



COA UNIT-2 Notes

Computer Organization & Architecture (Dr. A.P.J. Abdul Kalam Technical University)

Adders:

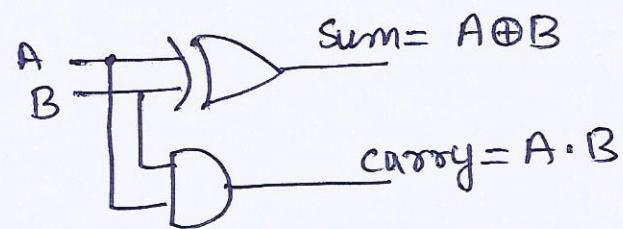
- Half Adders
- Full Adders
- n-bit Binary Parallel Adder
- Carry Look-ahead Adder (or Fast adder or speed adder).

Half Adder:

Half adder is a combinational circuit that adds two bits.

Truth Table of Half Adder:Circuit of Half adder

input		output	
A	B	Sum	carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

Full Adder:

Full adder is a combinational circuit that adds 3 bits.

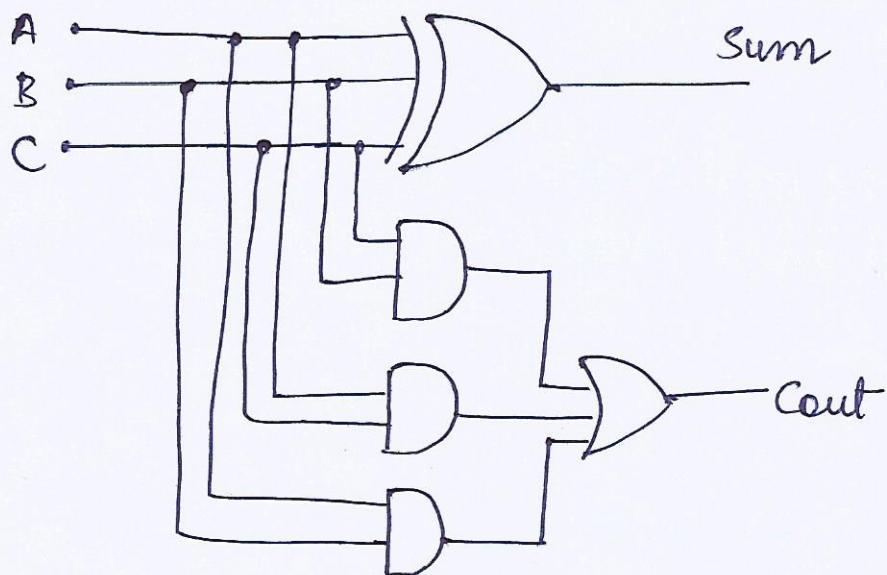
Truth Table:

A	B	C	sum	cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

$$\text{sum} = A \oplus B \oplus C$$

$$\text{cout} = AB + BC + AC$$

Logic diagram of Full adder:



$$C_{out} = \bar{A}\bar{B}C + A\bar{B}C + A\bar{B}\bar{C} + ABC$$

Properties :

$$= \bar{A}\bar{B}C + A\bar{B}C + A\bar{B}\bar{C} + ABC + ABC + ABC$$

$$\left\{ \because x+x=x \right\}$$

$$= BC[\bar{A}+A] + AC[\bar{B}+B] + AB[\bar{C}+C]$$

$$\left\{ x+\bar{x}=1 \right\}$$

$$= BC \cdot 1 + AC \cdot 1 + AB \cdot 1$$

$$\left\{ x \cdot 1 = x \right\}$$

$$C_{out} = BC + AC + AB$$

n-bit Binary Parallel Adder:

→ An n-bit binary parallel adder adds 2 n-bit numbers.

→ An n-bit binary adder uses n full adders.

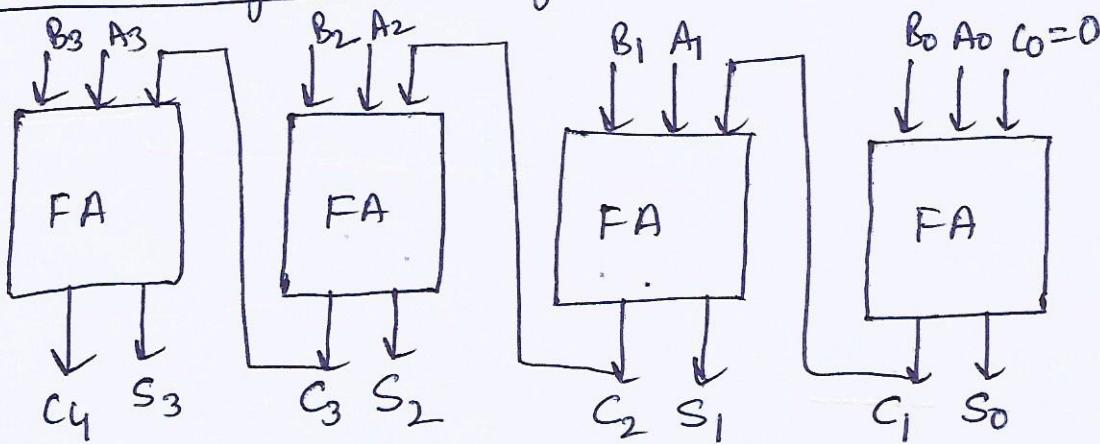
Example An 4-bit binary adder adds two 4-bit binary numbers

$$A = A_3 A_2 A_1 A_0$$

$$B = B_3 B_2 B_1 B_0$$

Addition

$$\begin{array}{r}
 & C_3 & C_2 & C_1 & C_0 = 0 & \Rightarrow \text{carry input} \\
 A_3 & A_2 & A_1 & A_0 & & \\
 + B_3 & B_2 & B_1 & B_0 & & \\
 \hline
 S_3 & S_2 & S_1 & S_0 & \Rightarrow \text{sum} \\
 C_4 & C_3 & C_2 & C_1 & \Rightarrow \text{carry output.}
 \end{array}$$

4-bit binary adder diagram.

A 4-bit binary parallel adder consists of 4 full adders in cascade, with output carry from one full adder, connected to the carry input of the next full adder.

Carry look ahead Adder:

Drawback of n-bit parallel adder -

- The carry output of one stage (or full adder) is connected to the carry input of the next higher stage (or full adder). This carry is called ripple carry.
- Therefore, the sum and carry can not be produced until the input carry occurs.
- This lead to a time delay in addition process. This delay is known as carry propagation delay.

Carry-Lookahead Adder:

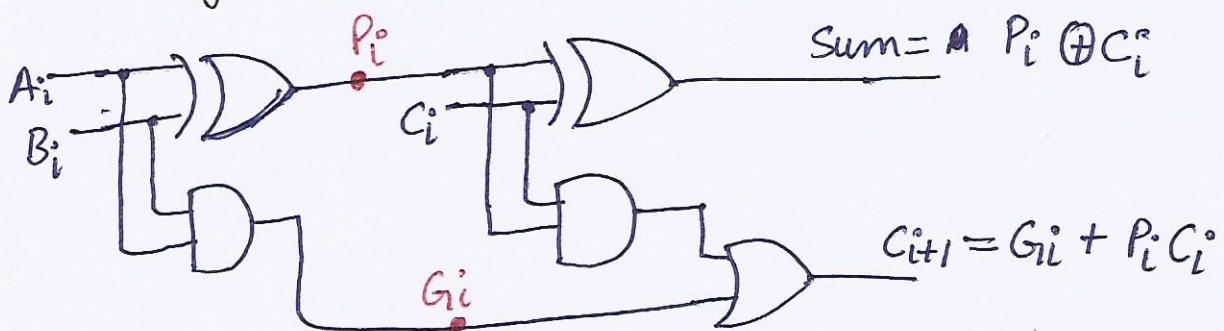
This method utilizes logic gates to look at the lower order bits of augend and addend to see if a higher order carry is to be generated or not

Carry look-ahead concept uses two functions

① Carry Propagate (P_i)

② Carry Generate (G_i)

Full adder using two half adders:



$$\text{Carry Propagation } P_i = A_i \oplus B_i$$

$$\text{Carry Generation } G_i = A_i \cdot B_i$$

Both RHS does not contain any C_i

The output $S_i = P_i \oplus C_i$

$$C_{i+1} = G_i + P_i C_i$$

we can create n-bit ~~look~~ carry look-ahead adder using these two equations.

Example

4 bit carry look-ahead adder:

For 4-bit, i is 0 to 3.

To remove the dependency of C_{i+1} on C_i , we need to express C_{i+1} into A_i, B_i and C_0 .

$$C_{i+1} = G_i + P_i C_i \quad \text{--- (1)}$$

Put

$$i=0 \quad C_{0+1} = G_{0+1} + P_0 C_0$$

$$C_1 = G_0 + P_0 C_0$$

$i=1$

$$C_{1+1} = G_1 + P_1 C_1$$

$$C_2 = G_1 + P_1 [G_0 + P_0 C_0]$$

$i=2$

$$C_{2+1} = G_2 + P_2 C_2$$

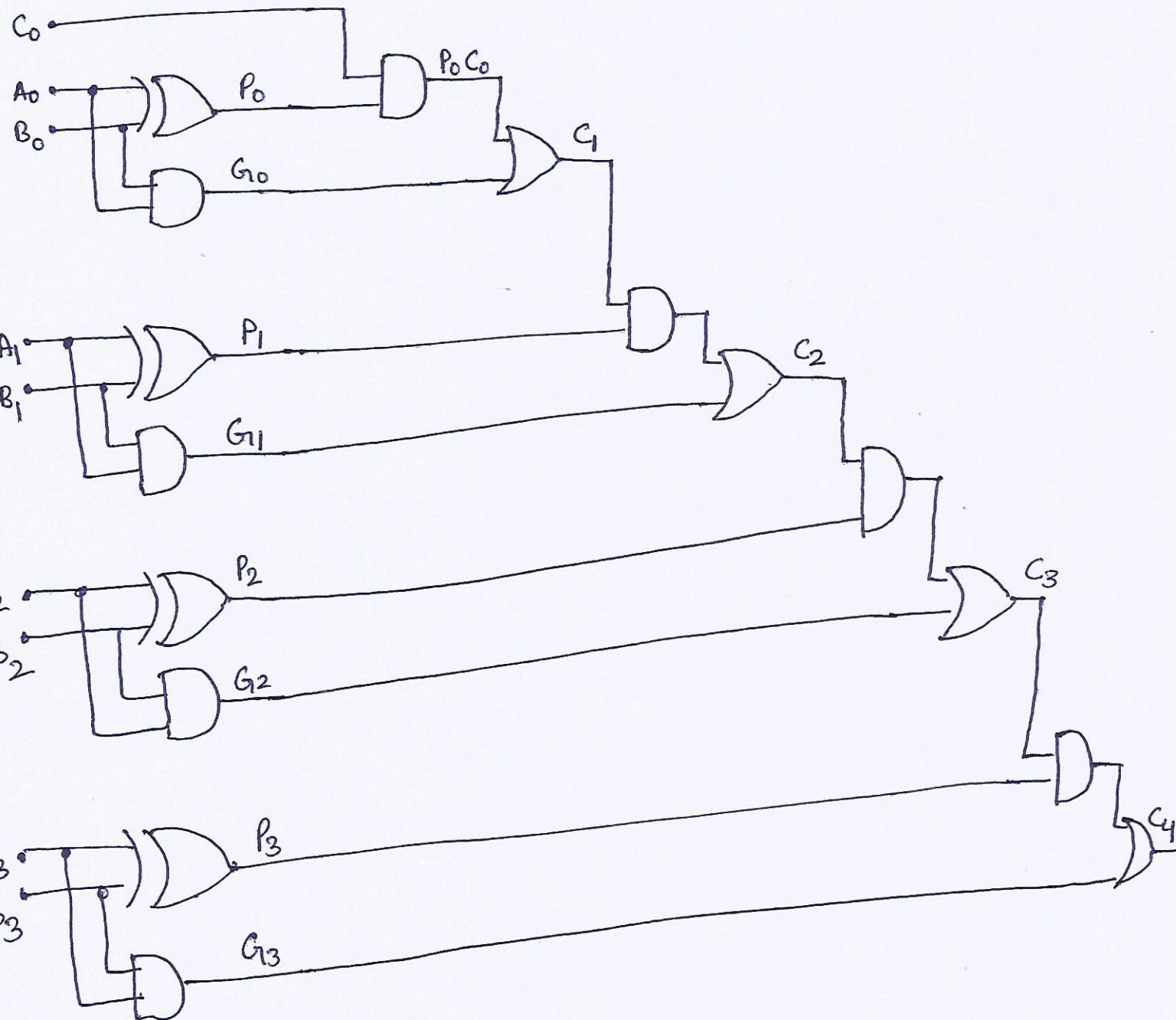
$$C_3 = G_2 + P_2 [G_1 + P_1 [G_0 + P_0 C_0]]$$

$i=3$

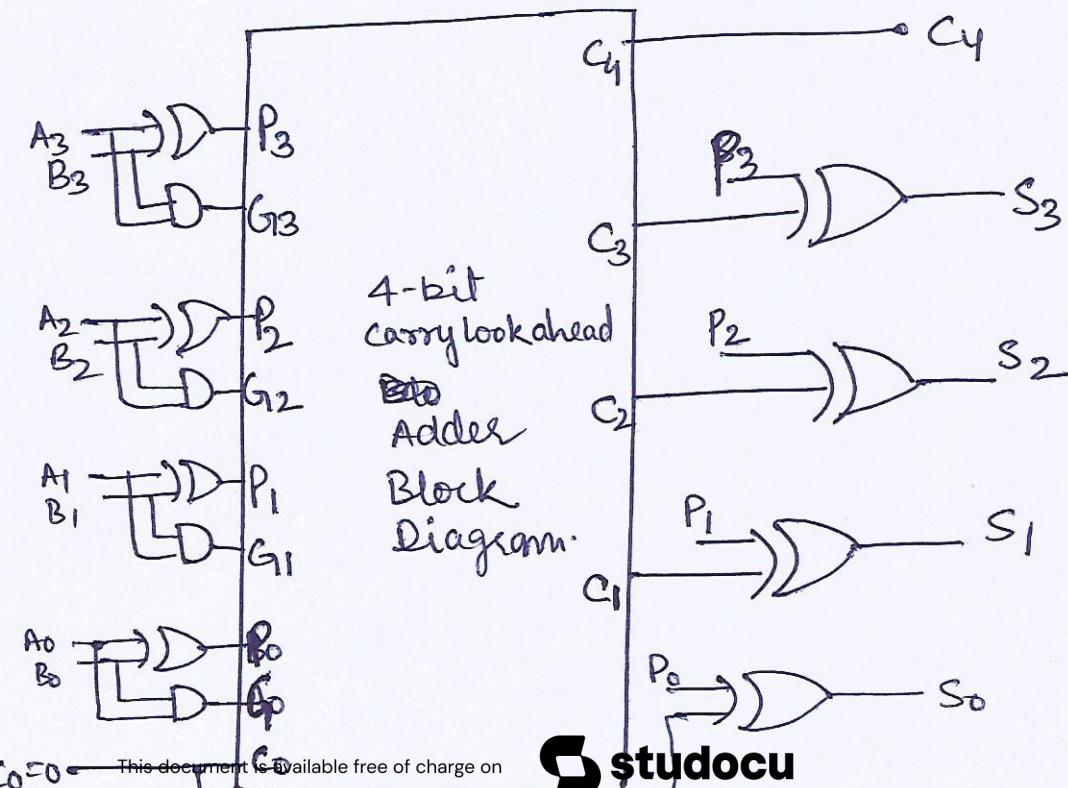
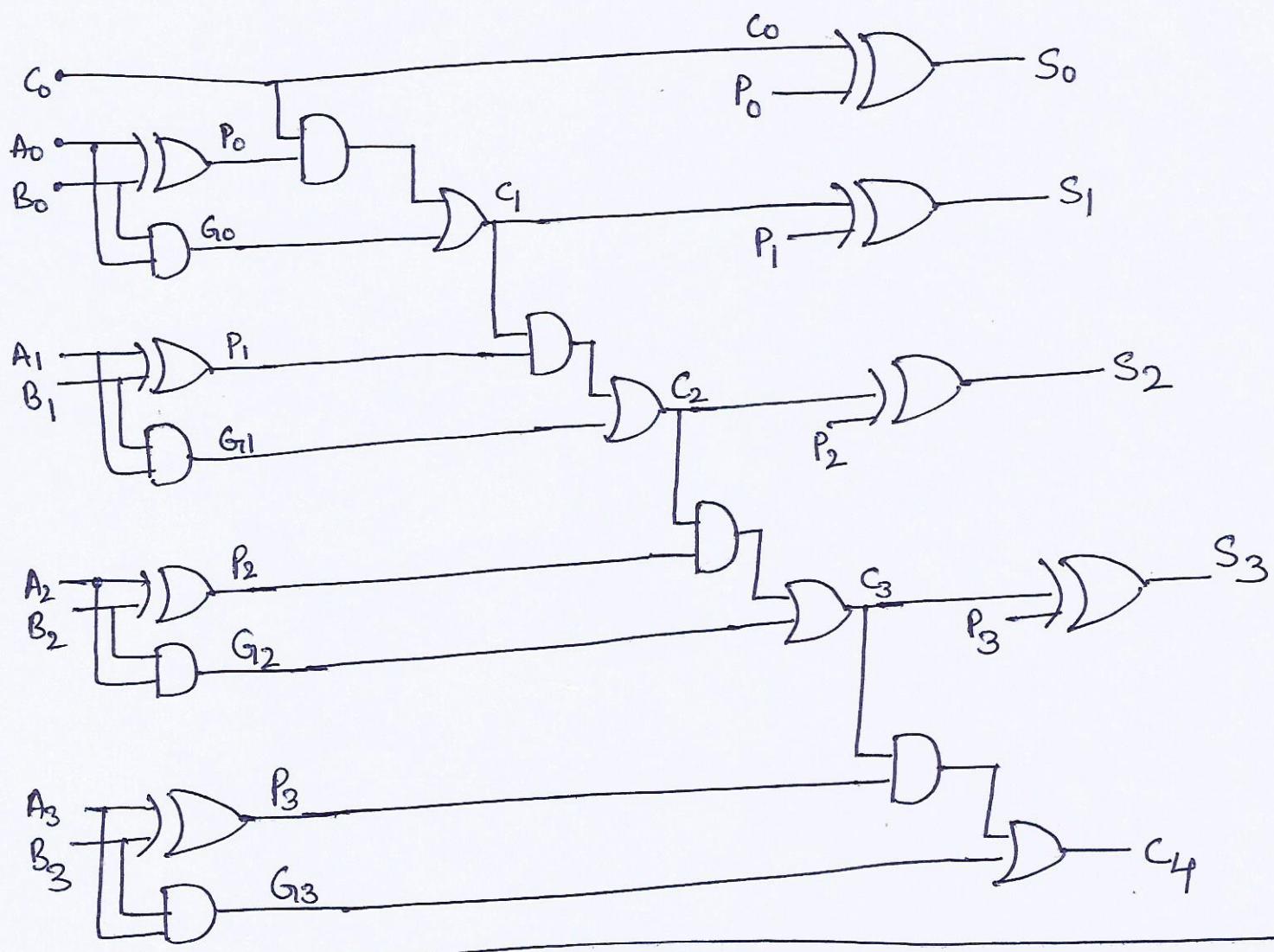
$$C_{3+i} = G_3 + P_3 C_3$$

$$C_4 = G_3 + P_3 [G_2 + P_2 [G_1 + P_1 [G_0 + P_0 C_0]]]$$

4-bit carry look-ahead ~~Adder~~ Carry Generation circuit

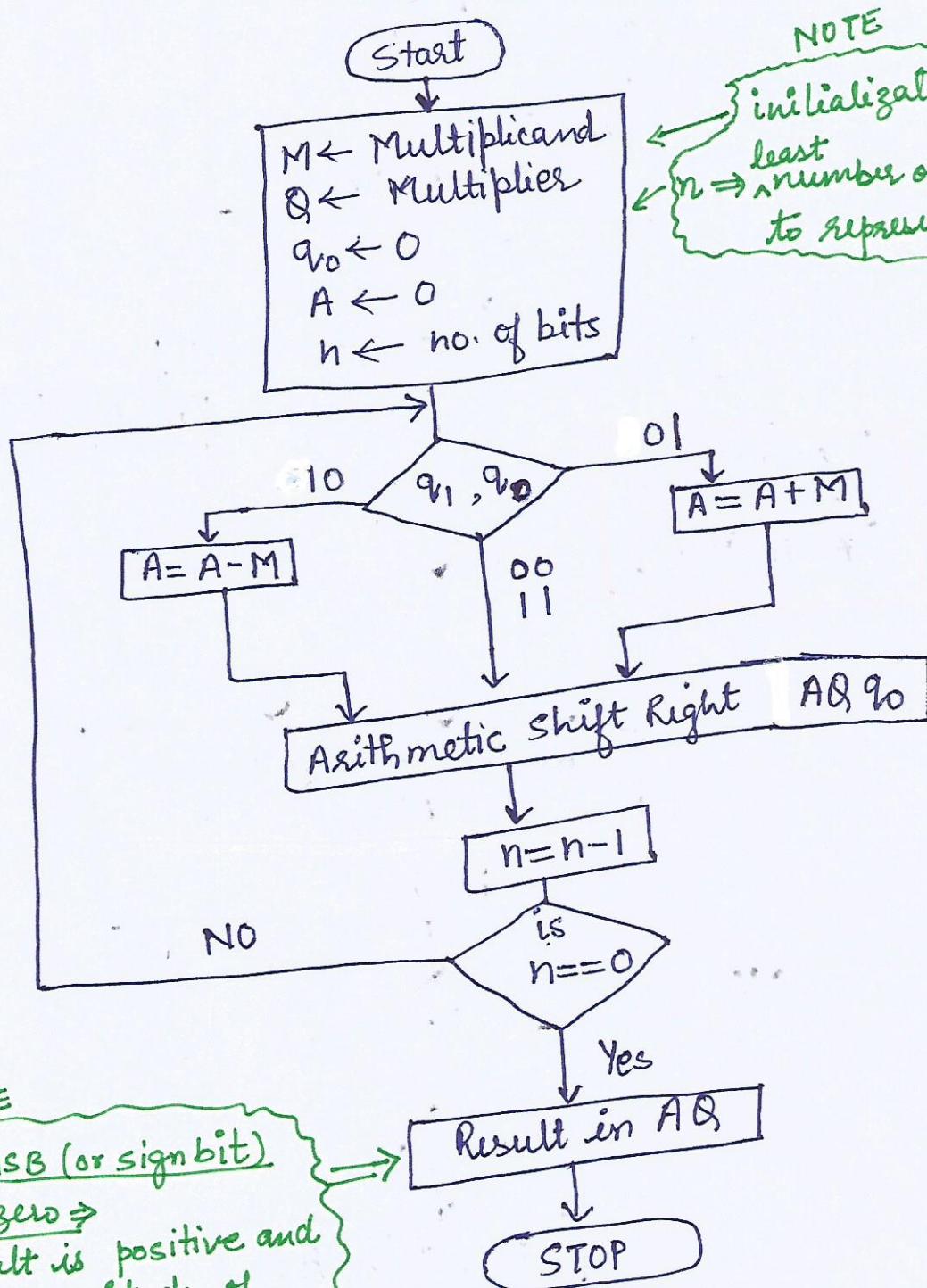


4 bit carry lookahead adder circuit:



Signed Multiplication : (Booth's Algorithm)

Flowchart:



NOTE

if MSB (or sign bit)
is zero \Rightarrow

Result is positive and
the magnitude of
the product.

if MSB (or sign bit 1)

Result is negative
take 2's complement to
get the magnitude and
place negative sign

NOTE
initialization
 $n \rightarrow$ least number of bits
to represent operands

Signed Booth Multiplication Example:

Example $(-7) \times (+3) = (-21)$

$$\begin{aligned} M &= (-7)_{10} = 2's(0111) \\ &= (1001)_2 \\ -M &= (0111)_2 \end{aligned}$$

Tracing Table

		$(+3)$			Action/Comment
	n	A	Q_{n-1}	q_0	
STEP①	4	0000	0011	0	initialization
		0111	0111	0	$A = A - M \Rightarrow A = A + (-M)$
		0011	1001	1	ASR AQq_0
STEP②	3	0011	1001	1	$n = n-1$
		0001	1100	1	ASR AQq_0
STEP③	2	0001	1100	1	$n = n-1$
		1010	1100	1	$A = A + M$
		1101	0110	0	ASR AQq_0
STEP④	1	1101	0110	0	$n = n-1$
		1110	1011	0	ASR AQq_0
	0	1110	1011	0	$n = n-1$

Result $\Rightarrow A Q$

11101011

↑ sign bit is one. So the result is negative, take 2's complement to get the magnitude and place negative sign.

Final result: $(-21)_{10}$

11101011
↓ 2's complement

00010100

↓ +1

00010101 $\Rightarrow (21)_{10}$

Array Multiplier:

Binary Multiplication is done by doing additions. Partial Products are calculated by multiplying the multiplicand by each bit of the multiplier and then summing the partial products.

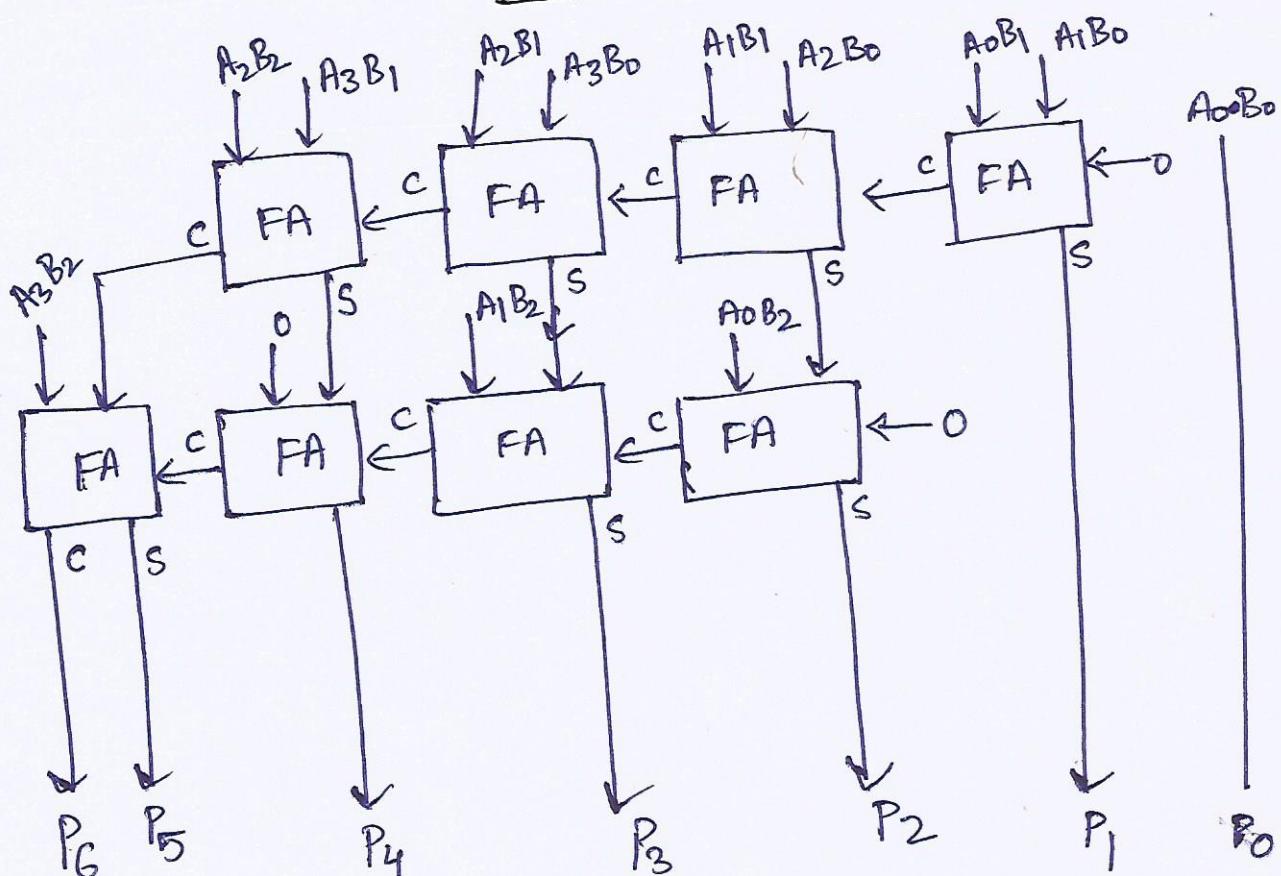
Ques: Design 4×3 binary or Array multiplier.

$$\text{Given } A \Rightarrow A_3 A_2 A_1 A_0$$

$$B \Rightarrow B_2 B_1 B_0$$

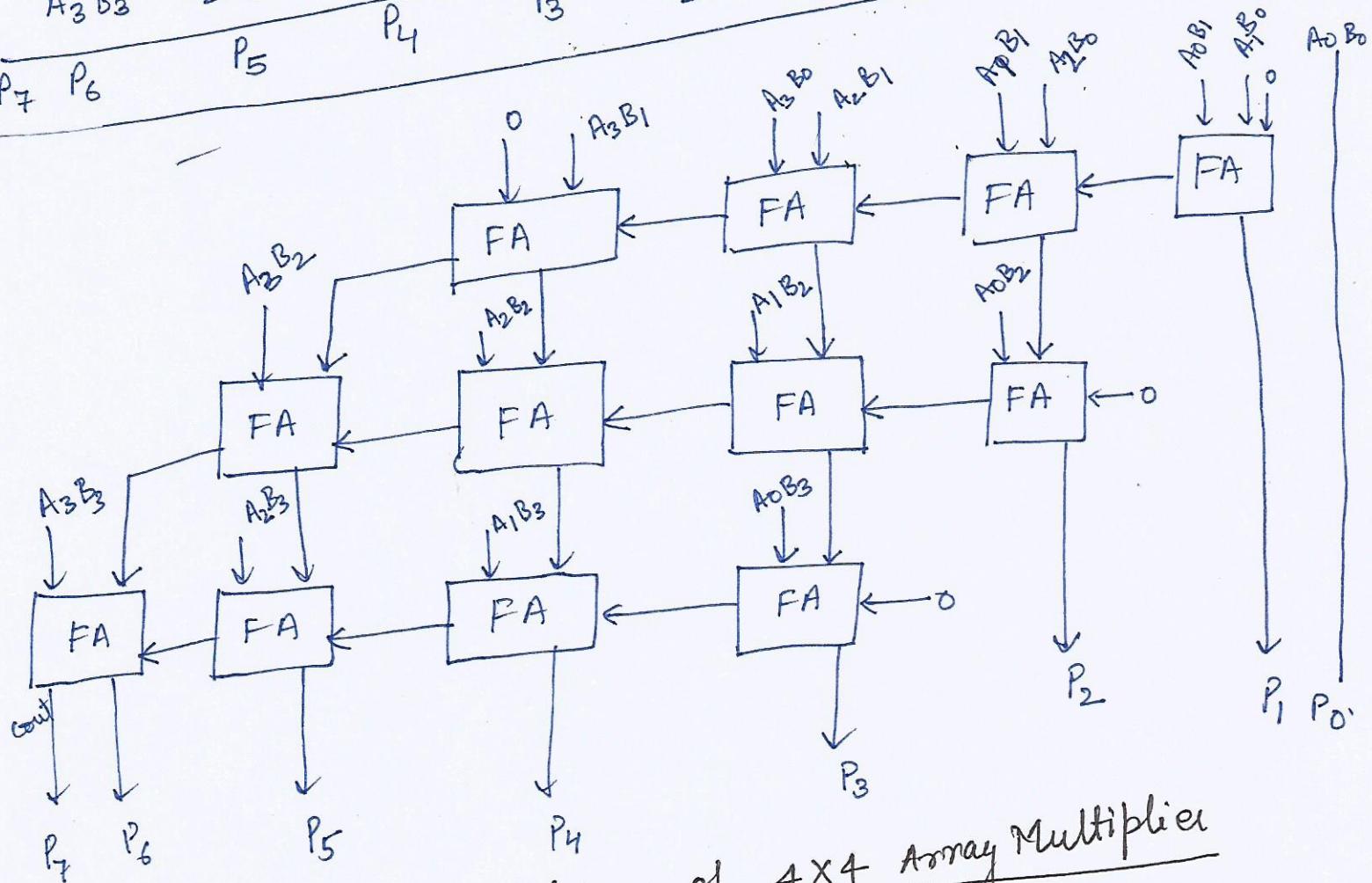
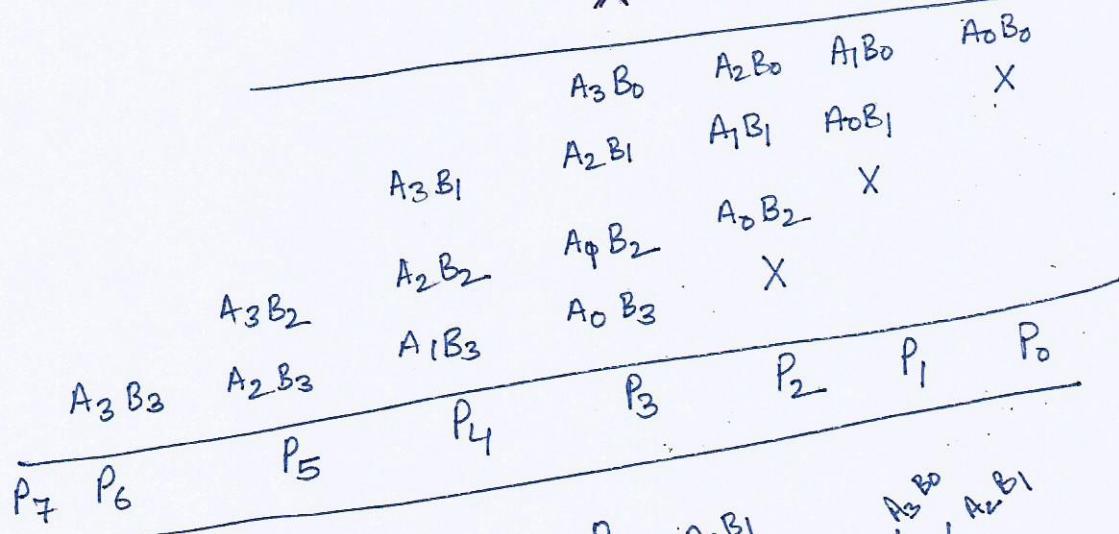
$$\text{Number of bits in product} \Rightarrow 4+3=7$$

$$\begin{array}{r}
 & A_3 & A_2 & A_1 & A_0 \\
 & \times & B_2 & B_1 & B_0 \\
 \hline
 & A_3 B_0 & A_2 B_0 & A_1 B_0 & A_0 B_0 \\
 & A_3 B_1 & A_2 B_1 & A_1 B_1 & A_0 B_1 \\
 & A_3 B_2 & A_2 B_2 & A_1 B_2 & A_0 B_2 \\
 \hline
 & P_6 & P_5 & P_4 & P_3 & P_2 & P_1 & P_0
 \end{array}$$



4x4 Binary Array Multiplier:

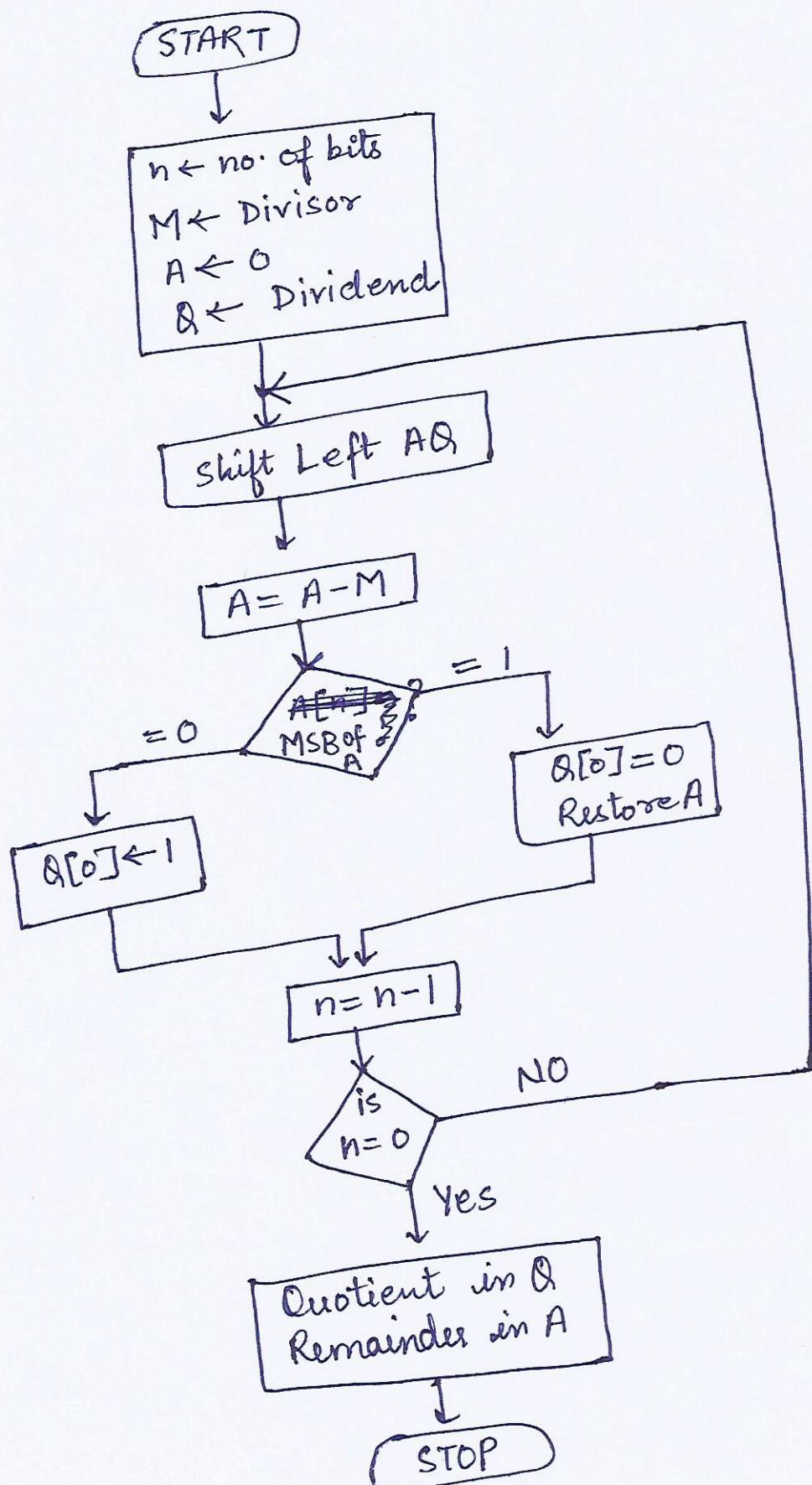
$$\begin{array}{r} A = A_3 \ A_2 \ A_1 \ A_0 \\ \times B = B_3 \ B_2 \ B_1 \ B_0 \end{array}$$



Circuit Diagram of 4x4 Array Multiplier

Division of unsigned Integer: Restoring Division Algorithm:

Flowchart:



Example: Divide 11 by 3.

(Q) Dividend = 11

(M) Divisor = 3

$M=3 \Rightarrow (00011)_2 \Rightarrow$ Number of bits in M = $n+1$

$Q=11 \Rightarrow (1011)_2$

$-M = (11101)_2$

Tracing Table

n	M	A	Q	Action / comment
4	00011	00000	1011	Initialization
		00001	011?	shift left AQ
		11110	011?	$A = A - M \Rightarrow 00001 + 11101 = 11110$
				$A[n] == 1$
3		00001	0110	$Q[0] \leftarrow 0$ Restore A $n = n-1$
		00010	110?	shift left AQ
2		11111	110?	$A = A - M$
		00010	1100	$Q[0] \leftarrow 0$, Restore A
		00010	1100	$n = n-1$
1		00101	100?	SL AQ
		00010	100?	$A = A - M$
		00010	1001	$Q[0] \leftarrow 1$
		00101	001?	$n = n-1$
		00010	001?	SL AQ
		00010	0011	$A = A - M$
				$Q[0] \leftarrow 1$

n	M	A	Q	Action / comment
11	00011	00010	0011	$n=n-1$

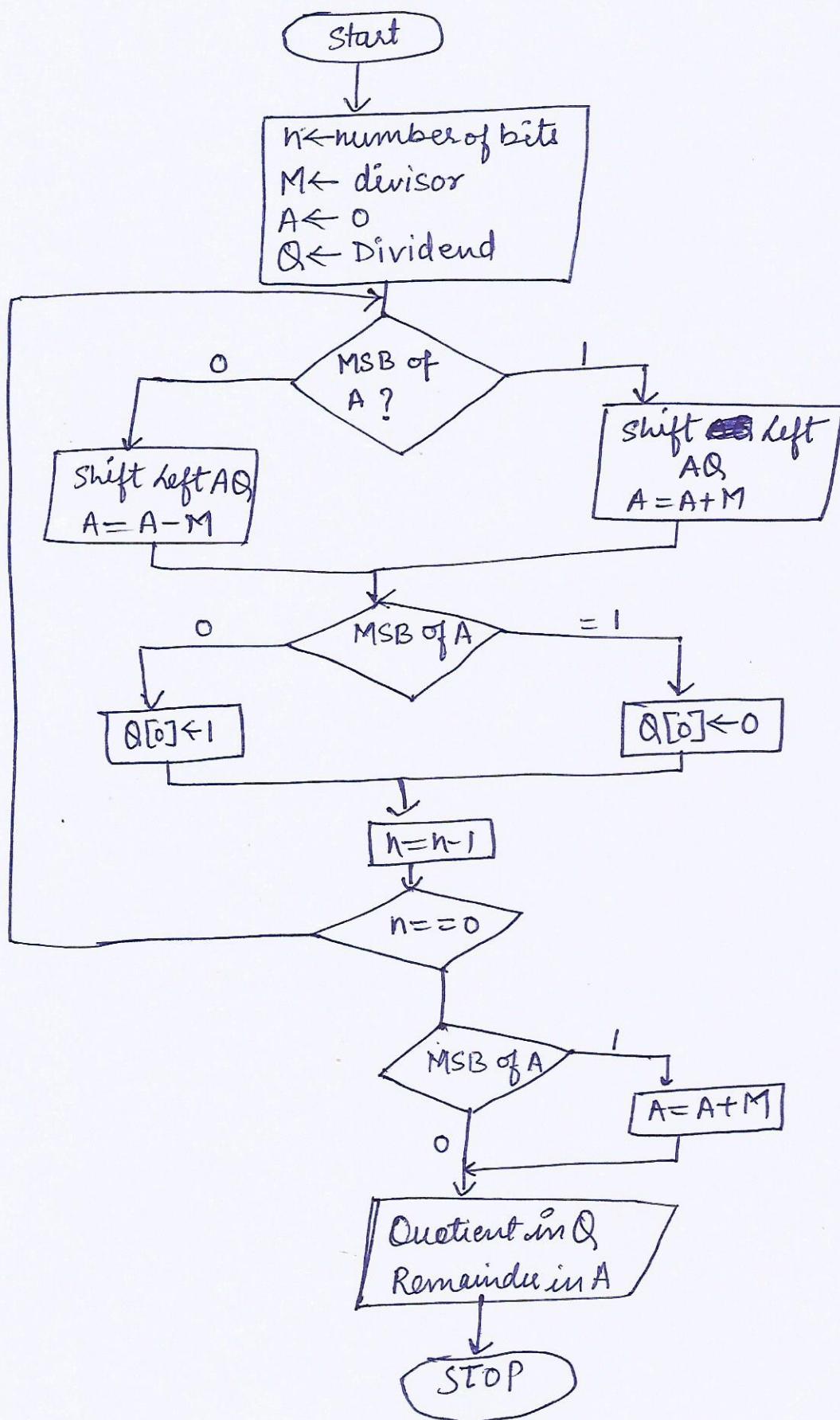
Now Quotient in $Q = 0011 \Rightarrow 3$

Remainder in $A = 00010 \Rightarrow 2$

11 (Dividend) / 3 (Divisor) \Rightarrow 3 Quotient
2 Remainder

Flow chart of Unsigned Integer Division

Non-Restoring Method



Example Divide 11 by 3.

Minalshi

(Q) Dividend = 11

Number of bits $\Rightarrow n = 4$

(M) Divisor = 3

Number of bits $M = n+1 = 5$

$$-M = 11101$$

Tracing Table:

<u>n</u>	<u>M</u>	<u>A</u>	<u>Q</u>	Action / comment
4	00011	00000	1011	Initialization
		00001	011?	SL A Q
		11110	011?	$A = A - M$
		11110	0110	$Q[0] \leftarrow 0$
		11110	0110	$n = n - 1$
3		11100	110?	SL A Q
		11111	110?	$A = A + M$
		11111	1100	$Q[0] \leftarrow 0$
2		11111	1100	$n = n - 1$
		11111	100?	SL A Q
		00010	100?	$A = A + M$
		00010	100?	$\begin{array}{r} 11111 \\ + 00011 \\ \hline 100010 \end{array}$
		00010	1001	↑ discard it
1		00101	001?	$Q[0] \leftarrow 1, n = n - 1$
		00010	001?	SL A Q
		00010	001?	$A = A - M$
		00010	0011	$Q[0] \leftarrow 1$
0 (Yes)		00010	0011	$n = n - 1$

MSB of A = 0 \Rightarrow then Result:

$$\text{Quotient in } Q = (0011)_2 = 3$$

$$\text{Remainder in } A = (00010)_2 = 2$$

Ques Divide 7 by 3 using Non-Restoring method.

$$\begin{array}{l} Q = \text{Dividend} = 7 \Rightarrow 111 \\ M = \text{Divisor} = 3 \Rightarrow 0011 \\ -M = 1101 \end{array} \quad] \quad n=3$$

0 → Subt
1 → Add.

<u>n</u>	A	Q	Action
3	0000	111	Initialization
	↑		
	0001	11?	
	+1101		SL AD, $A = A - M$
	1110	11?	$Q[0] \leftarrow 0$
	↑		
	1110	110	
2	1101	10?	Shift Left $A = A + M$
	0011		
carry discarded	0.000	101.	$Q[0] \leftarrow 1$
1	0001	01?	
	1101		$A = A - M$
	1110	010	$Q[0] \leftarrow 0$
0			
	↓	↓	
			Remainder Quotient

$$A \Rightarrow 1110$$

↑

$$A \Rightarrow 1110$$

$$\begin{array}{r} +M \Rightarrow +0011 \\ \hline 0001 \end{array}$$

$$\text{Remainder} = 1 \quad \text{Quotient} = 2$$

Floating Point Representation:

→ we can represent a floating point number in the form

$$\pm s \times B^{\pm E}$$

$B \Rightarrow$ Base, For binary $B=2$

⇒ This number can be stored in a binary word with three fields

- Sign (plus or minus)
- significand or Mantissa (S)
- Exponent E.

Sign	Biased Exponent	Mantissa or significand
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IEEE 754 Format:

- ① Single Precision (32 bit)
- ② Double Precision (64 bit)

	Signbit	Biased Exponent bits	Bias	Fraction bits
single Precision	1	8	127	23 bits
double Precision	1	11	1023	52 bits

Two functions for representing floating point numbers

→ Normalization

→ Biasing

Sign \Rightarrow	0 for positive numbers
	1 for negative numbers

A normalized number is one in which the most significant digit of the significand is one.

For base 2 representation

$$1.bbb\dots b \times 2^{E}$$

$$1.M \times 2^{E}$$

where $M \Rightarrow$ Mantissa / fraction / significand

Because the most significant bit is always one, it is unnecessary to store.

Given a number that is not normalized, the number may be normalized by shifting the radix point to the right of left most 1 bit and adjusting the exponent accordingly.

Biasing: Exponent is stored in biased representation.

In single precision case, 8 bits field yields the numbers 0 through 255.

with a bias of 127 ($2^7 - 1$), the true exponent values are in range -127 to +128

In double precision case, 11 bits field yields the numbers 0 through 2047 with a bias of 1023 ($2^{10} - 1$), the true exponent values are in range -1023 to +1024

Why Biasing: So that the bits in exponents can be treated as unsigned or nonnegative integers.

Instead of storing exponent in 2's complement, we can store it as unsigned number with the help of biasing

Ques: Represent $(-1234.125)_{10}$ into single precision and double precision representation.

Ans:- ① First convert 1234.125 into binary number.

$$(1234.125)_{10} = 10011010010.001$$

$$(1234.125)_{10} = 10011010010.001$$

Now Normalize this number \Rightarrow move the decimal point right to the leftmost 1.

$$\text{leftmost} = 1.0011010010001 \times 2^{+10}$$

single precision: 32 bit

single precision: 32 bit
g (ii) calculate biased exponent.

$$\begin{aligned}
 \text{Biased Exponent } E' &= E + \text{Bias} \\
 &= +10 + 12\text{f} \\
 &= 13\text{f} \\
 &= 10001001
 \end{aligned}$$

sign = negative = 1

sign	Biased Exponent	Mantissa (23)
1	10001001	00110100100010000000000000

Double Precision 64 bits

Double Precision 64 bits:
 (ii) calculate biased exponent $E' = E + \text{Bias}$
 $= 10 + 1023$

$$= 10 + 1023$$

$$= 1033$$

$$= 10000001001$$

sign	Exponent (E')	Mantissa (52)
1	10000001001	00110100100010000000- -0000

Ques Represent 000100110101.001101 using single precision and double precision.

Given number 000100110101.001101

Normalized Number $\Rightarrow 1.00110101001101 \times 2^{+9}$

Biased Exponent (SP): $E' = E + 127$

$$= 9 + 127$$

$$\Rightarrow 136$$

$$\Rightarrow 10001000$$

Biased Exponent (DP): $E' = E + 1023$

$$= 9 + 1023 \Rightarrow 1032$$

$$\Rightarrow 10000001000$$

single Precision

0	10001000	0011010100110100000000000
---	----------	---------------------------

Double Precision

0	10000001000100110101001101000---00
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Ques: Represent $-(0.0000010001100110)_2$ using single precision and double precision.

Normalize the number

$$\Rightarrow 1.00011001101 \times 2^{-6}$$

Biased Exponent in SP: $E' = E + \text{Bias} \Rightarrow -6 + 127$
 $\Rightarrow 121$
 $\Rightarrow 01111001$

Biased Exponent in DP: $E' \Rightarrow E + \text{Bias} \Rightarrow -6 + 1023$
 $\Rightarrow 1017$
 $\Rightarrow 0111111001$

single Precision \Rightarrow

1	0111001	000110011010000000000000
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Double Precision:-

1	0111111001	0001100110100000--000
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Floating-Point Addition and Subtraction:

- Four basic phases of the algorithm for addition and subtraction.

1. Check for zeros.
2. Align the significands
3. Add or subtract the significands
4. Normalize the result.

Floating Point Numbers	Addition and subtraction
$X = X_s \times B^{X_E}$ $Y = Y_s \times B^{Y_E}$	$X+Y = (X_s \times B^{X_E - Y_E} + Y_s) \times B^{Y_E}$ $X-Y = (X_s \times B^{X_E - Y_E} - Y_s) \times B^{Y_E}$ $\left. \begin{array}{l} X+Y = (X_s \times B^{X_E - Y_E} + Y_s) \times B^{Y_E} \\ X-Y = (X_s \times B^{X_E - Y_E} - Y_s) \times B^{Y_E} \end{array} \right\} X_E \leq Y_E$

Example:

$$X = 0.3 \times 10^2 = 30$$

$$Y = 0.2 \times 10^3 = 200$$

$$\begin{aligned}
 X+Y &= (0.3 \times 10^{2-3} + 0.2) \times 10^3 = (0.03 + 0.2) \times 10^3 \\
 &= 0.23 \times 10^3 \\
 &= 230
 \end{aligned}$$

$$\begin{aligned}
 X-Y &= (0.3 \times 10^{2-3} - 0.2) \times 10^3 = (0.03 - 0.2) \times 10^3 \\
 &= (-0.17) \times 10^3 \\
 &= -170
 \end{aligned}$$

Phase ① : Zero check: If either operand is 0, the other is reported as the result.

Phase ② : Significand alignment: To manipulate the numbers so that the two exponents are equal.

Alignment may be achieved by shifting either the smaller number to right (increasing its exponent) or shifting the larger number to left.

- often small number is picked for shifting

Phase ③ Addition: Next, two significands are added together taking into account their signs.

Phase ④ Normalization: This phase normalize the result.

Normalization consists of shifting significand digits left until the most significant digit is nonzero. Each shift causes a decrement of the exponent and thus could cause an exponent underflow. Finally, the result must be rounded off and the reported.

For floating-point addition and subtraction, it is necessary to ensure both the operands have the same exponent value. This may require shifting the radix point on one of the operands to achieve alignment.

A floating-point may produce one of these conditions:

- Exponent overflow
- Exponent Underflow
- Significand underflow
- Significand overflow

Exponent Overflow: A positive exponent exceeds the maximum possible exponent value. In some systems,

this may be designated as $+\infty$ or $-\infty$.

Exponent Underflow: A negative exponent is less than the minimum possible exponent value. (e.g. -200 is less than -127) This means the number is too small to be represented and it may be reported as zero.

Significand underflow: In the process of aligning significands digits may flow off the rightend of the significand. So, some form of rounding is required.

Significand overflow: The addition of two significands of the same sign may result in a carry out of the most significant bit. This can be fixed by realignment.

Flow chart for Floating Point Addition and Subtraction

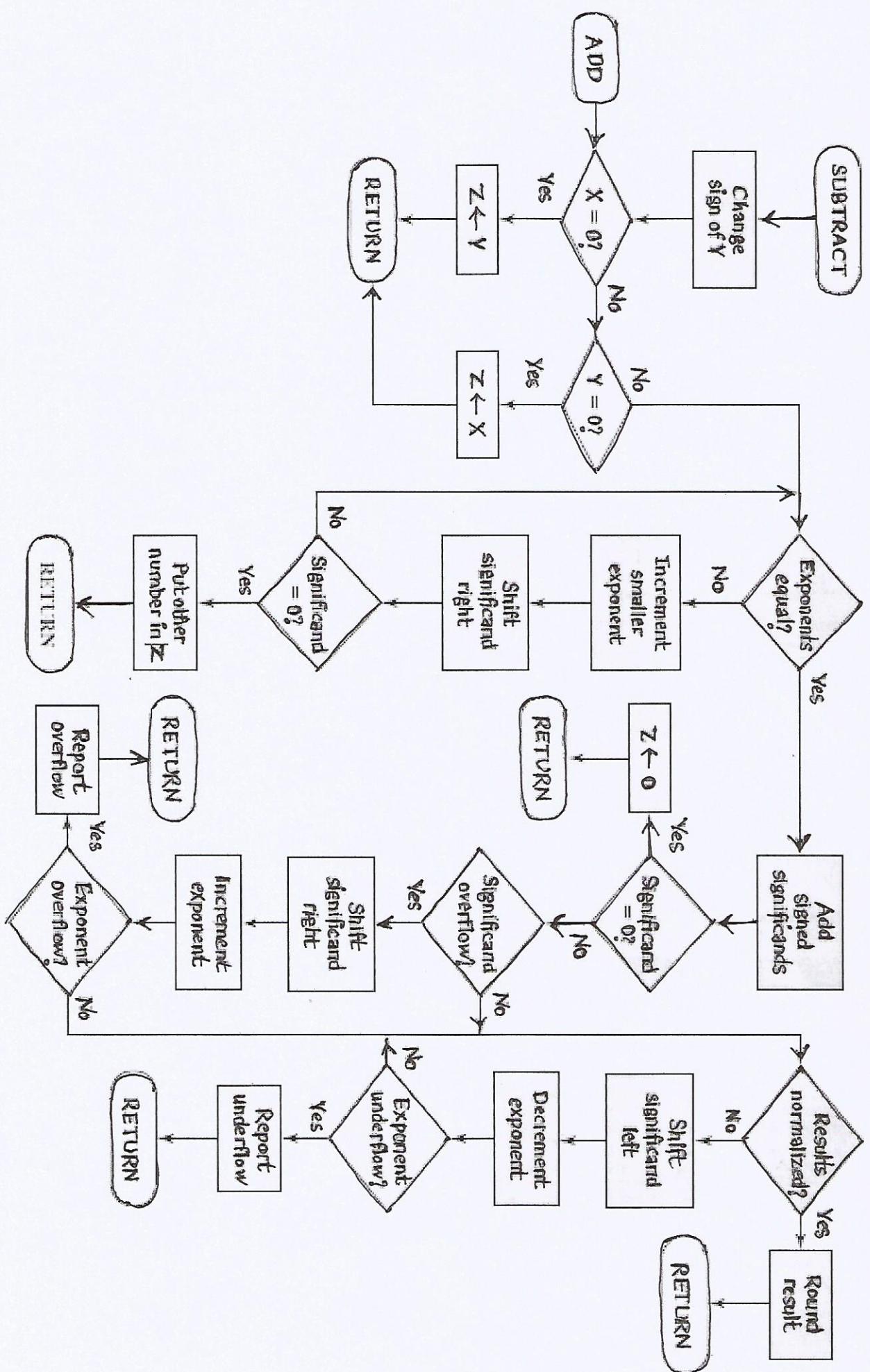


Figure 9.22 Floating-Point Addition and Subtraction ($Z \leftarrow Z \pm Y$)

Flow chart for Floating Multiplication

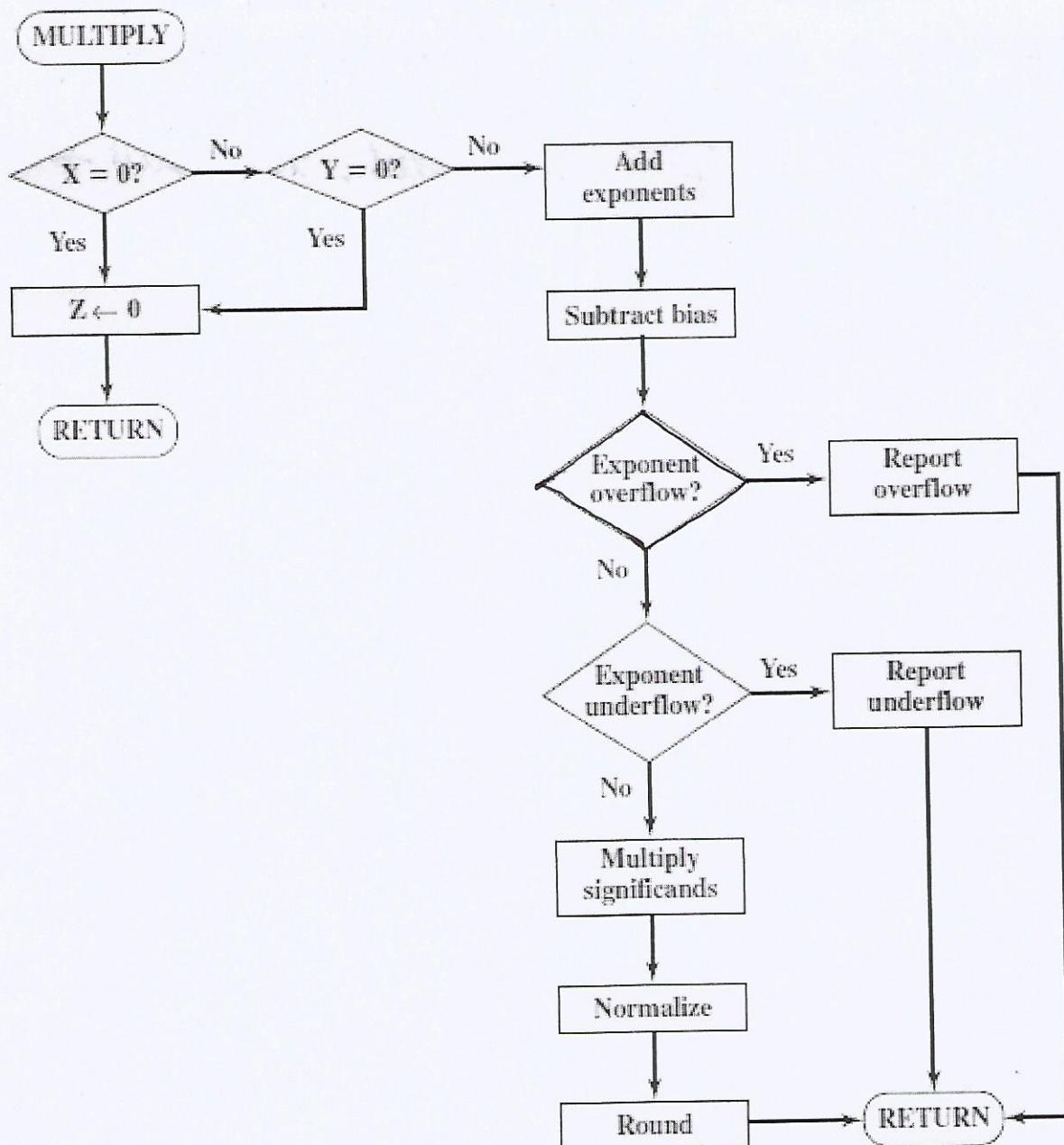


Figure 9.23 Floating-Point Multiplication ($Z \leftarrow X \times Y$)

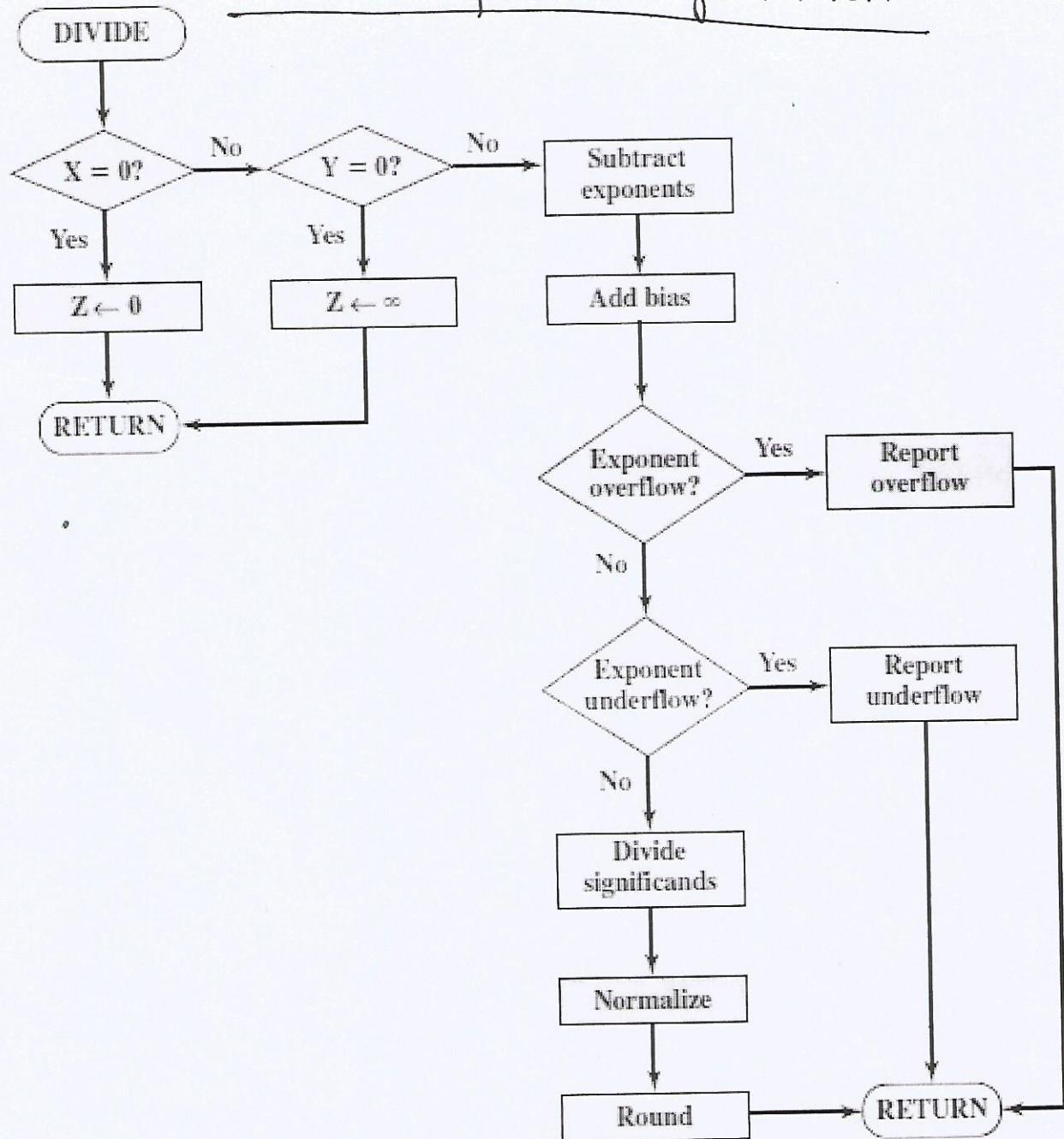
Floating-Point Multiplication:

- if either operand is 0(zero), zero is reported as result.
- The next, ~~step~~ step is to add the exponents. if the exponents are stored in biased form, the exponent sum would have doubled the bias. Thus, the bias must be subtracted from the sum. The result could be either an exponent underflow or exponent overflow, which would be reported.

- if the exponent of the product is within the range, the next step is to multiply the significand, taking into account their signs.
- After the product is calculated, the result is normalized and rounded.

flow3.jpg

Flow chart for Floating Division



Floating Point division:-

- Test for zero → if divisor is zero, an error is issued or result is set to infinity.
→ if dividend is zero, the result is zero.
- Next, the divisor exponent is subtracted from the dividend exponent. This removes bias, which must be added back in.
- Tests are made for exponent underflow or overflow.
- The next step is to divide the significands.
- Followed by normalization and rounding.

Logic Microoperations :

→ logic microoperations specify binary operations ~~for~~ for strings of bits stored in registers. These operations consider each bit of the register separately and treat them as binary variables.

For Example

- AND (\wedge)
- OR (\vee)
- NOT (\neg)
- NOR
- NAND
- X-OR (\oplus)
- X-NOR (\ominus)

$$R_1 = 1010$$

$$R_2 = 1100$$

$$\underline{\underline{AND}} \Rightarrow \begin{array}{r} 1010 \\ 1100 \\ \hline 1000 \end{array}$$

$$\underline{\underline{OR}} \Rightarrow \begin{array}{r} 1010 \\ 1100 \\ \hline 1110 \end{array}$$

$$\underline{\underline{NOT}} \Rightarrow \begin{array}{r} 1010 \\ \hline 0101 \end{array}$$

$$\underline{\underline{NOR}} : \begin{array}{r} 1010 \\ 1100 \\ \hline 0001 \end{array}$$

$$\underline{\underline{NAND}} : \begin{array}{r} 1010 \\ 1100 \\ \hline 0111 \end{array}$$

$$\underline{\underline{X-OR}} : \begin{array}{r} 1010 \\ 1100 \\ \hline 0110 \end{array}$$

$$\underline{\underline{X-NOR}} : \begin{array}{r} 1010 \\ 1100 \\ \hline 1001 \end{array}$$

X-NOR

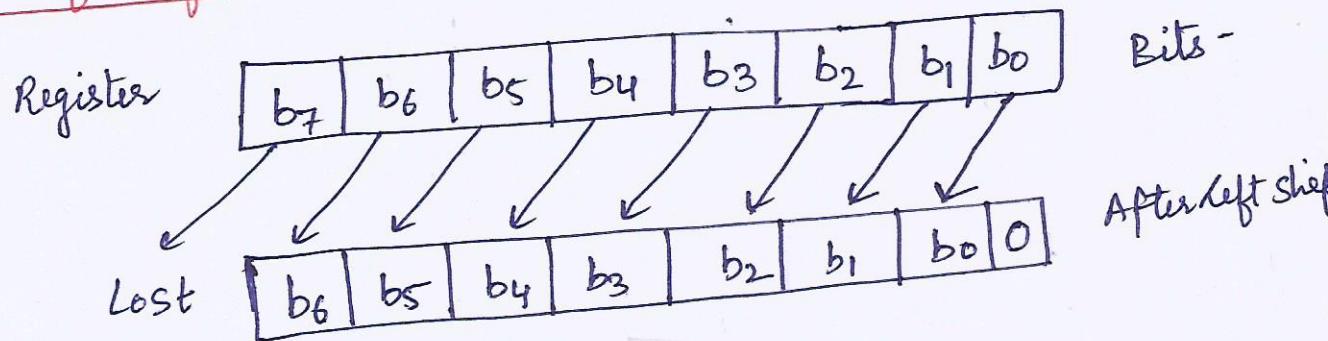
Shift Microoperations

→ Shift Micro operations are used for serial transfer of data.

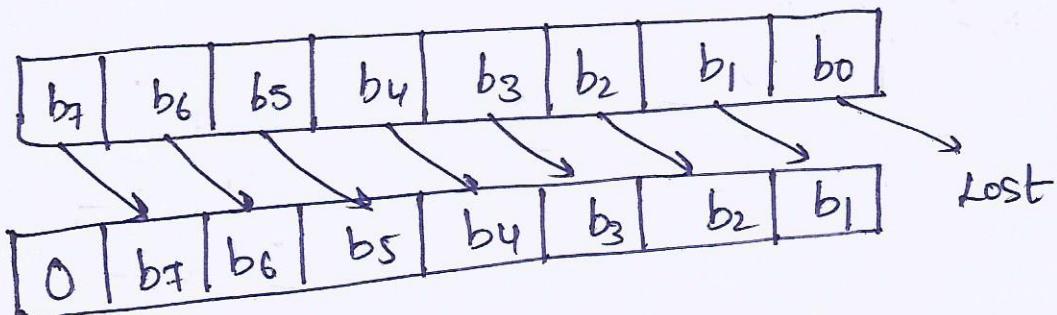
Types of Shift operations

- Logical shift —
 - shift left
 - shift right
- Circular shift or Rotate —
 - circular Right shift
 - circular left shift
- Arithmetic —
 - Arithmetic Right shift
 - Arithmetic left shift

Logical shift left:



Logical shift Right:



Example:

1 0 0 1 0 1 1 1
0 0 1 0 1 1 1 0 After shift left

1 0 0 1 0 1 1 1
0 1 0 0 1 0 1 1 After shift right

Circular Shift: (Also known as Rotate operation).

→ circulates the bits of the register around the two ends without loss of information.

circular left shift

b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀
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↓ After circular left shift

b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀	b ₇
----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------

Circular Right Shift

b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀
----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------

After ↓ circular Right shift

b ₀	b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁
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Example

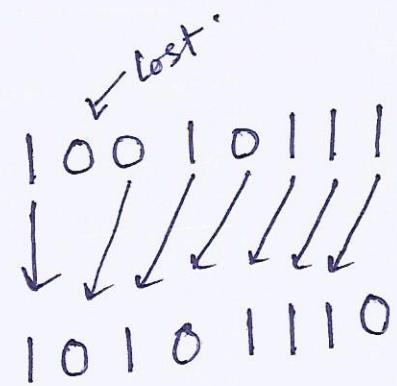
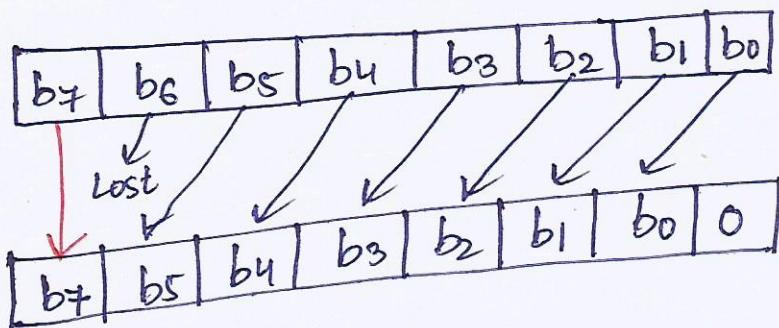
10010111
/ / / / / / /
00101111
Circular Left Shift

10010111
/ / / / / / /
11001011 → Circular Right Shift

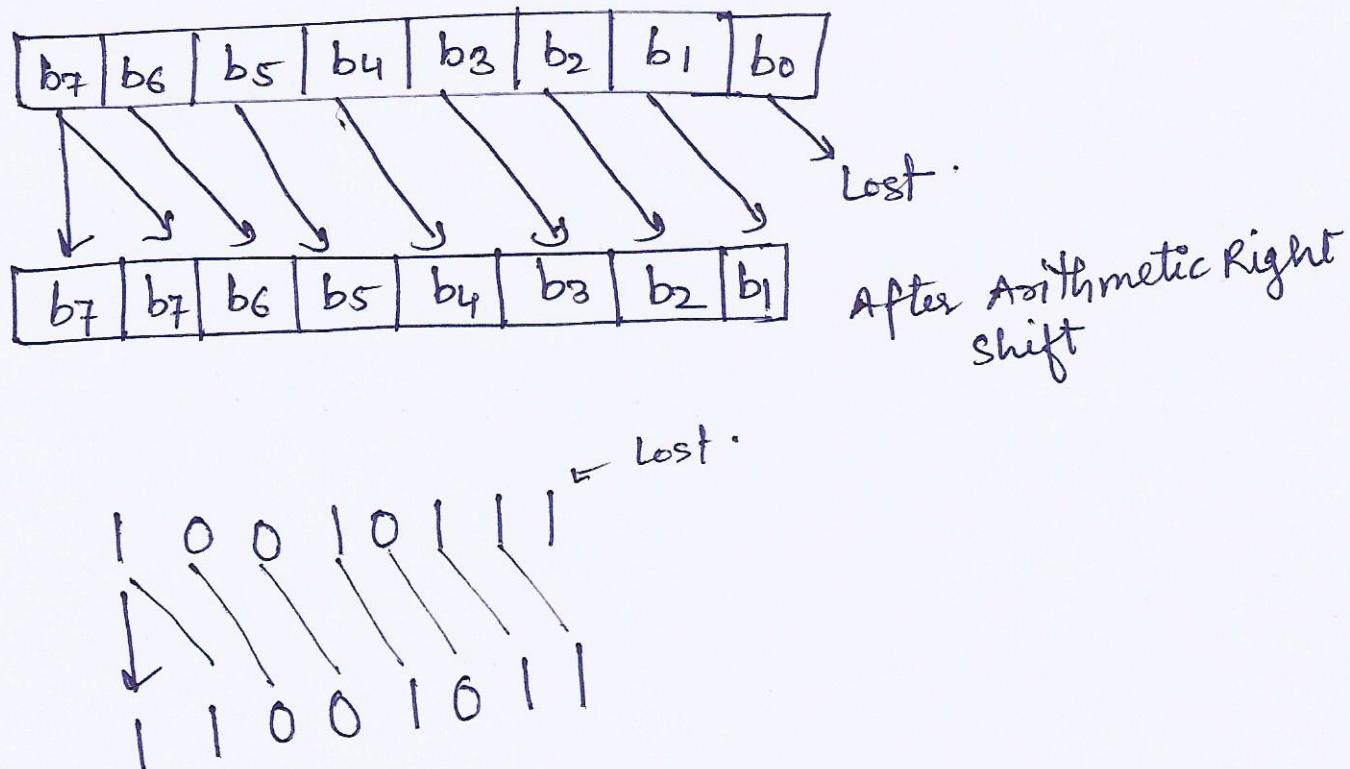
Arithmetic Shift:

- An Arithmetic shift operation shift a signed binary number to the left or right
- Arithmetic shift leaves the sign bit unchanged.

Arithmetic Left Shift



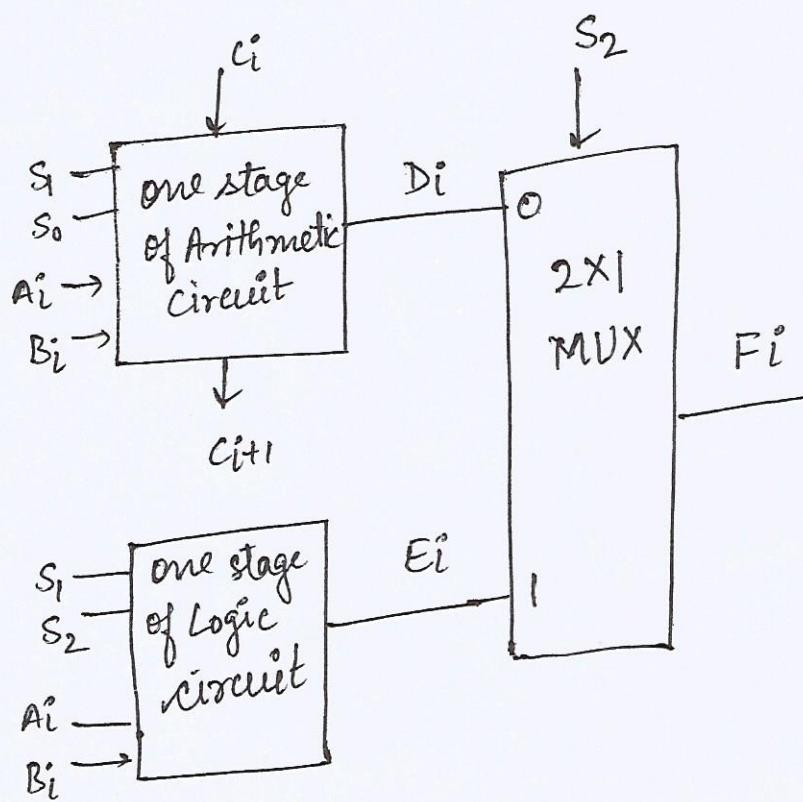
Arithmetic Right Shift:

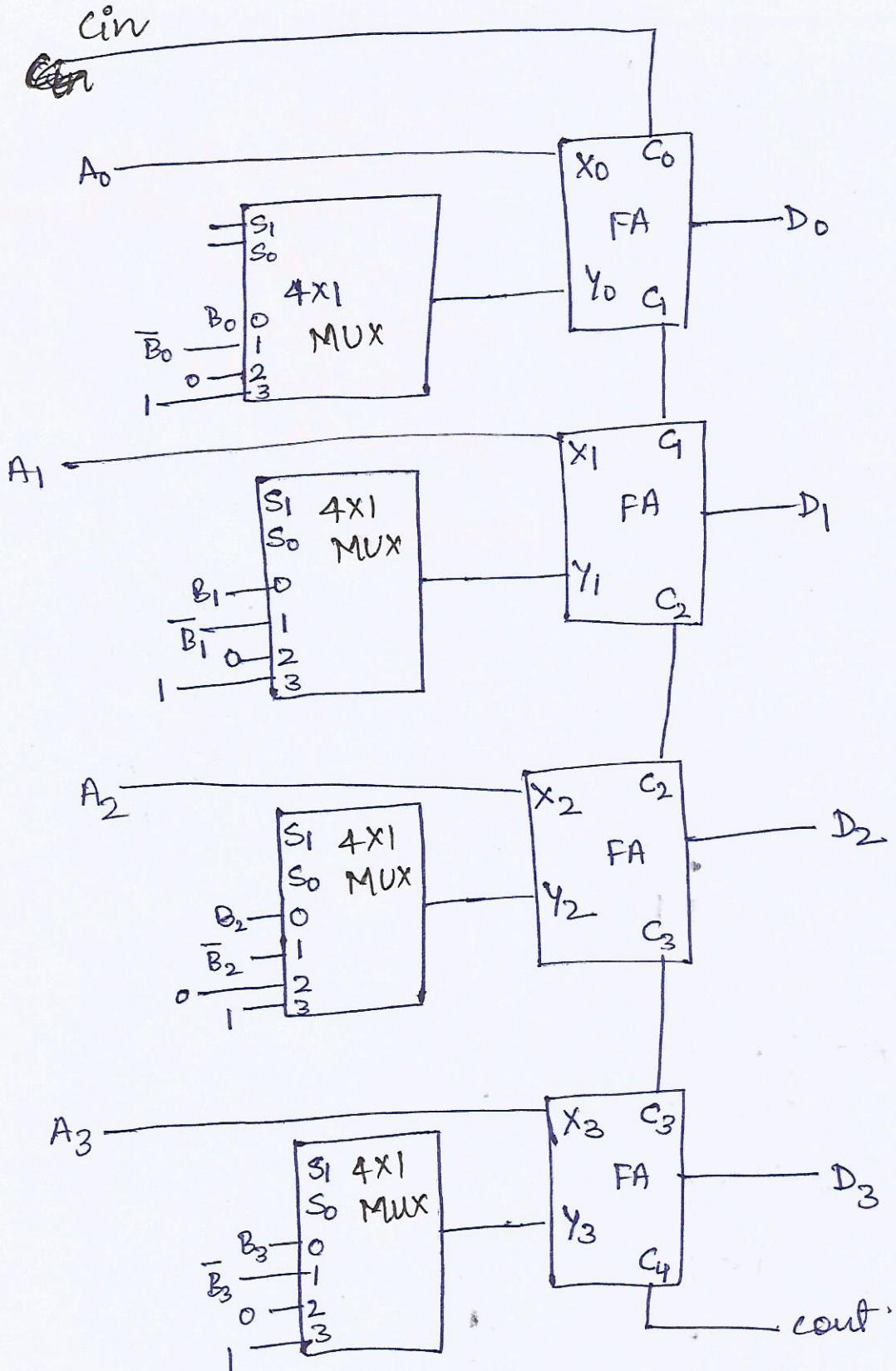


Arithmetic Logic Unit:

- CPU contains ALU, CU and Registers
- ALU is responsible for arithmetic and logical operations
- Basically ALU is a digital circuit that performs arithmetic operations like addition, subtraction, division and multiplication, and logical operations like AND, OR, XOR, NOT, etc.
- Types of ALU ⇒ ① Combinational ALU
② Sequential ALU.

combinational ALU's





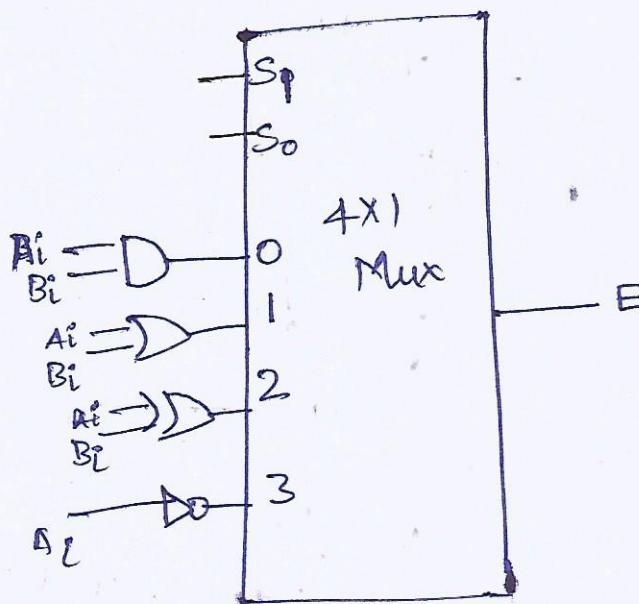
4 bit Arithmetic Circuit performing operations
like addition, subtraction, increment, decrement
and transfer etc.

Sq	S ₁	S ₀	Cin	Input	Output D = A + Y + Cin
0	0	0	0	B	D = A + B → Add
0	0	1	0	B	D = A + B + 1 → Add with carry
0	1	0	0	\bar{B}	D = A + \bar{B} → Subtract with Borrow
0	1	1	0	\bar{B}	D = A + \bar{B} + 1 → Subtract
1	0	0	0	0	D = A → Transfer A
1	0	1	0	0	D = A + 1 → Increment A
1	1	0	1	1	D = A - 1 → <u>Decrement A</u>
1	1	1	1	1	D = A → Transfer A

when S₁=1 and S₀=1 then
input is (B₃ B₂ B₁ B₀) 1111
which is 2's complement
of 0001 ⇒ It makes

$$\begin{array}{r}
 \begin{array}{cccc} A_3 & A_2 & A_1 & A_0 \\ +1 & 1 & 1 & 1 \end{array} \Rightarrow \begin{array}{c} A \\ +2\text{'s compl} \end{array} \\
 \hline
 \Rightarrow A + (-1) \\
 \Rightarrow A - 1
 \end{array}$$

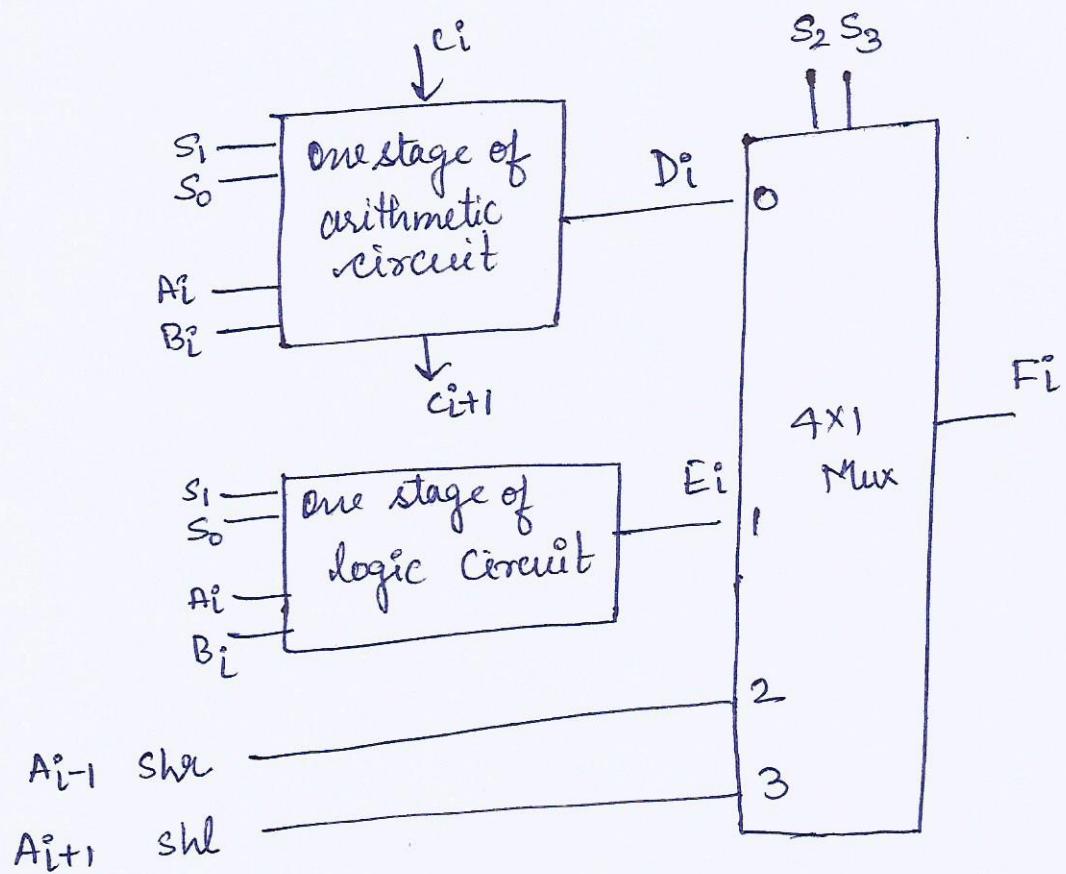
one stage of logic circuit:-



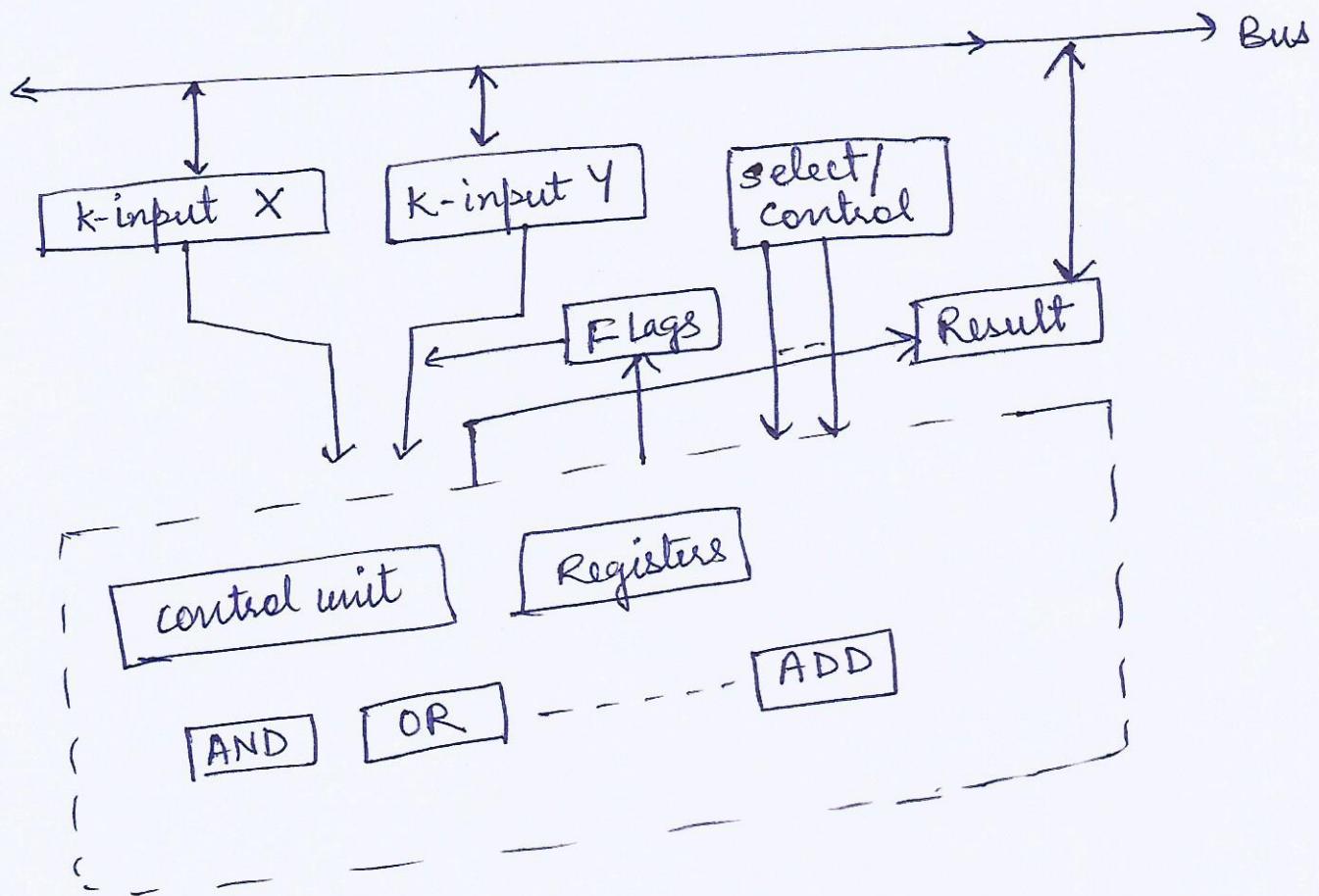
Function Table

S ₁	S ₀	Logic operation (E)
0	0	A _i AND B _i
0	1	A _i OR B _i
1	0	A _i XOR B _i
1	1	NOT A _i

One stage of arithmetic logic shift unit



Sequential logic circuit based ALU



- Two registers X and Y stores data or operands.
- select/control selects the appropriate arithmetic or logic operation which is performed over the data stored on register X and Y.
- After the execution of operation the result will be stored in result register.
- After the execution of operation flags may be set such as carry, zero result, positive or negative result, overflow, division by zero etc.