



Arrays: Mergesort

- Explore the Merge sort algorithm
- Understand the following aspects
 - Algorithm mechanism and pseudocode
 - Algorithm iterations on varying input
 - Algorithm time and space complexity



Merge sort

- Merge sort is a classic sort algorithm.
 - Optimal sort algorithm.
 - Used in system sort functions in major programming languages:
 - Java (sorting objects)
 - Python (stable sort)
- Algorithm scheme:
 - i. Divide given array into two halves.
 - ii. Sort each half recursively.
 - iii. Merge the two sorted halves.

Merge sort

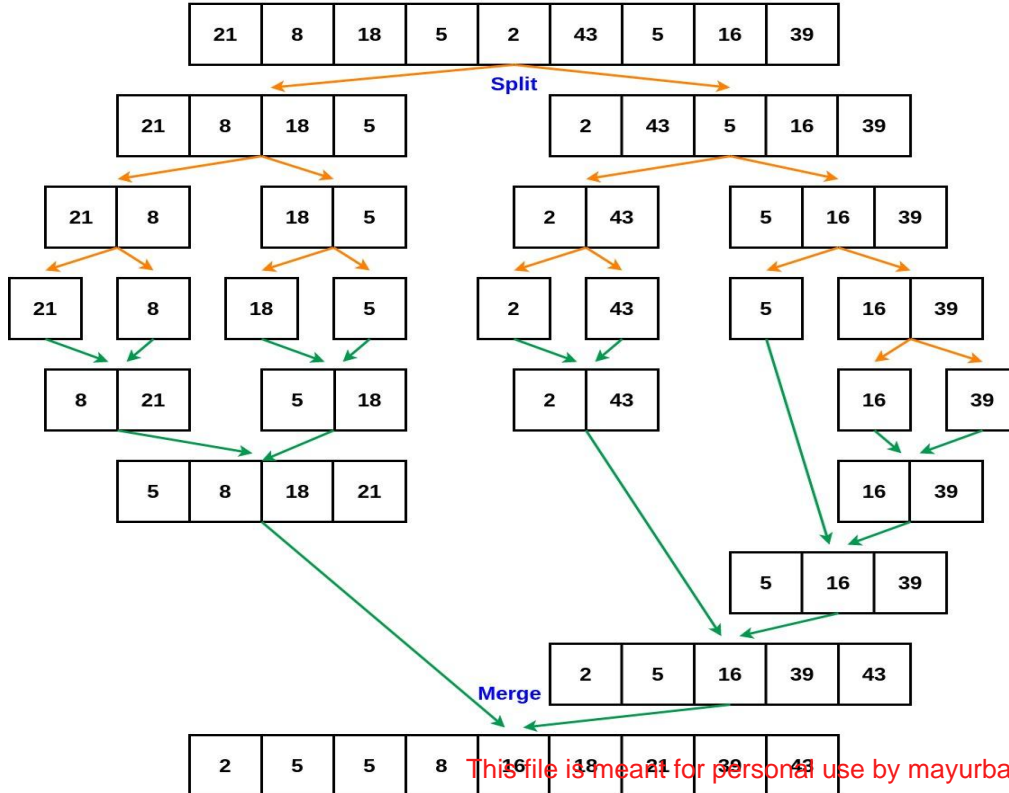
- Characteristics:
 - Merge sort is a comparison-based sort.
 - Regular Merge sort requires an auxiliary array for the merge operations.
 - In-place merge complicates algorithm, avoided.
- Comparison operation:
 - Defined as required for the data type:
 - Numbers
 - Strings
 - Objects: by attributes

Merge sort

- Merge sort uses the divide-and-conquer technique.
 - Divide the array into two halves.
 - Recursively sort the two halves.
- Iterative version of Merge sort also possible.

Merge sort

How Merge Sort Works



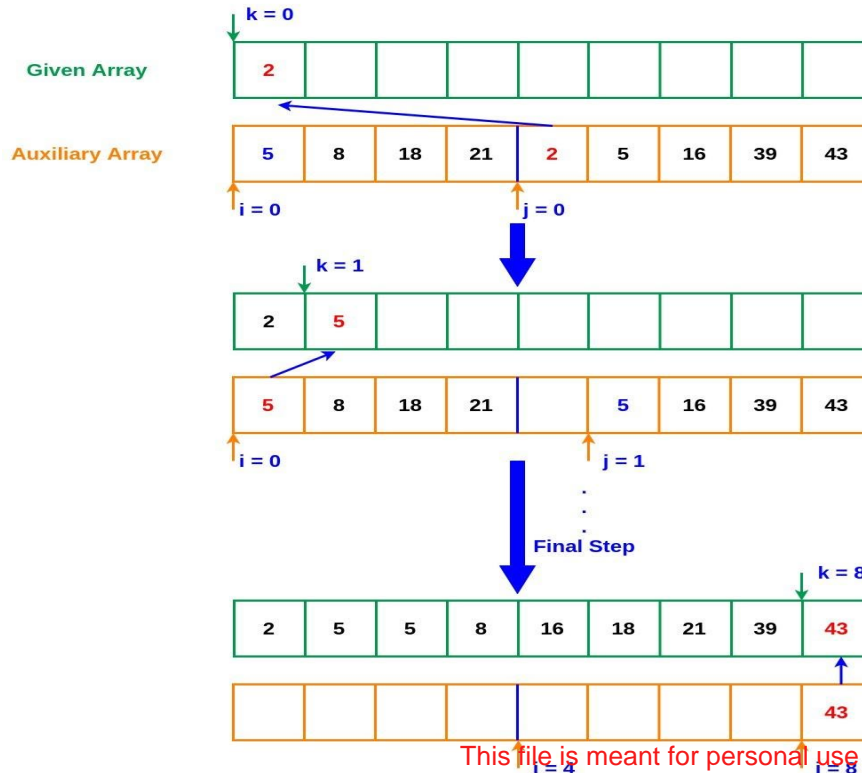
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Merge procedure

Merge Procedure



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Merge sort pseudocode

Code For Merge Sort

```
def merge(array, aux, lo, mid, hi):
    for k in range(lo, hi+1):

        i = lo
        j = mid+1

        for k in range(lo, hi+1):
            if i > mid:
                array[k] = aux[j]
                j+=1
            elif j > hi:
                array[k] = aux[i]
                i += 1
            elif aux[j] < aux[i]:
                array[k] = aux[j]
                j += 1
            else:
                array[k] = aux[i]
                i += 1
```

```
def merge_sort(array, aux, lo, hi):
    if hi <= lo:
        return

    mid = lo + (hi - lo) // 2

    merge_sort(array, aux, lo, mid)
    merge_sort(array, aux, mid+1, hi)

    merge(array, aux, lo, mid, hi)

def wrapper_sort(array):
    aux = []

    aux.extend(array)
    merge_sort(array, aux, 0, len(array)-1)
```

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Merge Sort Iterations On Array A (With Auxiliary Array AUX)

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Merge Sort Iterations On Pre-Sorted Array A (With Auxiliary Array AUX)

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Merge sort iterations

Merge Sort Iterations On Reverse-Sorted Array A (With Auxiliary Array AUX)

Array Indexes

	0	1	2	3	4	5	6	7	8
(VIOLET) Initial Array A	43	39	21	18	16	8	5	5	2
merge(A, AUX, 0, 0, 1)	39	43	21	18	16	8	5	5	2
merge(A, AUX, 2, 2, 3)	39	43	18	21	16	8	5	5	2
merge(A, AUX, 0, 1, 3)	18	21	39	43	16	8	5	5	2
merge(A, AUX, 4, 4, 5)	18	21	39	43	8	16	5	5	2

(GREEN)
Final
Sorted Array

	0	1	2	3	4	5	6	7	8
merge(A, AUX, 6, 6, 6)	18	21	39	43	8	16	5	5	2
merge(A, AUX, 7, 7, 8)	18	21	39	43	8	16	5	2	5
merge(A, AUX, 4, 6, 8)	18	21	39	43	2	5	5	8	16
merge(A, AUX, 0, 4, 8)	2	5	5	8	16	18	21	39	43

Merge sort: Complexity

- Analyzing time complexity:
 - Best case
 - Worst case
 - Average case
- Assume that the number of steps taken by the sort procedure (not merge), on an array of size N is $C(N)$.
- Mergesort has 2 steps:
 - Sort 2 sub-arrays recursively: $C(N/2) + C(N/2)$
 - Merge 2 sub-arrays of size $N/2$: N

Merge sort: Complexity

- Analyzing time complexity:
 - Simplifying assumption: $N = 2^k$, for some k
 - From the previous slide,

$$C(N) = 2C(N/2) + N$$

Dividing both sides by N ,

$$C(N) / N = C(N/2) / (N/2) + 1$$

$$C(N) / N = C(N/4) / (N/4) + 1 + 1 \quad (\text{Telescoping Property})$$

....

$$C(N) / N = C(N/N) / (N/N) + k \quad (\text{Remember, } N = 2^k)$$

$$C(N) / N = 0 + \log_2 N \quad (\text{Sorting 1 element!!})$$

$$C(N) = N \log_2 N$$

Merge sort: Complexity

- Analyzing time complexity:
 - The following can be demonstrated:
 - Worst case: $N \log_2 N$
 - Average case: $N \log_2 N$
 - Best case: $N \log_2 N$

Merge sort: Complexity

- Analyzing space complexity
 - Since the recursive sort and merge are executed throughout the sort:
 - Recursive sort: Constant space
 - Merge operation: Space proportional to N
 - Best case: space proportional to N
 - Worst case: space proportional to N
 - Average case: space proportional to N
 - The constant involved may change, from case to case.

We explored the Merge sort algorithm as follows:

- Algorithm mechanism and pseudocode
- Algorithm iterations on varying input
- Time and space complexity



Thank You