

**Lecture 16**

**Context Bounding Checkers I**

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# Huge Number of Thread Schedules

- ▶ Concurrent program with  $n$  threads where each thread has  $k$  instructions has

$$(n * k)! / (k!)^n \geq (n!)^k$$

interleavings

- ▶ Exponential in both  $n$  and  $k$ !
- ▶ Example: 5 threads with 5 instruction each

$$25! / 5!^5 = 6.2336074e+14$$

= 623 trillion interleavings

# Java Path Finder (JPF)

- ▶ Program checker for Java
- ▶ Properties to be verified
  - ▶ Program assertions
  - ▶ LTL properties
- ▶ Depth-first and breadth-first search, heuristics
  - ▶ Uses static analysis techniques to improve the efficiency of the search
- ▶ Requires a complete Java program
  - ▶ Cannot handle native code

# Combating State Space Explosion

- ▶ Symmetry reduction
  - ▶ Search equivalent states only once
- ▶ Partial order reduction
  - ▶ Do not search thread interleavings that generate equivalent behavior
- ▶ Static analyses
  - ▶ Reduce state space using static analyses
- ▶ User-provided restrictions
  - ▶ Manually bound variable domains, array sizes,...

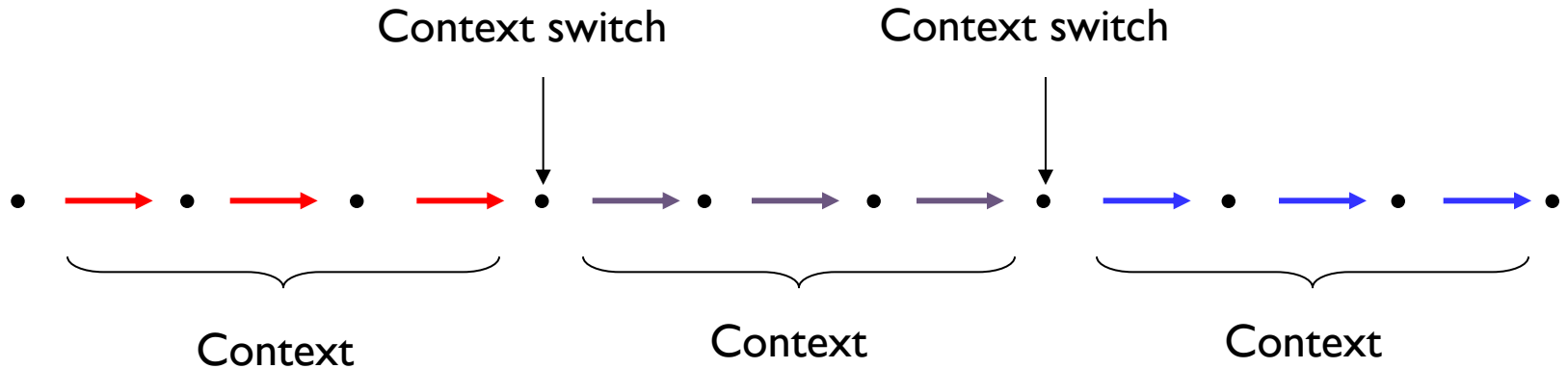
# This Time

- ▶ Context-bounded verification of concurrent programs
- ▶ Sequentialization of concurrent programs

# Context-Bounded Verification

slides acknowledgements: Shaz Qadeer, Madan Musuvathi

# Context-Bounded Verification



- ▶ Many subtle concurrency errors are manifested in executions with few context switches
- ▶ Analyze all executions with few context switches

# Context-Bounded Reachability Problem

- ▶ An execution is  $c$ -bounded if every thread has at most  $c$  contexts
- ▶ Does there exist a  $c$ -bounded execution from a state  $S$  to a state  $E$ ?



# CB Reachability is NP-Complete

## ► Membership in NP

- Witness is an initial state and  $n \cdot c$  sequences each of length at most  $|L \times G|$

$n$  = # of threads,  $c$  = # of contexts

$L$  = # of program locations,  $G$  = # of global states

## ► NP-hardness

- Reduction from the CIRCUIT-SAT problem

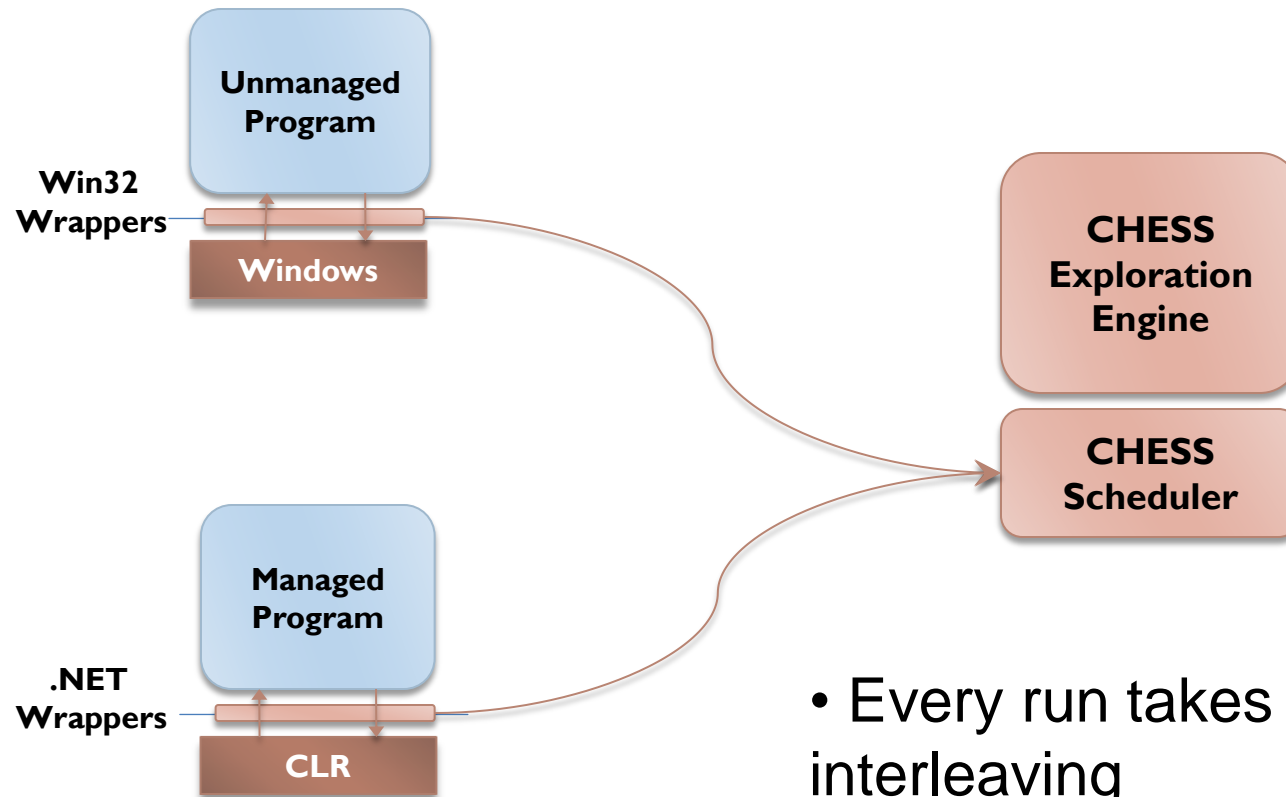
# Complexity of Safety Verification

	Unbounded	Context-bounded
Finite-state systems	PSPACE complete	NP-complete
Pushdown systems	Undecidable	NP-complete

# CHESS: Systematic Testing for Concurrency

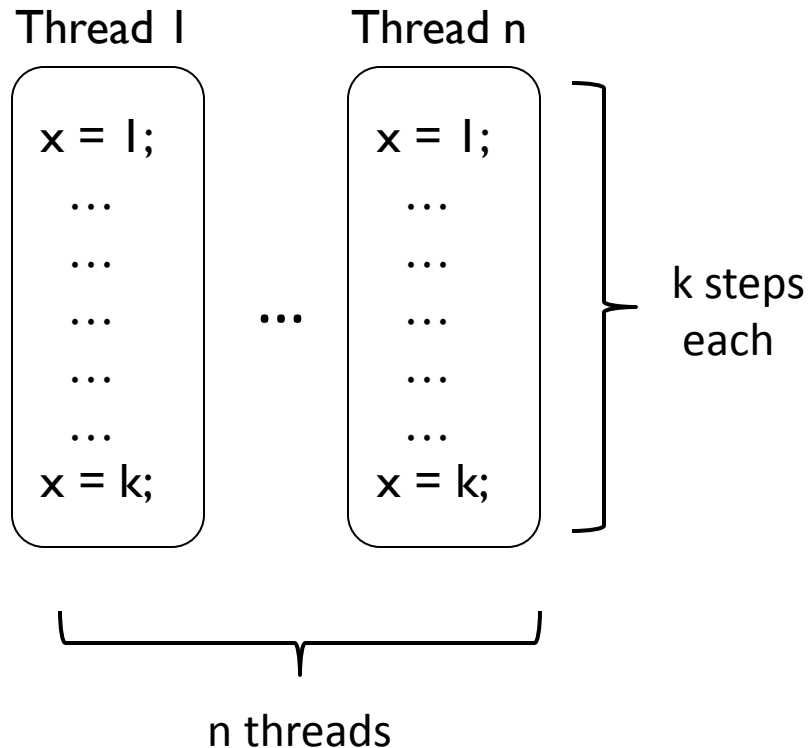
- ▶ CHESS is a user-mode scheduler
- ▶ Controls all scheduling nondeterminism
  - ▶ Replace the OS scheduler
- ▶ Guarantees:
  - ▶ Every program run takes a different thread interleaving
  - ▶ Reproduce the interleaving for every run

# CHESS Architecture



- Every run takes a different interleaving
- Reproduce the interleaving for every run

# State-Space Explosion

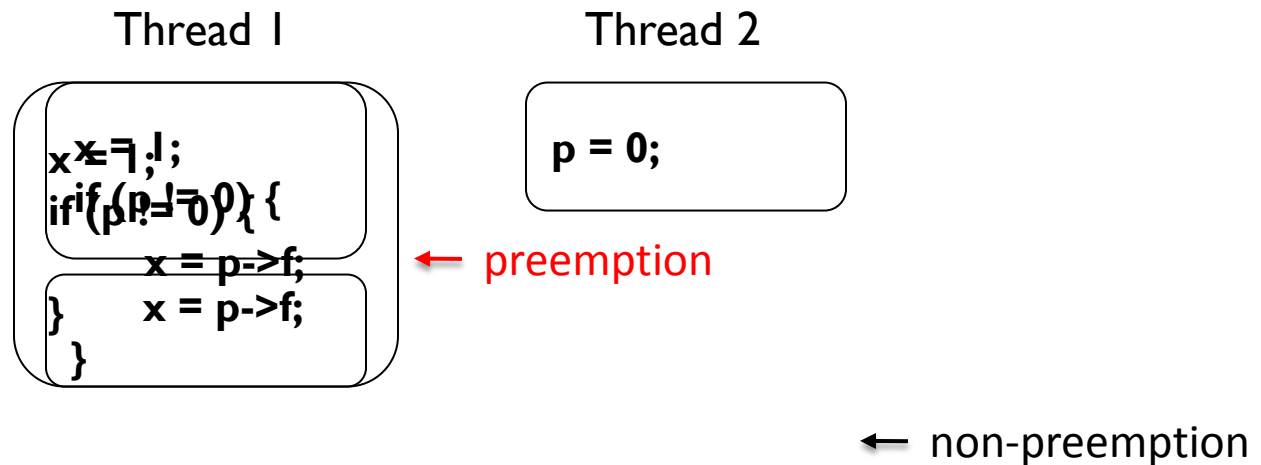


- ▶ Number of executions is  $O(n^{nk})$
- ▶ Exponential in both n and k
  - ▶ Typically:  $n < 10$ ,  $k > 1000$
- ▶ Limits scalability to large programs

**Goal:** Scale CHESSE to large programs (large k)

# Preemption-Bounding

- ▶ Prioritize executions with small # of preemptions
- ▶ Two kinds of context switches:
  - ▶ Preemptions – forced by the scheduler
    - ▶ E.g., time-slice expiration
  - ▶ Non-preemptions – a thread voluntarily yields
    - ▶ E.g., blocking on an unavailable lock, thread end

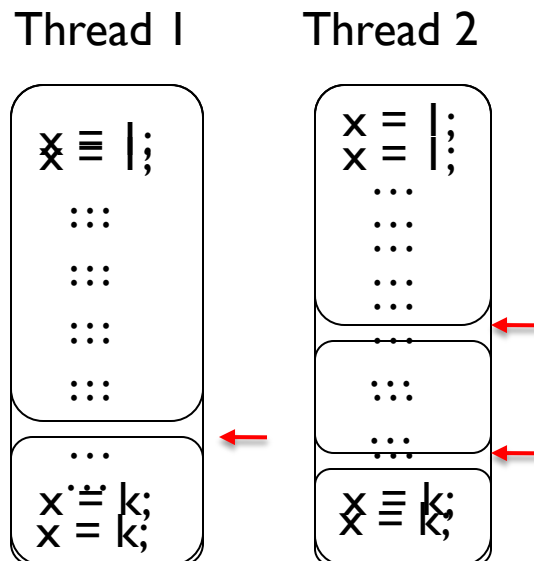


# Preemption-Bounding in CHESS

- ▶ The scheduler has a budget of  $c$  preemptions
  - ▶ Nondeterministically choose the preemption points
- ▶ Resort to non-preemptive scheduling after  $c$  preemptions
- ▶ Once all executions explored with  $c$  preemptions
  - ▶ Try with  $c+1$  preemptions

# Property 1: Polynomial Bound

- ▶ Terminating program with fixed inputs and deterministic threads
  - ▶  $n$  threads,  $k$  steps each,  $c$  preemptions
- ▶ Number of executions  $\leq {}_{nk}C_c * (n+c)!$   
 $= O((n^2k)^c * n!)$
- ▶ Exponential in  $n$  and  $c$ , but not in  $k$ !



- Choose  $c$  preemption points
- Permute  $n+c$  atomic blocks



## Property 2: Simple Error Traces

- ▶ Finds smallest number of preemptions to the error
- ▶ Number of preemptions better metric of error complexity than execution length

## Property 3: Coverage Metric

- ▶ If search terminates with preemption-bound of  $c$ , then any remaining error must require at least  $c+1$  preemptions
- ▶ Intuitive estimate for
  - ▶ The complexity of the bugs remaining in the program
  - ▶ The chance of their occurrence in practice

# Property 4: Many Bugs with Few Preemptions

Program	kLOC	Threads	Preemptions	Bugs
<b>Work-Stealing Queue</b>	1.3	3	2	3
<b>CDS</b>	6.2	3	2	1
<b>CCR</b>	9.3	3	2	2
<b>ConcRT</b>	16.5	4	3	4
<b>Dryad</b>	18.1	25	2	7
<b>APE</b>	18.9	4	2	4
<b>STM</b>	20.2	2	2	2
<b>PLINQ</b>	23.8	8	2	1
<b>TPL</b>	24.1	8	2	9

# Coverage vs Preemption-Bound

