - Rate Monotonic Scheduling
- EDF - Earliest Deadline First

- RMS is static priority driven pre-emptive approach.

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## Criteria

$$\frac{C_1}{P_1} + \frac{C_2}{P_2} + \frac{C_3}{P_3} + \dots + \frac{C_n}{P_n} = U \leq n(2^{\frac{1}{n}} - 1)$$

- C - Computation time , P - Process Period.

# RMS Implementation

```
int Period[] = {10,5,30,15};
int Execution[] = {2,1,5,2};
process *Array[N];
```

```
// Register signal and signal handler
signal(SIGINT, signal_callback_handler);
```

```
// Define the function to be called when ctrl-c (SIGINT) signal is sent to process
void signal_callback_handler(int signum)
{
    outputfile<<" Interrupt is Caught :- signal "<<signum<<endl;
    // Cleanup and close up stuff here
    outputfile<<" Interrupt is disabled. "<<endl;
    // Terminate program
    cout<<" Interrupt is Caught :- signal "<<signum<<endl;
    // Cleanup and close up stuff here
    cout<<" Interrupt is disabled. "<<endl;
}
```

```
Process 1 created successfully.
Process 2 created successfully.
Process 3 created successfully.
Process 4 created successfully.

CPU Utilization :-0.700000 Theoretical :- 0.756828
Process No. -> Waiting time -> Turn around time

ProcessID : 1 state : ready Tue Dec 6 23:38:29 2016
ProcessID : 2 state : ready Tue Dec 6 23:38:29 2016
ProcessID : 3 state : ready Tue Dec 6 23:38:29 2016
ProcessID : 4 state : ready Tue Dec 6 23:38:29 2016

ProcessID : 2 state : running Tue Dec 6 23:38:29 2016
ProcessID : 2 state : exit Tue Dec 6 23:38:30 2016

ProcessID : 2 -> 0 -> 1
ProcessID : 1 state : running Tue Dec 6 23:38:30 2016
ProcessID : 1 state : exit Tue Dec 6 23:38:32 2016

ProcessID : 1 -> 1 -> 3
ProcessID : 4 state : running Tue Dec 6 23:38:32 2016
ProcessID : 4 state : exit Tue Dec 6 23:38:34 2016
```

- The process with the (current) **closest deadline** is assigned the highest priority in the system and therefore it executes first.

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### Criteria

$$\sum_{i=1}^n \frac{C_i}{P_i} \leq 1 \quad (4)$$

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### Criteria

$$\sum_{i=1}^n \frac{C_i}{P_i} \leq 1 \quad (4)$$

- Hence, 100 % processor utilization is possible.



# EDF Implementation

```
int deadline[] = {15,12,9,8};  
int Execution[] = {4,3,5,2};  
int Arrivalttime[] = {0,0,2,5};
```

```
Process 1 created successfully.  
Process 2 created successfully.  
Process 3 created successfully.  
Process 4 created successfully.  
ProcessID : 1 state : ready Tue Dec 6 23:40:14 2016  
  
ProcessID : 2 state : ready Tue Dec 6 23:40:14 2016  
ProcessID : 2 state : running Tue Dec 6 23:40:14 2016  
ProcessID : 2 state : ready Tue Dec 6 23:40:16 2016  
ProcessID : 3 state : ready Tue Dec 6 23:40:16 2016  
ProcessID : 3 state : running Tue Dec 6 23:40:16 2016  
ProcessID : 3 state : ready Tue Dec 6 23:40:19 2016  
ProcessID : 4 state : ready Tue Dec 6 23:40:19 2016  
ProcessID : 4 state : running Tue Dec 6 23:40:19 2016  
ProcessID : 4 state : exit Tue Dec 6 23:40:21 2016  
ProcessID : 4 -> 2 -> 0 -> 2  
ProcessID : 3 state : running Tue Dec 6 23:40:21 2016  
ProcessID : 3 state : exit Tue Dec 6 23:40:23 2016  
ProcessID : 3 -> 5 -> 2 -> 7  
ProcessID : 2 state : running Tue Dec 6 23:40:23 2016
```

# EDF Deadline Missed

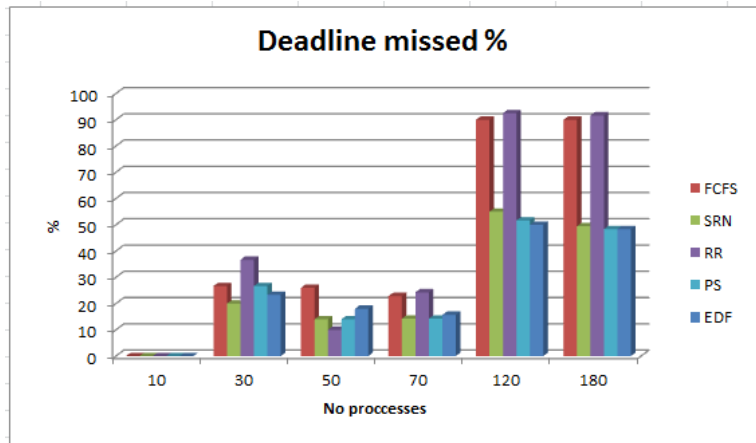


Figure : Deadline Miss with EDF

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# Conclusion

- We compared various CPU Scheduling Algorithms both static and real time and did their simulation in C++, after analyzing their results, we concluded that traditional scheduling algorithms like FCFS, RR, SJF, PS etc. missed lots of deadlines. Therefore we conclude that, these type of scheduling algorithms cannot be used for scheduling in Real time systems and the specialized real time algorithms like RMS and EDF are needed for process scheduling in Real time systems.

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