```
F1.
a)
         Xor = \a.\b.a (b False True) (b True False)
         Maj = \a.\b.\c.a (b True (c True False)) (b (c True False) False)
b)
         IsZero = \alpha.a (\x.True False x) True
         Invert = \a.a False True
         IsEven = \alpha.a (\xline x. Invert x) True
c)
     • Y = \{f.(\{x.x.x.x\}) (\{x.f.(x.x.x\})\}\}
      Next = h.\f.\x.f (h f x)
      Sum = \x.\y.x Next y
      AlmostFib = f.\x.(Eq x 0) 0 (
         (Eq x 1) 1 (
           (Sum (f (Sub x 1)) (f (Sub x 2)))
         )
       )
      Fib = Y AlmostFib
      AlmostDiv = f.\n.\c.(IsZero n) c (f (Sub n 3) (Next c))
      Div = Y AlmostDiv
      Div3 = \x.Div \times 0
d)
e)
         Pred = \n.\f.\x.n (\g.\h.h (g f)) (\u.x) (\u.u);
         Sub = \langle x. \rangle  Pred x
         Eq = \x.\y.(IsZero (Sub x y)) ((IsZero (Sub y x)) True False) False
         Leq = \x.\y. (IsZero (Sub x y) True False)
e)
f)
     • Less = \n.\m.(IsZero (Sub n m)) (IsZero (Sub m n) False True) False
         Div = \n.\mbox{\footnotement{M.}\acc.(Less n m)} (Pair acc n) (Div (Sub n m) m (Next acc))
2.
3.
4. a)
            Len = \list.\cnt.(IsNil list) cnt (Len (Tail list) (Next cnt))
            LenOfList = \list.Len list 0
  b)

    AlmostMap = \f.\first.\second.\acc.(IsNil first) acc ((IsNil second) acc

         (AlmostMap f (Tail first) (Tail second) (Cons (f (Head first) (Head second))
         acc)))
         Map2 = \f.\first.\second.Rev (AlmostMap f first second Nil)
  c)
          AlmostRev = \list.\acc.(IsNil list) acc (AlmostRev (Tail list) (Cons (Head
         list) acc))
            Rev = \x.AlmostRev \times Nil
          (\lambda x.xx)((\lambda y.y)(\lambda z.z))
     • ((\lambda y.y)(\lambda z.z)) \wedge ((\lambda y.y)(\lambda z.z))
               (\lambda x.x \ x)(\lambda z.z)
```

```
2. \overline{2} = \lambda f.\lambda x.f(fx)
        \overline{2}\ \overline{2} = (\lambda f.\lambda x.f(fx))(\lambda f'.\lambda x'.f'(f'x')) \underset{\beta}{\rightarrow} \lambda x.(\lambda f'.\lambda x'.f'(f'x'))((\lambda f'.\lambda x'.f'(f'x'))x) \underset{\beta}{\rightarrow} \lambda x.(\lambda x'.f'(f'x'))((\lambda f'.\lambda x'.f'(f'x'))x) \underset{\beta}
        \lambda x. (\lambda f'.\lambda x'.f'(f'x'))(\lambda y'.x(x\ y')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda y'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda x'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda x'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.(\lambda x'.x(x\ y'))((\lambda y'.x(x\ y'))x')) \underset{\beta}{\rightarrow} \lambda x. (\lambda x'.x(x\ y'))((\lambda y'.x(x\ y'))x')
        \lambda x.(\lambda x'.(\lambda x'.x(x\ x'))(x(x\ x'))) \xrightarrow{\beta} \tilde{\lambda x}.\lambda x'.(x(x\ (x(x\ x'))))
\overline{2}\,\overline{2}\,\overline{2}
\lambda > (\lambda f. \lambda x. f (f (f (f x)))) (\lambda f. \lambda x. f (f x))
\lambda > ((\lambda f. (\lambda x. (f (f (f (f x)))))) (\lambda f. (\lambda x. (f (f x)))))
\alpha > ((\lambda f. (\lambda x. (f (f (f (f x)))))) (\lambda X0. (\lambda X1. (X0 (X0 X1)))))
\beta > (\lambda x. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))))))
X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) x))))
\beta > (\lambda x. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))))))
(X1)))) (\lambda X1. (x (x X1))))))
\alpha > (\lambda x.~((\lambda X0.~(\lambda X1.~(X0~(X0~X1))))~((\lambda X0.~(\lambda X1.~(X0~(X0~X1))))~((\lambda X0.~(\lambda X1.~(X0~(X0~X1)))))
X1)))) (\lambda X2. (x (x X2)))))))
((\lambda X2. (x (x X2))) X1)))))
X2))) (x (x X1))))))
\alpha > (\lambda X3. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) (\lambda X1. ((\lambda X2. (X3 X1. (X0 X1))))))))
 (X3 X2))) (X3 (X3 X1)))))))
(X3 X1))))))))
\alpha > (\lambda X3. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) (\lambda X4. (X3 (X3 (X3 X3))))))
 (X3 X4))))))))
\beta > (\lambda X3. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) (\lambda X1. ((\lambda X4. (X3 (X3 (X3 (X3 X4))))) ((\lambda X4. (X3 (X3 X4))))))))
 (X3 (X3 (X3 X4)))) X1))))
\beta > (\lambda X3. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) (\lambda X1. ((\lambda X4. (X3 (X3 (X3 X4))))) (X3 (X3 (X3 X4)))))
(X3 (X3 X1))))))))
\alpha > (\lambda X5. ((\lambda X0. (\lambda X1. (X0 (X0 X1)))) (\lambda X1. ((\lambda X4. (X5 (X5 (X5 X4))))) (X5 (X5 (X5 X4))))))
 (X5 (X5 X1)))))))
X1)))))))))))
X6)))))))))))
(X5 (X5 (X5 (X5 (X5 X6)))))))) X1))))
(X5 (X5 (X5 (X5 X1))))))))))
\alpha > (\lambda X7. \ (\lambda X1. \ ((\lambda X6. \ (X7 \ X6))))))))))))))
 (X7 (X7 (X7 (X7 X1)))))))))
\beta > ( \lambda X7 . ( \lambda X1 . ( X7 )
X1)))))))))))))))))))
X1)))))))))))))))))
Здесь и далее m \equiv \overline{m} в зависимости от контекста и \stackrel{=}{=} \equiv =.
```

Индукцией покажем, что $nm = m^n$. Действительно,

$$nm = (\operatorname{Inc} (n-1))m = (\lambda f.\lambda x.(n-1)f(fx))m = \lambda x.(n-1)m(mx)$$

По предположению индукции

$$\lambda x.((n-1)m)(mx) = \lambda x.m^{n-1}(mx) = \overline{m^n}$$

Осталось показать, что для любых a и b верно.

$$\lambda f.a(bf) = \overline{ab}$$

Покажем это индукцией по a.

$$\lambda f.a(bf) = \lambda f.(\lambda x.(bf)^a x) = \lambda f.(\lambda x.(bf)^{a-1}((bf)x)) = \lambda f.(\lambda x.f^{b(a-1)}f^b x) = \overline{ab}$$

5.

$$S = \lambda x.\lambda y.\lambda z.xz(yz)$$

$$K = \lambda x. \lambda y. x$$

$$I = \lambda x.x$$

a)

 $\lambda x.x \ x = SII:$

$$\begin{array}{l} \lambda x.\lambda y.\lambda z.xz(yz)\ I\ I \underset{\beta}{\rightarrow} \lambda y.\lambda z.Iz(yz)\ I \underset{\beta}{\rightarrow} \lambda z.Iz(Iz) \underset{\beta}{\rightarrow} \lambda z.z(Iz) \underset{\beta}{\rightarrow} \lambda z.z\ z \\ \Omega = (\lambda x.x\ x)\ (\lambda x.x\ x) = (SII)\ (SII) \end{array}$$

b)

$$F = \lambda x. \lambda y. y \Leftrightarrow \lambda y. \lambda x. x = KI$$
:

$$\lambda x.\lambda y.x \ \lambda x.x \xrightarrow{\beta} \lambda y.\lambda x.x$$

$$\overline{1} = \lambda f. \lambda x. f \ x = SF$$
:

$$\lambda y.\lambda f.\lambda x.yx(fx) \ \lambda x'.\lambda y'.y' \underset{\beta}{\rightarrow} \lambda f.\lambda x.(\lambda x'.\lambda y'.y')x(fx) \underset{\beta}{\rightarrow} \lambda f.\lambda x.(\lambda y'.y')(fx) \underset{\beta}{\rightarrow} \lambda f.\lambda x.fx$$

c)

$$\lambda x.\lambda y.\lambda z.y = KK$$
:

$$\lambda x'.\lambda x.x' \ \lambda y.\lambda z.y \xrightarrow{\beta} \lambda x.\lambda y.\lambda z.y$$

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$$(\lambda x.xx)((\lambda y.y)(\lambda z.z)) \underset{\beta}{\rightarrow} (\lambda x.xx)(\lambda z.z)$$

$$(\lambda x.xx)((\lambda y.y)(\lambda z.z)) \underset{\beta}{\rightarrow} ((\lambda y.y)(\lambda z.z))((\lambda y.y)(\lambda z.z))$$

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a)

рассмотрим $\vdash A : \varphi$

доказательство типа А было получена по одному из правил:

Определение

Просто-типизированное лямбда-исчисление (по Карри). Типы: $\tau ::= \alpha | (\tau \to \tau)$. Язык: $\Gamma \vdash A : \varphi$

$$\frac{\Gamma, x : \varphi \vdash A : \psi}{\Gamma, x : \varphi \vdash X : \varphi} \; x \notin \Gamma \qquad \frac{\Gamma, x : \varphi \vdash A : \psi}{\Gamma \vdash \lambda x. A : \varphi \to \psi} \; x \notin \Gamma \qquad \frac{\Gamma \vdash A : \varphi \qquad \Gamma \vdash B : \varphi \to \psi}{\Gamma \vdash BA : \psi}$$

если по правилу 1, то это атомарная переменная, не имеет подвыражений если по правилу 2, то имеет вид $A=\lambda x.B$, и B имеет тип если по правилу 3, то A=BC, и B, и C имеют тип b)

пусть у Y или Ω есть тип

тогда у всех его подвыражений должен быть тип

в какой-то момент будет использовано правило 2:

$$\frac{\Gamma \vdash x: t_1 \to t_2 \quad \Gamma \vdash x: t_1}{\Gamma \vdash x\: x: t_2}$$

получается, у х бесконечный тип $t_1 = t_1 \to t_2 = (t_1 \to t_2) \to t_2...$

c)

newtype Mu a = Mu (Mu a -> a)
y f = (
$$h -> h$$
\$ Mu h) ($x -> f$. ($(Mu g) -> g) x $ x)$

newtype OmegaType = OmegaType { apply :: OmegaType -> OmegaType }

бесконечный тип

9a

$$\begin{split} \Gamma &= \{n: (\alpha \to \alpha) \to (\alpha \to \alpha), f: (\alpha \to \alpha), x: \alpha\} \\ &\frac{\frac{\overline{\Gamma \vdash x: \alpha} \quad \overline{\Gamma \vdash f: \alpha \to \alpha}}{\Gamma \vdash f: x: \alpha} \quad \frac{\overline{\Gamma \vdash n: (\alpha \to \alpha) \to (\alpha \to \alpha)} \quad \overline{\Gamma \vdash f: \alpha \to \alpha}}{\Gamma \vdash n f: (\alpha \to \alpha)} \\ &\frac{\Gamma &= \{f, n, x\} \vdash n f(f x): \alpha}{\{f, n\} \vdash \lambda x.n \ f(f x): (\alpha \to \alpha)} \\ &\frac{\overline{\{f, n\} \vdash \lambda x.n \ f(f x): (\alpha \to \alpha) \to (\alpha \to \alpha),}}{\{n\} \vdash \lambda f. \lambda x.n \ f(f x): (\alpha \to \alpha) \to (\alpha \to \alpha),} \end{split}$$