

Let f, g be the functions defined as $f(n) = 10^3 n \log n$ and $g(n) = \frac{n^2}{10^5 \log n}$.

Scegli una o più alternative:

- ☐ $f \in \Omega(g)$
- ☐ $f \in \Theta(g)$
- ☒ $f \in O(g)$

Nondeterministic Turing Machines:

Scegli una o più alternative:

- ☐ Can be simulated by deterministic TMs.
- ☒ If working in polynomial time, can be used to characterize **NP**
- ☐ Always work in polynomial time
- ☐ Are essential to define the complexity class **NP**

The universal Turing machine:

Scegli una o più alternative:

- ☒ Can simulate every Turing machine, with a polynomial overhead.
- ☐ Can simulate every Turing machine, but not itself
- ☐ Works in polynomial time.
- ☒ Is an essential ingredient of in the proof of existence of uncomputable problems.

Suppose a language \mathcal{L} is in **EXP** but not in **P**. Then:

Scegli una o più alternative:

- ☐ \mathcal{L} is necessarily **NP**-complete.
- ☒ The classes **NP** and **P** are different.
- ☒ There could be a nondeterministic polytime TM computing \mathcal{L}
- ☐ \mathcal{L} can be computed in polynomial time.

The notion of PAC-learnable concept class:

Scegli una o più alternative:

- ☐ Requires the output concept to have probability of error ε , in all cases
- ☒ Does not make any reference to the time complexity of the learning algorithm
- ☒ Needs to hold for every distribution **D** on the instance class.
- ☐ Cannot be reached when the underlying concept class is the one conjunctions of literals.

Problems by memory:

1. give a TM to decide $L = \text{set of strings for which if } 01 \text{ is present then it is followed by all zeroes}$
2. prove that the problem is in NP: check if a number is the sum of powers of 3 by giving a TM or pseudocode.
(asked to the professor, he said that 3^0 is not allowed as the problem would be trivial)
3. PP is the set of theorems expressed in the Principia Mathematica, published by Bertrand Russell in 1909-13. Do they fall in a complexity class? Motivate