

PHIL3110 - Assignment 1

Maxwell Bo

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Part A

Problem 1

$$1. \frac{\frac{Pa \vee Qa^{(1)} \quad (Pa \vee Qa) \rightarrow \perp}{\perp} \quad (1) (\neg I) \quad \frac{\perp}{\neg(Pa \vee Qa)} \quad (DM-1)}{\neg Pa \wedge \neg Qa}$$

De Morgan's Law 1

$$\frac{\frac{\frac{\varphi^{(1)}}{\varphi \vee \psi} \quad \neg(\varphi \vee \psi)}{\perp} \quad (1) (\neg I) \quad \frac{\perp}{\neg\varphi} \quad \frac{\frac{\frac{\psi^{(2)}}{\varphi \vee \psi} \quad \neg(\varphi \vee \psi)}{\perp} \quad (2) (\neg I) \quad \frac{\perp}{\neg\psi}}{\neg\varphi \wedge \neg\psi}$$

$$2. \frac{Qa \vee Ra^{(3)} \quad \frac{Qa^{(1)} \quad Qa \rightarrow Pa}{Pa} \quad \frac{Ra^{(2)} \quad Ra \rightarrow Pa}{Pa} \quad (1) (2) (\vee E)}{\frac{Pa}{(Qa \vee Ra) \rightarrow Pa} \quad (3) (\rightarrow I)}$$

$$3. \frac{\neg(\forall x Px \rightarrow \exists x Rx) \quad \frac{\frac{\forall x Px^{(1)} \quad \exists x Rx^{(2)}}{\forall x Px \wedge \exists x Rx} \quad \frac{\exists x Rx}{\forall x Px \rightarrow \exists x Rx} \quad (1) (\rightarrow I)}{\frac{\perp}{\neg\exists x Rx} \quad (2) (\neg I)}$$

Problem 2

1.

$$Q^{\mathcal{M}} = \{m_1\}$$

$$T^{\mathcal{M}} = \{\langle m_1, m_1 \rangle, \langle m_1, m_2 \rangle, \langle m_2, m_2 \rangle\}$$

2. Distressingly, \mathcal{L} does not define any constant symbols, nor does \mathcal{M} provide interpretations of constant symbols in \mathcal{M} .

Thus

$$\mathcal{M} \not\models \exists x \neg Txx$$

However assuming \mathcal{M}^+ , where \mathcal{M}^+ is the expanded model \mathcal{M} , where $m^{\mathcal{M}} = m$ for all $m \in M$, we see that

$$M \models \exists x \neg Txx$$

as

$$\langle m_3^{\mathcal{M}}, m_3^{\mathcal{M}} \rangle \notin T^{\mathcal{M}}$$

3. No, as \mathcal{M} does not define any constant symbols $\mathcal{M} \not\models \exists x \varphi$ for some arbitrary φ (as x will bind no constant symbols), and thus

$$\mathcal{M} \not\models \exists x \forall y (Qy \leftrightarrow Tyx)$$

Assuming \mathcal{M}^+ ,

$$\mathcal{M}^+ \not\models \exists x \forall y (Qy \leftrightarrow Tyx)$$

By fixing x to $m_1^{\mathcal{M}}$, we see that

$$\forall y \cdot y \in Q^{\mathcal{M}} \leftrightarrow \langle y, m_1^{\mathcal{M}} \rangle \in T^{\mathcal{M}}$$

as

$$\begin{aligned} m_1^{\mathcal{M}} &\in Q^{\mathcal{M}} \text{ and } \langle m_1^{\mathcal{M}}, m_1^{\mathcal{M}} \rangle \in T^{\mathcal{M}} \\ m_2^{\mathcal{M}} &\notin Q^{\mathcal{M}} \text{ and } \langle m_2^{\mathcal{M}}, m_1^{\mathcal{M}} \rangle \notin T^{\mathcal{M}} \\ m_3^{\mathcal{M}} &\notin Q^{\mathcal{M}} \text{ and } \langle m_3^{\mathcal{M}}, m_1^{\mathcal{M}} \rangle \notin T^{\mathcal{M}} \end{aligned}$$

Part B

Problem 3

1. (a) *BHMB* is not a good sandwich
 (b) $((B)JB)HB$ is a good sandwich
 (c) *BHBMBJM* is not a good sandwich
2. **Definition 1** *A sandwich is good iff there is some stage n such that sandwich $\in \text{Stage}(n)$ where:*
 - $\text{Stage}(0) = \{B\}$
 - $\text{Stage}(n+1)$ is the set of φ such that either:
 - (a) $\varphi \in \text{Stage}(n)$
 - (b) φ is of the form ψMB , ψHB , or ψJB , where $\psi \in \text{Stage}(n)$
3. Make A the set of all n such that for each $\psi \in \text{Stage}(n)$, ψ does not contain any two instances of the same ingredient adjacently.

Base

$0 \in A$. As $\text{Stage}(0) = \{B\}$, all $\varphi \in \text{Stage}(0)$ do not contain any two instances of the same ingredient adjacently. We can also see that all $\varphi \in \text{Stage}(0)$ are terminated by B

Induction Step

Suppose $n \in A$.

Lemma 1 *Since $n \in A$, and every $\varphi \in \text{Stage}(n)$ is of the form B , ψMB , ψHB , or ψJB , for $\psi \in \text{Stage}(m)$ for some $m < n$, all φ are terminated by B .*

Lemma 2 *Since $n \in A$, every $\varphi \in \text{Stage}(n)$ does not contain any two instances of the same ingredient adjacently.*

For each $\varphi \in \text{Stage}(n+1) \setminus \text{Stage}(n)$, recalling the definition for $\text{Stage}(n+1)$, φ is of the form ψMB , ψHB , or ψJB , for some $\psi \in \text{Stage}(n)$

Given

- all $\varphi \in \text{Stage}(n)$ are terminated by B (Lemma 1)
- all $\varphi \in \{MB, HB, JB\}$ don't begin with B , and are not themselves adjacent ingredients
- Lemma 2

all $\varphi \in \text{Stage}(n+1) \setminus \text{Stage}(n)$ do not contain two instances of the same ingredient adjacently.

Thus, $n+1 \in A$. Then by induction, we see that every $n \in A$.

Problem 4

1. ($\rightarrow I$) Suppose d_ψ is a derivation of $\Gamma, \varphi \vdash \psi$, with d_ψ being from stage n .

The ($\rightarrow I$) rule tells us that the $n+1^{\text{th}}$ stage contains a derivation d of $\Gamma, \varphi \vdash \varphi \rightarrow \psi$.

We must show $\Gamma, \varphi \models \varphi \rightarrow \psi$.

Since $d_\psi \in \text{Stage}_{\text{Der}}(n)$ we have $\Gamma, \varphi \models \psi$.

Let \mathcal{M} be a model which $\mathcal{M} \models \gamma$ for all $\gamma \in \Gamma \cup \varphi$.

Thus we have $\mathcal{M} \models \varphi$ and $\mathcal{M} \models \psi$, so $\mathcal{M} \models \varphi \rightarrow \psi$.

($\exists E$) Suppose:

- d_\exists is a derivation of $\Gamma \vdash \exists x\varphi(x)$
- d_ψ is a derivation of $\Delta, \varphi(a) \vdash \psi$, where a does not occur in Δ or $\exists x\varphi(x)$

where all are members of stage n . Then the ($\exists E$) rule tells us that the $n+1^{\text{th}}$ stage contains a derivation d of $\Gamma, \Delta, \varphi(a) \vdash \psi$.

We must show that $\Gamma, \Delta, \varphi(a) \models \psi$.

Since $d_\exists \in \text{Stage}_{\text{Der}}(n)$, we have $\Gamma \models \exists x\varphi(x)$.

Since $d_\psi \in \text{Stage}_{\text{Der}}(n)$, we have $\Delta, \varphi(a) \models \psi$.

Let \mathcal{M} be a model which $\mathcal{M} \models \gamma$ for all $\gamma \in \Delta \cup \varphi(a)$. Then $\mathcal{M} \models \psi$.

Suppose that $m = a^{\mathcal{M}}$. Now consider any model \mathcal{M}' which is exactly like \mathcal{M} except that $a^{\mathcal{M}'} = m' \neq a^{\mathcal{M}} = m$; i.e., we let $a^{\mathcal{M}'}$ be some arbitrary other m' from the domain M . Since $\Delta \cup \varphi(a)$ has a sentence with the constant symbol a in it, \mathcal{M}' only models ψ if it models $\varphi(a)$.

But the fact that all models \mathcal{M}' whose only difference from \mathcal{M} is their interpretation of the symbol a are such that $\mathcal{M}' \models \psi$ just means that:

$\exists m \in M \mathcal{M}' \models \exists x\varphi(x), \varphi(m) \Leftrightarrow \mathcal{M} \models \psi$.

2. Suppose $\psi := \chi \vee \delta \in \Delta$.

We claim that $\chi \vee \delta$ iff $\chi \in \Delta$ or $\delta \in \Delta$.

If $\chi \vee \delta \in \Delta$, then by Claim 120 and ($\vee I$) we have either χ or δ in Δ , and by ($\vee E$) we have χ and δ in Δ .

If either χ or δ are in Δ then Claim 120 and ($\vee I$) tell us that $\chi \vee \delta \in \Delta$.

$$\begin{aligned} \chi \vee \delta \in \Delta &\Leftrightarrow \chi \in \Delta \parallel \delta \in \Delta \\ &\Leftrightarrow \mathcal{M}^\Delta \models \chi \parallel \mathcal{M}^\Delta \models \delta \\ &\Leftrightarrow \mathcal{M}^\Delta \models \chi \vee \delta \end{aligned}$$

The first \Leftrightarrow is via our claim; the second \Leftrightarrow is by induction hypothesis; and the last \Leftrightarrow is from the \models definition.

3. Suppose $\psi := \forall x\chi(x)$.

Then we claim $\forall x\chi(x) \in \Delta$ iff there is some constant symbol $a \in \mathcal{L}(C)$ such that $\chi(a) \in \Delta$.

Suppose $\forall x\chi(x) \in \Delta$. Then by Claim 120 and ($\forall E$), there is some $a \in \mathcal{L}(C)$ such that $\chi(a) \in \Delta$.

On the other hand, if for some some $a \in \mathcal{L}(C)$, $\chi(a) \in \Delta$, then by Claim 120 and $(\forall I)$, $\forall x \chi(x) \in \Delta$.

Then we have

$$\begin{aligned} \forall x \chi(x) \in \Delta &\Leftrightarrow \text{there is some } a \in M^\Delta \text{ such that } \chi(a) \in \Delta \\ &\Leftrightarrow \text{there is some } a \in M^\Delta \text{ such that } \mathcal{M}^\Delta \models \chi(a) \\ &\Leftrightarrow \mathcal{M}^\Delta \models \forall x \chi(x) \end{aligned}$$

The first \Leftrightarrow is via our claim; the second \Leftrightarrow is by induction hypothesis; and the final \Leftrightarrow is via the \models definition.