

Assignment 3: Derivation

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1. (a) n is a **value** parameter. m is a **result** parameter.
- (b) $inv \triangleq lrun(A, n_0, m)$
- (c) Let

$$\begin{aligned} pre(A, n) &\triangleq lrun(A, n, n+1) \\ post(A, n, m) &\triangleq mrun(A, n_0, m) \end{aligned}$$

by the specification of the procedure, where inv , pre and $post$ implicitly capture parameters (A, n, m) from the frame.

$$\begin{aligned} &n, m : [pre, post] \\ \sqsubseteq &\quad \{\text{Composition: middle predicate is } inv\} \\ &n, m : [pre, inv]; \quad n, m : [inv, post] \\ \sqsubseteq &\quad \{\text{Assignment: } pre \Rightarrow inv[m \setminus n + 1]\} \\ &m := n + 1; \quad n, m : [inv, post] \\ \therefore & \end{aligned}$$

$$\begin{aligned} inv[m \setminus n + 1] &\equiv lrun(A, n_0, m)[m \setminus n + 1] \\ &\equiv lrun(A, n_0, n + 1) \end{aligned}$$

and

$$lrun(A, n, n + 1) \Rightarrow lrun(A, n_0, n + 1)$$

Let

$$guard \triangleq m < A.len \wedge A_{n_0} \neq A_m$$

$$\begin{aligned} \sqsubseteq &\quad \{\text{Strengthen post: } inv \wedge \neg guard \Rightarrow post\} \\ &m := n + 1; \quad n, m : [inv, inv \wedge \neg guard] \\ \therefore & \end{aligned}$$

$$\begin{aligned} &inv \wedge \neg guard \Rightarrow post \\ \equiv &\quad lrun(A, n_0, m) \wedge \neg(m < A.len \wedge A_{n_0} \neq A_m) \Rightarrow mrun(A, n_0, m) \\ \equiv &\quad lrun(A, n_0, m) \wedge \neg(m < A.len \wedge A_{n_0} \neq A_m) \Rightarrow lrun(A, n_0, m) \wedge (m < A.len \Rightarrow A_{n_0} \neq A_m) \\ \equiv &\quad lrun(A, n_0, m) \wedge (\neg(m < A.len) \vee \neg(A_{n_0} \neq A_m)) \Rightarrow lrun(A, n_0, m) \wedge (\neg(m < A.len) \vee (A_{n_0} \neq A_m)) \end{aligned}$$