Advanced R Companion

R4DS Reading Group

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Chapter 1

Welcome

A companion to Advanced R and supplement to Advanced R Solutions Recordings for the accompanied presentation for each chapter can be seen here.

Chapter 2

Names and Values

2.2 Binding basics

Should we care about R internals?

Guided by this blogpost, we can use the .Internal function to inspect metadata associated with our objects:

```
x <- list(1:5)
.Internal(inspect(x))
# > @0x000001b6a4af9fc8 19 VECSXP g0c1 [NAM(7)] (len=1, tl=0)
# > @0x000001b6a3321d90 13 INTSXP g0c0 [NAM(7)] 1 : 5 (compact)
```

- @0x000001b6a4af9fc8 address (memory location)
- 19 VECSXP type (full list here)
- g0 garbage collector info token
- c1 size of object (small vector)
- $\bullet\,$ NAM(7) named value of the object (if greater than one copy on modify)
- len=1 length of object
- tl true length of object
- small snippet of the data

```
{
    x <- list(1:5)
    .Internal(inspect(x))
}

# < @0x000001b6a4b4a3a0 19 VECSXP g0c1 [NAM(1)] (len=1, tl=0)
# < @0x000001b6a3a814d0 13 INTSXP g0c0 [NAM(7)] 1 : 5 (compact)
```

It is of note here that without curly brackets we have to use copy-on-modify, but within curly brackes we can use copy-in-place because NAM(1)

2.3 Copy-on-modify

copy-on-modify vs copy-in-place: is one more preferable in certain situations?

modify in place only happens when objects with a single binding get a special performance optimization and to environments.

2.2.2 Exercises

Question 3 digs into the syntactically valid names created when using read.csv(), but what is the difference between quotation and backticks?

If we create an example csv

```
example2223 <- tibble(
   `if` = c(1,2,3),
   `_1234` = c(4,5,6),
   `column 1` = c(7,8,9)
)
write.csv(example2223, "example2223.csv", row.names = FALSE)</pre>
```

Import using adjusted column names to be syntactically valid:

Import using non-adjusted column names

```
read.csv(file = "example2223.csv", check.names = FALSE)
```

Import using the tidyverse where names are not adjusted

```
df_non_syntactic_name <- read_csv(file = "example2223.csv")</pre>
```

```
## Parsed with column specification:
## cols(
## `if` = col_double(),
## `_1234` = col_double(),
## `column 1` = col_double()
## )
```

However I really don't understand the difference between backticks and quotation marks. For example when I select a column in the case of non-syntactic in the tidyverse I can use quotation marks or backticks

```
df_non_syntactic_name %>% select("if")
```

But in base R, I can do this with quotation marks, but not backticks:

```
df__non_syntactic_name["if"]
```

Error in `[.default`(df__non_syntactic_name, `if`) : invalid subscript type 'special'

According to ?Quotes backticks are used for "non-standard variable names" but why in base R they don't work to select columns but in the tidyverse they work to select variables?

The easiest way to think about this is that backticks refer to objects while quotation marks refer to strings. dplyr::select() accepts object references as well as string references, while base R subsetting is done with a string or integer position.

2.3.2 Function calls

Can we go over and break down figure in 2.3.2

When you create this function:

```
crazyfunction <- function(eh) {eh}</pre>
```

eh doesn't exist in memory at this point.

```
x \leftarrow c(1,2,3)
```

x exists in memory.

```
z <- crazyfunction(x)
```

z now points at x, and eh still doesn't exist (except metaphorically in Canada). eh was created and exists WHILE crazyfunction() was being run, but doesn't get saved to the global environment, so after the function is run you can't see its memory reference.

The round brackets (eh) list the arguments, the curly brackets {eh} define the operation that it's doing - and you're assigning it to crazyfunction.

R functions automatically return the result of the last expression so when you call that object (the argument eh) it returns the value of that argument. This is called **implicit returns**

2.3.3 Lists

Checking the address for a list and its copy we see they share the same references:

```
11 <- list(1,2,3)
12 <- 11
identical(lobstr::ref(l1),lobstr::ref(l2))

## [1] TRUE
lobstr::obj_addr(l1[[1]])

## [1] "0x7ffd85e056b0"
lobstr::obj_addr(l2[[1]])</pre>
```

```
## [1] "0x7ffd85e056b0"
```

But why isn't this the case for their subsets? Using obj_addr they have different addresses, but when we look at their references they are the same

```
lobstr::obj_addr(l1[1])
```

```
## [1] "0x7ffd89f79a30"
```

```
lobstr::ref(l1[1])

## [1:0x7ffd89f4d148] <list>
## [2:0x7ffd85e056b0] <dbl>

lobstr::obj_addr(l2[1])

## [1] "0x7ffd91436fa0"

identical(lobstr::obj_addr(l1[1]), lobstr::obj_addr(l2[1]))
```

[1] FALSE

This is because using singular brackets wraps the value 1 in a new list that is created on the fly which will have a unique address. We can use double brackets to confirm our mental model that the sublists are also identical:

```
identical(lobstr::obj_addr(l1[[1]]), lobstr::obj_addr(l2[[1]]))
```

```
## [1] TRUE
```

What's the difference between these 2 addresses <0x55d53fa975b8> and 0x55d53fa975b8?

Nothing - it has to do with the printing method:

```
x <- c(1, 2, 3)
print(tracemem(x))

## [1] "<0x7ffd90e60b08>"

cat(tracemem(x))

## <0x7ffd90e60b08>
lobstr::obj_addr(x)
```

```
## [1] "0x7ffd90e60b08"
```

When would you prefer a deep copy of a list to a shallow copy? Is this something to consider when writing functions or package development or is this more something that's optimized behind the scenes?

Automagical!

2.3.5 Character vectors

Is there a way to clear the "global string pool"?

According to this post it doesn't look like you can directly, but clearing all references to a string that's in the global string pool clears that string from the pool, eventually

2.3.6.2 Exercise

When we look at tracemem when we modify x from an integer to numeric, x is assigned to three objects. The first is the integer, and the third numeric - so what's the intermediate type?

```
x <- c(1L, 2L, 3L)
obj_addr(x)
tracemem(x)
x[[3]] <- 4</pre>
```

```
[1] "0x7f84b7fe2c88"
[1] "<0x7f84b7fe2c88>"
tracemem[0x7f84b7fe2c88 -> 0x7f84b7fe5288]:
tracemem[0x7f84b7fe5288 -> 0x7f84bc0817c8]:
```

What is 0x7f84b7fe5288 when the intermediate x <- c(1L, 2L, 4) is impossible?

When we assign the new value as an integer there is no intermediate step. This probably means c(1,2, NA) is the intermediate step; creating an intermediate vector that's the same length of the final product with NA values at all locations that are new or to be changed

```
x <- c(1L, 2L, 3L)
obj_addr(x)

## [1] "0x7ffd8a1b4d88"

tracemem(x)

## [1] "<0x7ffd8a1b4d88>"
```

tracemem[0x7ffd8a1b4d88 -> 0x7ffd91204e08]: eval eval withVisible withCallingHandlers handle timing_

You can dig into the C code running this:

```
pryr::show_c_source(.Internal("[<-"))</pre>
```

2.4.1 Object size

x[[3]] < -4L

If I have two vectors, one 1:10 and another c(1:10, 10), intuitively, I would expect the size of the second vector to be greater than the size of the first. However, it seems to be the other way round, why?

```
x1 <- 1:10
x2 <- rep(1:10, 10)
lobstr::obj_size(x1)</pre>
```

680 B

```
lobstr::obj_size(x2)
```

448 B

If we start with the following three vectors:

```
x1 <- c(1L, 2L, 3L, 4L, 5L, 6L, 7L, 8L, 9L, 10L)
x2 <- 1:10
x3 <- rep(1:10, 10)
lobstr::obj_sizes(x1, x2, x3)</pre>
```

```
## * 96 B
## * 680 B
## * 448 B
```

Intuitively, we would have expected x1 < x2 < x3 but this is not the case. It appears that the rep() function coerces a double into integer and hence optimizes on space. Using:, R internally uses ALTREP.

ALTREP would actually be more efficient if the numbers represented were significantly large, say 1e7.

```
x4 <- 1:1e7
x5 <- x4
```

```
x5[1] <- 1L
lobstr::obj_sizes(x4, x5)</pre>
```

```
## * 680 B
## * 40,000,048 B
```

Now, the size of x4 is significantly lower than that of x5. This seems to indicate that ALTREP becomes super efficient as the vector size is increased.

2.5.1 Modify-in-place

"When it comes to bindings, R can currently only count 0, 1, or many. That means that if an object has two bindings, and one goes away, the reference count does not go back to 1: one less than many is still many. In turn, this means that R will make copies when it sometimes doesn't need to."

Can we come up with an example of this? It seems really theoretical right now.

First you need to switch your Environment tab to something other than global in RStudio!

```
Now we can create a vector:

v <- c(1, 2, 3)
(old_address <- lobstr::obj_addr(v))

## [1] "0x7ffd728f54c8"

Changing a value within it changes its address:
v[[3]] <- 4
(new_address <- lobstr::obj_addr(v))

## [1] "0x7ffd9162d748"

old_address == new_address

## [1] FALSE

We can assign the modified vector to a new name, where y and v point to the same thing.
y <- v
(y_address <- lobstr::obj_addr(y))

## [1] "0x7ffd9162d748"

(v_address <- lobstr::obj_addr(v))
```

```
(v_address <- lobstr::obj_addr(v))
## [1] "0x7ffd9162d748"
y_address == v_address
## [1] TRUE</pre>
```

[1] 11toE

Now if we modify v it won't point to the same thing as y:

```
v[[3]] <- 3
(y_address <- lobstr::obj_addr(y))
```

```
## [1] "0x7ffd9162d748"
(v_address <- lobstr::obj_addr(v))
```

```
## [1] "0x7ffd90e6bba8"
```

```
y_address == v_address
```

[1] FALSE

But if we now change y to look like v, the original address, in theory editing y should occur in place, but it doesn't - the "count does not go back to one"!

```
y[[3]] <- 3
(new_y_address <- lobstr::obj_addr(y))</pre>
```

```
## [1] "0x7ffd90e71278"

new_y_address == y_address
```

```
## [1] FALSE
```

Can we break down this code a bit more? I'd like to really understand when and how it's copying three times. As of R 4.0 it's now copied twice, the 3rd copy that's external to the function is now eliminated!!

```
# dataframe of 5 columns of numbers
x <- data.frame(matrix(runif(5 * 1e4), ncol = 5))
# median number for each column
medians <- vapply(x, median, numeric(1))

# subtract the median of each column from each value in the column
for (i in seq_along(medians)) {
    x[[i]] <- x[[i]] - medians[[i]]
}</pre>
```

```
cat(tracemem(x), "\n")
```

```
<0x7fdc99a6f9a8>
```

```
for (i in 1:5) {
  x[[i]] <- x[[i]] - medians[[i]]
}</pre>
```

```
tracemem[0x7fdc99a6f9a8 -> 0x7fdc9de83e38]:
tracemem[0x7fdc9de83e38 -> 0x7fdc9de83ea8]: [[<-.data.frame [[<-
tracemem[0x7fdc9de83ea8 -> 0x7fdc9de83f18]: [[<-.data.frame [[<-
tracemem[0x7fdc9de83f18 -> 0x7fdc9de83f88]:
tracemem[0x7fdc9de83f88 -> 0x7fdc9de83ff8]: [[<-.data.frame [[<-
tracemem[0x7fdc9de83ff8 -> 0x7fdc9de84068]: [[<-.data.frame [[<-
tracemem[0x7fdc9de83ff8 -> 0x7fdc9de84068]:
tracemem[0x7fdc9de84068 -> 0x7fdc9de840d8]:
tracemem[0x7fdc9de84068 -> 0x7fdc9de84148]: [[<-.data.frame [[<-
tracemem[0x7fdc9de84148 -> 0x7fdc9de84148]: [[<-.data.frame [[<-
tracemem[0x7fdc9de84148 -> 0x7fdc9de84228]:
tracemem[0x7fdc9de84228 -> 0x7fdc9de84228]:
tracemem[0x7fdc9de84228 -> 0x7fdc9de84298]: [[<-.data.frame [[<-
tracemem[0x7fdc9de84298 -> 0x7fdc9de84308]: [[<-.data.frame [[<-
tracemem[0x7fdc9de84308 -> 0x7fdc9de843e8]: [[<-.data.frame [[<-
tracemem[0x7fdc9de843e8 -> 0x7fdc9de84458]: [[<-.data.frame [[<-
tracemem[0x7fdc9de843e8] -> 0x7fdc9de84458]: [[<-.data.frame [[<-
tracemem[0x7fdc9de843e8] -> 0x7fdc9de84458]: [[<-.data.frame [[<-
```

When we run tracemem on the for loop above we see each column is copied twice followed by the [[<-.data.frame [[<-, the stack trace showing exactly where the duplication occurred.

So what is [[<-.data.frame? It's a function! By looking at '?[[<-.data.frame we see this is used to "extract or replace subsets of data frames."

When we write x[[i]] <- value, it's really shorthand for calling the function [[<-.data.frame with inputs x, i, and value.

Now let's step into the call of this base function by running debug(``[[<-.data.frame``):

```
debug(`[[<-.data.frame`)</pre>
```

and once inside, use tracemem() to find where the new values are assigned to the column:

```
function (x, i, j, value)
{
  if (!all(names(sys.call()) %in% c("", "value")))
    warning("named arguments are discouraged")
  cl <- oldClass(x)
# this is where another copy of x is made!
  class(x) <- NULL</pre>
```

tracemem[0x7fdc9d852a18 -> 0x7fdc9c99cc08]:

```
nrows <- .row_names_info(x, 2L)</pre>
  if (is.atomic(value) && !is.null(names(value)))
    names(value) <- NULL</pre>
  if (nargs() < 4L) {</pre>
    nc <- length(x)</pre>
    if (!is.null(value)) {
      N <- NROW(value)
      if (N > nrows)
         stop(sprintf(ngettext(N, "replacement has %d row, data has %d",
           "replacement has %d rows, data has %d"), N,
           nrows), domain = NA)
      if (N < nrows)</pre>
         if (N > OL && (nrows \%\N == OL) && length(dim(value)) <=
           value <- rep(value, length.out = nrows)</pre>
         else stop(sprintf(ngettext(N, "replacement has %d row, data has %d",
           "replacement has %d rows, data has %d"), N,
           nrows), domain = NA)
    }
    x[[i]] \leftarrow value
    if (length(x) > nc) {
      nc <- length(x)</pre>
      if (names(x)[nc] == "")
        names(x)[nc] <- paste0("V", nc)</pre>
      names(x) <- make.unique(names(x))</pre>
    class(x) <- cl</pre>
    return(x)
  if (missing(i) || missing(j))
    stop("only valid calls are x[[j]] <- value or x[[i,j]] <- value")</pre>
  rows <- attr(x, "row.names")</pre>
  nvars <- length(x)</pre>
  if (n <- is.character(i)) {</pre>
    ii <- match(i, rows)</pre>
    n <- sum(new.rows <- is.na(ii))</pre>
  if (n > OL) {
```

```
ii[new.rows] <- seq.int(from = nrows + 1L, length.out = n)</pre>
      new.rows <- i[new.rows]</pre>
    }
    i <- ii
  }
  if (all(i \ge 0L) \&\& (nn < max(i)) > nrows) {
    if (n == 0L) {
      nrr <- (nrows + 1L):nn</pre>
      if (inherits(value, "data.frame") && (dim(value)[1L]) >=
        length(nrr)) {
        new.rows <- attr(value, "row.names")[seq_len(nrr)]</pre>
        repl <- duplicated(new.rows) | match(new.rows,</pre>
          rows, OL)
        if (any(repl))
          new.rows[repl] <- nrr[repl]</pre>
      else new.rows <- nrr</pre>
    }
    x <- xpdrows.data.frame(x, rows, new.rows)
    rows <- attr(x, "row.names")</pre>
    nrows <- length(rows)</pre>
  iseq <- seq_len(nrows)[i]</pre>
  if (anyNA(iseq))
    stop("non-existent rows not allowed")
  if (is.character(j)) {
    if ("" %in% j)
      stop("column name \"\" cannot match any column")
    jseq <- match(j, names(x))</pre>
    if (anyNA(jseq))
      stop(gettextf("replacing element in non-existent column: %s",
        j[is.na(jseq)]), domain = NA)
  else if (is.logical(j) || min(j) < OL)</pre>
    jseq <- seq_along(x)[j]</pre>
  else {
    jseq <- j
    if (max(jseq) > nvars)
      stop(gettextf("replacing element in non-existent column: %s",
        jseq[jseq > nvars]), domain = NA)
  if (length(iseq) > 1L || length(jseq) > 1L)
    stop("only a single element should be replaced")
  x[[jseq]][[iseq]] <- value
  # here is where x is copied again!
  class(x) <- cl</pre>
# tracemem[0x7fdc992ae9d8 -> 0x7fdc9be55258]:
```

}

Thus seeing exactly where the three as of R 4.0: two! copies are happening.

х

Chapter 3

Vectors

3.2.1 Scalars

Can you have NA in vector Hell yeah!

3.2.3 Missing values

NA is a 'sentinel' value for explicit missingness - what does 'sentinel' mean?

A sentinel value (also referred to as a flag value, trip value, rogue value, signal value, or dummy data) is a special value in the context of an algorithm which uses its presence as a condition of termination. Also worth noting two NAs are not equal to each other! For instance, in C++ there's a special character to identify the end of a string I think another example of a sentinel value might be in surveys where you sometimes see missing data or N/A coded as 999, or 9999 (or maybe just 9)

Another example of a sentinel value might be in surveys where you sometimes see missing data or N/A coded as 999, or 9999 (or maybe just 9). The possible values in a column of data might be:

```
factor(c(1,1,1,1,2,3,3,4,4,9), levels = c(1,2,3,4,9))
```

Sentinels are typically employed in situations where it's easier/preferable to have a collection of values of the same type - represented internally using the same conventions and requiring the same amount of memory - but you also need a way to indicate a special circumstance. So like in the case of survey data you may, for example, see a variable indicating that an individual is 999 years old but the correct way to interpret that is that the data was not collected.

3.2.4 Testing and coercion

Why does the book warn us against using is.vector(), is.atomic() and is.numeric()? [read docs]

- is.atomic will also return true if NULL
- is.numeric tests if integer or double NOT factor, Date, POSIXt, difftime
- is.vector will return false if it has attributes other than names

3.3.1 Setting Attributes

Working in the medical field I have to import SAS files a lot where the column names have to adhere to specific abbreviations so they're given a label attribute for their full name. What are some other common

uses for attributes?

Hypothesis test attributes!

3.3.2 setNames

We can use setNames to apply different values to each element in a vector. How do we do this for our own custom attribute? The code below does NOT work!

```
my_vector <- c(
    structure(1, x = "firstatt_1"),
    structure(2, x = "firstatt_2"),
    structure(3, x = "firstatt_3")
)

my_vector <- setNames(my_vector, c("name_1", "name_2", "name_3"))

# mental model: shouldn't this should return $names and $x?
attributes(my_vector)</pre>
```

```
## $names
## [1] "name_1" "name_2" "name_3"
```

As soon as you instantiate a vector the attributes are lost. BUT we can store it as a list *within* the vector to keep them! We can create a custom attribute function and use that with map to add a list inside our dataframe:

```
custom_attr <- function(x, my_attr) {
   attr(x, "x") <- my_attr
   return(x)
}

as_tb <-
   tibble(
   one = c(1,2,3),
   x = c("att_1", "att_2", "att_3"),
   with_arr = map2(one, x, ~custom_attr(.x, .y))
)

as_tb$with_arr</pre>
```

```
## [[1]]
## [1] 1
## attr(,"x")
## [1] "att_1"
##
## [[2]]
## [1] 2
## attr(,"x")
## [1] "att_2"
##
## [[3]]
## [1] 3
## attr(,"x")
## [1] "att_3"
```

3.3.3 Dimensions

Because NROW and NCOL don't return NULL on a one dimensional vector they just seem to me as a more flexible option. When do you *have* to use ncol and nrow?

It may be better practice to always use NROW and NCOL!

As long as the number of rows matches the data frame, it's also possible to have a matrix or array as a column of a data frame. (This requires a slight extension to our definition of a data frame: it's not the length() of each column that must be equal, but the NROW().)

```
df \leftarrow data.frame(x = c(1,2,3,4,5),
                  y = c(letters[1:5])
length(df$y) == NROW(df$y)
## [1] TRUE
What's an example of where length() != NROW()
The case of a matrix!
my_matrix <- 1:6</pre>
dim(my_matrix) \leftarrow c(3,2)
my_matrix
##
         [,1] [,2]
## [1,]
           1
## [2,]
## [3,]
            3
length(my_matrix) == NROW(my_matrix) # 6 != 3
## [1] FALSE
length(my_matrix) == NCOL(my_matrix) # 6 != 2
```

```
## [1] FALSE
```

The length of the matrix is 6, and if we manipulate the dimensions of the matrix we see that the NROW is 3 and and NCOL is 2.

3.4 S3 atomic vectors

How is data type typeof() different from class()?

Classes are built on top of base types - they're like special, more specific kinds of types. In fact, if a class isn't specified then class() will default to either the implicit class or typeof.

So Date, POSIXct, and difftime are specific kinds of doubles, falling under its umbrella.

```
lubridate::is.Date(Sys.Date())

## [1] TRUE

is.double(Sys.Date())

## [1] TRUE

lubridate::is.POSIXct(Sys.time())

## [1] TRUE
```

```
is.double(Sys.time())
## [1] TRUE
lubridate::is.difftime(as.difftime(c("0:3:20", "11:23:15")))
## [1] TRUE
is.double(as.difftime(c("0:3:20", "11:23:15")))
## [1] TRUE
But then why does my_factor fail to be recognized under its more general integer umbrella?
my_factor <- factor(c("a", "b", "c"))
is.factor(my_factor)
## [1] TRUE
is.integer(my_factor)
## [1] FALSE
XXX</pre>
```

3.4.2 Dates

Why are dates calculated from January 1st, 1970?

Unix counts time in seconds since its official "birthday," – called "epoch" in computing terms – which is Jan. 1, 1970. This article explains that the early Unix engineers picked that date arbitrarily, because they needed to set a uniform date for the start of time, and New Year's Day, 1970, seemed most convenient.

3.5.1 Lists

When should you be using list() instead of c()

It's really contingent on the use case. In the case of adding custom classes it's worth noting that those are lost once you c() those objects together!

3.6.8 Data frames and tibbles

What does 'lazy' mean in terms of as_tibble?

Technically lazy evaluation means that expressions are not evaluated when they are bound to variables, but their evaluation is deferred until their results are needed by other computations. In this context though we think Hadley just meant that it's treated as a character if it "looks and smells like a character".

The solution manual gives the answer and notes:

```
df_coltypes <- data.frame(
    a = c("a", "b"),
    b = c(TRUE, FALSE),
    c = c(1L, OL),
    d = c(1.5, 2),
    e = c("one" = 1, "two" = 2),
    g = factor(c("f1", "f2")),
    stringsAsFactors = FALSE</pre>
```

```
as.matrix(df_coltypes)
```

```
## a b c d e g
## one "a" "TRUE" "1" "1.5" "1" "f1"
## two "b" "FALSE" "0" "2.0" "2" "f2"
```

"Note that format() is applied to the characters, which gives surprising results: TRUE is transformed to "TRUE" (starting with a space!)."

...But where is the format() call happening? I don't see a space!

After running debug(as.matrix(df_coltypes)) and going down a rabbit hole we found this is a bug that has been addressed! See issue here

Conclusion

How does vectorization make your code faster

Taking the example from Efficient R Programming:

VECTORIZED:

```
sum(log(x))
```

NON-VECTORIZED:

```
s <- 0
for(x0 in x) {
   s <- s + log(x0)
}</pre>
```

The vectorized code is faster because it obeys the golden rule of R programming: "access the underlying C/Fortran routines as quickly as possible; the fewer functions calls required to achieve this, the better".

- Vectorized Version:
 - 1. sum [called once]
 - 2. log [called once]
- Non-vectorized:
 - 1. + [called length(x) times]
 - 2. log [called length(x) times]

In the vectorised version, there are two primitive function calls: one to log (which performs length(x) steps in the C level) and one to sum (which performs x updates in the C level). So you end up doing a similar number of operations at C level regardless of the route.

In the non-vectorised form you are passing the logic back and forth between R and C many many times and this is why the non-vectorised form is much slower.

A vectorized function calls primitives directly, but a loop calls each function length(x) times, and there are 1 + length(x) assignments to s. Theres on the order of 3x primitive function calls in the non-vectorised form!!

Resources:

• Check out Jenny Brian's slides

• Great article by Noam Ross

Putting it all together in a single diagram:



Chapter 4

Subsetting

4.1 Introduction

"There are three subsetting operators [. [[, \$. What is the distinction between an operator and a function? When you look up the help page it brings up the same page for all three extraction methods. What are their distinctions and do their definitions change based on what you're subsetting? Can we make a table?

```
[
[[
$
ATOMIC
RETURNS VECTOR WITH ONE ELEMENT
SAME AS [
NOPE!
LIST
```

RETURNS A LIST

RETURNS SINGLE ELEMENT FROM WITHIN LIST

RETURN SINGLE ELEMENT FROM LIST [CAN ONLY USE WHEN LIST VECTOR HAS A NAME]

MATRIX

RETURNS A VECTOR

RETURNS A VECTOR OR SINGLE VALUE

NOPE!

DATA FRAME

RETURNS A VECTOR OR DATA FRAME

RETURNS VECTOR/LIST/MATRIX OR SINGLE VALUE

RETURNS VECTOR/LIST/MATRIX USING COLUMN NAME

TIBBLE

RETURNS A TIBBLE

RETURNS A VECTOR OR SINGLE VALUE

RETURNS THE STR OF THE COLUMN - TIBBLE/LIST/MATRIX

If we think of everything as sets (which have the properties of 0,1, or many elements), if the set has 1 element it only contains itself and NULL subsets. Before you subset using [or [[count the elements in the set. If it has zero elements you are done, if it has one element [will return itself - to go further you need to use [[to return its contents. If there is more than one element in the set then [will return those elements. You can read more about subsetting here

4.2.1 Selecting multiple elements

Why is numeric(0) "helpful for test data?"

This is more of a general comment that one should make sure one's code doesn't crash with vectors of zero length (or data frames with zero rows)

Why is subsetting with factors "not a good idea"

Hadley's notes seem to say subsetting with factors uses the "integer vector of levels" - and if they all have the same level, it'll just return the first argument. Subsetting a factor vector leaves the factor levels behind unless you explicitly drop the unused levels

4.2.2 lists

We've been talking about $\$ as a shorthand for [[. Using the example list $x \leftarrow list(1:3, "a", 4:6)$ can we use x\$1 as shorthand for x[[1]]?

The "shorthand" refers to using the name of the vector to extract the vector. If we give 1:3 a name such as test = 1:3

```
x <- list(named_vector = 1:3, "a", 4:6)
x[[1]] == x$named_vector</pre>
```

```
## [1] TRUE TRUE TRUE
```

As such, \$ is a shorthand for x[["name_of_vector"]] and not shorthand for x[[index]]

4.3.1 [[

The book states:

While you must use [[when working with lists, I'd also recommend using it with atomic vectors whenever you want to extract a single value. For example, instead of writing:

```
for (i in 2:length(x)) {
  out[i] <- fun(x[i], out[i - 1])
}</pre>
```

It's better to write

```
for (i in 2:length(x)) {
  out[[i]] <- fun(x[[i]], out[[i - 1]])
}</pre>
```

Why? Can we see this in action by giving x, out, and fun real life values?

If we have a vector

```
df_x <- c("Advanced", "R", "Book", "Club")
We can use [ or [[ to extract the third element of df_x
df_x[3]
## [1] "Book"

df_x[[3]]
## [1] "Book"
But in the case where we want to extract an element from a list [ and [[ no longer give us the same results
df_x <- list(A = "Advanced", B = "R", C = "Book", D = "Club")

df_x[3]
## $C
## [1] "Book"
df_x[3]]</pre>
```

Because using [[returns "one element of this vector" in both cases, it makes sense to default to [[instead of [since it will reliably return a single element.

4.3.5 Exercise

[1] "Book"

The question asks to describe the upper.tri function - let's dig into it!

```
x \leftarrow outer(1:5, 1:5, FUN = "*")
upper.tri(x)
         [,1]
               [,2]
                     [,3]
                           [,4]
                                 [,5]
## [1,] FALSE TRUE TRUE
                          TRUE
                                 TRUE
## [2,] FALSE FALSE TRUE
                          TRUE
                                 TRUE
## [3,] FALSE FALSE FALSE
                          TRUE
## [4,] FALSE FALSE FALSE FALSE
                                 TRUE
## [5,] FALSE FALSE FALSE FALSE
```

We see that it returns the upper triangle of the matrix. But I wanted to walk through how this function actually works and what is meant in the solution manual by leveraging $.row(dim(x)) \le .col(dim(x))$.

The function .row() and .col() return a matrix of integers indicating their row number

```
.row(dim(x))
         [,1] [,2] [,3] [,4] [,5]
##
## [1,]
            1
                 1
                       1
                            1
## [2,]
            2
                 2
                       2
                            2
                                  2
## [3,]
            3
                 3
                       3
                            3
                                  3
## [4,]
            4
                 4
                       4
                            4
                                  4
## [5,]
            5
                                  5
.col(dim(x))
         [,1] [,2] [,3] [,4] [,5]
## [1,]
                 2
            1
                       3
## [2,]
                                  5
            1
                 2
                       3
                            4
                                  5
## [3,]
            1
                 2
                       3
## [4,]
            1
                 2
                       3
                                  5
                       3
                                  5
## [5,]
            1
.row(dim(x)) \leftarrow .col(dim(x))
          [,1]
                [,2]
                       [,3]
                             [,4] [,5]
                             TRUE TRUE
## [1,] TRUE
                TRUE
                       TRUE
## [2,] FALSE TRUE
                       TRUE
                             TRUE TRUE
## [3,] FALSE FALSE
                       TRUE
                             TRUE TRUE
## [4,] FALSE FALSE FALSE
                             TRUE TRUE
## [5,] FALSE FALSE FALSE FALSE TRUE
```

Is there a high level meaning to a . before function? Does this refer to internal functions? [see: ?row vs ?.row]

Objects in the global environment prefixed with . are hidden in the R (and RStudio) environment panes - so functions prefixed as such are not visible unless you do ls(all=TRUE). Read more here and (here)[https://stackoverflow.com/questions/7526467/what-does-the-dot-mean-in-r-personal-preference-naming-convention-or-more]

4.3.3 Missing and OOB

Let's walk through examples of each

LOGICAL ATOMIC

```
c(TRUE, FALSE)[[0]] # zero length
# attempt to select less than one element in get1index <real>
c(TRUE, FALSE)[[4]] # out of bounds
# subscript out of bounds
c(TRUE, FALSE)[[NA]] # missing
# subscript out of bounds
```

LIST

```
list(1:3, NULL)[[0]] # zero length
# attempt to select less than one element in get1index <real>
list(1:3, NULL)[[3]] # out of bounds
# subscript out of bounds
```

```
list(1:3, NULL)[[NA]] # missing
# NULL
```

NULL

```
NULL[[0]] # zero length
# NULL
NULL[[1]] # out of bounds
# NULL
NULL[[NA]] # missing
# NULL
```

4.5.8 Logical subsetting

"Remember to use the vector Boolean operators & and |, not the short-circuiting scalar operators && and ||, which are more useful inside if statements."

Can we go over the difference between & and && (and | vs | |) I use brute force to figure out which ones I need...

&& and | | only ever return a single (scalar, length-1 vector) TRUE or FALSE value, whereas | and & return a vector after doing element-by-element comparisons.

The only place in R you routinely use a scalar TRUE/FALSE value is in the conditional of an if statement, so you'll often see && or || used in idioms like: if (length(x) > 0 && any(is.na(x))) { do.something()}

In most other instances you'll be working with vectors and use & and | instead.

Using && or || results in some unexpected behavior - which could be a big performance gain in some cases:

- ullet | | will not evaluate the second argument when the first is TRUE
- && will not evaluate the second argument when the first is ${\tt FALSE}$

```
true_one <- function() { print("true_one evaluated."); TRUE}</pre>
true two <- function() { print("true two evaluated."); TRUE}</pre>
# arguments are evaluated lazily. Unexpected behavior can result:
c(T, true_one()) && c(T, true_two())
## [1] "true_one evaluated."
## [1] "true_two evaluated."
## [1] TRUE
c(T, true_one()) && c(F, true_two())
## [1] "true_one evaluated."
## [1] "true_two evaluated."
## [1] FALSE
c(F, true_one()) && c(T, true_two())
## [1] "true_one evaluated."
## [1] FALSE
c(F, true_one()) && c(F, true_two())
## [1] "true_one evaluated."
```

```
## [1] FALSE
c(T, true_one()) || c(T, true_two())
## [1] "true_one evaluated."
## [1] TRUE
c(T, true_one()) || c(F, true_two())
## [1] "true_one evaluated."
## [1] TRUE
c(F, true_one()) || c(T, true_two())
## [1] "true_one evaluated."
## [1] "true_two evaluated."
## [1] TRUE
c(F, true_one()) || c(F, true_two())
## [1] "true_one evaluated."
## [1] "true_one evaluated."
## [1] "true_one evaluated."
## [1] "true_two evaluated."
## [1] "true_two evaluated."
```

4.5.8 Boolean algebra

Read more about Special Primatives here

The unwhich() function takes a boolean and turns it into a numeric - would this ever be useful? How?

XXX

"x[-which(y)] is not equivalent to x[!y]: if y is all FALSE, which(y) will be integer(0) and -integer(0) is still integer(0), so you'll get no values, instead of all values."

Can we come up with an example for this plugging in values for x and y

```
c(TRUE, FALSE)[-which(FALSE)]

## logical(0)
c(TRUE, FALSE)[!FALSE]

## [1] TRUE FALSE
```

Chapter 5

Control Flow

5.2.2 Vectorised if

```
Why does if else print "Out of range" once but if else prints it twice?
if (sample(1:10, 1) == 1) print("In range") else print("Out of range")
[1] "Out of range"
ifelse((sample(1:10, 1) == 1), print("In range"), print("Out of range"))
[1] "Out of range"
[1] "Out of range"
var <- if (sample(1:10, 1) == 1) print("In range")</pre>
var
NULL
var <- ifelse((sample(1:10, 1) == 1), print("In range"), print("Out of range"))</pre>
## [1] "In range"
var
## [1] "In range"
ifelse explicitly returns its result, whereas if invisibly returns it!
ifelse(c(TRUE, FALSE, TRUE), 1:2, 3)
## [1] 1 3 1
Honestly, ifelse() is weird. Try this too:
ifelse(c(TRUE, FALSE, TRUE), 1:10, 3)
## [1] 1 3 3
5.3 Loops
Can the body of the loop change the set?
my_set <- c(1, 20, 99)
```

```
for (i in my_set){
   if (i==1){
      my_set[9]= 20
    }
   print("hello")
   print(my_set)
}

## [1] "hello"
## [1] 1 20 99 NA NA NA NA NA 20
## [1] "hello"
## [1] 1 20 99 NA NA NA NA NA 20
## [1] "hello"
## [1] "hello"
## [1] 1 20 99 NA NA NA NA NA 20
```

5.2.3 switch statement

Looks like you can't!

"Closely related to if is the switch() statement. It's a compact, special purpose equivalent"

What is meant here by "special purpose"? Can we come up with a case where you can't substitute if for switch or the other way around? Use switch. Is it safe to say to switch when you have character inputs (as the book suggests) and use case_when or if for numerics?

Switch is special in that only ONE value can be true, as in the case from our presentation the shiny input can only ever be ONE of the strings on the left it cannot be multiple.

5.3.1 common pitfalls

What does the book mean by leveraging vector for preallocating the loop output? How is this different from creating an empty list to instantiate the loop?

```
means <- c(1, 50, 20)

out_vector <- vector("list", length(means))
for (i in 1:length(means)) {
   out_vector[[i]] <- rnorm(10, means[[i]])
}

out_list <- list()
for (i in 1:length(means)) {
   out_list[[i]] <- rnorm(10, means[[i]])
}</pre>
```

By preallocating the length of the out_vector we're leveraging modify-in-place rather than copy-on-modify.

The book warns against using length(x) and suggests using seq_along instead. Is there any downside to using seq_along or a case when you'd prefer length(x) over seq_along(x)? I can't think of any downsides to using it!

We have yet to find a downside but should look into this further!

5.3.3.1 Exercises

```
x <- numeric()
out <- vector("list", length(x))
for (i in 1:length(x)) {
  out[i] <- x[i] ^ 2
}
out</pre>
```

```
## [[1]]
## [1] NA
```

I understand that this works because we can count down in a loop - so the first iteration x[1] == NA, and the second x[2] == numeric(0) but where does this value go? Is it just evaluated but not stored since R counts from 1?

This question is actually the bigger question, "Can you assign something to index 0 in R?" and it doesn't seem that you can.

Assignment to a valid index

```
mylist = list()
mylist[1] <- c("something")
mylist[1]
## [[1]]</pre>
```

Assignment to [0]

[1] "something"

```
mylist = list()
mylist[0] <- c("something")
mylist[0]</pre>
```

```
## list()
```

It's interesting that it's syntactically correct to assign to mylist[0] but it does nothing!

Empty index

```
mylist = list()
mylist[1]
```

```
## [[1]]
## NULL
```

R defaultly accesses the first layer using [and says there are no elements in the first layer of mylist and returns NULL. But when you do mylist[[1]] R says Out of Bounds because the first layer of mylist has been called and there is no layer at all. That's why R throws error in [[case.

Impossible Index [[

```
mylist = list()
mylist[[0]]
```

Error in mylist[[0]] : attempt to select less than one element in get1index <real>

[[isolates one element whereas [can be used to get subsets of varying sizes. [[gets precisely one element, no more no less, or it fails. So [[0]]<- cannot work because index zero cannot hold a value.

Impossible Index [[<-

```
mylist = list()
mylist[[0]] <- "something"</pre>
```

Error in mylist[[0]] <- "something" : attempt to select less than one element in integerOneIndex

Undefined name

Selection using an undefined - but possible name - gives NULL

```
mylist = list()
mylist[["undefined_name"]]
```

Out of Bounds

Selection using a possible - but not currently allocated - index gives an out of bounds error

```
mylist = list()
mylist[[10]]
```

Error in mylist[[10]] : subscript out of bounds

Conclusion

Can we put these concepts together in diagram form? Let's work on improving these schematics!

Let's first put functions into words:

- If condition then run code, else quit
- A vectorized version of: if condition then run code, else run other code
- For every element in list of elements do what is in between the curly braces
- While this condition is true do this
- Repeat this until something tells you to break out



We can visualize how the ifelse logic on a single element above will operate on each element within a vector:



For instance if we can run:

```
ifelse(c(TRUE, FALSE, TRUE), 1:2, 3)
```

[1] 1 3 1

Lets break down what's happening:

Create a test answer:

```
test <- c(TRUE, FALSE, TRUE)
yes <- 1:2
no <- 3
```

which indices in the test are TRUE and which are FALSE

```
yes_idx <- which(test) # 1,3
no_idx <- which(!test) # 2</pre>
```

set up our answer

```
answer <- test # T, F, T
```

grow the yes and no answers to the length of the test (input)

```
yes_final <- rep(yes, length.out = length(test))
no_final <- rep(no, length.out = length(test))</pre>
```

fill the answer with yes or no from the enlarged yes/no answers

```
answer[yes_idx] <- yes_final[yes_idx] # 1,1
answer[no_idx] <- no_final[no_idx] # 3</pre>
```

return our final answer:

answer

```
## [1] 1 3 1
```

Another example: we can run

```
ifelse(c(TRUE, FALSE, TRUE, FALSE, TRUE), 1:10, "FALSE")
```

we can see that if else places the numbers in 1:10 based on their index where our condition is to TRUE and inserts the string "FALSE" whenever the condition is FALSE

Complexity

Cyclomatic complexity is a software metric used to indicate the complexity of a program. It is computed using the control flow graph of the program: the nodes of the graph correspond to indivisible groups of commands of a program, and a directed edge connects two nodes if the second command might be executed immediately after the first command.

We can use the package cyclocomp to look at the cyclomatic complexity of functions. Looking at each function from the Chapter 5 presentation:

```
controlflow_functions <- list(
   if_func = if_func,
   if_else_func = if_else_func,
   ifelse_func = ifelse_func,
   casewhen_func = casewhen_func,
   switch_func = switch_func,
   for_func = for_func,
   while_func = while_func,
   repeat_func = repeat_func)</pre>
```

We see that the for loop was our most complex function and while had a complexity of 3. The rest of our functions had a complexity of 1.

As Colin Fay states:

"Splitting a complex function into smaller functions is not a magic solution because (A) the global complexity of the app is not lowered by splitting things into pieces (just local complexity), and (B) the deeper the call stack, the harder it can be to debug."

Chapter 6

Functions

6.2.2 Primitives

So if you are familiar with C can you just write a function in C in R? What does that process look like? I think this is a bigger question of digging into the relationship between C and R.

Primitives are part of R core, and can only be written by the R-core team. At its heart, R is the set of primitive C functions underneath it.

You can use Rcpp to include C++ code in your R code, but these aren't Primitives. There are also other techniques which we'll likely see covered in later chapters. Here's an example using Rcpp.

```
Rcpp::cppFunction('#include<string>
  std::string IPA() {
  std::string val = "IPAs suck";
  return val;
  }')
val <- IPA()
val</pre>
```

[1] "IPAs suck"

Are there any non-base primitives? If so how is that possible!

XXX

6.2.5.1 Exercises

This question is flagged as "started" let's try to complete it!

1. **Q**: Given a name, like "mean", match.fun() lets you find a function. Given a function, can you find its name? Why doesn't that make sense in R?

A: A name can only point to a single object, but an object can be pointed to by 0, 1, or many names. What are names of the functions in the following block?

```
function(x) sd(x) / mean(x)

## function(x) sd(x) / mean(x)

f1 <- function(x) (x - min(x)) / (max(x) - min(x))
f2 <- f1
f3 <- f1</pre>
```

There isn't a 1 to 1 mapping between functions and names in R. Multiple names may point to the same function as we see for f1, f2, and f3. Also, each function has its own environment so it's possible that two functions might have the same "code" but are not the same because they have different environments (or closures). Lastly, anonymous functions don't have names so we'd have no way to look these up.

We could find the names of our functions if they are all in the global environment using body(x) = body(y)

```
names(which(sapply(mget(lsf.str()), function(x) body(x) == body(function(x) (x - min(x)) / (max(x) - min(x))
```

```
[1] f1 f2 f3
```

But that's just deparsing the body into a string and comparing the values. So if you want to think of two functions as being equal if their deparsed body strings as the same, then that's technically possible but that is just like searching for every variable that has the value of 5 [possible but not efficient].

The main point is that name -> object is a one way (non-unique) look up in R. There's no efficient way to go backwards. This is true for all values, not just functions.

6.3 Function composition

When comparing nested, intermediate, and piping functions, it looks like Hadley flips the order of f() and g() between bullet points

It does look like he does that!

```
f <- function(z) {
  cat("g is:", z)
}

g <- function(x) {
  x * 2
}</pre>
```

6.0.1 Nested

```
f(g(2))
```

g is: 4

6.0.2 Intermediate

This is written in the book as $y \leftarrow f(x)$; g(y) but should be flipped to $y \leftarrow g(x)$; f(y) if we are to follow the nested example

```
y <- g(2)
f(y)
```

g is: 4

6.0.3 Piping

This also needs to be flipped from x %>% f() %>% g() to x %>% g() %>% f()

```
2 %% g() %%% f()
```

```
## g is: 4
```

6.4 Lexical scoping

"The scoping rules use a parse-time, rather than a run-time structure"? What is "parse-time" and "run-time"? How do they differ?

parse-time is when the function gets defined: when the formals and body get set. run-time is when it actually gets called. This function doesn't get past parse-time because of the syntax error

```
get_state <- function(in_df, state_name){</pre>
  out_df % in_df[in_df$state == state_name, ]
  return(out_df)
}
get_state <- function(in_df, state_name){</pre>
  out_df % in_df[in_df$state == state_name, ]
Error: unexpected input in:
"get_state <- function(in_df, state_name){</pre>
  out_df % in_df[in_df$state == state_name, ]"
  return(out_df)
Error: object 'out_df' not found
}
Error: unexpected '}' in "}"
This function will get parsed successfully but could fail at run at run-time if the input data frame doesn't
have a column named state:
get_state <- function(in_df, state_name){</pre>
  out_df <- in_df[in_df$state == state_name, ]</pre>
  return(out_df)
get_state(iris, 'setosa')
## [1] Sepal.Length Sepal.Width Petal.Length Petal.Width Species
## <0 rows> (or 0-length row.names)
```

At R's build-time, if you want to run a function from a package that isn't loaded it will not throw an error but at run-time it will if the required package is not loaded:

Without dplyr this will fail

```
func_1(test_tbl, x1)
```

Error in select({: could not find function "select"

This will work:

```
library(dplyr)
#>
#> Attaching package: 'dplyr'
#> The following objects are masked from 'package:stats':
#>
#> filter, lag
#> The following objects are masked from 'package:base':
#>
#> intersect, setdiff, setequal, union
func_1(test_tbl, x1)
```

```
## # A tibble: 5 x 1
## x1
## < dbl>
## 1 0.385
## 2 0.188
## 3 0.245
## 4 0.747
## 5 0.271
```

6.4.3 A fresh start

How would we change this code so that the second call of g11() is 2?

```
g11 <- function() {
  if (!exists("a")) {
    a <- 1
  } else {
    a <- a + 1
  }
  a
}</pre>
```

```
## [1] 1
g11 <- function() {
   if (!exists("a")) {
      a <- 1
   } else {
      a <- a + 1
   }
   a <<- a
   a
}
g11()</pre>
```

```
## [1] 1
```

6.5 Lazy evaluation

"This allows you to do things like include potentially expensive computations in function arguments that will only be evaluated if needed"

Does anyone have an example of this? We discussed a function that will only perform expensive tasks given the context of the function perhaps?

Maybe a situation where we can give a function default arguments where sample is a stand in for longer expensive functions like different fancy modeling techniques? We can workshop this...

```
mega_model <- function(values, x = sample(1:100, 10), y = sample(200:300, 10), z = sample(300:400, 10)
) {
    dplyr::case_when(
        is.numeric(values) ~ x,
        is.character(values) ~ y,
        TRUE ~ z
    )
}
mega_model(c("a", "b", "c"))</pre>
```

[1] 201 237 203 279 225 242 286 260 243 264

6.5.1 Promises

Can we discuss the order that this happening in? Is it that Calculating... is printed, then x*2 then x*2 again? I am still reading this as: h03(double(20), double(20)) which is an incorrect mental model because the message is only printed once...

```
double <- function(x) {
  message("Calculating...")
  x * 2
}
h03 <- function(x) {
  c(x, x)
}</pre>
```

Explain what's happeining here below in words, and restructure the promise image to make more sense

```
double <- function(y) {
  message("Calculating...")
  cat("double before\n")
  print(pryr::promise_info(y))
  force(y)
  cat("double after\n")
  print(pryr::promise_info(y))
  y * 2
}
h03 <- function(x) {
  cat(paste0("h03 before\n"))
  print(pryr::promise_info(x))
  force(x)
  cat("h03 after\n")
  print(pryr::promise_info(x))
  c(x, x)</pre>
```

##

```
double(h03(20))
## Calculating...
## double before
## $code
## h03(20)
##
## $env
## <environment: R_GlobalEnv>
##
## $evaled
## [1] FALSE
## $value
## NULL
##
## h03 before
## $code
## [1] 20
##
## $env
## <environment: R_GlobalEnv>
## $evaled
## [1] FALSE
##
## $value
## NULL
##
## h03 after
## $code
## [1] 20
##
## $env
## NULL
## $evaled
## [1] TRUE
##
## $value
## [1] 20
##
## double after
## $code
## h03(20)
##
## $env
## NULL
##
## $evaled
## [1] TRUE
```

```
## $value
## [1] 20 20
## [1] 40 40
```



```
plop <- function(a, b) a * 10
plop(2, var_doesnt_exist)</pre>
```

```
## [1] 20
```

var_doesnt_exist is a promise within g, we use the promises within g when we call f but f never uses its second argument so this runs without a problem. When would we want to leverage this behavior?

The unevaluated <code>var_doesnt_exist</code> doesn't exist , but we can use <code>substitute</code> to to get the expression out of a promise! If we modify our function we can play with the expression contained in b:

```
plop <- function(a, b) {
  cat("You entered", deparse(substitute(b)), "as `b` \n")
  a * 10
}
plop(a = 2, b = var_doesnt_exist)</pre>
```

You entered var_doesnt_exist as `b`

[1] 20

We can even evaluate b and use it to create a dplyr like pull function:

```
plop <- function(a, b) {
   eval(substitute(b), envir = a)
}
plop(iris, Species)[1:10]</pre>
```

[1] setosa setosa setosa setosa setosa setosa setosa setosa setosa ## Levels: setosa versicolor virginica

6.5.2 Default arguments

I don't quite understand why x = ls() is different from ls() here; aren't we still assigning x = ls() but without specifying x?

```
h05 <- function(x = ls()) {
  a <- 1
```

[1] "h05"

```
X
}
# this makes sense to me
h05()
## [1] "a" "x"
# how is this different from above?
h05(ls())
    [1] "a"
##
                                 "answer"
                                                           "as_tb"
   [4] "casewhen_func"
                                 "controlflow_functions"
                                                           "custom_attr"
    [7] "df"
                                 "df_coltypes"
                                                           "df_non_syntactic_name"
## [10] "df x"
                                 "double"
                                                           "example2223"
                                 "f1"
                                                           "f2"
```

```
## [13] "f"
## [16] "f3"
                                                            "func_1"
                                  "for_func"
## [19] "g"
                                  "g11"
                                                            "get_state"
                                  "h05"
                                                            "i"
## [22] "h03"
## [25] "if_else_func"
                                  "if func"
                                                            "ifelse_func"
## [28] "IPA"
                                  "11"
                                                            "12"
                                  "medians"
                                                            "mega_model"
## [31] "means"
                                  "my_matrix"
                                                            "my_set"
## [34] "my_factor"
                                  "mylist"
                                                            "new_address"
## [37] "my_vector"
## [40] "new_y_address"
                                  "no"
                                                            "no_final"
                                  "old_address"
                                                            "out"
## [43] "no_idx"
## [46] "out_list"
                                  "out_vector"
                                                            "plop"
## [49] "repeat_func"
                                  "subsetting_table"
                                                            "switch_func"
                                  "test_tbl"
                                                            "true_one"
## [52] "test"
## [55] "true_two"
                                  "v"
                                                            "v_address"
## [58] "val"
                                  "var"
                                                            "while_func"
## [61] "x"
                                  "x1"
                                                            "x2"
                                  "x4"
                                                            "x5"
## [64] "x3"
## [67] "y"
                                  "y_address"
                                                            "yes"
## [70] "yes_final"
                                  "yes_idx"
```

The difference is where the promise is created. ls() is always evaluated inside h05 when x is evaluated. The difference is the environment. When ls() is provided as an explicit parameter, x is a promise whose environment is the global environment. When ls() is a default parameter, it is evaluated in the local environment where it is used.

Hypothesis: does nesting ls() in h05 first evaluate ls() then evaluate h05()?

```
library(magrittr)
h05 <- function(x = {y <- 4;ls()}) {
    a <- 1
    x
}
h05()

[1] "a" "x" "y"
ls()</pre>
```

```
{y <- 4;ls()} %>% h05()

[1] "h05" "y"

ls()

[1] "h05" "y"

h05({x <- 5;ls()})

[1] "h05" "x" "y"

ls()

[1] "h05" "x" "y"</pre>
```

Notice in all of the latter calls, a is not returned - so it's not evaluating 1s() inside of the function.

6.5.4.3 Exercise

I understand this problem is showing us an example of name masking (the function doesn't need to use the y = 0 argument because it gets y = 0 argument because y = 0 argument because y = 0 argument because y = 0 argument y = 0 argum

```
y <- 10
f1 <- function(x = {y <- 1; 2}, y = 0) {
  c(x, y)
}
f1()</pre>
```

[1] 2 1

The curly brackets are an expression, which can be read as

```
{
    y <- 1
    2
}</pre>
```

[1] 2

This is returning 2 and setting 1 to y. The colon can be read as a new line in the expression. x is called inside the function and overwrites the argument value of y

```
f1 <- function(x = {y <- 1; 2}, y = 0) {
    # At this point, neither x nor y have been needed or evaluated. This is "lazy".
    cat(y)
    # "Summon the current state of y".
    # R looks first in the function env - is there a y there? No
    # R then looks at the arguments - is there a y there? Yes, it's 0 -> print out zero
    # If R had not found y in the arguments, then it would look in the parent environment of the function
    # That's where it would find y = NULL - but since it already found a default arg, it already stopped
    cat(x)
    # "Summon the current state of x"
    # x is an expression that first sets y to equal 1 and then returns the number 2 -> print out 2
    c(x, # "Summon the current state of x" - x is still the expression that sets y to 1 and then x to 2
    y) # "Summon the current state of y" - y was set to 1, so y is 1 here.
}
f1()
```

```
## 02
## [1] 2 1
```

Compare to:

```
f1 <- function(x = {y <- 1; 2}, y = 0) {
   c(y, # "Summon the current state of y" - y has not yet been set to anything else, so it returns zero
      x) # "Summon the current state of x" - x is still the expression that sets y to 1 and then x to 2
      # However, y has already been returned - so even though y is now set to 1, it's called after the
}
x = NULL
y = NULL
f1()</pre>
```

[1] 0 2

What is happening here: - The default value of x is "assign 1 to y then return 2 implicitly." - The default value of y is 0. - x and y aren't their defaults until they're referenced and there isn't a value, so when you invoke the question x is the default, but y is never the default. - When you get to c(x, x) is evaluated. Now the value of y is 1 instead of its default value! - When you get to c(x, y), y is now 1, so the return is 2 1

Using the original f1 function, if we write f1(x <- 5) we get 5 0, When you make that call, the function's x argument gets set to x <- 5 (instead of the default). When you get to c(x), the x <- 5 call gets evaluated in the calling environment (global, most likely, unless you're calling it from inside a function or something).

To see where x gets assigned, try this:

```
f2 <- function(x = {y <- 1; 2}, y = 0) {
    # c(x, y)
    "Nothing to see here."
}
f2(x <- 5)
## [1] "Nothing to see here."
x</pre>
```

```
## Error: object 'x' not found
```

Since x is never used in this version, the $x = \{x < -5\}$ promise never gets evaluated, so x never gets set in the calling environment. But if you do the same thing with f1, x is now 5 in the calling environment.

Also note that calling the <- function returns the value (the second argument) invisibly, so y <- $\{x <- 5\}$ assigns 5 to both y and x. I wouldn't recommend ever doing this on purpose, but it's useful to know for debugging weird cases.

A piece that **ALMOST** confused me was that the function's default value only ever "exists" in the function's environment, not in the calling environment, so the original case doesn't change y to 1 globally. But f1({y <- 1; 2}) **WILL** change y globally... but does not change the value of y inside the function.

6.5.4.4 Exercise

I know this isn't exactly needed to answer the question, but how do we access a function that has methods? For instance - here I want to dig into the hist function using hist

hist

```
## function (x, ...)
## UseMethod("hist")
## <bytecode: 0x7ffd834fbdc8>
```

```
## <environment: namespace:graphics>
```

does not give me the actual contents of the actual function....

We need to access is using hist. <method>

hist.default

6.6 dot dot dot

"(See also rlang::list2() to support splicing and to silently ignore trailing commas..." Can we come up with a simple use case for list2 here? The docs use list2(a = 1, a = 2, b = 3, b = 4, 5, 6) but how is this different from list?

```
identical(
  rlang::list2(a = 1, a = 2, b = 3, b = 4, 5, 6) ,
  list(a = 1, a = 2, b = 3, b = 4, 5, 6)
)
```

[1] TRUE

list2 is most helpful when we need to force environment variables with data variables. We can see this by creating a function that takes a variable number of arguments:

```
library(rlang)

numeric_list <- function(...) {
   dots <- list(...)
   num <- as.numeric(dots)
   set_names(num, names(dots))
}

numeric_list2 <- function(...) {
   dots <- list2(...)
   num <- as.numeric(dots)
   set_names(num, names(dots))
}

numeric_list(1, 2, 3)</pre>
```

```
## [1] 1 2 3
numeric_list2(1, 2, 3)
```

[1] 1 2 3

The main difference with list(...) is that list2(...) enables the !!! syntax to splice lists:

"lapply() uses ... to pass na.rm on to mean()" Um, how?

```
x <- list(2, 3)
numeric_list2(1, !!! x, 4)
```

```
## [1] 1 2 3 4
numeric_list(1, !!! x, 4)
```

```
numeric_list(1, !!! x, 4)
Error in !x : invalid argument type
```

```
x <- list(c(1, 3, NA), c(4, NA, 6))
str(lapply(x, mean, na.rm = TRUE))</pre>
```

```
## List of 2
## $ : num 2
## $ : num 5
```

An lapply takes on two main arguments: what you want to loop over and the function to apply to each element. By including ... lapply allows you to supply additional arguments which will be passed to the function inside the lapply. In this case, na.rm = TRUE is being applied to mean every time it's being called in the loop.

6.6.1.2 Exercise

I tried running browser(plot(1:10, col = "red")) to peek under the hood but only got Called from: top level in the console. What am I missing?

We can use debugonce!

```
debugonce(plot())
```

6.7.4 Exit handlers

"Always set add = TRUE when using on.exit() If you don't, each call to on.exit() will overwrite the previous exit handler." What does this mean?

add = TRUE is important when you have more than one on.exit function!

```
j08 <- function() {
  on.exit(message("a"))
  on.exit(message("b"), add=TRUE)
}
j08()</pre>
```

a

b

Can we go over this code? How does it not change your working directory after you run the function

```
cleanup <- function(dir, code) {
  old_dir <- setwd(dir)
  on.exit(setwd(old_dir), add = TRUE)

  old_opt <- options(stringsAsFactors = FALSE)
  on.exit(options(old_opt), add = TRUE)
}

cleanup("~")
getwd()</pre>
```

[1] "/Users/mayagans/Documents/bookclub-Advanced_R/QandA"

The behavior of setwd "changing the working directory" is actually a side effect of the function - it invisibly returns the previous working directory as the value of the function (potentially for the exact purpose demonstrated). We can use this within our on.exit function to change back to the prior working directory!

If on.exit fails will it continue onto the next on.exit so long as add == TRUE ξ 'on.exit fails it'll continue onto the next one

```
f <- function() {
  on.exit(stop("Error"))
  on.exit(message("yay, still called."), add = TRUE)
}
f()</pre>
```

```
Error in f() : Error
yay, still called.
```

6.7.5.4 Exercise

This question is flagged as "started" let's try to complete it! Hadley comments in the repo: "I think I'm more interested in supplying a path vs. a logical value here".

Q: How does the chdir parameter of source() compare to with_dir()? Why might you prefer one approach to the other? The with_dir() approach was given in the book as

```
in_dir <- function(dir, code) {
  old <- setwd(dir)
  on.exit(setwd(old))

force(code)
}</pre>
```

A: with_dir() takes a path to a working directory as an argument. First the working directory is changed accordingly. on.exit() ensures that the modification to the working directory are reset to the initial value when the function exits.

with_dir gives you the flexibility to change the path to wherever you want (maybe a parent-level folder) whereas source(chdir=TRUE) changes the path to "where that file lives specifically".

Given the following file structure:

```
bookclub-Advanced R
    .gitignore
    AdvancedR. Rproj
    README.html
    README.md
 -data
      .Rhistory
      beer reviews.csv.gz
      beer reviews.data
      import data.R
 -images
      controlflow.png
      copy on modify fig1.png
      copy on modify fig2.png
      median benchmark.png
      r4ds.png
      vectors.png
```

Imagine I want to run import_data.R, but it needs to reference images/controlflow.png. we can do this by setting the wd to advRbookclub:

```
with_dir("/advRbookclub/data/import_data.R","/advRbookclub")
```

Or we can use:

```
source("/advRbookclub/data/import_data.R", chdir=TRUE)
```

but then, we'd need to include something like setwd(here::here()) in import_data.R so that it goes back to AdvancedR.Rproj and sets the working directory there.

In conclusion:

- source is a base R function so reduces dependencies. Once set, you could use setwd("..") assuming you can have some confidence that it's part of a repository or something?
- with_dir is exported from the withr package but is a more fine-tuned control by passing the specific folder name as opposed to a boolean TRUE | FALSE.

6.7.5.5 Exercise

Can we go over the source code of capture.output and capture.output2?

There were several new terms here to me when going over this function:

- stderr error output
- stdout normal output

```
#!/usr/bin/env Rscript
# redirect.r -- experiments with warning and error messages
# Get any arguments (and ignore them)
arguments <- commandArgs(trailingOnly=TRUE)</pre>
# Default
write("1) write() to stderr", stderr())
write("1) write() to stdout", stdout())
warning("1) warning()")
# Ignore all warnings
options(warn = -1)
write("2) write() to stderr", stderr())
write("2) write() to stdout", stdout())
warning("2) warning()")
# Send all STDERR to STDOUT using sink()
options(warn = 0) # default setting
sink(stdout(), type="message")
write("3) write() to stderr", stderr())
write("3) write() to stdout", stdout())
warning("3) warning()")
# Send all STDOUT to STDERR using sink()
sink(NULL, type="message") # default setting
sink(stderr(), type="output")
write("4) write() to stderr", stderr())
write("4) write() to stdout", stdout())
warning("4) warning()")
# Send messages and output to separate files
msg <- file("message.Rout", open="wt")</pre>
out <- file("output.Rout", open="wt")</pre>
```

```
sink(msg, type="message")
sink(out, type="output")
write("5) write() to stderr", stderr())
write("5) write() to stdout", stdout())
warning("5) warning()")
```

• sink diverts R output to a connection.

```
zz <- file("all.Rout", open = "wt")
sink(zz)
sink(zz, type = "message")
try(log("a"))
## revert output back to the console -- only then access the file!
sink(type = "message")
sink()
file.show("all.Rout")</pre>
```

- textConnection allows R character vectors to be read as if they were being read from a text file.
 - They can capture R output to a character vector
 - They can be used to create a new character object or append to an existing one in the user's workspace.
 - At all times the complete lines output to the connection are available in the R object. Closing the connection writes any remaining output to a final element of the character vector.

```
function (..., file = NULL, append = FALSE, type = c("output",
                                                     "message"), split = FALSE)
  # Capture dots
  # [-1L] removes the list()
  args <- substitute(list(cat("a", "b", "c", sep = "\n")))[-1L]
  # match
  type <- match.arg(type)</pre>
  # set default return value
  rval <- NULL
  # set default closer
  closeit <- TRUE</pre>
  # if file is null, then
  if (is.null(file))
    # If file is null, then create a write-only text connection object which will
    # save to the variable rval in the execution environment (local = TRUE).
    # see https://biostatmatt.com/R/R-conn-ints.pdf for gritty info on connections (and sinks)
   file <- textConnection("rval", "w", local = TRUE)</pre>
  else if (is.character(file))
    # if "file" is a character vector, then interpret it as a filename. Open a
    # file connection in either append or write mode, depending on the value of
    # "append"
   file <- file(file, if (append)
      "a"
      else "w")
  else if (inherits(file, "connection")) {
    # if "file" is already a connection object, check if it is open. If not, open it
    # in append mode, if specified, otherwise in write mode.
    # inherits refers to the S3 class system.
```

```
# Browse[2]> class(file)
  # [1] "textConnection" "connection"
 if (!isOpen(file))
    open(file, if (append)
      "a"
      else "w")
  # if the connection is already open, don't close it in this function.
 else closeit <- FALSE</pre>
}
# if you get here, then you misspecified "file"
else stop("'file' must be NULL, a character string or a connection")
# sink all output of type "type" into the connection "file". If you would like
# the output to continue to its original source, then "split" it.
# by default, messages (messages, warnings, errors) go to stderr and
# everything else to stdout.
sink(file, type = type, split = split)
on.exit({
  # on exit, call sink with the same arguments and without "file" being specified.
  # this will cause the sink from the line before to terminate.
 sink(type = type, split = split)
  # Close the connection (always, unless "file" was provided as
 # an already open connection)
 if (closeit) close(file)
})
# store the calling environment in pf. i.e. pf refers to the environment in
# which capture.output was called.
pf <- parent.frame()</pre>
# define a local function which will evaluate its sole argument (expr) in the
# parent frame.
evalVis <- function(expr) withVisible(eval(expr, pf))</pre>
# for each argument collected in the dot dot.
# use split = TRUE to help you debug at this point. When you try to poke around
# with the sink applied, you are rightfully stymied because all output is going
# to the sink connection!
for (i in seq_along(args)) {
 # store the argument in expr
 expr <- args[[i]]</pre>
  # based on the mode of the expr, evaluate it.
 tmp <- switch(mode(expr),</pre>
                expression = lapply(expr, evalVis),
                call = ,
                name = list(evalVis(expr)),
                stop("bad argument"))
  # print any visible values output during evaluation.
  # This print will be collected by the sink we set up earlier and saved to
  # the file connection.
 for (item in tmp) if (item$visible)
   print(item$value)
# calling on.exit with no arguments will clear the exit handler.
# we are doing this because the on.exit was designed to safeguard the sink
```

```
# and files being closed in the case of an exception when evaluating the
  # passed in arguments.
  on.exit()
  # undo the sink
  sink(type = type, split = split)
  # close the file if necessary
  if (closeit)
    close(file)
  # return the captured output or null invisibly otherwise.
  if (is.null(rval))
    invisible(NULL)
  else rval
capture.output2 <- function(code) {</pre>
 temp <- tempfile()</pre>
  on.exit(file.remove(temp), add = TRUE)
  sink(temp)
  on.exit(sink(), add = TRUE)
 force(code)
 readLines(temp)
}
```

The second function will always sink output to a temporary file, and then return the results by reading the file back in (and returning a character vector). It uses two exit handlers, one to clean up the temporary file, and one to remove the sink.

6.8.4 Replacement functions

Can we put into words the translation for

We can dig into the source code, but the jist is that in order to implement these complex assignments:

- 1. Copy x into a temporary variable *temp*
- 2. [<-(names(*tmp*), 2, "two") modifies the second element of the names of *temp*,
- 3. names<-(*tmp* assigns step 2 to *temp* names
- 4. Clean up by removing the temp variable

6.8.6.3 Exercise

This question is flagged as "started" let's try to complete it!

Q: Explain why the following code fails:

A: First, let's define x and recall the definition of modify() from the textbook:

```
x <- 1:3

`modify<-` <- function(x, position, value) {
   x[position] <- value
   x
}</pre>
```

R internally transforms the code and the transformed code reproduces the error above.

```
get("x") <- `modify<-`(get("x"), 1, 10)
#> Error in get("x") <- `modify<-`(get("x"), 1, 10) :
#> target of assignment expands to non-language object
```

The error occurs during the assignment, because no corresponding replacement function, i.e. get<- exists for get(). To confirm this we can reproduce the error via the following simple example.

```
get("x") \leftarrow 2
#> Error in get("x") \leftarrow 2: target of assignment expands to non-language object
```

I don't really see why this needs to be expanded upon....

Chapter 7

Environments

7.2.3 Parents

What's the relationship between base, empty enviornment, and global enviornment?

The working environment is the environment whose names would currently mask names in any other environment. Note that in the book, the term current environment is used not working environment. At any given time there is an environment that is most immediately-visible i.e. its names will mask the same names in any other environment. This is the current environment. The identity of the current environment can change e.g. when entering a function. It is often the case that the current environment is the global environment i.e. typically when working interactively.

Are functions the only mechanism for changing the current environment? (I believe the answer is yes.)

Consider the following from the book:

The current environment, or current_env() is the environment in which code is currently executing. When you're experimenting interactively, that's usually the global environment, or global_env().

The use of the word "usually" in this sentence raises the possibility that there's an exception. That exception could possibly mean

- 1. There's a way to change the environment in interactive mode without going into a function
- 2. You can experiment interactively within a function or
- 3. Something I'm ignorant of.

I know that you can step-through a function in the debug browser so maybe that's what he means?

But then the very next line in the text says:

The global environment is sometimes called your "workspace", as it's where all interactive (i.e. outside of a function) computation takes place.

It both i) defines interactive specifically to exclude the inside of a function and ii) makes a stronger claim, "all interactive computation" vs "usually".

If you put a breakpoint or browser() call inside a function, then you would be working interactively in a function environment.

Does hierarchy of environments have a computer science name (linked list?)?

The hierarchy of environments is not a linked list, it is a directed graph.

Clarification: the global environment's parent is all the packages loaded within the environment, not the empty environment, right? How do we check this? parent.env(global)?

Each package becomes a parent to the global environment. We can inspect local environments using search_envs()

search_envs()

```
[[1]] $ <env: global>
   [[2]] $ <env: package:sloop>
   [[3]] $ <env: package:memoise>
  [[4]] $ <env: package:gt>
## [[5]] $ <env: package:openintro>
##
  [[6]] $ <env: package:rlang>
   [[7]] $ <env: package:cyclocomp>
## [[8]] $ <env: package:lobstr>
## [[9]] $ <env: package:kableExtra>
## [[10]] $ <env: package:forcats>
## [[11]] $ <env: package:stringr>
## [[12]] $ <env: package:dplyr>
## [[13]] $ <env: package:purrr>
## [[14]] $ <env: package:readr>
## [[15]] $ <env: package:tidyr>
## [[16]] $ <env: package:tibble>
## [[17]] $ <env: package:ggplot2>
## [[18]] $ <env: package:tidyverse>
## [[19]] $ <env: package:stats>
## [[20]] $ <env: package:graphics>
## ... and 6 more environments
```

7.2.4 Super assignment

Let's expand on the concept of super assignment

```
x <- 0
f <- function() {
    x <- 2
    x <<- 1
    x
}</pre>
```

```
## [1] 2
x
```

[1] 1

Note that the assignment inside f is local, but super assignment "never creates a variable in the current environment" so it modifies the global x and not the local x.

If you want to break your brain a bit, check out:

```
x <- 0
f <- function() {
    x <- x
    x <<- x + 1
    x</pre>
```

```
f()

## [1] 0

x

## [1] 1

f()

## [1] 1

x

## [1] 2

f()

## [1] 2

x

## [1] 3
```

7.2.5 Getting and setting

"But you can't use [[with numeric indices, and you can't use [:"

It makes sense to me that you can't use numeric indexes because objects in an environment aren't ordered, but why can't you use [? The solutions manual states: "The second option would return two objects at the same time. What data structure would they be contained inside?

[returns an object of the same type, if we were to apply this to environments we'd have an environment returning an environment.

7.2.6 Advanced bindings

Hadley mentions delayed bindings are used when autoloading datasets with packages - can we find an example of this? How is this different from including LazyData: true in your description file?

The LazyData: true entry in DESCRIPTION just informs that delayed binding should be used. autoload does use delayed bindings [from source code: do.call("delayedAssign", list(name, newcall, .GlobalEnv, .AutoloadEnv))] but lazydata that gets used by packages is ultimately implemented as an internal function. If you scan through View(loadNamespace):

```
if (file.exists(paste0(dbbase, ".rdb")))
    lazyLoad(dbbase, env)
    dbbase <- file.path(pkgpath, "data", "Rdata")
    if (file.exists(paste0(dbbase, ".rdb")))
        lazyLoad(dbbase, .getNamespaceInfo(env, "lazydata"))</pre>
```

and lazyLoad eventually calls .Internal(makeLazy(vars, vals, expr, db, envir)) which leads you to https://github.com/wch/r-source/blob/726bce63825844715860d35fdf76539445529f52/src/main/builtin.c# L103 and ultimately the delayed binding is realized as a promise with defineVar(name, mkPROMISE(expr0, eenv), aenv);

7.3 Recursing over environments

How can we re-write where so that it returns all functions with the same name?

```
where2 <- function(in_name, env = caller_env()) {</pre>
  all_functions <- ""
  index <-1
  while (!identical(env, empty_env())) {
    # if success
    if (env_has(env, in_name)) {
      all_functions[index] <- env_name(env)</pre>
      index <- index + 1</pre>
      #return()
    }
    # inspect parent
    env <- env_parent(env)</pre>
  # base case - I'm missing this I think
  return(all_functions)
# load dplyr so you have two filters
library(dplyr)
where2("filter")
## [1] "package:dplyr" "package:stats"
We can also create a function that emulates pryr's where:
where3 <- function(name, env = parent.frame(), found = character()) {</pre>
  if (identical(env, emptyenv()))
    return(found)
  else if (exists(name, env, inherits = FALSE)) {
    Recall(name, parent.env(env), c(found, environmentName(env)))
 } else {
    Recall(name, parent.env(env), found)
```

```
## [1] "package:dplyr" "package:stats"
```

7.3.1.2 Exercises

where3("filter")

}

I understood the recursion in the prior example, but what is inherits doing here? Can we go through this line for line and discuss what is happening in this function?

```
fget <- function(name, env = caller_env(), inherits = TRUE) {
    # Base case
    if (env_has(env, name)) {
        obj <- env_get(env, name)

        if (is.function(obj)) {
            return(obj)
        }
    }</pre>
```

```
if (identical(env, emptyenv()) || !inherits) {
    stop("Could not find function called \"", name, "\"", call. = FALSE)
}

# Recursive Case
fget(name, env_parent(env))
}

# Test
mean <- 10
fget("mean", inherits = TRUE)

## function (x, ...)
## UseMethod("mean")
## <bytecode: 0x7ffd86d0c9b0>
## <environment: namespace:base>
```

Inherits is an argument that stops the function from performing the recursive action of looking into the parent environment for the name. Inherits is acting like a valve. If false, and the name wasnt found in the current environment, then stop because the next expression searches the parent of env. In the case environment is the empty environment then stop because the empty environment doesnt have a parent.

7.4.2 Function environment

I found this section a little confusing. Can we go over the second figure in the section? Where does x live? g points to x but x is in the global environment? Can we come up with our own example for a function being bound to the global environment but accessing variables from its own environment? (I think this is what the second figure in the section is trying to display)

```
y <- 1
e <- env()
e$g <- function() 1
e$g</pre>
```

```
## function() 1
```

This can be seen in our where2 example in the recursion section!

7.4.3 Namespaces

"Every binding in the package environment is also found in the namespace environment; this ensures every function can use every other function in the package. But some bindings only occur in the namespace environment. These are known as internal or non-exported objects, which make it possible to hide internal implementation details from the user."

When you're developing a package does the namespace environment just come with your package for free when you build it or you need to create both your package and its namespace env?

How do you create functions that exist only in the namespace environment?

When you are building the package, everything in the NAMESPACE file gets generated by roxygen2. Attaching a package puts the package in the search path of namespaces. So when you run library(package) or require(package) it creates the namespace environment. Package developers control what names are available by exporting to namespace (and otherwise you can call non-exported with :::).

For example let's say you want to extend forcats to use ordered factors to lump the tails. In order to get it to work, we need to use some forcats helper functions. We could just copy the code for the functions and added them to our scripts, but for prototyping it's easy enough to just use :::

```
fct_lump_ordered <- function(f, n, prop, q, w = NULL,</pre>
                                other_level_low = "Other Low",
                                other_level_high = "Other High",
                                ties.method = c("min", "average", "first", "last", "random", "max")) {
  f <- check ordered(f)</pre>
  w <- forcats:::check_weights(w, length(f))</pre>
  ties.method <- match.arg(ties.method)</pre>
  levels <- levels(f)</pre>
  if (is.null(w)) {
    count <- as.vector(table(f))</pre>
    total <- length(f)
  } else {
    count <- as.vector(tapply(w, f, FUN = sum))</pre>
    total <- sum(w)</pre>
  }
  if (all(missing(n), missing(prop), missing(q))) {
    lump <- forcats:::in_smallest(count)</pre>
    lump <- lump_range(!lump)</pre>
    new_levels <- ifelse(lump == -1L, other_level_low, ifelse(lump == 1L, other_level_high, levels))</pre>
  } else if (!missing(n) ) {
    if (n < 0) {
      rank <- rank(count, ties = ties.method)</pre>
      n <- -n
    } else {
      rank <- rank(-count, ties = ties.method)</pre>
    if (sum(rank > n) <= 1) {
      return(f)
    }
    lump <- lump_range(rank <= n)</pre>
    new_levels <- ifelse(lump == -1L, other_level_low, ifelse(lump == 1L, other_level_high, levels))</pre>
  } else if (!missing(prop)) {
    prop_n <- count/total</pre>
    if(prop < 0) {
      lump <- lump_range(prop_n <= -prop)</pre>
      new_levels <- ifelse(lump == -1L, other_level_low, ifelse(lump == 1L, other_level_high, levels))</pre>
    } else {
      if (sum(prop_n <= prop) <= 1) {</pre>
        return(f)
      lump <- lump_range(prop_n >= prop)
      new_levels <- ifelse(lump == -1L, other_level_low, ifelse(lump == 1L, other_level_high, levels))
  } else if (!missing(q)) {
    cdf <- cumsum(count)/sum(count)</pre>
    lump \leftarrow lump\_range(cdf >= q[1] \& cdf \leftarrow= q[2])
    new_levels <- ifelse(lump == -1L, other_level_low, ifelse(lump == 1L, other_level_high, levels))</pre>
  }
  if (other_level_low %in% new_levels && other_level_high %in% new_levels) {
    f <- forcats::lvls_revalue(f, new_levels)</pre>
```

```
forcats::fct_relevel(f, other_level_low)
  forcats::fct_relevel(f, other_level_high, after = Inf)
} else if (other_level_low %in% new_levels) {
  f <- forcats::lvls_revalue(f, new_levels)
    forcats::fct_relevel(f, other_level_low)
} else if (other_level_high %in% new_levels) {
  f <- forcats::lvls_revalue(f, new_levels)
    forcats::fct_relevel(f, other_level_high, after = Inf)
}
else {
  f
}</pre>
```

How do conflicted packages identify duplicate function names and print them out

We can use conflict_scout!

```
conflicted::conflict_scout
```

```
## function (pkgs = NULL)
## {
##
       pkgs <- pkgs %||% pkgs_attached()</pre>
       objs <- lapply(pkgs, pkg_ls)</pre>
##
##
       names(objs) <- pkgs
##
       index <- invert(objs)</pre>
##
       potential <- Filter(function(x) length(x) > 1, index)
       unique <- Map(unique_obj, names(potential), potential)
##
       conflicts <- Filter(function(x) length(x) > 1, unique)
##
       conflicts <- map2(names(conflicts), conflicts, superset principle)</pre>
##
##
       conflicts <- map2(names(conflicts), conflicts, drop_moved)</pre>
##
       for (fun in ls(prefs)) {
##
            if (!has_name(conflicts, fun))
##
           conflicts[[fun]] <- prefs_resolve(fun, conflicts[[fun]])</pre>
##
##
##
       conflicts <- compact(conflicts)</pre>
##
       new_conflict_report(conflicts)
## }
## <bytecode: 0x7ffd8c069af0>
## <environment: namespace:conflicted>
```

7.5 Call stacks

What exactly is a frame?

A frame is a singular step within the CST - in thr following example, f, g, and h are each frames

```
f <- function(x) {
   g(x = 2)
}
g <- function(x) {
   h(x = 3)
}
h <- function(x) {
   stop()</pre>
```

}

7.6 Data Structures

Can we discuss what's happening in this function? Since it's similar to the setwd function from last week could we build on this function to include on.exit()?

```
my_env <- new.env(parent = emptyenv())
my_env$a <- 1

get_a <- function() {
    my_env$a
}

set_a <- function(value) {
    old <- my_env$a
    my_env$a <- value
    invisible(old)
}</pre>
```

We can use those two functions like we did getting and setting the work directory inside another function - using the old value within on.exit to reset my_env\$a outside of the function: Hadley suggests returning invisible old so that you don't need to explicitly call get_a, you can just assign the output of set_a

```
set_a <- function(value) {
  old <- my_env$a
  my_env$a <- value
  invisible(old)
}

do_thing_where_env_a_is_value(value, code) {
  myoldenv <- set_a(value) # sets a to the NEW value and returns the OLD value for storage
  on.exit(set_a(myoldenv), add=TRUE)
  force(code)
}</pre>
```

The bullet point mentions hashmaps but I'm still really unclear on what that is. What are hash tables and are they related?

They are sort of like fast look-up tables! It may be impossible to show with a small example. To demonstrate that the lookup is fast you'd need to store a lot of stuff in it first. You'd also probably need an alternative implementation of a lookup table to compare it to. We could probably implement a hash table relatively easily* using environments. Demonstrating that using environments is a good way to do it seems difficult.

• looking at a data structures textbook (Cormen, Lieserson & Rivest), a hash table is expected to provide users with fast functions for insert, search, and delete.

Chapter 8

Conditions

8.3 Ignoring conditions

I can't quite think of a time where try is more appropriate than tryCatch - does anyone have an example?

It seems that try is just a wrapper for tryCatch and you can use it whenever you'd use try

When would you actually use suppressWarning? Maybe when loading libraries? I created an example for suppressMessages, but does someone have a better, practical use case?

As a function user, it's common to suppress warnings when loading in data and using readr and dplyr-but not so much as a function creator

8.4.2 Exiting handlers

```
tryCatch(
  message = function(cnd) "There",
  {
    message("Here")
    stop("This code is never run!")
  }
)
#> [1] "There"
```

The protected code is evaluated in the environment of tryCatch(), but the handler code is not, because the handlers are functions.

Clarification question what is the "handler code" and "protected code"?

The protected code is inside the {} and the handler code is message = function(cnd) "There"

In the tryCatch example where we have finally print "Thank God for Beer" I find it interesting that this is printed **before** the code inside the tryCatch. Can anyone explain why?

From the help it's, "expression to be evaluated before returning or exiting." That means it's the final thing that happens inside the tryCatch, but it happens before the return (of NA or the string), so it happens first.

8.4.3 Calling handlers

How would you define bubbling up?

Bubbling up: By default, a condition will continue to propagate to parent handlers, all the way up to the default handler (or an exiting handler, if provided)

I would consider what is being described there "bubbling up".

Why is this message executed once per message in the function?

```
withCallingHandlers(
  message = function(cnd) cat("Caught a message!\n"),
  {
    message("Someone there?")
    message("Why, yes!")
  }
)
```

```
#> Caught a message!
#> Someone there?
#> Caught a message!
#> Why, yes!
```

withCallingHandlers could be understood as: "for each {message}, do x"

The following prints the message once:

```
withCallingHandlers(
  message = function(cnd) cat("Caught a message!\n"),
  {
    warning("Someone there?")
    message("Why, yes!")
  }
)
```

```
## Warning in withCallingHandlers(message = function(cnd) cat("Caught a message!
## \n"), : Someone there?
## Caught a message!
## Why, yes!
```

and it comes in the warning message, so it's not really a "direct" print

Oh wait, the handler code is named message, so is it possible that the two message calls in the protected block are actually calling both base::message and the message - named handler?

The return value of a calling handler is ignored because the code continues to execute after the handler completes; where would the return value go? That means that calling handlers are only useful for their side-effects.

Can we come up with an example for this masking? I think seeing it will help me understand...

This just means that you cannot capture the return value of your handler (but you sort of can)

```
f <- function() {
  my_lovin <- NULL
  withCallingHandlers(
    message = function(cnd) {my_lovin <<- "not this time"},
    {
       never_ever_gunna_get_it = message("No, you're never gonna get it")
    }
  )
  my_lovin</pre>
```

```
}
f()
## No, you're never gonna get it
## [1] "not this time"
How does muffling differ from suppressWarnings?
muffling allows for an over-ride/replacement of messages, while suppress just quiets everything:
fn <- function() {</pre>
  inform("Beware!", "my_particular_msg")
  inform("On your guard!")
  "foobar"
}
Now we can use a new definition of my_particular_msg to replace "Beware"
with_handlers(fn(),
  my_particular_msg = calling(function(cnd) {
    inform("Dealt with this particular message")
    cnd muffle(cnd)
  })
## Dealt with this particular message
## On your guard!
## [1] "foobar"
whereas suppressMessages just returns foobar:
suppressMessages(fn())
## [1] "foobar"
```

8.4.5.2 Exercises

In the example

```
show_condition <- function(code) {
  tryCatch(
    error = function(cnd) "error",
    warning = function(cnd) "warning",
    message = function(cnd) "message",
    {
      code
      NULL
    }
  )
}</pre>
```

The first three calls to show_condition make sense to me, and I even understand that the first time the code is evaluated inside a tryCatch it exists (that's why it returns message) but how are you supplying 3 arguments to a function that just takes on one argument, code? Is that what the {} are for?

```
show_condition({
  10
```

b

```
message("?")
warning("?!")
})
```

{} let us execute multiple lines of code!

8.4.5.3 Exercises

I couldn't follow the manual's answer for what's happening here, can we come up with our own answer for what's happening here in words?

```
withCallingHandlers(
  message = function(cnd) message("b"),
  withCallingHandlers(
    message = function(cnd) message("a"),
    message("c")
  )
)
## b
## a
```

The first call to withcallinghandlers adds a condition handler for conditions with class "message" to the handler stack (not sure if its actually a stack?) and then executes the second withcallinghandlers which adds another condition handler for conditions with class "message" to the handler stack (see above), and then executes the code message("c").

What happens next is a chain reaction of handlers. The call message("c") is handled by the inner handler, which then calls message("a"), but message("a") is caught by the outer handler, and so it outputs b first.

Then the inner handler resolves itself and outputs a and then that condition (having not been muffled) "bubbles up" to the outer handler, which calls message("b") again producing the second b and again since the message wasn't muffled, the original condition message("c") "bubbles up" to the top where it is evaluated producing the output c

Another exercise: Guess the output of these two functions:

```
i <- 1
withCallingHandlers(
  message = function(cnd) {message(paste0(i, ". b")); i <<- i+1},
  withCallingHandlers(
    message = function(cnd) {message(paste0(i, ". a")); i <<- i+1},
    {message(paste0(i, ". c"))}
)</pre>
```

```
## 1. b
## 1. a
## 3. b
## 1. c
```

```
i <- 1
withCallingHandlers(
    message = function(cnd) {i <<- i+1; message(paste0(i, ". b"))},
    withCallingHandlers(
        message = function(cnd) {i <<- i+1; message(paste0(i, ". a"))},
        {message(paste0(i, ". c"))}
    )

## 3. b

## 2. a

## 4. b

## 1. c</pre>
```

8.6.3 Resignal

Can we go over what is happening here?

```
warning2error <- function(expr) {
  withCallingHandlers(
    warning = function(cnd) abort(conditionMessage(cnd)),
    expr
)
}
warning2error({
  x <- 2 ^ 4
    warn("Hello")
})</pre>
```

Error: Hello

The function warning2error captures an expression which is evaluated by withCallingHandlers where you have defined a handler for warning conditions. The handler captures the condition end raised by warn which is structure(list(message = "Hello"), class = c("warning", "condition")

The function conditionMessage is an s3 generic which evaluates to conditionMessage.condition which simply accesses cnd\$message this is then the input to abort which raises an error with the message "Hello"

8.6.4 Record

Why are we using cnd_muffle here?

```
catch_cnds <- function(expr) {
  conds <- list()
  add_cond <- function(cnd) {
    conds <<- append(conds, list(cnd))
    cnd_muffle(cnd)
}

withCallingHandlers(
  message = add_cond,
  warning = add_cond,</pre>
```

```
expr
)

conds
}

catch_cnds({
  inform("a")
  warn("b")
  inform("c")
})
```

If we remove cnd_muffle we see that a, b, and c are printed to the console prior to getting the conds output.

Would it be possible in the second Record example to create a function that doesn't require us to put the abort statement at the end, just ignoring it? Or because abort is an exiting handler it needs to be last?

```
catch_cnds <- function(expr) {</pre>
  conds <- list()</pre>
  add_cond <- function(cnd) {</pre>
    conds <<- append(conds, list(cnd))</pre>
    cnd_muffle(cnd)
  tryCatch(
    error = function(cnd) {
      conds <<- append(conds, list(cnd))</pre>
    },
    withCallingHandlers(
      message = add cond,
      warning = add_cond,
      expr
  )
  conds
}
catch_cnds({
  abort("a")
  inform("b")
  warn("c")
})
```

I'm not positive, but I don't think so without mucking around in the C code. You "can" using try instead of tryCatch, but I can't think of a way to have it let you try each line in your passed in expression for example

What is signal and what is it doing here?

```
log <- function(message, level = c("info", "error", "fatal")) {
    # if we remove match.arg and just use level
    # the signal returns infoerrorfatal as a single string
    # we need match.arg to find the selected level
    # and it defaults to the first if none are selected
    level <- match.arg(level)
    signal(message, "log", level = level)</pre>
```

```
}
```

Signal is a general function that calls abort, inform or warn. It has the same signature as those functions except for the class argument that is necessary for signal but NULL by default for each of the others

```
signal <- function(message, class, ..., .subclass) {
  if (!missing(.subclass)) {
    deprecate_subclass(.subclass)
  }
  message <- collapse_cnd_message(message)
  cnd <- cnd(class, ..., message = message)
  cnd_signal(cnd)
}</pre>
```

If you create a condition object by hand, and signal it with signalCondition(), cnd_muffle() will not work. Instead you need to call it with a muffle restart defined, like this: withRestarts(signalCondition(cond), muffle = function() NULL)

Where does this code go given the prior example?

We would replace the code cnd_muffle(cnd):

```
ignore_log_levels <- function(expr, levels) {
    withCallingHandlers(
    log = function(cnd) {
        if (cnd$level %in% levels) {
            # cnd_muffle(cnd)
            withRestarts(signalCondition(cnd), muffle = function() NULL)
        }
    },
    expr
)
}
record_log(ignore_log_levels(log("Hello"), "warning"))</pre>
```

[info] "Hello"

8.6.6.2 Exercises

Calling handlers are called in the context of the call that signaled the condition. Exiting handlers are called in the context of the call to tryCatch().

What exactly does this mean? tryCatch evaluates what we were calling the protected code first and calling handlers execute the handling code first? Can we make a simple example?

- tryCatch is a project manager who oversees everything and then personally hands over the end product
- with Calling Handlers writes some procedures/guidelines and assumes everyone has enough information to get their jobs done

Comparing:

```
f <- function() g()
g <- function() h()
h <- function() message("!")
withCallingHandlers(f(), message = function(cnd) {
  lobstr::cst()</pre>
```

```
cnd_muffle(cnd)
})
##
     1. +-base::withCallingHandlers(...)
##
     2. +-global::f()
##
##
     3. | \-global::g()
##
     4. |
            \-global::h()
              \-base::message("!")
##
                +-base::withRestarts(...)
##
     6. I
                | \-base:::withOneRestart(expr, restarts[[1L]])
##
                     \-base:::doWithOneRestart(return(expr), restart)
##
     8. 1
##
     9. |
                \-base::signalCondition(cond)
    10. \-(function (cnd) ...
##
          \-lobstr::cst()
```

withCallingHandlers is run to completion before f is called/put onto the stack. While f is in progress, g needs to be called so g is put on top of the stack, then h etc. Eventually they're all completed and taken off the stack (in reverse order). Then the handler (function (cnd) ... is called

While

tryCatch is still on the stack when the handler is called i.e. value, tryCatchOne, tryCatchList and tryCatch are all still in line to be completed

Calling handlers are called in the context of the call that signaled the condition, in this case, f(). So the message handler returns a value to the environment where f() is a meaningful call. Exiting handlers are called in the context of the call to tryCatch(). The exiting handler returns to an ongoing tryCatch() so it can do whatever it needs to do.

8.6.6.4 Exercises

There's no way to break out of the function because we're capturing the interrupt that you'd usually use!

What does this mean? You can't stop the function if you set an interrupt argument inside a tryCatch?

That's exactly it - this was a warning not to use inturrupt in your functions!

Slides

What environment(s) these restarts are called within, or if that even applies or matters

Note: for this example I set my Environment panel in Rstudio to "Manual Refresh Only" (the curly arrow menu) while running this to make sure it wasn't doing anything to confuse me.

```
warning = function(w) { print(paste('warning:', w )); browser() },
                                # error print the error and send us to browser
                               error=function(e) { print(paste('e:',e )); browser()} ) {
  print(paste("big expensive step we don't want to repeat for x:",x))
  z <- x # the "expensive operation"
  print("Main function caller env:")
  print(rlang::caller_env())
  print("Main function current env:")
  print(rlang::current env())
  print("Main function parent:")
  print(rlang::env_parent(rlang::current_env(), 1))
  # second function on z that isn't expensive but could potentially error
  withRestarts(
    withRestarts(
      withCallingHandlers(
          print("withCallingHandlers caller_env:")
          print(rlang::caller_env())
          print("withCallingHandlers current_env:")
          print(rlang::current_env())
          print("withCallingHandlers parent:")
          print(rlang::env_parent(rlang::current_env(), 1))
          print(paste("attempt cheap operation for z:",z))
          return(log(z))
        },
        warning = warning,
        error = error
      ),
      force_positive = function() {
        z <<- -z
        print("force_positive caller_env:")
        print(rlang::caller_env())
        print("force_positive current_env:")
        print(rlang::current_env())
        print("force_positive parent:")
        print(rlang::env_parent(rlang::current_env(), 1))
      }
   ),
    set_to_one = function() {z <<- 1}</pre>
  )
expensive function(-1)
```

You can run it yourself (and then invokeRestart("force_positive")) to see the results, but to summarize:

- The withCallingHandlers part is executing in exactly the same environment as the function.
- The force_positive restart is executing in its own new environment.
- The parent of that environment is the execution environment of the main function.

The caller_env for force_positive, though, is still a bit of a mystery. Some further poking found that its grandparent is the base package's namespace, but I don't grok why. Maybe because it's the browser() environment? Is that a thing?

Are there any default restarts include in base R so that you could invokeRestart("XXX") or is the restart

always user defined?

From the warning documentation:

While a warning is being processed, a muffleWarning restart is available. If this restart is invoked with invokeRestart, then warning returns immediately.

We can also use the computeRestarts function which lists all available "default" restart functions:

```
computeRestarts()
```

```
[[1]]
<restart: abort >
```

It seems computeRestarts() is able to find abort even without {rlang} attached! Which means it surely would be able to find some function from the base packages if one existed. But it does not, so that leads me to believe that they aren't any beyond muffleWarning. it seems, in general, you need to specify a function name for invokeRestart()

What environment(s) these restarts are called within, or if that even applies or matters

Given this little example:

use the restart from the midlevel function

```
# low level function
# if text isn't correct then abort with class malformed_text
is_correct <- function(text) {</pre>
  print('parent in low: '); print(env_parent())
  print('current in low: '); print(env parent())
  if (text != "CORRECT") {
    msg <- pasteO(text, " needs to be 'CORRECT'")</pre>
    abort(
      "malformed_text",
      message = msg,
      text = text
    )
  }
  return(text)
}
# medium level function
# apply low level function to all text in a list
# if it doesn't pass give it the restart skip_text
# which will change that entry to NA
replace_not_correct <- function(all_text) {</pre>
  # f <- function(x) { print(current_env()); is_correct(x) }</pre>
  lapply(all_text, function(text) {
    withRestarts(
      # f(text),
      { print('parent in mid: '); print(env_parent()); print('current in mid: '); print(current_env())
      skip_text = function(e){ NA; print('parent in mid skip: '); print(env_parent()); print('current
    )
 })
}
# res <- replace_not_correct(list('A', 'B', 'C'))</pre>
# res <- replace_not_correct(list('CORRECT', 'A'))</pre>
# res
# high level function
# if you have an error of type malformed_text,
```

```
analyze_text <- function(all_text) {</pre>
  withCallingHandlers(
    malformed_text = function(e) invokeRestart("skip_text"),
    lapply(all_text, replace_not_correct)
  )
analyze_text(list("CORRECT", "NOT", "CORRECT"))
## [1] "parent in mid: "
## <environment: 0x7ffd899f7890>
## [1] "current in mid: "
## <environment: 0x7ffd899f8038>
## [1] "parent in low: "
## <environment: R_GlobalEnv>
## [1] "current in low: "
## <environment: R_GlobalEnv>
## [1] "parent in mid: "
## <environment: 0x7ffd899e8e00>
## [1] "current in mid: "
## <environment: 0x7ffd88b447b0>
## [1] "parent in low: "
## <environment: R_GlobalEnv>
## [1] "current in low: "
## <environment: R GlobalEnv>
## [1] "parent in mid skip: "
## <environment: 0x7ffd88b447b0>
## [1] "current in mid skip: "
## <environment: 0x7ffd8d6e2c58>
## [1] "parent in mid: "
## <environment: 0x7ffd8d6d0c48>
## [1] "current in mid: "
## <environment: 0x7ffd8d6d04d8>
## [1] "parent in low: "
## <environment: R_GlobalEnv>
## [1] "current in low: "
## <environment: R_GlobalEnv>
## [[1]]
## [[1]][[1]]
## [1] "CORRECT"
##
##
## [[2]]
## [[2]][[1]]
## <environment: 0x7ffd8d6e2c58>
##
##
## [[3]]
## [[3]][[1]]
## [1] "CORRECT"
```

So, the parent environment of the handler (is that the right term?) called skip_text in the mid-level function is the environment of the expr part of withRestarts (i.e. the first parameter), also in the mid-level function.

Are there any default restarts include in base R so that you could invokeRestart("XXX") or is the restart

always user defined?

From the warning documentation:

While a warning is being processed, a muffleWarning restart is available. If this restart is invoked with invokeRestart, then warning returns immediately.

We can also use the computeRestarts function which lists all available "default" restart functions:

```
computeRestarts()
```

```
[[1]]
<restart: abort >
```

It seems computeRestarts() is able to find abort even without {rlang} attached! Which means it surely would be able to find some function from the base packages if one existed. But it does not, so that leads me to believe that they aren't any beyond muffleWarning. it seems, in general, you need to specify a function name for invokeRestart()

Let's revisit my example from the talk: can we build on this to use tidyeval so that the user can write beer_states %>% beer_mean(state) [I think we need to use .data and note how state is given as an object so we need to use tidyeval to suppress it's evaluation...]

```
beer_states <- readr::read_csv(</pre>
  'https://raw.githubusercontent.com/rfordatascience/tidytuesday/master/data/2020/2020-03-31/beer_state
beer_mean <- function(.data, x) {</pre>
  column_name <- deparse(substitute(x))</pre>
  msg <- glue::glue("Can't calculate mean, {column_name} is not numeric")
  if (!is.numeric(.data[[column name]])) {
    abort(
      message = msg,
      arg = column_name,
      data = .data
    )
  } else {
    mean(which(!is.na(.data[[column_name]])[.data[[column_name]]]))
  }
}
beer_states %>% beer_mean(barrels)
```

[1] 806.4551

Chapter 9

Functionals

9.2.2 Anonymous functions and shortcuts

We saw you'll get an error if you try to map over elements that don't exist, and can use .default to override that. Is this related to tryCatch in some way? Can we look at the map source code for .default? And how would we overcome this error if we were to use base R's lapply(x, 'two')

```
x <- list(
    list(one = "a"),
    list(two = "b"),
    list(three = "c")
)

map_chr(x, 'two', .default = NA)

## [1] NA "b" NA
lapply(x, 'two')</pre>
```

Error in get(as.character(FUN), mode = "function", envir = envir) : object 'two' of mode 'function' was

Within ${\tt purrr}$ there's the function

```
find_extract_default <- function(.null, .default) {
   if (!missing(.null)) {
      .null
   } else {
      .default
   }
}</pre>
```

So it doesn't seem to be a conditional, but rather a way to deal with missing errors

```
as.character(lapply(x,`[[`,"two"))
```

```
[1] "NULL" "b" "NULL"
```

9.2.6.4 Exercise

In order to extract p-values the solution manual suggests using map_dbl but could we use pluck to get these values? If so how?

```
trials <- map(1:100, ~ t.test(rpois(10, 10), rpois(10, 7)))
# tibble(p_value = map_dbl(trials, "p.value"))
map_dbl(trials, pluck, "p.value")
     [1] 1.154984e-01 2.519610e-02 2.631330e-01 8.875650e-02 8.587759e-01
##
     [6] 1.404054e-02 2.160124e-02 2.101519e-02 7.674028e-03 3.496454e-03
##
    [11] 2.711314e-02 2.645881e-02 2.914382e-03 1.520664e-01 7.761718e-03
##
    [16] 3.337838e-03 8.654453e-03 4.051869e-02 5.017592e-01 2.676743e-02
   [21] 1.304179e-01 2.499865e-01 6.243841e-02 4.119358e-01 1.499279e-02
##
##
    [26] 5.443831e-02 4.493478e-02 1.896964e-01 1.214582e-01 1.413273e-01
##
   [31] 5.760040e-01 6.869600e-02 3.171074e-03 8.103523e-02 1.281316e-03
   [36] 7.465622e-02 2.809172e-02 8.225027e-02 9.319087e-01 3.932204e-03
##
    [41] 8.830920e-04 1.284256e-01 3.941924e-01 5.206744e-02 7.305352e-02
##
    [46] 6.208799e-02 3.365046e-02 1.882240e-02 2.276767e-01 5.833807e-01
##
##
   [51] 2.335841e-01 8.015547e-03 7.011201e-04 2.027645e-01 4.813626e-03
    [56] 9.746532e-03 3.957417e-01 5.995076e-03 1.968747e-01 1.426307e-02
##
    [61] 1.243559e-02 1.214715e-03 9.943176e-02 7.864401e-04 1.301142e-02
    [66] 7.650938e-05 2.759159e-05 3.299408e-02 2.620086e-01 1.162073e-01
##
   [71] 2.532421e-02 5.117365e-02 6.821415e-04 4.509707e-03 1.137253e-01
##
##
   [76] 3.300675e-03 4.560727e-01 8.740656e-03 4.565355e-04 3.652628e-01
##
    [81] 3.185187e-01 1.393712e-02 2.801115e-01 3.695600e-03 1.715129e-02
##
    [86] 2.625241e-02 1.650533e-02 1.395559e-02 4.065866e-03 1.001851e-06
   [91] 2.331682e-02 1.814688e-01 1.939516e-02 6.705850e-02 8.826917e-03
##
   [96] 1.774541e-01 8.273213e-03 6.387857e-02 3.115453e-02 2.798558e-02
```

9.2.6.5 Exercise

Can we make this work?

```
x <- list(
  list(1, c(3, 9)),
  list(c(3, 6), 7, c(4, 7, 6))
)

triple <- function(x) x * 3
# map(x, map, .f = triple)</pre>
```

Specifying .f there replaces the second argument of the top level map function, so you'd be doing map(.x = x, .f = triple, map) which is not what you mean. What you want here is:

```
map_depth(x, 2, triple)
```

```
## [[1]]
## [[1]][[1]]
## [1] 3
##
## [[1]][[2]]
## [1] 9 27
##
##
## [[2]][[1]]
## [[2]][[1]]
## [1] 9 18
##
##
## [[2]][[2]]
```

```
## [1] 21
##
## [[2]][[3]]
## [1] 12 21 18
```

[[2]]

[[3]]

two is now 12

9.4.1 Same type of output as input: modify()

When using modify we now get the warning: Warning message: modify() is deprecated as of rlang 0.4.0. Vector tools are now out of scope for rlang to make it a more focused package.

How should we rewrite this example?

```
data.frame(
  x = 1:3,
  y = 6:4
) %>%
  modify( ~ .x * 2)
...
```

rlang::modify!= purrr::modify()

9.4.5 Any number of inputs: pmap() and friends

I want to use pmap to map over a vector, but the metadata for the functions other arguments is elsewhere - how would I combine these using pmap? The book says to name the metadata columns the same name as your function which I did but this still doesn't work?

```
my_data <- 1:3
metadata <- tribble(</pre>
  ~id, ~mult, ~adder,
 "one", 2,
           3,
 "two",
                    6,
  "three", 4,
the_function <- function(vec, id, mult, adder) {</pre>
 glue::glue("{id} is now {vec * mult + adder}")
# my_data doesn't change but we want to map over the metadata
\# x = string
# y = multiplier
\# z = adder
# pmap(list(metadata), ~the_function(vec = my_data))
pmap(c(list(my_data), metadata), the_function)
## [[1]]
## one is now 7
```

three is now 19

9.4.6.2 Exercise

I see how we can use iwalk and walk2 for writing to multiple files, but the question asks about disadvantages to this - what would they be?

```
cyls <- split(mtcars, mtcars$cyl)
paths <- file.path(temp, paste0("cyl-", names(cyls), ".csv"))
walk2(cyls, paths, write.csv)

mtcars %>%
   split(mtcars$cyl) %>%
   set_names(~ file.path(temp, paste0("cyl-", .x, ".csv"))) %>%
   iwalk(~ write.csv(.x, .y))
```

A readability problem mostly, as well as implicitly using names(cyls) as opposed to explicitly like it's used in paths

9.5.4 Multiple inputs

Can we think of a simple example for reduce()?

A Fibonacci function!

[1] 0 1 1 2 3 5 8 13 21 34 55

9.6.3.2 Exercise

I understand that if (length(x) == 1L) return(x[[1L]]) covers the case like $simple_reduce(1, +)$ but what's the deal with if (length(x) == 0L) return(default) and what exactly is default?

```
simple_reduce <- function(x, f, default) {
    # when would you use reduce on something length 0?
    if (length(x) == 0L) return(default)
    if (length(x) == 1L) return(x[[1L]])

    out <- x[[1]]
    for (i in seq(2, length(x))) {
       out <- f(out, x[[i]])
    }
    out
}</pre>
```

Default is the user supplied second number to add by, and we can use integer(0) to work with a length(0) vector.

```
simple_reduce(integer(0), `+`, default = 0L)
```

[1] 0

You would want to make sure your reduce is covered in the case of length zero so that your program doesn't crash when a user passes in an empty vector.

9.7.1 Matrices and arrays

How would you tidyverse the rowSums function?

```
## [1] 0 0 2
pmap_dbl(x, ~sum(c(...) > 2))
```

```
## [1] 0 0 2
```

I'm not quite sure what idempotent means but Hadley warns that a2d and a1 aren't the same, but isn't that just because we're using 1 which is row wise, and not 2? What is the warning he is heeding us against here?

```
a2d <- matrix(1:20, nrow = 5)
a1 <- apply(a2d, 1, identity)
identical(a2d, a1)</pre>
```

```
## [1] FALSE
```

I think he is just trying to drive home the idea that you might expect no change by apply-ing an operator like identity which leaves its input unchanged, however what it instead returns is the transpose. So perhaps it would have been better to more specifically say that apply(..., MARGIN = 1) isn't idempotent.

9.7.3.2 Exercise

What's an example of using eapply (iterates over the (named) elements of an environment)?

```
baseenv() %>% eapply(is.primitive) %>% unlist %>% which %>% names
```

Can we come up with an example for rapply allowing us to apply a function to only a specified class? Does something like this exist within purrr?

Answer from this community post:

9.0.0.1 rapply

```
"character",
   how = "replace"
  )
}
x %>% as_integer_rapply() %>% str()
## List of 3
## $ :List of 2
    ..$ a: int 1
## ..$ b: num 2
## $ c: int 3
## $ d: num 4
9.0.0.2 map
as_integer_recursive_map <- function(x) {</pre>
 x %>%
    map_if(is.character, as.integer) %>%
    map_if(is.list, as_integer_recursive_map)
}
x %>% as_integer_recursive_map() %>% str()
## List of 3
## $ :List of 2
    ..$ a: int 1
##
##
   ..$ b: num 2
## $ c: int 3
## $ d: num 4
9.0.0.3 Comparing the two:
microbenchmark::microbenchmark(
  as_integer_rapply(x),
  as_integer_recursive_map(x),
  times = 100
)
## Unit: microseconds
##
                           expr
                                     min
                                                lq
                                                        mean
                                                                median
                                                                               uq
##
           as_integer_rapply(x)
                                   9.772
                                           13.2935
                                                     38.3984
                                                               20.2105
## as_integer_recursive_map(x) 1015.665 1088.8885 1221.8427 1152.3090 1266.9110
        max neval
##
## 1831.531
              100
## 2114.119
              100
as_integer_safe <- function(y) {</pre>
  if (is.character(y) && all(grepl("^{\t}, y))) {
    as.integer(y)
  } else {
    у
  }
}
```

```
x[[1]][[3]] <- "dog"
rapply(x, f = as_integer_safe, how = "replace")
## [[1]]
## [[1]]$a
## [1] 1
##
## [[1]]$b
## [1] 2
##
## [[1]][[3]]
## [1] "dog"
##
##
## $c
## [1] 3
## $d
## [1] 4
```

Chapter 10

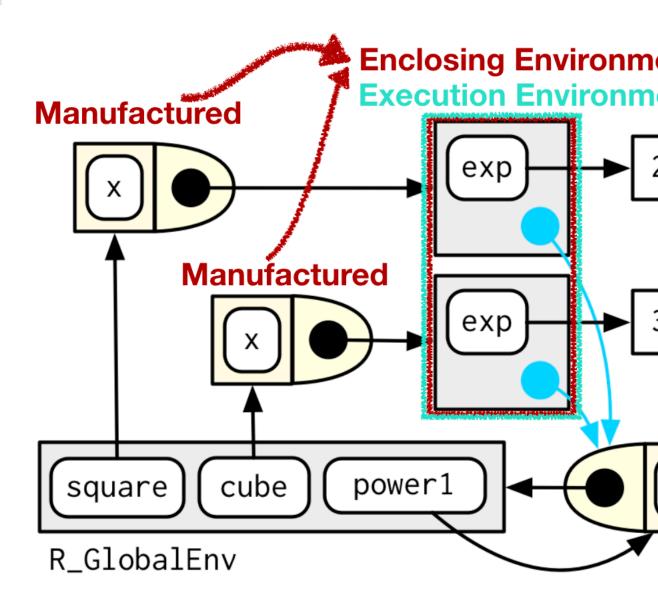
Function factories

10.2 Factory fundamentals

Let's really solidify this statement in our own words using the power1 and square/cube example:

The enclosing environment of the manufactured function is an execution environment of the function factory.

The enclosing environment of the manufactured function is an executatory.



- enclosing environment (this is the function environment like square and cube) [in red]
- execution environment (this is the same environment as above but from the "perspective" of power1) [in green]

square and cube have their own enclosing environments where x is different for each function - but their parent, the environment of power1, is the same.

Because power1 is a function, it has its own execution environment which is usually ephemeral but in this

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case it is the environment of both square and cube – and that is how the manufactured functions have access to exp.

10.2.3 Forcing evaluation

I think this force section is trying to convey how we should use force on every argument in the factory's execution environment (here just exp) because you want to avoid lazy eval in the manufactured function's enclosing environment (like how we want to force(exp = 2) for square) but how does it not just become force(exp = 3) when we set x < -3?

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#>

#> #> \$value
NULL

Without force, exp doesn't get evaluated when power is called. square is a promise to create a function to raise its input to the exp'th power. Once square is called, the promise evaluates, and square becomes a function to raise its input to whatever exp was that first time square was called. That's what makes it dangerous; what square does depends on the order in which things are called. Dangers and confusions like this are part of why function factories don't get used that much, despite the fact that they're really quite cool.

```
suppressMessages(requireNamespace("pryr"))
# make sure this isn't already defined
rm(x)
power2 <- function(exp) {</pre>
  cat(paste0("outer execution environment: ", capture.output(environment()), "\n"))
  cat(paste0("outer enclosure: ", capture.output(parent.env(environment())), "\n"))
  cat(paste0("exp promise info: \n", paste0("\t", capture.output(pryr::promise_info(exp)), collapse = "
  out <- function(x_) {</pre>
    cat(paste0("inner execution environment: ", capture.output(environment()), "\n"))
    cat(paste0("inner enclosure: ", capture.output(parent.env(environment())), "\n"))
    cat(paste0("exp promise info: \n", paste0("\t", capture.output(pryr::promise_info(exp)), collapse =
  }
}
square <- power2(x)
#> outer execution environment: <environment: 0x0000000015812568>
#> outer enclosure: <environment: R_GlobalEnv>
#> exp promise info:
#>
  $code
#>
   x
#>
#>
    $env
#>
    <environment: R_GlobalEnv>
#>
    $evaled
#>
   [1] FALSE
#>
#>
```

#> \$env
#> NULL

```
square(4)
#> inner execution environment: <environment: 0x0000000012aee2a8>
#> inner enclosure: <environment: 0x000000015812568>
#> exp promise info:
#> $code
#> x
#>
#>
  $env
#> <environment: R_GlobalEnv>
#>
#> $evaled
  [1] FALSE
#>
#>
#> $value
#> NULL
#>
#> Error in square(4): object 'x' not found
The exp promise is looking for x in the enclosure of power2 (global) and can't find it.
x <- 2
square(4)
#> inner execution environment: <environment: 0x0000000017c915d0>
#> inner enclosure: <environment: 0x000000015812568>
#> exp promise info:
#> $code
#>
#>
#> $env
#>
   <environment: R_GlobalEnv>
#>
#> $evaled
#> [1] FALSE
#>
#> $value
#> NULL
#> Warning in square(4): restarting interrupted promise evaluation
#> [1] 16
Notice, that the exp promise still wasn't evaled until the last line, x_ ^ exp. The warning is caused by our
previous failed attempt to evaluate the promise.
x <- 3
square(5)
#> inner execution environment: <environment: 0x0000000017d9b270>
#> inner enclosure: <environment: 0x000000015812568>
#> exp promise info:
#> $code
#>
#>
```

```
#>
#> $evaled
#> [1] TRUE
#>
#> $value
#> [1] 2
#>
#> [1] 25
# exp was evaled already, so no further warning is produced, and our change to x has no effect.
rm(x)
square <- power2(x)</pre>
#> outer execution environment: <environment: 0x0000000017e96e78>
#> outer enclosure: <environment: R_GlobalEnv>
#> exp promise info:
#> $code
#> x
#>
#> $env
#> <environment: R_GlobalEnv>
#>
#> $evaled
#> [1] FALSE
#>
#> $value
#> NULL
#>
environment(square)$x <- 2</pre>
square(6)
#> inner execution environment: <environment: 0x000000017f8a530>
#> inner enclosure: <environment: 0x000000017e96e78>
#> exp promise info:
#> $code
#> x
#>
#> $env
#> <environment: R_GlobalEnv>
#>
#> $evaled
#> [1] FALSE
#>
#> $value
#> NULL
#>
#> Error in square(6): object 'x' not found
Defining x in the function environment doesn't work, because the promise isn't looking for x there
power2 <- function(exp) function(x_) x_ ^ exp</pre>
# wasn't sure this would work without braces...neat
square <- power2(x)</pre>
environment(square)$exp <- 2</pre>
square(6)
```

#> \$code

```
[1] 36
square <- power2(x)
exp <- 2
square(6)</pre>
```

Error in square(6): object 'x' not found

This doesn't work because we need to define x within environment(power2)\$exp <- 2

10.1.1 USING FORCE

```
rm(exp)
rm(x, power2, square)
power2 <- function(exp) {</pre>
  cat(paste0("outer execution environment: ", capture.output(environment()), "\n"))
  cat(paste0("outer enclosure: ", capture.output(parent.env(environment())), "\n"))
  cat(paste0("exp promise info before: \n", paste0("\t", capture.output(pryr::promise_info(exp)), collage
  force(exp)
  cat(paste0("exp promise info after: \n", paste0("\t", capture.output(pryr::promise_info(exp)), collap
  out <- function(x_) {</pre>
    cat(paste0("inner execution environment: ", capture.output(environment()), "\n"))
    cat(paste0("inner enclosure: ", capture.output(parent.env(environment())), "\n"))
   cat(paste0("exp promise info: \n", paste0("\t", capture.output(pryr::promise_info(exp)), collapse =
   x_ ^ exp
 }
}
square <- power2(x)</pre>
#> outer execution environment: <environment: 0x00000000184e4e90>
#> outer enclosure: <environment: R_GlobalEnv>
#> exp promise info before:
#> $code
#> x
#>
#> $env
#> <environment: R_GlobalEnv>
#>
#> $evaled
#> [1] FALSE
#>
#> $value
#> NULL
#> Error in force(exp): object 'x' not found
# error because x isn't defined
x <- 2
square <- power2(x)</pre>
#> outer execution environment: <environment: 0x0000000018cafaf8>
#> outer enclosure: <environment: R_GlobalEnv>
#> exp promise info before:
```

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```
#>
    х
#>
#>
#>
    <environment: R_GlobalEnv>
#>
    $evaled
#>
#>
    [1] FALSE
#>
#>
    $value
    NULL
#>
#>
    exp promise info after:
    $code
#>
#>
    х
#>
#>
    $env
#>
    NULL
#>
#>
    $evaled
#>
    [1] TRUE
#>
#>
    $value
#>
    [1] 2
#>
square(4)
#> inner execution environment: <environment: 0x0000000018d55b48>
#> inner enclosure: <environment: 0x000000018cafaf8>
#> exp promise info:
#>
    $code
#>
    х
#>
#>
    $env
#>
    NULL
#>
#>
    $evaled
#>
    [1] TRUE
#>
#>
    $value
#>
    [1] 2
#> [1] 16
x <- 3
square(4)
```

10.2.5 garbage collection

Does {factory} handle the issue of temporary objects in the manufactured function not getting garbage collected?

We don't need to use rm because the execution environment of the factory remains ephemeral (like any other function call). Rather than returning a function wrapped in the execution environment, it returns a function in the caller environment, allowing normal garbage collection to clean up the temporary objects created during creation of the manufactured functions.

10.2.6.2 Exercise

Base R contains two function factories, approxfun() and ecdf(). Read their documentation and experiment to figure out what the functions do and what they return.

approxfun

approxfun creates a function (i.e. mathematical formula) that can be used to find the linear interpolation for any given point, based on the data that was passed into the factory-function at the time it was passed into the function

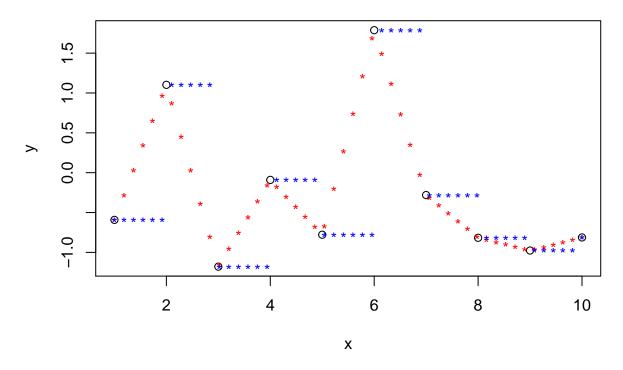
```
dat <- data.frame(x=1:10, y=(1:10)^2)
dat</pre>
```

```
##
           У
## 1
       1
           1
## 2
           4
## 3
       3
          9
      4 16
## 5
       5 25
## 6
       6
          36
## 7
       7
         49
## 8
          64
## 9
       9 81
## 10 10 100
fn <- approxfun(dat$x,dat$y)</pre>
fn(4.5)
```

```
## [1] 20.5
x <- 1:10
y <- rnorm(10)
plot(x, y, main = "approx(.) and approxfun(.)")
points(approx(x, y), col = 2, pch = "*")
points(approx(x, y, method = "constant"), col = 4, pch = "*")</pre>
```

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approx(.) and approxfun(.)



ECDF

```
x <- rnorm(3) # 3 random numbers
Fn <- ecdf(x) # create distribution function for each x
Fn(x) # apply the distribution function to each number</pre>
```

[1] 0.6666667 0.3333333 1.0000000

10.2.6.3 Exercise

How would we make this function using {factory}

```
pick <- function(i) {
  force(i)
  function(x) x[[i]]
}
library(factory)
pick <- build_factory(
    .factory_fun = function(x) x[[i]],
    i
)</pre>
```

10.3.1 Labelling

What is the "factory" happening here? That we're supplying a function as the label argument? Is summarize also a factory then?

It isn't that the ggplot2 functions are factories, it's that they're a great place to use factories. They're

actually functionals when you pass in a function, since they're using a function as an argument!

10.3.4 Exercises

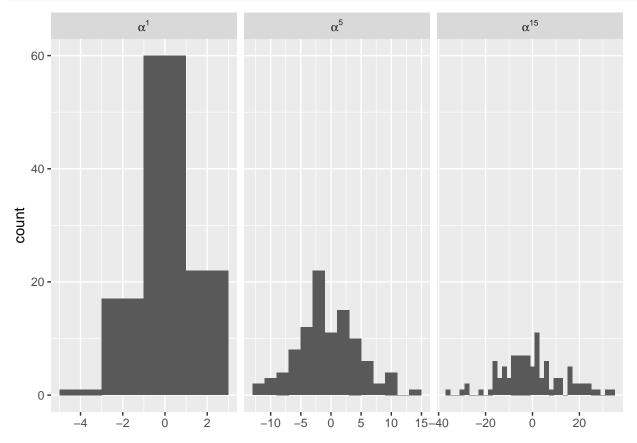
Can we use code to answer the question: compare and contrast $ggplot2::label_bquote()$ with $scales::number_format()$

Data:

```
sd <- c(1, 5, 15)
n <- 100
df <- data.frame(x = rnorm(3 * n, sd = sd), sd = rep(sd, n))</pre>
```

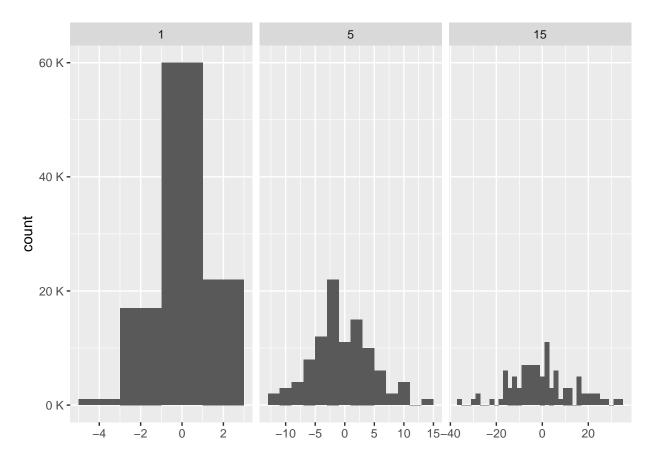
Using label_bquote using plotMath explanation.....

```
ggplot(df, aes(x)) +
  geom_histogram(binwidth = 2) +
  facet_wrap(~ sd, scales = "free_x", labeller = label_bquote(alpha ^ .(sd))) +
  labs(x = NULL)
```



Using number_format will force the computation of the parameters, formatting numbers explanation.....

```
ggplot(df, aes(x)) +
  geom_histogram(binwidth = 2) +
  facet_wrap(~ sd, scales = "free_x") +
  labs(x = NULL) +
  scale_y_continuous(
  labels = scales::number_format(scale = 1, suffix = " K")
)
```



10.2 10.5 factories and functionals

Can we {factory} this?

```
names <- list(
  square = 2,
  cube = 3,
  root = 1/2,
  cuberoot = 1/3,
  reciprocal = -1
)
funs <- purrr::map(names, power1)</pre>
```

```
power4 <- build_factory(
    fun = function(x) {
        x ^ exponent
    },
    exponent
)

power_names <- list(
    square = 2,
    cube = 3,
    root = 1/2,
    cuberoot = 1/3,
    reciprocal = -1
)

powers <- purrr::map(power_names, power4)</pre>
```

Chapter 11

Function operators

11.1 Introduction

A function operator is a function that takes one (or more) functions as input and returns a function as output

...How exactly does this definition differ from a function factory?

Term	Required Input	Optional Output	Output
Functionals Function Factory Function Operator	Function	Vector	Vector
	NA	Vector,Function	Function
	Function	Function	Function

Function operators are closely related to function factories; indeed they're just a function factory that takes a function as input

Can anyone confirm that, by definition, function operators can only take functions as inputs (and not any kind of vector)? Book quote: "A function operator is a function that takes one (or more) functions as input and returns a function as output." If this is true, then something like the following would only be a function factory, not both a function factory and a function operator:

```
sleepy <- function(f, n){
  force(f)
  function(...){
    cat(
      glue::glue('Sleeping for {n} second{ifelse(n != 1, "s", "")}.'),
      sep = '\n'
    )
    Sys.sleep(n)
    f(...)
}
sleepy_print <- sleepy(print, 1.5)
sleepy_print('hello world')</pre>
```

```
## Sleeping for 1.5 seconds.
## [1] "hello world"
```

Function factories are "any function that returns a function" and Hadley uses "simple" examples where you

pass in a numeric or character to get back a customized function (pass a number to a power function to get a function dedicated to cubes)

Function operators are factories where you pass in a function and modify its behaviour a little, without actually building up the logic for a function from scratch. The operator does not supply the actual meat of the logic. examples: safely, silently, rate-limity etc

11.2.1 purrr::safely

It seems like safely is a condition wrapper - are all conditions a kind of function operator?

safely/possibly/quietly are function operators that are condition-wrappers, but tryCatch/withCallingHandlers are not function operators (they catch the outputs and do X thing but don't actually return a function)

How would we apply safely to find which dataframe in a list of dataframes is causing an error?

```
library(tidyverse)
random_starwars <- function() {</pre>
  starwars %>%
    sample n(50) %>%
    select(name, height,mass)
}
broken_starwars <- function(){</pre>
  starwars %>%
  sample_n(50) %>%
  select(name, height, mass) %>%
  mutate(height = as.character(height))
}
list_data <- list(a = random_starwars(),</pre>
                  b = random starwars(),
                  c = broken starwars(),
                  d = random starwars())
map_dfr_safely <- function(.x, .f, ... ) {</pre>
  results <- map(.x, safely(.f), ...)
  list(result = bind rows(!!!map(results, "result")), error = map(results, "error"))
}
map_dfr_safely(list_data, semi_join, filter_starwars, by = c("name", "height"))
## Warning in tbl_vars_dispatch(x): restarting interrupted promise evaluation
## Warning in tbl_vars_dispatch(x): restarting interrupted promise evaluation
## Warning in tbl_vars_dispatch(x): restarting interrupted promise evaluation
## $result
## # A tibble: 0 x 0
##
## $error
## $error$a
## <simpleError in tbl_vars_dispatch(x): object 'filter_starwars' not found>
##
```

```
## $error$b
## <simpleError in tbl_vars_dispatch(x): object 'filter_starwars' not found>
##
## $error$c
## <simpleError in tbl_vars_dispatch(x): object 'filter_starwars' not found>
##
## $error$d
## <simpleError in tbl_vars_dispatch(x): object 'filter_starwars' not found>
```

11.2.2 memoise

How do you check how much memory is allocated to memoise's caching?

```
random_starwars <- function(...){
    starwars %>%
        sample_n(3) %>%
        select(name)}
memoised_starwars <- memoise(random_starwars)
a <- random_starwars()
b <- memoised_starwars(1)
c <- memoised_starwars(2)
d <- memoised_starwars(3)

lobstr::obj_size(get("_cache",environment(memoised_starwars)))</pre>
```

16,248 B

Could we create a memoise wrapper that clears the cache when that's a certain number?

```
random_starwars <- function(...){</pre>
  dplyr::starwars %>%
    dplyr::sample_n(3) %>%
    dplyr::select(name)
cache_memory_size <- function(f) {</pre>
  cache <- get("_cache", environment(f))</pre>
  lobstr::obj_size(cache)
capped_memoise <- function(..., .cache_size) {</pre>
  force(.cache_size)
  .self_ref <- memoise::memoise(...)</pre>
  f <- function(...) {</pre>
    env <- parent.env(environment())</pre>
    cache_size <- cache_memory_size(env[["_self_ref"]])</pre>
    if (cache_size > env[["_cache_size"]]) {
      message(paste0("Clearing cache [0",cache_size,"]\n"))
      memoise::forget(env[["_self_ref"]])
    env[["_self_ref"]](...)
  assign("_cache_size", .cache_size, environment(.self_ref))
  assign("_self_ref", .self_ref, environment(.self_ref))
  environment(f) <- environment(.self_ref)</pre>
  f
```

17,128 B

```
capped_memoised_starwars <- capped_memoise(random_starwars, .cache_size = 30000)</pre>
for (i in 1:100) {
  capped_memoised_starwars(i)
  print(cache_memory_size(capped_memoised_starwars))
## 14,448 B
## 15,328 B
## 16,232 B
## 17,152 B
## 18,056 B
## 18,960 B
## 19,864 B
## 20,640 B
## 21,480 B
## 22,376 B
## 23,272 B
## 24,112 B
## 24,944 B
## 25,656 B
## 26,488 B
## 27,376 B
## 28,216 B
## 28,928 B
## 29,768 B
## 30,608 B
## Clearing cache [@30608]
## 14,440 B
## 15,336 B
## 16,240 B
## 17,136 B
## 18,032 B
## 18,928 B
## 19,760 B
## 20,696 B
## 21,592 B
## 22,432 B
## 23,208 B
## 24,064 B
## 24,912 B
## 25,744 B
## 26,520 B
## 27,288 B
## 28,128 B
## 28,984 B
## 29,840 B
## 30,608 B
## Clearing cache [@30608]
## 14,432 B
## 15,336 B
## 16,232 B
```

```
## 18,064 B
## 18,968 B
## 19,864 B
## 20,704 B
## 21,472 B
## 22,376 B
## 23,208 B
## 23,984 B
## 24,760 B
## 25,536 B
## 26,328 B
## 27,216 B
## 28,048 B
## 28,816 B
## 29,648 B
## 30,536 B
## Clearing cache [030536]
## 14,440 B
## 15,336 B
## 16,168 B
## 17,000 B
## 17,896 B
## 18,736 B
## 19,568 B
## 20,392 B
## 21,304 B
## 22,208 B
## 22,920 B
## 23,808 B
## 24,728 B
## 25,496 B
## 26,336 B
## 27,048 B
## 27,904 B
## 28,680 B
## 29,456 B
## 30,288 B
## Clearing cache [@30288]
## 14,448 B
## 15,336 B
## 16,224 B
## 17,072 B
## 17,968 B
## 18,744 B
## 19,656 B
## 20,432 B
## 21,256 B
## 22,160 B
## 22,992 B
## 23,832 B
## 24,720 B
```

25,624 B

```
## 26,464 B
## 27,288 B
## 28,128 B
## 29,048 B
## 29,952 B
## 30,728 B
Can we come up with simple example for these function operators?
x <- list(
  c(0.512, 0.165, 0.717),
  c(0.064, 0.781, 0.427),
  c(0.890, 0.785, 0.495),
  "oops"
f <- function(...) {sum(...)}</pre>
possibly
map_dbl(x, possibly(f, NA))
## [1] 1.394 1.272 2.170
                              NA
quietly
Quetly doesn't work on errors, only warnings:
map(x, quietly(function(x) tryCatch(f(x), error = function(e) warning(e)))
```

$auto_browser$

```
map_dbl(x, auto_browse(f))
```

11.2.3.1 Exercises

Vectorize() provides a convenient and concise notation to iterate over multiple arguments, but has some major drawbacks that mean you generally shouldn't use it.

What are these drawbacks the solution manual is referring to?

11.0.0.1 Dean's Way

```
library(tidyverse)
patient_name <- function(path) {
    path_list <- str_split(path, "/") %>% unlist()
    paste(path_list[length(path_list) - 1], path_list[length(path_list)], sep = "_")
}

# Vectorize it with Vectorize
patient_name_v <- Vectorize(patient_name)
patient_name_v(path = c("some/path/abc/001.txt", "another/directory/xyz/002.txt"))

## some/path/abc/001.txt another/directory/xyz/002.txt
## "abc_001.txt" "xyz_002.txt"</pre>
```

11.0.0.2 Jim's Way

```
patient_name_best <- function(path) {</pre>
  pasteO(basename(dirname(path)), "_", basename(path))
```

1. SIMPLIFY logical or character string; attempt to reduce the result to a vector, matrix or higher dimensional array; see the simplify argument of sapply. This means the type of the function output depends on the input.

```
patient_name_v("some/path/abc/001.txt")
## some/path/abc/001.txt
           "abc_001.txt"
patient_name_v(character())
## named list()
patient_name_v(NULL)
```

- ## list()
 - $2. \ \, \text{Vectorize does not generate functions with easily inspect-able code}$
 - 3. Vectorize functions use do.call(), which can have unexpected performance consequences
 - 4. Vectorize does not actually make your code execute faster

Chapter 12

Base Types

Introduction

encapsulation: the user doesn't need to worry about details of an object because they are encapsulated behind a standard interface.

What is meant here by encapsulation?

Random forests, glm, neural nets are implemented differently internally. But to predict the class for a new sample based on a fitted model, you just use the 'predict' generic. You don't need to know how prediction is performed, just that 'predict(my_fit)' makes it happen

This makes [RC] harder to reason about, but allows them to solve problems that are difficult to solve in the functional OOP style of S3 and S4.

Maybe this question is better answered when we finish this section but what are these scenarios in which RC is a better choice than S3 and S4?

12.3 Base types

What is a symbol? How is this different from a character?

typeof(quote(a))

[1] "symbol"

symbol is the difference between: "a" and a. One is a string and the other represents an object in the environment. A quoted object is not the same as a string

Chapter 13

S3

13.2 Basics

What's a real use case for unclass?

I wanted to pass a paletteer palette to leaflet::colorNumeric which ends up calling the S3 method leaflet:::toPaletteFunc but that only has methods defined for character function and matrix. The paletteer output is a character vector with class color and so colorNumeric(palette, domain) failed. Simple solution: colorNumeric(unclass(palette), domain)

13.3 Classes

Can we come up with a simple example for a constructor, validator, and helper?

IPA Constructor

Create the custom class

```
new_ipa <- function(x = double()) {
  structure(x, class = "ipa")
}</pre>
```

IPA Validator

Function to make sure you can use your custom class

```
validate_ipa <- function(x = double()) {
   stopifnot(is.numeric(x))
   # has to be at least 15 bitterness
   stopifnot(x >= 15)
   x
}
```

IPA Helper

Calls both the validator and the constructor

```
ipa <- function(x) {
  x <- as.numeric(x)
  # Both helper and valditor</pre>
```

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```
x <- new_ipa(x)
res <- validate_ipa(x)
res
}

13.0.0.1 Test it out!
Fails as expected:
ipa(5)

Error in validate_ipa(x) : x >= 15 is not TRUE

But this works!
ipa(50)

## [1] 50
## attr(,"class")
## [1] "ipa"
```

13.3.4.3 Exercises

Can we go over the factor source code, commenting it line for line?

```
# using the example:
x <- c("one", "two", "three", "one")</pre>
factor <- function (x = character(), levels, labels = levels, exclude = NA, ordered = is.ordered(x), nm
   # if character vector is null set to character()
  if (is.null(x))
      x <- character()
  \# nx - names of x
  # in our case NULL
    nx \leftarrow names(x)
    # we didn't supply levels so it's missing
    if (missing(levels)) {
      # set y to the unique names found in the vector
      # nmax by default is NA but this is an argument
      # for unique setting the max number of unique values
        y <- unique(x, nmax = nmax)
        # get the order of the unique values
        # defaults to alphabetical
        ind <- order(y)</pre>
        # put our vector in order using only unique values
        levels <- unique(as.character(y)[ind])</pre>
    # because the function arg is a promise
    # we need to eval now to use it? I think?
    force(ordered) # FALSE
    # if x isn't a character string make it one
    if (!is.character(x))
        x <- as.character(x)
    # match the number of levels with NAs
```

```
# then if they are NA it will resolve to TRUE
    # only return the TRUE levels
    # this is how we get rid of NAs!
    levels <- levels[is.na(match(levels, exclude))]</pre>
    # get the ordered levels
    f <- match(x, levels)</pre>
    # if the names aren't null give f those names
    if (!is.null(nx))
        names(f) \leftarrow nx
    # if missing levels make levels back into a character string?
    # why?
    if (missing(labels)) {
        levels(f) <- as.character(levels)</pre>
    # otherwise get the legenth of your labels
    # if its the same as the levels
    else {
        nlab <- length(labels)</pre>
        if (nlab == length(levels)) {
            # xlevls is labels as character
            # then we get the unique values of that
            # and set that to xlevs
            nlevs <- unique(xlevs <- as.character(labels))</pre>
            # this is null
            at <- attributes(f)
            # attr levels
            at$levels <- nlevs
            # match the levels with ordered levels
            # and put the levels back in numeric order?
            # I think?
            f <- match(xlevs, nlevs)[f]</pre>
            # now set the attributes to the numeric levels
            attributes(f) <- at
        }
        # if you only have one lavel
        else if (nlab == 1L)
            # add numerics to the end of the label?
            levels(f) <- pasteO(labels, seq_along(levels))</pre>
            # otherwise invalid length
        else stop(gettextf("invalid 'labels'; length %d should be 1 or %d",
            nlab, length(levels)), domain = NA)
    }
    # add ordered and factor to f's class list
    class(f) <- c(if (ordered) "ordered", "factor")</pre>
    f
}
factor(x)
```

13.3.4.4 Exercises

Now that we've gone through the source code, I don't see where the contrasts attibute comes into play? We can look at the C() function but where is this used within factor()?

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It is a little misleading to say that factors "have an ... attribute", where I really think he means that "some functions might check if that attribute has been assigned".

13.4.1 Generics

In the example:

```
my_new_generic <- function(x) {
   UseMethod("my_new_generic")
}</pre>
```

The book states: "If you wonder why we have to repeat my_new_generic twice, think back to Section 6.2.3" Can we go over why?

I think he's referring to the fact that you have to use a string representation inside the UseMethod call.

I'm guessing part of the answer to this is baked into this warning on the UseMethod function docs:

Warning: Prior to R 2.1.0 UseMethod accepted a call with no arguments and tried to deduce the generic from the context.

Basically, you're specifying the name of the function and then telling R to go find this generic function's methods. You can't self-reference the name of the generic after it's been assigned to itself?

Maybe the name doesn't exist!

```
test.default <- function() {
    "test"
}

# not sure why you would want to do this, but you can
(function() {
    UseMethod("test")
})()</pre>
```

[1] "test"

Also because you might have assigned your generic to a new name, how would it know to use the true "generic" name.

```
superuglygenericfunction <- function() {
   UseMethod("superuglygenericfunction")
}
superuglygenericfunction.default <- function() {
   "ugh"
}
nice <- function() {
   UseMethod("superuglygenericfunction")
}
nice()</pre>
```

```
## [1] "ugh"

#> [1] "ugh"
```

13.4.4.5 Exercises

Can we explain what's happening here in our own words I was a little confused with the solutions manual? g(x) uses 1 from the .default method but 10 from the generic? Why?

```
g <- function(x) {
    x <- 10
    y <- 10
    UseMethod("g")
}

g.default <- function(x) c(x = x, y = y)

x <- 1
    y <- 1
    g(x)</pre>
```

х у 1 10

UseMethod creates a new function call with arguments matched as they came in to the generic. Any local variables defined before the call to UseMethod are retained (unlike S).

When you call g(x) even though you appear to be ignoring it in your definition of g and assigning x <-10, when you call UseMethod it will basically match the x in the g.default formals to the x that was passed in. The y isn't matched to any argument and is retained.

Tyler Grant Smith:m: 13 minutes ago r g <- function(x) { x <- 10 y <- 10 UseMethod("g") } g.default <- function(x = NULL) c(x = x, y = y) x <- 1 y <- 1 # this is interesting behavior g() #> y #> 10 g(x) #> x y #> 1 10

13.5 Object styles

Unfortunately, describing the appropriate use of each of these object styles is beyond the scope of this book.

Can we use the vctrs package in a concrete example to address some of these fringe cases and reasons why length(x) is causing an issue?

An object style is a pattern for structuring your class. You can use any structure you want but these 4 patterns get used very often. Date and factor are both vector style for example: the core data for objects of both types are stored in a single vector. And there are some other context-appropriate things denoted by attributes in each case.

I don't think length is an issue but when you move away from the vanilla vector style object then the meaning changes. For example, under the hood, a data frame is really a list of vectors. Each of the vectors (columns) has the same length. So if you take the length of a data frame you get the length of the list i.e. the number of columns.

```
length(mtcars)
```

```
## [1] 11
```

The result of calling length on a data frame isn't necessarily intuitive: most people would probably say that the number of rows is a more natural value for the length of a data frame.

This is the same for dates: which are actually of length 3 when classes and 9 when unclassed.

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```
x <- as.POSIXlt(ISOdatetime(2020, 1, 1, 0, 0, 1:3))
length(x)

## [1] 3
length(unclass(x))

## [1] 11</pre>
```

13.6 Inheritance

Can we make our own example of a method with inheritance?

Creating an IBU calculator for beer, we can use the .ipa method for a beer of class .wipa:

```
library(sloop)
\# IBUs are calculated as onces x alpha acid x boil time / constant
# we're going to use a different constant for IPAs and Stouts
# generic
IBU_calculator <- function(x) {</pre>
  UseMethod("IBU_calculator")
}
# default
IBU_calculator.default <- function(x){</pre>
  stop("Can only calculate IBU for beers, not ",class(x))
# ipa
IBU_calculator.ipa <- function(x) {</pre>
  1 * x * 30 / 7.25
}
# stout
IBU_calculator.stout <- function(x) {</pre>
  1.5 * x * 30 / 6
}
#### some beers ####
ipa_1 <- structure(</pre>
  15, class = c("ipa","beer")
# define a wine to see default dispatch
wine_1 <- structure(3,class = c("pinot","wine"))</pre>
# wipa "inherits" class ipa so we can use that method!
wipa_1 <- structure(20, class = c("wipa", "ipa", "beer"))</pre>
s3_dispatch(IBU_calculator(ipa_1))
```

=> IBU_calculator.ipa

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```
IBU_calculator.beer
   * IBU_calculator.default
s3_dispatch(IBU_calculator(wipa_1))
      IBU_calculator.wipa
## => IBU_calculator.ipa
      IBU calculator.beer
## * IBU_calculator.default
s3_dispatch(IBU_calculator(wine_1))
##
      IBU_calculator.pinot
##
      {\tt IBU\_calculator.wine}
## => IBU_calculator.default
13.1
        13.4.3.3
Can we go over why this doesn't return a1?
generic2 <- function(x) UseMethod("generic2")</pre>
generic2.a1 <- function(x) "a1"</pre>
generic2.a2 <- function(x) "a2"</pre>
generic2.b <- function(x) {</pre>
  class(x) <- "a1"
  NextMethod()
generic2(structure(list(), class = c("b", "a2")))
[1] "a2"
generic2 <- function(x) UseMethod("generic2")</pre>
generic2.a1 <- function(x) { cat("generic2.a1\n"); cat(class(x)) }</pre>
generic2.a2 <- function(x) { cat("generic2.a2\n"); cat(class(x)) }</pre>
generic2.b <- function(x) {</pre>
  cat("generic2.b\n")
  cat("Regardless of what we do, NextMethod will search for the following",
      "classes:\n", paste0(.Class, collapse = ", "), "\n")
  class(x) <- "a1"
  NextMethod()
}
x <- structure(1:1e8, class = c("b", "a2"))
generic2(x)
## generic2.b
## Regardless of what we do, NextMethod will search for the following classes:
## b, a2
## generic2.a2
```

13.7.1 S3 and Base Types

a1

WHY isn't the dispatch of mean(x1) determined by its type???

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```
x1 <- 1:5
class(x1)
[1] "integer"
s3_dispatch(mean(x1))
   mean.integer
   mean.numeric
=> mean.default
methods (mean)
## [1] mean.Date
                         mean.default
                                          mean.difftime
                                                            mean.POSIXct
## [5] mean.POSIX1t
                         mean.quosure*
                                          mean.vctrs_vctr*
## see '?methods' for accessing help and source code
```

mean doesn't actually have a specific method defined in base R for int or numeric. the default method handles all base types, type checking along the way:

mean.default

```
## function (x, trim = 0, na.rm = FALSE, ...)
## {
##
       if (!is.numeric(x) && !is.complex(x) && !is.logical(x)) {
##
           warning("argument is not numeric or logical: returning NA")
           return(NA_real_)
##
##
       if (na.rm)
##
##
           x \leftarrow x[!is.na(x)]
##
       if (!is.numeric(trim) || length(trim) != 1L)
           stop("'trim' must be numeric of length one")
##
       n <- length(x)
##
       if (trim > 0 && n) {
##
           if (is.complex(x))
##
##
                stop("trimmed means are not defined for complex data")
##
           if (anyNA(x))
##
               return(NA_real_)
##
           if (trim >= 0.5)
##
               return(stats::median(x, na.rm = FALSE))
           lo \leftarrow floor(n * trim) + 1
##
##
           hi <- n + 1 - lo
##
           x <- sort.int(x, partial = unique(c(lo, hi)))[lo:hi]</pre>
##
##
       .Internal(mean(x))
## }
## <bytecode: 0x7ffd8b1ec9a8>
## <environment: namespace:base>
```

13.7.3 Internal Generics

Defining a single group generic for your class overrides the default behaviour for all of the members of the group.

Does this mean that if I define abs for my class it overrides the behaviour of sign, sqrt and all the others in the Math group?

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Essentially, these are all (loosely) math-related functions, so R understands how to do these for anything it can understand as a numeric (somehow). If you're adding a class that isn't handled in one of these numeric-ish ways, you'll need to tell R how to handle the mathematics. You can't make a group like "Math", "Ops", "Summary", or "Complex" and expect R to know what it's doing.

If you define abs for your class it won't know how to do the other things. Suppose you defined your own min.POSIXct then when you call min(Sys.time()) it will encounter your function, evaluate, and then that is the end. Otherwise, min would call the min Primitive which is implemented in C as found here and if you trace through that you will see that is where the GroupDispatch is called.

13.7.4 Double Dispatch

What is happening here:

```
Ops.foo <- function(...) print("Called twice for some reason?")
x <- structure(2, class = "foo")
x > 0
```

- [1] "Called twice for some reason?"
- [1] "Called twice for some reason?"

Double dispatch – where it's calling the same method on x and zero. The inequality needs to convert both to the same class.

The book mentions that we can use vctrs::vec_arith to avoid the aboce, but does that pertain to this case? I tried: vctrs::vec_arith(">", x, 0) But I think my understanding is still off...

The vec_arith generic will first dispatch on class(x) So, in this case it tries to find vec_arith.foo but we haven't defined that yet, so it instead finds vec_arith.default which signals the error.

```
x <- structure(2, class = "foo")
x > 0
[1] TRUE
```

```
vec_arith(">", x, 0)
```

Error: <foo> > <double> is not permitted
vec_arith will dispatch to this first, because it finds the "foo" generic

```
vec_arith.foo <- function(op, x, y, ...) {
    # now it will dispath on y, which in this case is a numeric.
    UseMethod("vec_arith.foo", y)
}
vec_arith(">", x, 0)
```

Error in UseMethod("vec_arith.foo", y): no applicable method for 'vec_arith.foo' applied to an object of

Uh oh another error, and this is super confusing because none of our objects are class double. (typeof(x) is "double"). The problem is that vec_arith.foo further dispatches on the second argument and it is looking for vec_arith.numeric or vec_arith.default, but we haven't defined that yet.

```
vec_arith.foo.numeric <- function(op, x, y, ...) {
  vctrs::vec_arith_base(op, x, y, ...)
}
vec_arith(">", x, 0)
```

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```
vec_arith(">", 0, x)
```

Error: <double> > <foo> is not permitted

now vec_arith will first dispatch on the 0 and scope to vec_arith.numeric which will then try to dispatch to vec_arith.numeric.foo which we havent defined yet.

```
vec_arith.numeric.foo <- function(op, x, y, ...) {
  vctrs::vec_arith_base(op, x, y, ...)
}
vec_arith(">", 0, x)
```

[1] FALSE

What is happening here?!

```
Math.bar <- function(x) {</pre>
  print("Math bar")
  NextMethod(.Generic)
}
Math.foo <- function(x) {</pre>
  print("Math foo")
  NextMethod(.Generic)
}
x <- 1
class(x) <- c("foo", "bar", class(x))</pre>
abs(x)
`abs.foo` <- function(x) {</pre>
  print("Generic foo")
  NextMethod()
}
abs(x)
`abs.bar` <- function(x) {</pre>
  print("Generic bar")
  NextMethod()
}
abs(x)
rm(abs.foo)
# infinite recursion continually calling abs.bar, only calling Math.foo once
rm(Math.foo)
# back to working
abs(x)
```