Assignment 3: Derivation

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- 1. (a) n is a value parameter. m is a result parameter.
 - (b) $inv \triangleq lrun(A, n_0, m)$
 - (c) Let

$$pre(A, n) \triangleq lrun(A, n, n + 1)$$

 $inv(A, n, m) \triangleq lrun(A, n_0, m)$
 $post(A, n, m) \triangleq mrun(A, n_0, m)$

by 1(b), and the specification of the procedure. inv, pre and post implicitly capture variables (A, n, m) as parameters from the frame, for syntactic convience. s.t.

 $\begin{array}{l} n,m:[pre,\,post] \\ & \subseteq \quad \{ \text{Composition: middle predicate is } inv \} \\ & n,m:[pre,\,inv]; \quad n,m:[inv,\,post] \\ & \subseteq \quad \{ \text{Assignment: } pre \Rrightarrow inv[m\backslash n+1] \} \\ & m:=n+1; \quad n,m:[inv,\,post] \end{array}$

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$$inv[m \backslash n + 1] \equiv lrun(A, n_0, m)[m \backslash n + 1]$$

 $\equiv lrun(A, n_0, n + 1)$

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$$lrun(A, n, n + 1) \implies lrun(A, n_0, n + 1)$$

Let

$$guard \triangleq m < A.len \land A_{n_0} = A_m$$

s.t.

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$$\begin{array}{ll} & inv \land \neg guard \; \Rrightarrow \; post \\ \equiv \; & \{ \text{Expansion of definitions} \} \\ & lrun(A, n_0, m) \land \neg (m < A.len \land A_{n_0} = A_m) \; \Rrightarrow \; mrun(A, n_0, m) \\ \equiv \; & \{ \text{Expansion of functions} \} \\ & lrun(A, n_0, m) \land \neg (m < A.len \land A_{n_0} = A_m) \; \Rrightarrow \; lrun(A, n_0, m) \land (m < A.len \Rightarrow A_{n_0} \neq A_m) \\ \equiv \; & \{ \text{De Morgan's law - negation of conjunction} \} \end{array}$$

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\begin{array}{l} \operatorname{lrun}(A, n_0, m) \wedge (\neg (m < A.len) \vee \neg (A_{n_0} = A_m)) \; \Rrightarrow \; \operatorname{lrun}(A, n_0, m) \wedge (m < A.len \Rightarrow A_{n_0} \neq A_m) \\ \equiv \; \{P \Rightarrow Q \; \equiv \; \neg P \vee Q\} \\ \operatorname{lrun}(A, n_0, m) \wedge (\neg (m < A.len) \vee \neg (A_{n_0} = A_m)) \; \Rrightarrow \; \operatorname{lrun}(A, n_0, m) \wedge (\neg (m < A.len) \vee (A_{n_0} \neq A_m)) \\ \equiv \; \{\} \\ \operatorname{true} \end{array}
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