

COMP4500

Assignment 1

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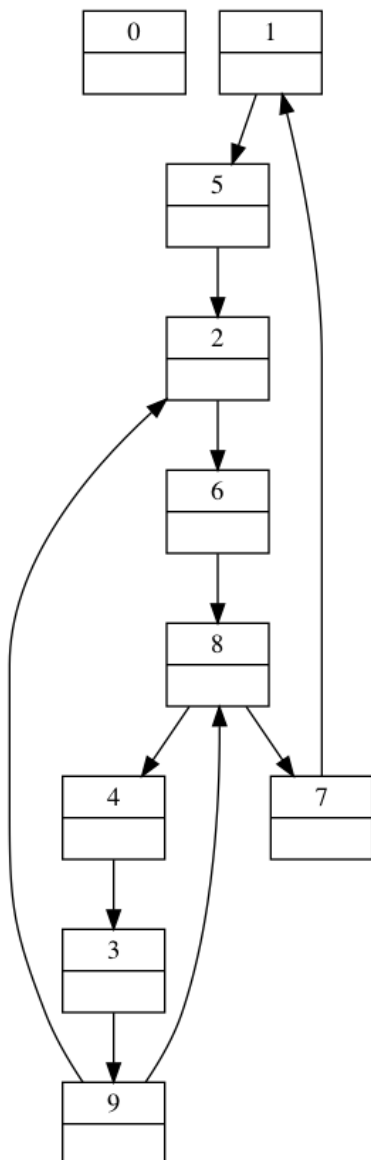
September 17, 2018

Question 1: Constructing SNI and directed graph

(a) Creating your SNI

My initial input number was 984392687152. My resulting SNI was the same, 984392687152.

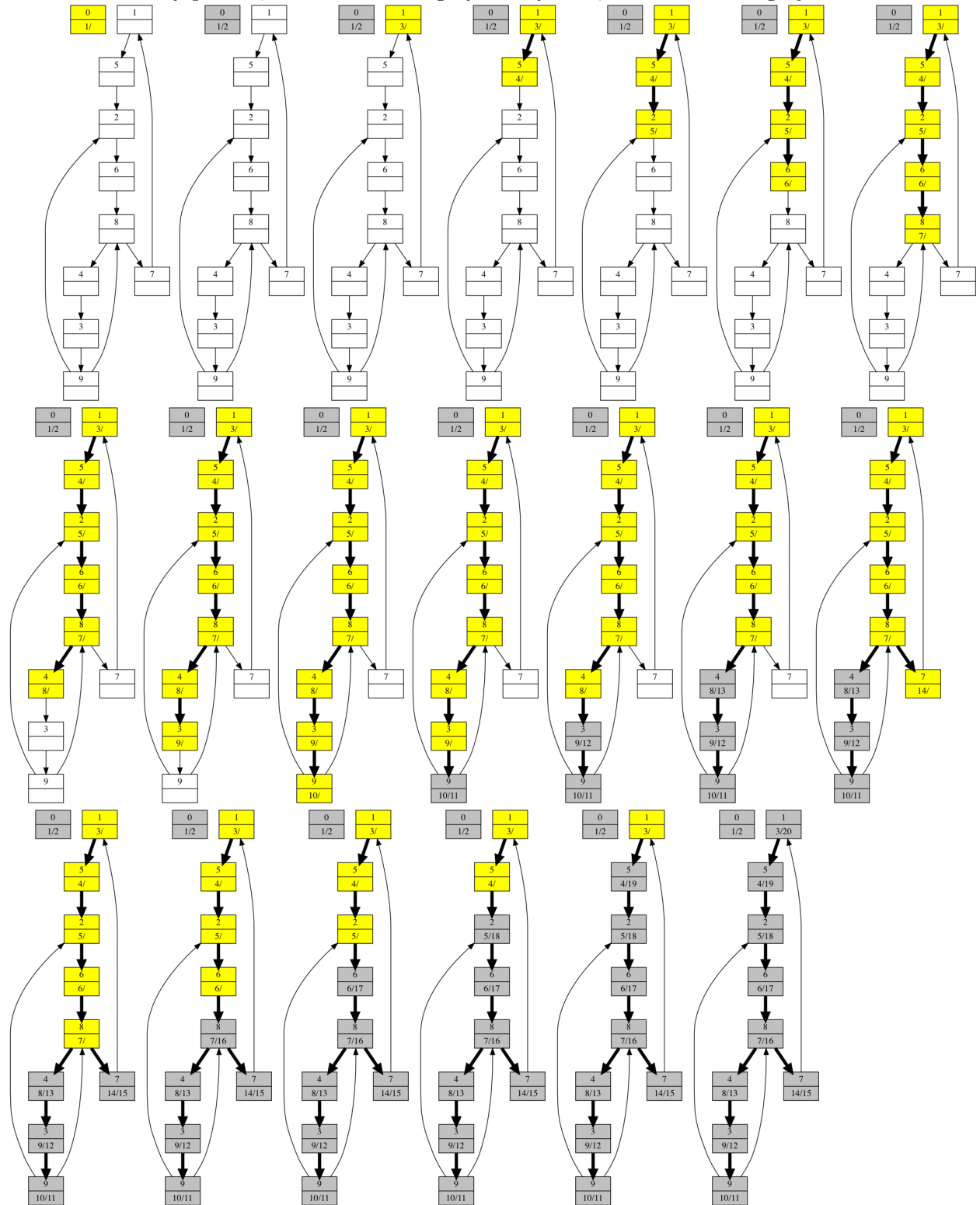
(b) Graph S



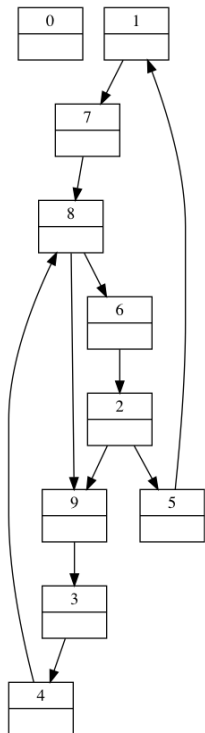
Question 2: Strongly connected components

(a) Step 1 of the SCC algorithm using S as input

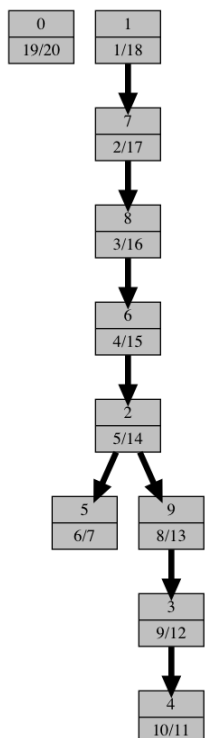
To be nicer to my printer, I've substituted gray with yellow, and black with gray.



(b) Step 2 of the SCC algorithm - S^T



(c) Step 3 and 4 of the SCC algorithm



The tree rooted with 1 was constructed first, and the tree rooted with 0 was constructed second.

Question 3: Design and implement a solution

Question 4: Worst-case complexity analysis

(a) and (b) combined

Rather than treating each \mathcal{P} as a node, and each \mathcal{D} as an edge in a graph, this algorithm treats each \mathcal{D} as a node.

A valid edge between two \mathcal{D} constitutes a two element sub-path of a possible route from P_s to P_d .

```
def findMinimumCost(locations, source, ts, destination, td, deliveries):
    sourceToDeliveries: HashMap<Location, HashSet<Delivery>> =
        new HashMap().onLookupFail(new HashSet())
    destinationToDeliveries: HashMap<Location, HashSet<Delivery>> =
        new HashMap().onLookupFail(new HashSet())

    # D iterations
    # O(1) loop body
    # O(P) HashSet construction cost, due to handshake lemma
    # Overall: O(D + P)
    # Worst-case: |P| = 2 * |D|,
    #   if each delivery bridges two unique locations
    # Overall: O(D)
    for delivery in deliveries:
        # foreach P, we construct a HashSet, if the P lookup fails
        sourceDeliveries: HashSet<Delivery> =
            sourceToDeliveries[delivery.source] # O(1)
        destinationDeliveries: HashSet<Delivery> =
            destinationToDeliveries[delivery.destination] # O(1)

        sourceDeliveries.add(delivery) # O(1)
        destinationDeliveries.add(delivery) # O(1)

    adjacency: HashMap<Delivery, HashSet<Delivery>> =
        new HashMap().onLookupFail(new HashSet())

    # D iterations
    # Worst-case: O(D) loop body,
    #   if successive locations always depart after predecessor arrival
    # Overall: O(D^2)
    for delivery in deliveries:
        candidateNeighbours: HashSet<Delivery> = sourceToDeliveries[delivery.destination]
            # worst-case O(D)
            .filter(d -> delivery.arrival() <= d.departure())

        adjacency[delivery] = candidateNeighbours # O(1)

    # O(D) due to implementation constraints
    sources: HashSet<Delivery> = sourceToDeliveries[source]

    # djikstras is \Theta(V * lg V + E * lg V),
    # As E is worst-case O(V^2), substituting, we get
    # \Theta(V * lg V + V^2 * lg V)
    # As our Ds are vertexes in our use of Djikstras algorithm, we get
    # Overall: \Theta(D * lg D + D^2 * lg D), which is
    # Overall: \Theta(D^2 * lg D)
    dijkstra(G = adjacency, sources)

    # O(|D|-1)
    minimumCost: Int = destinationToDeliveries[destination].minBy(d -> d.d).d

    return cost == inf ? -1 : cost
```

(b)

The time complexity of Dijkstra's algorithm, $\Theta(\mathcal{D}^2 \lg \mathcal{D})$, dominates all other parts of our algorithm ($O(\mathcal{D})$, $O(\mathcal{D}^2)$, $O(\mathcal{D})$, $O(|\mathcal{D}| - 1)$), as well as providing the tightest upper and lower-bound.

Thus, the time complexity is $\Theta(\mathcal{D}^2 \lg \mathcal{D})$ which describes both an asymptotic upper and lower bound on the worst case time complexity of this algorithm.

We assume that `HashSet` and `HashMap` never degenerate into $O(n)$ puts, gets, adds and nexts, and choose the $O(1)$ best case scenario. We use a Binary Heap in the implementation of Dijkstra's, with $O(\lg V)$ `Extract-Mins` and `Decrease-Keys`.

A \mathcal{D} that describes a complete graph of valid delivery sequencing would exploit the worst-case complexity of this algorithm.