# Square Root Decomposition Ideas

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# 1 Ideas

• Update Query:  $k \ v, \ a[k] \leftarrow v$ Range Frequency:  $l \ r \ x$ , count Frequency of x in range [l, r]

#### Implementation Details:

- Preprocess the array into blocks of size  $\sqrt{n}$ .
- Maintain a frequency map for each block.
- For a query, combine results from relevant blocks.
- Update Query:  $k \ v, \ a[k] \leftarrow v$ Query:  $l \ r \ c$ , count elements in range [l, r] greater than or equal to c.

#### Implementation Details:

- Preprocess the array into blocks of Ordered Multiset of size  $\sqrt{n}$ .
- Update Query: 0 a b, set power of hole a to value b.

  Jump Query: 1 a, find the last hole visited and the total number of jumps before the ball leaves the row.

#### Implementation Details:

- Preprocess the array of N holes into blocks of size  $\sqrt{N}$ .
- For each hole i, precompute a pair of values: (1) the next hole the ball lands in *outside* its current block, and (2) the number of jumps it takes to get there.
- A jump query can then be answered in  $O(\sqrt{N})$  time by chaining these precomputed block-to-block jumps.
- An update only requires recomputing the values for the single block that was changed, also in  $O(\sqrt{N})$  time.

```
void process() {
                                   (int i = n - 1; i >= 0; i--) {
  int idx = i / SQ, r = min(n - 1, (i / SQ + 1) * SQ - 1LL);
  if (i + arr[i] > r) {
    jumps[i] = {i, 1};
}
 3
 4
                                       else
 6
                                            jumps[i] = jumps[i + arr[i]];
jumps[i].second++;
 8
                         }
10
                }
11
                void update(int idx, ll val) {
    arr[idx] = val;
    int blk_idx = idx / SQ;
    int l = blk_idx * SQ, r = min(n - 1, (blk_idx + 1) * SQ - 1LL);
    for (int i = r; i >= 1; i--) {
        if (i + arr[i] > r) {
            jumps[i] = {i, 1};
        } else {
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14
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17
18
19
                                       20
21
23
                         }
24
26
                pair < int, int > query(int 1) {
    pair < int, int > ans = {0, 0};
    int idx = 1;
    while (idx < n) {</pre>
27
29
30
                                   ans.first = jumps[idx].first;
ans.second += jumps[idx].second;
idx = jumps[idx].first + arr[jumps[idx].first];
32
33
                          return ans:
```

Range Update: L R X, add X to heights of hills in range [L, R].
 Jump Query: i k, find the final hill after making k jumps starting from hill i. A jump is to the nearest hill on the right that is strictly taller and within a distance of 100.

#### Implementation Details:

- Divide the array into blocks of size  $\sqrt{N}$ . Use a lazy array (e.g., lazy\_add) for efficient range updates on full blocks.
- For each hill *i*, precompute: the next immediate jump (next\_jump), the final hill reached if jumping only within the block (block\_exit\_node), and the number of jumps to get there (jumps\_to\_exit).
- This preprocessing is done efficiently in  $O(\sqrt{N})$  per block by iterating backwards with a stack to find the next greater element.
- A jump query uses the precomputed data to skip across blocks and performs a small linear scan for inter-block jumps, leading to an overall  $O(\sqrt{N})$  query time.

```
#include <bits/stdc++.h>
2
         using namespace std;
         using 11 = long long;
3
         const int MAX_N = 100005, MAX_SQRT_N = 320, MAX_JUMP_DIST = 100;
5
6
         int n, q, block_size;
         11 heights[MAX_N], lazy_add[MAX_SQRT_N];
         int next_jump[MAX_N], block_exit_node[MAX_N], jumps_to_exit[MAX_N];
9
10
         void rebuild block(int block idx) {
             int start_idx = block_idx * block_size;
int end_idx = min((block_idx + 1) * block_size, n);
11
             if (lazy_add[block_idx] != 0) {
                  for (int i = start_idx; i < end_idx; ++i) heights[i] += lazy_add[block_idx];</pre>
15
```

```
16
                   lazy_add[block_idx] = 0;
              }
17
18
              stack<int> s;
19
              for (int i = end_idx - 1; i >= start_idx; --i) {
   while (!s.empty() && (s.top() - i > MAX_JUMP_DIST || heights[s.top()] <= heights[i
20
21
                        ])) {
22
                        s.pop();
23
                   }
24
                   if (!s.empty()) {
25
                        next_jump[i] = s.top();
block_exit_node[i] = block_exit_node[next_jump[i]];
jumps_to_exit[i] = jumps_to_exit[next_jump[i]] + 1;
26
27
                   } else {
29
30
                        next_jump[i] = i;
                        block_exit_node[i] = i;
31
                        jumps_to_exit[i] = 0;
32
33
34
                   s.push(i);
35
              }
36
         }
37
         void update_range(int 1, int r, 11 val) {
   int start_block = 1 / block_size;
38
39
              int end_block = r / block_size;
40
41
              if (start_block == end_block) {
    for (int i = 1; i <= r; ++i) heights[i] += val;</pre>
42
43
                   rebuild_block(start_block);
44
45
                   return:
46
47
48
              for (int i = 1; i < (start_block + 1) * block_size; ++i) heights[i] += val;</pre>
49
              rebuild_block(start_block);
50
              for (int i = start_block + 1; i < end_block; ++i) lazy_add[i] += val;</pre>
51
52
              for (int i = end_block * block_size; i <= r; ++i) heights[i] += val;</pre>
53
54
              rebuild_block(end_block);
55
56
         int query_jumps(int start_idx, int k) {
   int current_hill = start_idx;
57
58
59
                   // Case 1: We don't have enough jumps to take the full precomputed path // We must take a single, precomputed step instead.
61
62
                   if (jumps_to_exit[current_hill] > 0 && k < jumps_to_exit[current_hill]) {</pre>
63
                        current_hill = next_jump[current_hill];
64
65
                        continue;
66
67
                   }
68
                   // Case 2: We have enough jumps to use the precomputed path.
if (jumps_to_exit[current_hill] > 0) {
69
70
                        k -= jumps_to_exit[current_hill];
current_hill = block_exit_node[current_hill];
71
73
                   }
74
                   // Case 3: We are now at a block's exit point and still have jumps left
75
76
                   // We must perform one jump by scanning manually. if (k > 0) {
77
78
                        bool jumped = false;
                        11 current_true_height = heights[current_hill] + lazy_add[current_hill /
79
                             block_size];
80
                        81
                             11 next_true_height = heights[j] + lazy_add[j / block_size];
82
                             if (next_true_height > current_true_height)
83
                                  current_hill = j;
84
85
                                  jumped = true;
86
87
                                  break;
88
90
                        if (!jumped) break; // No further jump possible, we are stuck.
91
                   }
92
              return current_hill + 1;
93
94
          int main() {
97
              ios_base::sync_with_stdio(false), cin.tie(NULL);
98
              cin >> n >> q;
99
```

```
block_size = static_cast < int > (sqrt(n));
if (block_size == 0) block_size = 1;
100
101
102
                       for (int i = 0; i < n; ++i) {
    cin >> heights[i];
103
104
                       }
105
106
                       for (int i = (n - 1) / block_size; i >= 0; --i) {
    rebuild_block(i);
107
108
109
                      while (q--) {
   int type;
   cin >> type;
   if (type == 1) {
      int i, k;
      cin >> i >> k;
      cout << query_jumps(i - 1, k) << " \n";
   l else {</pre>
110
111
112
113
114
115
\frac{116}{117}
                              } else { int 1, r;
118
119
                                     ll x;
cin >> 1 >> r >> x;
update_range(1 - 1, r - 1, x);
120
121
122
                              }
123
124
125
                       return 0;
126
```