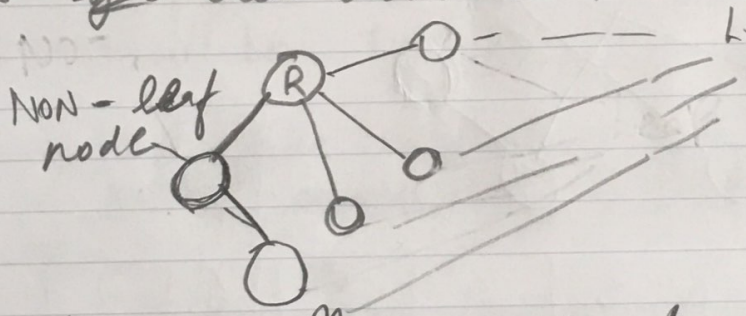
Exam 1

Q1

i. R, will have the maximum height. A height of any node is the number of edges from one node to its leaf. So, if R is the root then all the nodes are connected to it so, R will have the number of edges, hence the maximum height.

ii There is not enough information to prove this. Because to find a depth of a node in a tree you would calculate the length from the Root R, to L. In this case L is a random leaf node and the origin problem states "One root node R & Several leaf and One non leaf"

This tree would look like



According to this graph L is a "Random" leaf node so if is not for certain it connected to "One - non leaf node" or part of one of the "Several" leaf nodes. So we cannot come to the conclusion that 'L' has the maximum depth.

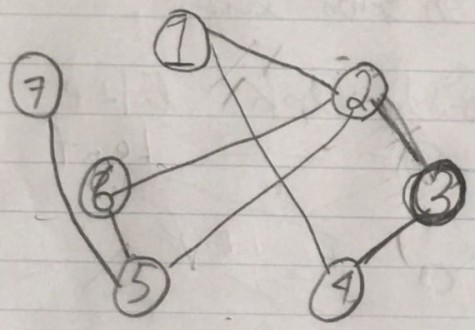
iii The height of a tree can be calculated recursively. We can do this by first calculating the left & right Subtrees heights.

To do this

(a) Size of right subtree
(b) Size of left subtree
(c) The height will be the size of right + left & add it with (1) this will give the height of tree recursively.

the best algorithm to use for this is a stack algorithm.

Q2) 1-4, 1-2, 2-3, 2-5, 2-6, 4-3, 5-6, 5-7
Undirected, 10 nodes



1)

Nodes	1	2	3	4	5	6	7
1	0	1	0	1	0	0	0
2	1	0	1	0	1	1	0
3	0	1	0	1	0	0	0
4	1	0	1	0	0	0	0
5	0	1	0	0	0	1	1
6	0	1	0	0	1	0	0
7	0	0	0	0	1	0	0

2) The time complexity for adjacency matrix is $O(n^2)$ since it's a undirected graph.

(2) The time complexity for the matrix is $O(n^2)$
Since this is a undirected graph it still is $O(n^2)$

3) This has the same reasoning as the last
the time complexity will be $O(n^2)$

Q3 1 million integers, in array randomly.
Each cell → 1 nano sec

(1) We would need to take each cell an multi by it by 1 million
Worst case is = 1 million * 1 nano sec.

(2) It would not be any different if an arry got sorted as long as the time & no. of integers do not Change it will be 1 million * 1 nano sec.

(3) Yes, just like with most serches the binary search will be better we can improve it by using $O(\log n)$

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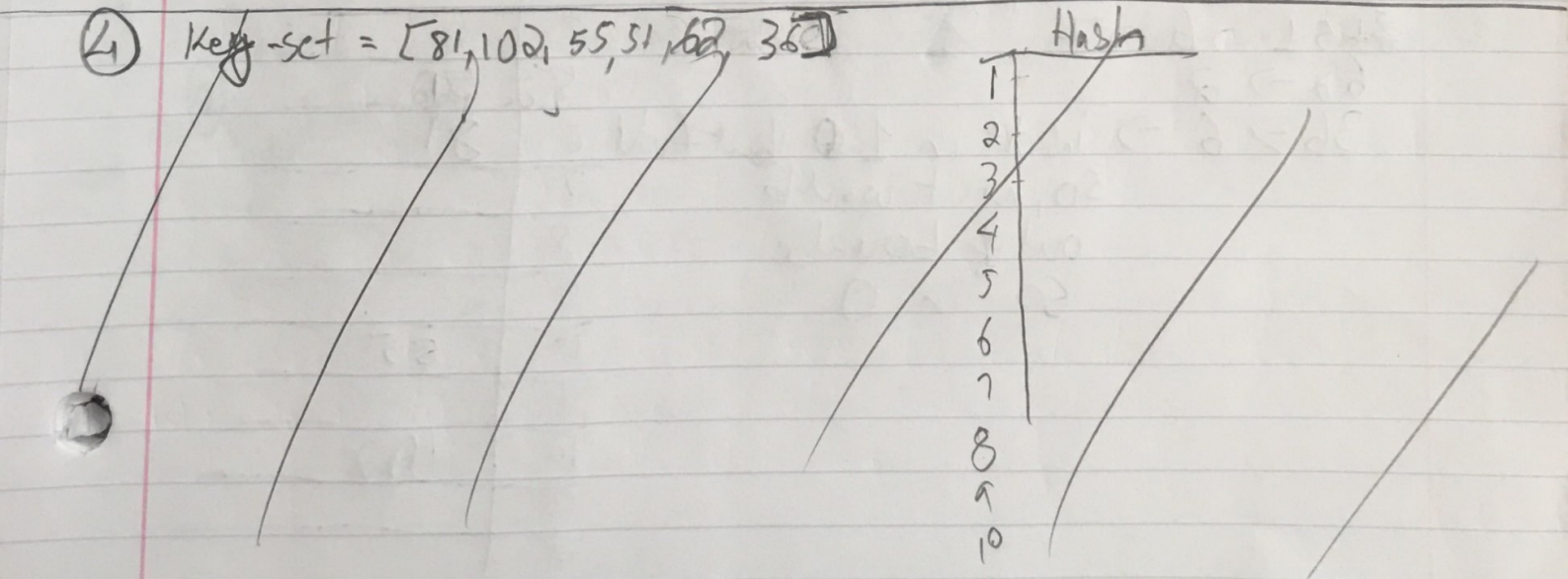
(4) Key-set = [81, 102, 55, 51, 62, 36]



$S(K) = S(K) \% 15$

(24) [81, 102, 55, 51, 62, 36]

(i) Lin Probing

So 81 % 15 → 6
102 % 15 → 12
55 % 15 → 10
51 % 15 → 6 but full so 7
62 % 15 → 2
36 % 15 → 6 but full so 8

	Table
1	
2	62
3	
4	
5	
6	81
7	51
8	36
9	
10	55
11	
12	100
13	
14	
15	

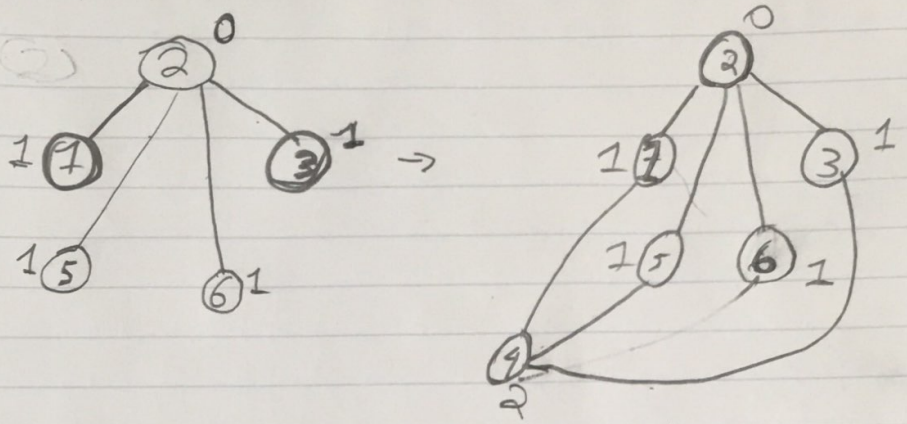
This is 102

ii) Quad Probing

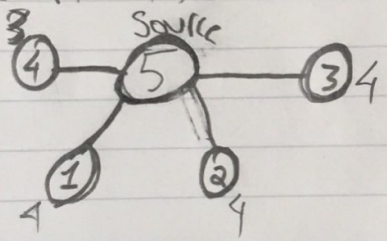
81 → 6
102 → 12
55 → 10
51 → 6 → 7 insert in
62 → 2
36 → 6 → insert in 10 but full
so, next would
out of bounds
so in 0

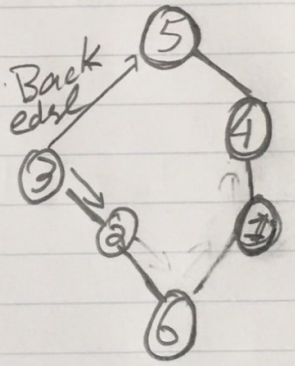
	Table
1	36
2	62
3	
4	
5	51
6	81
7	
8	
9	
10	55
11	
12	102
13	
14	
15	

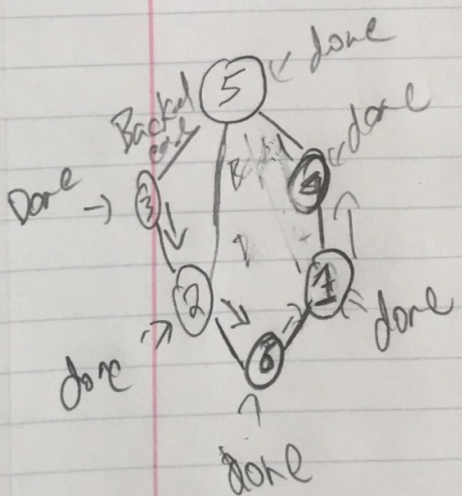
Q5 (1) Breadth First Search



(6) Depth First Search







Back edges = (3,5) (6,3), (1,5), (2,3), (4,3)