

编译原理模拟试题

Final exam demo

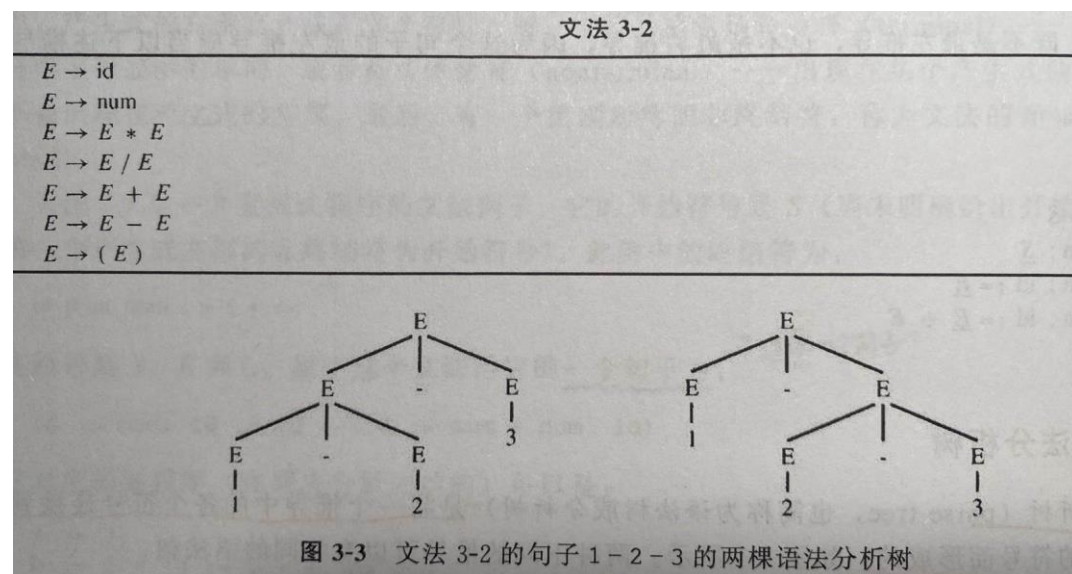
1、正则表达式和有限自动机

正则表达式可以转换为 NFA 和 DFA，也可以写出对应的 CFG。已知字母表{a,b,c}，任意构思一个正则式及其对应的 NFA 和 DFA，以及 CFG，并给出 NFA 和 DFA 的状态数。（考核标准中，NFA 和 DFA 必须具有起始状态和接受状态，DFA 出项必须包含字母表中的每个符号。CFG 必须是标准的 BNF 形式）

2、CFG 的改写

a) 将二义性文法改造成为无二义性文法。

Rewrite the grammar into an unambiguous grammar.



b) 对以下文法提取左因子，消除左递归。

Left factor and eliminate left recursive

参考以前作业

3、自顶向下的分析

a. 计算下面文法的 nullable、FIRST 和 FOLLOW 集合：

$S \rightarrow u B D z$
 $B \rightarrow B v$
 $B \rightarrow w$
 $D \rightarrow E F$
 $E \rightarrow y$
 $E \rightarrow$
 $F \rightarrow x$
 $F \rightarrow$

b. 构造 LL(1)分析表。

c. 给出证据说明该文法不是 LL(1)文法。

d. 尽可能少地修改该文法使它成为一个接收相同语言的 LL(1)文法。

a. Calculate nullable, FIRST, and FOLLOW for this grammar:

$$\begin{aligned} S &\rightarrow u B D z \\ B &\rightarrow B v \\ B &\rightarrow w \\ D &\rightarrow E F \\ E &\rightarrow y \\ E &\rightarrow \\ F &\rightarrow x \\ F &\rightarrow \end{aligned}$$

b. Construct the LL(1) parsing table.

c. Give evidence that this grammar is not LL(1).

d. Modify the grammar as little as possible to make an LL(1) grammar that accepts the same language.

4、Bottom-Up Parsers

3.11 构造下面这个文法的 LR(0) 状态，然后确定该文法是否为 SLR 文法。

$$\begin{array}{ll} 0 & S \rightarrow B \$ \\ 1 & B \rightarrow \text{id } P \\ 2 & B \rightarrow \text{id } (E] \\ 3 & P \rightarrow \\ 4 & P \rightarrow (E) \\ 5 & E \rightarrow B \\ 6 & E \rightarrow B , E \end{array}$$

Construct the LR(0) states for this grammar, and then determine whether it is an SLR grammar.

$$\begin{array}{ll} 0 & S \rightarrow B \$ \\ 1 & B \rightarrow \text{id } P \\ 2 & B \rightarrow \text{id } (E] \\ 3 & P \rightarrow \\ 4 & P \rightarrow (E) \\ 5 & E \rightarrow B \\ 6 & E \rightarrow B , E \end{array}$$

5、Syntax-Directed Translation

参考书上习题 6.6

Consider the following grammar for integer binary trees (in linearized form):

$$btree \rightarrow (\text{number } btree \text{ } btree) \text{ nil}$$

Write a **YACC specification** to check that a binary tree is ordered, that is, that the values of the numbers of the first subtree are \leq the value of the current number and the values of all the numbers of the second subtree are \geq the value of the current number. For example, $(2 (1 \text{ nil nil}) (3 \text{ nil nil}))$ is ordered, but $(1 (2 \text{ nil nil}) (3 \text{ nil nil}))$ is not. You may use the following attributes: **int num**, **bool order**. Use bison syntax: **\$i.num** refers to the **num** attribute of the *i*th symbol of the production and **\$\$num** refers to the num attribute of the production's result. Do not use global variables.