The Quintessential Quandary Guide

Mike Bond

July 6, 2022

1 Introduction

Quandary is a language that combines elements of functional languages (especially Scheme/Lisp) and imperative languages (especially Java).

Mike Bond created Quandary in Summer 2019 to use in programming languages classes (specifically CSE 3341 and 6341 at Ohio State University), as part of an effort to better connect the technical material and implementation projects.

This guide tries to cover everything related to Quandary, especially how to use Quandary as a programming language (by writing Quandary programs and running them using the Quandary reference interpreter) and how to implement a Quandary interpreter (by extending the Quandary skeleton interpreter and testing the modified interpreter using the grading script and grading test cases).

2 Getting Started

Unless stated otherwise, commands in this guide should be run from the root Quandary-Public directory. Some commands may only work if your shell is bash. To change your shell to bash, run

chsh -s /bin/bash

from any directory.

Complete the steps in this section in order.

2.1 Choose a Platform

Linux. Linux is your best bet. Everything should just work.

macOS. macOS is a solid second choice. The Quandary scripts will work mostly fine, except for one issue:

In order for the grading script (grade.sh) to work on macOS, you may need to install the correct variant of the realpath command using the following command (run from any directory):

brew install coreutils

If brew command fails, you'll need to first get Homebrew (visit http://brew.sh).

Windows. If you have the misfortune of using Windows, you can still make it work. You can either

- run Linux in an OS virtual machine such as VirtualBox;
- view and edit the code on Windows, but build and run the interpreter on stdlinux; or
- run Windows Subsystem for Linux (WSL).

If you use WSL, here's an issue you may encounter at some point:

You may need to run the following commands on the quandary scripts to fix errors regarding trailing r characters:

```
sed -i 's/\r$//' ref/quandary
sed -i 's/\r$//' skeleton/quandary
```

Some people have found that they need to run the commands on other files such as the test case files (.dat files) and TODO.

2.2 Get Quandary

Clone Quandary-Public by running this command from the directory where you want to contain the Quandary-Public directory:

```
git clone git@github.com:mdbond/Quandary-Public.git
```

Quandary-Public contains the following directories and files:

- ref/ contains the Quandary reference interpreter
- skeleton/ contains the Quandary skeleton interpreter
- examples/ contains Quandary code examples
- grading/ contains grading script and test cases
- quandary.pdf is this document

You should view and edit the Quandary code—particularly the skeleton interpreter—in an IDE. Not sure which IDE to use? Use Visual Studio Code. For VSC to understand the skeleton interpreter (e.g., to support navigating and detecting errors), you'll want to open the skeleton directory as the root folder in VSC.

Note that you'll still need to *build* the skeleton interpreter using the Makefile (see below). But you can probably configure VSC to run the Makefile every time you change a source file.

2.3 Install Java

To run the reference interpreter, you'll need a Java virtual machine (JVM), i.e., the java command. To build the skeleton interpreter, you'll need the Java compiler, i.e., the javac command. To install both tools, install the Java Development Kit (JDK).

You may already have the JDK. You can skip this part if the javac command already works.

Not sure which Java implementation to install? Install OpenJDK.

If you're on Ubuntu, you can install the OpenJDK JDK version 17 using this command:

```
sudo apt install openjdk-17-jdk
```

If you're on stdlinux, run the following (just once in your life):

```
subscribe JDK-CURRENT
```

Then log out and back in.

Folks have reported needing a JDK version of at least 11 or 12.

2.4 Get the Reference Interpreter Working

Run the reference interpreter, which should print usage information:

```
$ ref/quandary
Expected format: quandary [OPTIONS] QUANDARY_PROGRAM_FILE INTEGER_ARGUMENT
Options:
   -gc (MarkSweep|MarkSweepVerbose|RefCount|Explicit|NoGC)
   -heapsize BYTES
   -ct TIMEOUT_IN_SECONDS
BYTES must be a multiple of the word size (8)
Quandary process returned 0
```

Run the reference interpreter with a Quandary program and argument:

```
$ ref/quandary examples/primes2.q 20
Interpreter returned (2 . (3 . (5 . (7 . (11 . (13 . (17 . (19 . nil)))))))
Quandary process returned 0
```

2.5 Get the Skeleton Interpreter Working

Dependencies. Before building the Quandary skeleton interpreter, you need to download and extract CUP and JFlex and set environment variables that the skeleton's Makefile is expecting.

Download JFlex and CUP using the following URLs:

- JFlex 1.7.0: https://jflex.de/release/jflex-1.7.0.tar.gz
- CUP 0.11b-20160615: http://www2.cs.tum.edu/projects/cup/releases/java-cup-bin-11b-20160615. tar.gz

Then extract them and set JFLEX_DIR and CUP_DIR to point to their locations. For example,

```
tar -zxf jflex-1.7.0.tar.gz --directory $HOME
mkdir -p $HOME/cup && tar -zxf java-cup-bin-11b-20160615.tar.gz --directory $HOME/cup
```

which should put JFlex in \$HOME/jflex-1.7.0 and put CUP's two JAR files in \$HOME/cup.

You can of course put JFlex and CUP in other places if you like.

Set the environment variables JFLEX_DIR and CUP_DIR to the locations of JFlex and CUP, respectively. Use absolute, not relative, paths for JFLEX_DIR and CUP_DIR. For example, if JFlex and CUP are in \$HOME/jflex-1.7.0 and \$HOME/cup, respectively, and your shell is bash:

```
export JFLEX_DIR=$HOME/jflex-1.7.0
export CUP_DIR=$HOME/cup
```

Even better: Add these commands to your \$HOME/.bashrc, and they'll be set every time you open a terminal.

Build the skeleton interpreter. Run make in the skeleton directory:

```
(cd skeleton && make)
```

which will run make in the skeleton directory and later return to the parent directory (Quandary-Public).

Run the skeleton interpreter. Now you can run

2.6 What to Do Next

Write your own Quandary programs and run them with the reference interpreter. Modify the Quandary skeleton in order to implement the projects described in Section ??. Run the grading script (Section ??) to evaluate your modified skeleton. Read the rest of this document before attempting the projects.

3 Understanding the Skeleton Interpreter

4 Academic Integrity

You'll implement the projects by modifying the skeleton interpreter. You'll want to save your code somewhere like in a GitHub repository. However, you must store your code in a *private* repository. Storing your code in a public repository, or making your code public in any other way, during or after the semester, is a violation of academic integrity. And of course don't share or show your interpreter source code to anyone either.

5 Troubleshooting and Suggestions

5.1 Suggestions for Modifying the Skeleton Interpreter

Don't modify files that are generated automatically by JFlex or CUP. To see which files aren't generated automatically, run make clean to eliminate generated files.

5.2 Asking for Help

If you can't figure out the answer or find it in this document or on Piazza, the best ways to ask for help are (in order from most to least recommended):

- Make a public post on Piazza.
- Attend instructor or TA office hours.
- Ask in class.

If you're having technical difficulties like getting a weird error, post as much information as possible. If you encounter problems running a script, run it prefaced with bash -x For example:

bash -x skeleton/quandary examples/primes2.q 20

Another example:

bash -x grading/grade.sh skeleton/myproject.tgz ref/quandary grading/calc-public.dat examples

and then post the full output of the command along with other information including your platform.

5.3 Finding Bugs in the Reference Interpreter

Students who find a bug in the reference interpreter will receive \$20. You must be the first to make a public Piazza post demonstrating the bug.

Often you'll find an issue and not know whether you're doing something wrong or you've found a reference interpreter bug. Don't worry—just make a public Piazza post explaining the issue (you don't need to know you've found a reference interpreter bug to get credit for it).

5.4 Implementation Language

Because the skeleton is written in Java, it's natural to extend the skeleton to implement your interpreter in Java. However, you don't have to extend the skeleton, and you don't even have to use Java—you can use C/C++ or Rust if you like (using any other language requires prior instructor approval).

If you're interested in porting the skeleton to C/C++, note that JFlex and CUP have close equivalents for C/C++: The JFlex manual (https://jflex.de/manual.html; "Porting from lex/flex") says that the input file Scanner.jflex is similar in format to the format expected by the C/C++ flex tool. Likewise, Java CUP is based on the C/C++ tool YACC, and they use similar input file formats. So you can probably port Scanner.jflex to flex, and Parser.cup file to YACC.

6 Grading and Submitting Your Interpreter

6.1 Generating a Submission

To generate a file containing your interpreter, run

(cd skeleton; make submission)

which will run make clean and generate myproject.tgz.

6.2 The Grading Script and Test Cases

grade.sh SUBMISSION_TGZ REF_IMPL TESTCASE_LIST TESTCASE_DIR

where

- SUBMISSION_TGZ is the .tgz being submitted
- REF_IMPL is the Quandary reference interpreter script
- TESTCASE_LIST is a file that specifies a list of test cases; each test case is on its own line and has the following format:

POINTS PROGRAM INPUT

where

- POINTS is the number of points the test case is worth
- PROGRAM is the file containing the Quandary program (must be located in TESTCASE_DIR)
- INPUT is the integer input to the program
- TESTCASE_DIR is the location of the program files listed in TESTCASE_LIST

Here's an example:

grading/grade.sh skeleton/myproject.tgz ref/quandary grading/calc-public.dat examples

The public test case files (*-public.dat) contain example test cases for each project. The point values are arbitrary/meaningless.

TODO: Note that isrefint.q should FAIL.

6.3 Submitting Your Project

Upload your .tgz to Carmen by the deadline.

7 Quandary Subsets (i.e., Quandary Projects)

Add some interesting variants

8 Implementation Suggestions and Hints

See last semester's Piazza posts

9 Quandary Language and Runtime Specification

9.1 Syntax and Semantics

Colors denote productions used only for heap (including the Q and Ref types), concurrency, and mutation. Later, the list of built-in functions uses the same color coding.

```
\langle program \rangle ::= \langle funcDefList \rangle
\langle \mathit{funcDefList} \rangle ::= \langle \mathit{funcDef} \rangle \ \langle \mathit{funcDefList} \rangle
\langle funcDef \rangle ::= \langle varDecl \rangle \ (\langle formalDeclList \rangle) \ \{\langle stmtList \rangle\}
\langle varDecl \rangle ::= \langle type \rangle IDENT
                                                                                   // Variables and functions are immutable by default
                 | mutable \(\langle type \rangle \) IDENT // Mutable vars can be updated; mutable funcs can perform updates
                                                                                                                               // 64-bit signed integer
\langle type \rangle ::= int
                                                     // Reference to a heap object with left and right fields of type Q; or nil
            Ref
                                                                                                                      // Super type of int and Ref
\langle formalDeclList \rangle ::= \langle neFormalDeclList \rangle
\langle neFormalDeclList \rangle ::= \langle varDecl \rangle, \langle neFormalDeclList \rangle
                               |\langle varDecl \rangle|
\langle stmtList \rangle ::= \langle stmt \rangle \langle stmtList \rangle
\langle stmt \rangle ::= \langle varDecl \rangle = \langle expr \rangle;
                                                                                                                 // Declare and initialize variable
            | IDENT = \langle expr \rangle;
                                                                 // Update to already-declared-and-initialized (mutable) variable
                if ( \langle cond \rangle ) \langle stmt \rangle
                if (\langle cond \rangle) \langle stmt \rangle else \langle stmt \rangle
                while (\langle cond \rangle) \langle stmt \rangle
                                                                                                                      // Pointless without mutation
                IDENT ( \langle exprList \rangle );
                                                                                                             // IDENT must be mutable function
               free \langle expr \rangle;
                                                                          // Frees memory iff explicit memory management enabled
                print \langle expr \rangle;
                                                                                           // Prints evaluated value followed by a newline
                return \langle expr \rangle;
                \{ \langle stmtList \rangle \}
\langle exprList \rangle ::= \langle neExprList \rangle
\langle neExprList \rangle ::= \langle expr \rangle , \langle neExprList \rangle
                      |\langle expr\rangle|
```

```
// Special constant value of type Ref
\langle expr \rangle ::= nil
                                                                                                                                 // 64-bit signed integer of type int
                   INTCONST
                   IDENT
                   -\langle expr \rangle
                                                                                                    // Need for explicit downcast from {\tt Q} to {\tt int} or {\tt Ref}
                  ( \langle type \rangle ) \langle expr \rangle
                   IDENT ( \langle exprList \rangle )
                   \langle binaryExpr \rangle
                                                           // Evaluates the left and right sides of the binary expression concurrently
                   [ \langle binaryExpr \rangle ]
               | (\langle expr \rangle)
\langle binaryExpr \rangle ::= \langle expr \rangle + \langle expr \rangle
                           |\langle expr \rangle - \langle expr \rangle
                             \langle expr \rangle * \langle expr \rangle
\langle expr \rangle . \langle expr \rangle
                                                                                                  // Evaluates to a Ref referencing a new heap object
\langle cond \rangle ::= \langle expr \rangle \leq \langle expr \rangle
                    \langle expr \rangle >= \langle expr \rangle
                                                                                                                                    // For comparing int values only
                    \langle expr \rangle == \langle expr \rangle
                    \langle expr \rangle != \langle expr \rangle
                                                                                                                                    // For comparing int values only
                    \langle expr \rangle < \langle expr \rangle
                    \langle expr \rangle > \langle expr \rangle
                    \langle cond \rangle && \langle cond \rangle
                    \langle cond \rangle \mid \mid \langle cond \rangle
                   ! \langle cond \rangle
                    (\langle cond \rangle)
```

Lexical analysis: An IDENT is a sequence of letters, digits, and underscores such that the first character is not a digit.

If an INTCONST exceeds the bounds of a 64-bit signed integer, the interpreter's behavior is undefined.

Quandary's syntax is case sensitive.

Quandary allows Java/C/C++-style "block" comments /* like this */

9.2 Precedence and dangling else

Precedence of operators in high-to-low order:

```
    Expressions in parentheses (()) or brackets ([])
    - used as a unary operator and ( \langle type \rangle ) (cast operator)
    *
    - used as a binary operator and +
    .
    <=, >=, ==, !=, <, and >
    !
    && and | |
```

All operators are left assocative.

Dangling else ambiguity is resolved by matching an else with the nearest if statement allowed by the grammar.

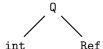
9.3 Static type checking

The Quandary interpreter checks the following rules prior to executing the program.

Declarations: A program must not define a function with the same name as another function, including the built-in functions. A program must only call functions defined in the program or built-in functions. A program must define a function named main that takes a single argument of type int.

A function must not declare a variable with the same name as a variable that has been defined earlier (including as a parameter) in the same or an outer/containing lexical scope (demarcated by curly braces, i.e., {}, or by being a single conditional statement in an if/else statement or while loop). An expression may only access variables declared within the same or an outer/containing lexical scope. Thus for any variable name v at any program point, either v can be accessed or declared, but never both.

Types and conversions: All $\langle expr \rangle$ evaluation—including function actuals, return values, and free statements—must be statically type-checked as much as possible, according to the following type hierarchy:



Upcasts and same-casts are permitted to be either implicit (silent) or explicit. Downcasts must be explicit, i.e., $|\langle expr \rangle ::= (\langle type \rangle) \langle expr \rangle|$, and are checked at run time.

Function calls must have the same number of actuals as the function definition's number of formals.

Immutability: Variables and functions are *immutable* unless declared as mutable. An immutable variable must not be the assigned-to variable in an assignment statement, i.e., statement (stmt) := IDENT = statement (stmt);

An immutable function's body must not contain calls to mutable functions (including built-in mutable functions).

A call *statement* may only call a mutable function.

Miscellaneous: Every function must be statically guaranteed to return a value. The interpreter's static checking may verify this property by simply checking that the function's last statement is a **return** statement (and reporting an error if not).

A function may contain return statements that make code statically unreachable. In general, statically unreachable code is not erroneous.

9.4 Built-in functions

Q left(Ref r) - Returns the left field of the object referenced by r

Q right(Ref r) - Returns the right field of the object referenced by r

int isAtom(Q x) - Returns 1 if x's value is nil or an int; returns 0 otherwise (if x's value is a non-nil Ref)

int isNil(Q x) - Returns 1 if x's value is nil; returns 0 otherwise (if x's value is an int or a nonnil Ref)

mutable int setLeft(Ref r, Q value) - Sets the left field of the object referenced by r to value, and returns 1

mutable int setRight(Ref r, Q value) - Sets the right field of the object referenced by r to value,
and returns 1

```
mutable int acq(Ref r) - Acquires the lock of the object referenced by r and returns 1
```

mutable int rel(Ref r) - Releases the lock of the object referenced by r and returns 1

int randomInt(int n) – Returns a random int in [0, n)

9.5 Language semantics (including dynamic type checking) and operation of the interpreter

The interpreter executes the defined function called main and passes a command-line parameter as main's argument:

```
$ ./quandary
Expected format: quandary [OPTIONS] QUANDARY_PROGRAM_FILE INTEGER_ARGUMENT
Options:
   -gc (MarkSweep|MarkSweepVerbose|RefCount|Explicit|NoGC)
   -heapsize BYTES
BYTES must be a multiple of the word size (8)
```

The interpreter prints the return value of main in the following way:

```
Interpreter returned ((5 . nil) . (-87 . (9 . 3)))
```

Incorrect command-line parameters, including QUANDARY_PROGRAM_FILE not being found, have undefined behavior (i.e., the interpreter may fail in any way).

Function calls: Function call semantics are pass-by-value.

Order of evaluation: The interpreter evaluates expressions in left-to-right order, i.e., it evaluates the left side of (non-concurrent) binary expressions before the right side, and it evaluates function call actual expressions in left-to-right order.

Binary boolean operators (&& and ||) use short-circuit evaluation.

Dynamic type checking: The interpreter should check executed type downcasts at run time and report a fatal dynamic type checking error on a type downcast failure.

Heap mutation: A new heap object's left and right fields are each initialized to an int or Ref value, and must remain as either an int or Ref value, respectively, for the duration of the execution. Thus the interpreter should fail with a dynamic type checking error if the setLeft() or setRight() function attempts to overwrite an int slot with a Ref value, or a Ref slot with an int value.

The purpose of this restriction is to avoid the implementation challenge of updating both the value and associated type metadata atomically (which is an issue if implementing objects using "raw" memory).

nil dereference: Calling left(), right(), setLeft(), setRight(), acq(), or rel() with a first argument evaluating to nil should cause a fatal nil dereference error at run time.

Memory management: An execution should report an "out of memory" error if and only if the non-freed memory exceeds the specified maximum heap size.

The interpreter potentially supports explicit memory management and mark–sweep and reference counting garbage collection (and optionally others as well, e.g., semi-space). See command-line arguments above.

Explicit memory management only: An execution that accesses a freed object has undefined semantics. An execution that performs double-free on a reference has undefined semantics. An execution that tries to free nil has undefined semantics.

Trace-based garbage collection only: An evaluation of an allocation expression ($\langle binaryExpr \rangle ::= \langle expr \rangle$) performs trace-based GC when and only when the non-freed memory exceeds the specified maximum heap size. Trace-based GC frees objects that are transitively unreachable from the roots (functions' local variables and intermediate values). Implementing support for stopping multiple threads at GC-safe points is not required; if trace-based GC is triggered when multiple threads are active, the interpreter has undefined behavior (but ideally it will report an error, to help with debugging).

Concurrency: A concurrently evaluated binary expression [$\langle binaryExpr \rangle$] evaluates the left and right child expressions in two new concurrent threads (i.e., thread fork), and waits for both threads to finish (i.e., thread join). Every thread that is not blocked eventually makes progress.

Thread fork and join and lock acquire and release are synchronization operations that induce happens-before edges. Conflicting accesses unordered by happens-before constitute a data race.

An execution of a program with a data race has undefined semantics. An execution in which a thread performs a rel() of a lock it does not hold, has undefined semantics.

Error checking: To help with grading, the interpreter *process* should return one of the following error codes as appropriate:

- 0 success
- 1 lexical analysis or parsing error
- 2 static checking error
- 3 dynamic type checking error
- 4 nil dereference error
- 5 Quandary heap out-of-memory error

The interpreter script (quandary) should print this return code. Specifically, the script should handle executions as follows.

1. For a non-erroneous, terminated program, the script should print the following as its last two lines:

Interpreter returned RETURN_VALUE_OF_MAIN Quandary process returned 0

where RETURN_VALUE_OF_MAIN is return value of the Quandary program's main function.

Printing anything or nothing before that is fine.

- 2. For a non-erroneous, non-terminating execution (e.g., a program execution with an infinite loop),¹ the script should not terminate. Printing anything or nothing is fine.
- 3. For an execution that should return error code ERROR_CODE, the script should print the following as its last line:

¹Execution with unbounded call depth has undefined behavior.

Quandary process returned ERROR_CODE

Printing anything or nothing before that is fine.

4. For an execution that has undefined behavior, any behavior and output is fine (including uncaught exceptions). An interpreter can safely assume that programs and inputs with undefined behavior will not be executed.

If the *interpreter program itself* runs out of stack memory, runs out of heap memory, or allocates too many threads, then behavior is undefined (any behavior is acceptable). For a reasonable Quandary input program, the interpreter should succeed if given enough stack memory, heap memory, and thread count limit.

9.6 Implementing the interpreter

An interpreter written in Java or C++ should allocate heap objects into raw memory (represented by a primitive array in Java, for example), and assume that raw memory provides only low-level load, store, and compare-and-set operations. When writing the interpreter, use the provided Heap class to emulate raw memory.