

# Assignment 1: Background theory

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1. (a)  $y : [ \text{true}, (x = 0 \Rightarrow y = 0) \wedge (x \neq 0 \Rightarrow y = \frac{y_0}{x}) ]$   
 (b)  $y : [ \text{true}, (x = 0 \Rightarrow \text{true}) \wedge (x \neq 0 \Rightarrow y = \frac{y_0}{x}) ]$   
 (c) TODO
2.  $x, y : [\text{true}, x = z^2 \wedge y = z^4]$   
 $\sqsubseteq \{ \text{Composition} \}$   
 $x, y : [\text{true}, x = z^2]; x, y : [x = z^2, x = z^2 \wedge y = z^4]$   
 $\sqsubseteq \{ \text{Assignment: } \text{true} \Rightarrow x = z^2[x \setminus z^2] \}$   
 $x = z^2; x, y : [x = z^2, x = z^2 \wedge y = z^4]$   
 $\sqsubseteq \{ \text{Assignment: } x = z^2 \Rightarrow x = z^2 \wedge y = z^4[y \setminus x^2] \}$   
 $x := z^2; y := x^2$

3. (a) Assuming

$$\begin{aligned} wp(y := 10, \text{true}) &\equiv \text{true}[y \setminus 10] \\ &\equiv \text{true} \end{aligned}$$

we can conclude that

$$\begin{aligned} wp(\text{if } (x > 0 \vee y < 10) \rightarrow y := 10 \text{ fi}, \text{true}) &\equiv (x > 0 \vee y < 10) \wedge \\ &\quad ((x > 0 \vee y < 10) \rightarrow wp(y := 10, \text{true})) \\ &\equiv (x > 0 \vee y < 10) \wedge \text{true} \\ &\equiv (x > 0 \vee y < 10) \end{aligned}$$

As  $y < 10 \Rightarrow (x > 0 \vee y < 10)$ , the Hoare Triple is true.

- (b) Assuming

$$\begin{aligned} wp(x := x + y, P[x \setminus x + y]) &\equiv (P[x \setminus x + y])[x \setminus x + y] \\ &\equiv P[x \setminus (x + y) + y] \\ &\equiv P[x \setminus x + 2y] \end{aligned}$$

Using this

$$\{ \text{true}[x \setminus x + 2y] \} x := x + y \{ \text{true}[x \setminus x + y] \}$$

holds but

$$\{ (x > 8)[x \setminus x + 2y] \} x := x + y \{ (x > 8)[x \setminus x + y] \} \equiv \{ x + 2y > 8 \} x := x + y \{ x + y > 8 \}$$

does not. Therefore, the choice of  $P$  determines the validity of the Hoare Triple, due to  $P$  having arbitrary. Assuming that  $P$  is quantified over all possible predicates, the Hoare Triple does not hold

4. (a)  $y : [y < 10, y > 0]$   
 $\sqsubseteq \{ \text{Selection: } y < 10 \Rightarrow (x > 0 \vee y < 10) \}$   
**if**  $(x > 0 \vee y < 10) \rightarrow y : [(x > 0 \vee y < 10) \wedge (y < 10), y > 0]$  **fi**  
 $\sqsubseteq \{ \text{Absorption 1: } (x > 0 \vee y < 10) \wedge (y < 10) = y < 10 \}$   
**if**  $(x > 0 \vee y < 10) \rightarrow y : [y < 10, y > 0]$  **fi**  
 $\sqsubseteq \{ \text{Assignment: } y < 10 \Rightarrow y > 0[y \setminus 10] \}$   
**if**  $(x > 0 \vee y < 10) \rightarrow y := 10$  **fi**
- (b)  $y : [y < 10, y > 0]$   
 $\not\sqsubseteq \{ \text{Selection: } y < 10 \not\Rightarrow ((x > 0) \wedge (y < 10)) \}$   
**if**  $((x > 0) \wedge (y < 10)) \rightarrow y : [((x > 0) \wedge (y < 10)) \wedge (y < 10), y > 0]$  **fi**

5.

$\forall S, B \cdot (\text{repeat } S \text{ until } B \equiv S; \text{ do } \neg B \rightarrow S \text{ od})$

$\therefore$

$w : [P, Q]$   
 $\sqsubseteq \{ \text{Composition} \}$   
 $w : [P, I]; w : [I, Q]$   
 $\sqsubseteq \{ \text{Strengthen Postcondition: } I \wedge \neg(\neg B) \Rightarrow Q \}$   
 $w : [P, I]; w : P[I, I \wedge \neg(\neg B)]$   
 $\sqsubseteq \{ \text{Repetition} \}$   
 $w : [P, I]; \text{do } (\neg B) \rightarrow w : [I \wedge \neg B, I \wedge (0 \leq V < V_0)] \text{ od}$

$w : [P, I]$  and  $w : [I \wedge \neg B, I \wedge (0 \leq V < V_0)]$  can both refine to the same program,  $S$ .

Thus,  $P \Rightarrow I \wedge \neg B$ , so such that for an arbitrary  $S$  with precondition  $P$ ,  $S$  can preserve the requirements of the **do** loop.

Given both  $I \wedge \neg(\neg B) \Rightarrow Q$  and  $P \Rightarrow I \wedge \neg B$ , we can formulate

if  $P \Rightarrow I \wedge \neg B$  then  $w : [P, I \wedge B] \sqsubseteq \text{repeat } w : [I \wedge \neg B, I \wedge (0 \leq V < V_0)] \text{ until } B$   
 where  $I$  is a loop invariant, and  $V$  is a loop variant