

Bootstrapping of Commodity Forward Prices

Jake C. Fowler

February 2023

THIS DOCUMENT IS CURRENTLY WORK IN PROGRESS

1 Introduction

This document presents an algorithm for the bootstrapping commodity forwards, futures and swap contracts.

It is applicable for linear commodity contracts contingent on either a continuous delivery of physical commodity, financial payoff based on the average of an index throughout the contract delivery period.

The rest of this document will refer to just forward prices and curves, even though the same methodology can be applied to futures and swap instruments as well.

Also the document refers to the *delivery period* of the contracts, in place of *delivery or fixing period* where *fixing period* is the set of time periods from which the average price of an index is used to calculate the payoff of a swap contract.

1.1 Bootstrapping Definition and Uses

The process of bootstrapping is to take a set of prices of forward and derive a piecewise flat forward curve consistent with these prices. In the typical case, the delivery periods of the input forward contracts will be overlapping, so bootstrapping can be seen as an algorithm for producing a set of forward prices for contiguous delivery periods. In this way, bootstrapping should in the most general case be seen as an interpolation method.

One example of the use of a bootstrapped curve is for it to be loaded into an ETRM (Energy Trading Risk Management) system for the valuation of a portfolio of vanilla trades. The bootstrapping creates an unambiguous price for each delivery period in the lowest granularity which a commodity trades, which has practical benefits over just using the raw forward prices of traded contracts.

A piecewise flat bootstrapped curve can also be used for the conservative valuation of a physical asset with embedded optionality.

Finally bootstrapping can be used to derive the inputs to an interpolation method which is then used to calculate a smoothed higher granularity forward curve.

2 No-Arbitrage Forward Price Condition

This section describes the no-arbitrage relationship between the set of prices of forward contracts for a consecutive delivery period and price of a single lower granularity contract which spans the whole delivery period of these. It is this relationship which forms the basis of the bootstrapping algorithm.

Denote:

- F as a forward price of the *big* contract. The notation does not contain the concept of the time when this is observed, but it can just be assumed that the current valuation date is earlier than t_s the start of the delivery period.
- f_i as the forward price for the i th *small* contract a higher granularity period than F has, with $i = 1 \dots n$
- D_i as the discount factor from the settlement date of contract f_i to the valuation date. We assume that the delivery period of f_i is short enough for the whole contract to be settled on a single date.
- w_i is the weighting of the delivery period of contract i , relative to the delivery period that forward price F corresponds to.

To show the no-arbitrage relationship, construct a portfolio which is long one unit of the *big* contract and short w_i units of each *small* contract. This portfolio is essentially flat, does not require and upfront cash, and so should have zero PV. If the PV were positive then a long position would present an arbitrage, as would a negative PV presents an arbitrage from taking a short position in the portfolio. Expressing the PV as the discounted cash flow from take the portfolio to settlement:

$$\sum_{i=1}^n (F - f_i) w_i D_i = 0 \quad (1)$$

Which can be rearranged to:

$$F = \frac{\sum_{i=1}^n f_i w_i D_i}{\sum_{i=1}^n w_i D_i} \quad (2)$$

This is the no-arbitrage relationship between the *big* contract and its associated *small* contracts, and can be seen as a weighted average.

3 Bootstrapping Algorithm

3.1 Linear System

It is important to note for later discussion in this paper that the above solution will be *least squares* in two different senses. Firstly, should matrix \mathbf{A} have rank less than m , as would be the case where there is some redundancy in the input contracts, $\hat{\mathbf{b}}$ will not necessarily be equal to the vector of input contracts \mathbf{b} . It will however be the vector in the column space of \mathbf{A} which has the lowest sum of element-wise squared difference to \mathbf{b} .

Secondly, should matrix \mathbf{A} have rank less than it's number of columns, as well typically be the case in practice, the linear system $\mathbf{Ax} = \mathbf{b}$ will have infinite solutions, and $\hat{\mathbf{x}}$ will be chosen as the smallest solution, i.e. the one with the lowest sum of squared elements. At first it seems that this second *least squares* property is a desirable one, given we are looking for a piecewise flat output curve. Think of the case of bootstrapping a single contract to the forward prices of the periods within it's delivery period. A vector with the same price for all sub-periods as the single input contract price will be also equal to the vector of prices with lowest sum of squared elements. The section below however shows that the least squares solution is not necessarily a good one for the bootstrapping problem, and introduces an adjustment which improves the results.

3.2 The Problem of Partially Overlapping Contracts

The bootstrapping algorithm presented in the previous section has some cases, where it does not perform well. Take an example of two input contracts, the first with delivery covering the first three periods: t_1, t_2, t_3 . The second contract has delivery covering two periods: t_3, t_4 . If the prices of these two contracts is both 10.0, then the vector of bootstrapped price is calculated as $[8.0 \ 8.0 \ 14.0 \ 6.0]^T$. Intuitively this result comes as a surprise, one would expect that the bootstrapped price of all four periods will equal 10.0, the price of both input contracts.

It can easily be verified that these prices average back to the input contracts over just the delivery periods of each input forward. The next check that these results are not erroneous is to calculate the euclidean length of the bootstrapped price vector and compare it to the length of the intuitively expected flat price vector of 4 10's. The actual bootstrapped vector is of length 360.0, which is indeed lower than the length of flat price vector of 400.0. Where matrix \mathbf{A} is underdetermined, and hence has infinite solutions, we know that the bootstrapper will return the single solution with the smallest length. Although the above check does not confirm that there has not been a mistake (i.e. a smaller solution vector exists), it does bring into question whether the criteria of returning the price vector of lowest length is indeed desirable.

The weighting matrix is as follows:

$$\mathbf{A} = \begin{bmatrix} \frac{1}{3} & \frac{1}{3} & \frac{1}{3} & 0 \\ 0 & 0 & 0.5 & 0.5 \end{bmatrix} \quad (3)$$

With pseudo-inverse:

$$\mathbf{A}^+ = \begin{bmatrix} 1.2 & -0.4 \\ 1.2 & -0.4 \\ 0.6 & 0.8 \\ -0.6 & 1.2 \end{bmatrix} \quad (4)$$

To verify that the problem seen above will also be the case with an arbitrary forward price of the input contracts equal to c , we can calculate the derived bootstrapped prices algebraically:

$$\mathbf{x} = \mathbf{A}^+ \mathbf{b} = \begin{bmatrix} 0.8c \\ 0.8c \\ 1.4c \\ 0.6c \end{bmatrix} \quad (5)$$

This confirms that whatever the prices of the two input forward contracts, the bootstrapped contracts will not be constant and equal to the input contract price.

The above example is contrived and simplified for ease of explanation. Cases seen with real market data being bootstrapped with a similar structure, where two forward contract one another, but neither is completely overlapped by the other. In such cases the input forward prices are not all equal, so the identification of the bootstrapper results being problem is not the same as the above. example. The problem seen with real market data is that sections of the bootstrapped curve vary wildly from the input prices.

Such cases are referred to in this rest of this paper as *partially overlapping contracts* and an examples seen in real market data when a contract for a one week delivery partially overlaps with a month.

As such, *partially overlapping contracts* should not merely be thought of as edge cases, and the bootstrapping algorithm needs to be made robust to such cases.

3.3 Introduction of Nullspace Component to Solution

The previous section has shown an example where the least squares solution does not give one we desire, due to a deviation from the prices of the input contracts where partially overlapping contracts are present. To rectify this, we can note that in the problematic scenario, and in the typical case of running the bootstrapping algorithm, the rank of matrix \mathbf{A} will be less than it's number of columns n , the number of bootstrapped prices. In such cases there will be infinite solutions, and a general form of these solutions found using the fact that any vector in the nullspace of \mathbf{A} can be added to any particular solution.

If the solution which is closest to $\mathbf{0}$ is not the good, then we need to define a new vector \mathbf{x}^{target} , the solution which should be closest to will be choosen. In the case of the problematic example where all input contracts have the same price, then \mathbf{x}^{target} with all elements equal to this price seems like the obvious choice. See the section below for a discussion on choosing \mathbf{x}^{target} .

Defined \mathbf{K} as a matrix with columns containing the orthonormal basis of the nullspace of \mathbf{A} , and error vector as $\mathbf{e} = \mathbf{x}^{target} - \mathbf{x}$. We want to find a linear combination of the columns of \mathbf{K} which can be added to \mathbf{x} to move it closest to \mathbf{x}^{target} . This corresponds to find the least-squares solution of \mathbf{c} to the following system:

$$\mathbf{K}\mathbf{c} = \mathbf{e} \quad (6)$$

The least-squares solution can be found be solving the Normal Equations:

$$\mathbf{K}^T \mathbf{K} \mathbf{c} = \mathbf{K}^T \mathbf{e} \quad (7)$$

As \mathbf{K} is orthonormal $\mathbf{K}^T \mathbf{K} = \mathbf{I}$, hence \mathbf{c} is calculated as:

$$\mathbf{c} = \mathbf{K}^T \mathbf{e} \quad (8)$$

Or, substituting for \mathbf{e} and multiplying out the \mathbf{K}^T term:

$$\mathbf{c} = \mathbf{K}^T \mathbf{x}^{target} - \mathbf{K}^T \mathbf{x} \quad (9)$$

\mathbf{x} is in the Row Space of \mathbf{A} (see Strang (1998)), hence orthogonal to any vector in the nullspace of \mathbf{A} , so $\mathbf{K}^T \mathbf{x} = \mathbf{0}$ and the expression for \mathbf{c} can be simplified:

$$\mathbf{c} = \mathbf{K}^T \mathbf{x}^{target} \quad (10)$$

Using this result, our nullspace adjusted bootstrapped prices can be written as

$$\mathbf{x}^* = \mathbf{x} + \mathbf{K}\mathbf{c} \quad (11)$$

Substituting for \mathbf{x} and \mathbf{c} :

$$\mathbf{x}^* = \mathbf{A}^+ \mathbf{b} + \mathbf{K} \mathbf{K}^T \mathbf{x}^{target} \quad (12)$$

3.4 Example of Using Nullspace Component

This section present

3.5 Choosing The Target Vector

In the case where all input forward prices are the same, choosing \mathbf{x}^{target} is easy, but what about the more realistic case when all prices are different? This is somewhat arbitrary, but intuition says that each element of \mathbf{x}^{target} should be a function of the prices of all input contracts which span the delivery period each element of \mathbf{x}^{target} represents the bootstrapped price of. Some experimentation has shown that taking the price of contract with minimum length of delivery period yields good results. In the case where there is more than one contract has the same minimum length, which in practice would be unusual, then the price of the contract which is delivers earliest should be used.

4 Future Work

More research into choice of the target vector can be done, including an analysis of both the intuitive and mathematical reasoning behind the logic for calculating \mathbf{x}^{target} presented in this paper, as well as alternative strategies. The results of varying \mathbf{x}^{target} should be analysed with real market data to assess usage in practice.

The whole approach to bootstrapping in this paper should also be compared to algorithms which perform both bootstrapping and interpolation, such as Benth et al. (2007). In the case of partially overlapping contracts, the inclusion of the nullspace component and target vector presented in this paper is used to impose a shape on the resulting bootstrapped prices, but one could argue that using a spline is more suitable due to the smooth structure of a spline giving a more desirable shape. A piecewise flat curve could then be averaged back from the result of the spline, and the use of a bootstrapper to create the input to another spline algorithm is no longer needed. Two downside of using a combined bootstrapping and interpolation algorithm can be seen. The first is that introduces additional complication in the case where only a piecewise flat curve is required from the bootstrapping. The second is that a higher order polynomial is required to ensure that the number of unknowns (polynomial coefficients) is at least as large as the number of constraints in the linear system solved in the spline calculation. This is because the number of piecewise polynomials being solved for can vary when the input contracts are allowed to overlap. The higher the order of the piecewise polynomials, the more tendency it has to oscillate. The author has observed undesirable oscillations when using a 4th order polynomial, as used in Benth et al. (2007), when interpolating natural gas and power forward curves. In practice it is unlikely that a polynomial should change it's slope more than once when interpolating commodity forward curve data.

The handling of shaping constraints which are collinear with the forward price constraints (and prior shaping constraints) in an efficient manner is an area which requires more work. For example if the forward prices for Q1-23,

Jan-23, and Feb-23 are all present, then the shaping ratio or spread from Q1-2023 to Mar-23 is redundant because the 3 prices already specify the Mar-23 forward price. Mathematically, the addition of the shaping constraint row to matrix \mathbf{A} will not change the rank of \mathbf{A} . In such situations the shaping row will never be required and ideally the bootstrapping algorithm should detect this and not append the shaping constraint row to \mathbf{A} . The calculation could involve fully recalculating the rank of \mathbf{A} after each constraint row has been appended, or whether each new row is a linear combination of existing rows, but this could be expensive. There should be a more efficient algorithm which for each new row makes use of the calculations used to check previous rows. Some research into rank-revealing QR factorisations might yield an approach.

Similarly, checking whether forward price constraints are collinear with other forward price constraints could be desired, although the case for discarding redundant forward price constraints is less clear cut than it is for shaping constraints. The algorithm caller would have to be aware of the dependency of the results on the order in which the forward prices and shaping factors are fed in. To make the results deterministic, an ordering first step would be required.

Such handling of constraint collinearity would be required for a truly formulaic bootstrapping algorithm.

References

- F. E. Benth, S. Koekebakker, and F. Ollmar. Extracting and applying smooth forward curves from average-based commodity contracts with seasonal variation. *The Journal of Derivatives*, 15:52–66, 2007.
- Gilbert Strang. *Introduction to Linear Algebra: 2nd Edition*. Wellesley-Cambridge Press, 1998.