

CS330 Assignment 3

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Condition variable

Defined struct `cond_t` which includes a name and sleeplock as variables. Used this struct for implementing condition variables. Implemented `cond_init`, `cond_wait`, `cond_signal` and `cond_broadcast` functions. `Cond_wait` used `condsleep` function written in `proc.c`, while `cond.broadcast` used a predefined function `wakeup` and `cond_signal` used an edited version `wakeupone` of the same `wakeup` function. Spinlock in `wakeup` is converted to sleeplock so that process sleeps while it has to wait.

Semaphore

Similarly a struct `semaphore` is defined that contains an integer variable to store the semaphore value, a condition variable and sleeplock. `Semaphore.c` contains `sem_init`, `sem_wait` and `sem_post`. `Sem_init` initialises condition variable by `cond_init`, and initialises sleeplock using `initsleeplock` functions. The semaphore value is initialised by input parameter. `Sem_wait` allows semaphore number of processes to run and while next processes waits until `sem_post` is called by any of the process. `Sem_post` increments the value by 1 and wakes up a sleeping process using `cond_signal`.

System calls

1 Barrier

We have used an array of condition variable to implement Barrier. The barrier array is an integer array which stores -1 when not initialised or freed, and number of processes called after initialised.

1.1 `barrier_alloc`

Initialises the first free barrier and initialises its value 0 from -1, initialise corresponding condition variable and returns the id of the barrier that is allocated.

1.2 `barrier`

Uses a global lock for printing so that no other process can print when 1 process is printing may be that process be of different barrier id. Uses the sleep lock of condition variable so that a process of same barrier waits until all the process of the same barrier enters first.

1.3 `barrier_free`

Sets the value of freed barrier again to -1.

2 Condition Producer Consumer

Created `buffer_elem` struct having 2 condition variables insert and delete. A sleeplock, an integer storing value and a flag named full.

2.1 `buffer_cond_init`

Initialises tail and head to 0.also initialises inser, delete, print sleeplock that is common for all processes. Initialises the values of the `buffer_elem`.

2.2 `cond_produce`

Producer produces if the `buffer_elem` that the tail points to is empty, and while producing uses inserlock of that particular `buffer_elem`.

2.3 `cond_consume`

Similar process as of producer but vice versa.

3 Semaphore Producer Consumer

Same function as of contion producer and consumer using emaphores.

3.1 `buffer_sem_init`

Initialises `nextp`, `nextc` the producer and consumer pointers to 0. Initialises 4 semaphoers namely - `sem_prod=0`, `sem_cons=0`, `sem_empty=size of buffer`, `sem_full=0`

3.2 `sem_produce`

Produces if the buffer element is empty using `sem_prod` semaphore. Every time `nextp` is incremented cyclically. checking empty is done by semaphore and posted full when produced.

3.3 `sem_consume`

Produces if the buffer element is empty using `sem_prod` semaphore. Method similar to producer just viceversa.

Observation

Semaphore takes more time to implement as the number increases.