Square Root Decomposition Ideas

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1 Ideas

• Update Query: $k \ v, \ a[k] \leftarrow v$ Range Frequency: l r x, count Frequency of x in range [l, r]

Implementation Details:

- Preprocess the array into blocks of size \sqrt{n} .
- Maintain a frequency map for each block.
- For a query, combine results from relevant blocks.
- Update Query: $k \ v, \ a[k] \leftarrow v$ Query: l r c, count elements in range [l, r] greater than or equal to c.

Implementation Details:

- Preprocess the array into blocks of Ordered Multiset of size \sqrt{n} .
- Update Query: 0 a b, set power of hole a to value b. Jump Query: 1 a, find the last hole visited and the total number of jumps before the ball leaves the row.

Implementation Details:

- Preprocess the array of N holes into blocks of size \sqrt{N} .
- For each hole i, precompute a pair of values: (1) the next hole the ball lands in *outside* its current block, and (2) the number of jumps it takes to get there.
- A jump query can then be answered in $O(\sqrt{N})$ time by chaining these precomputed blockto-block jumps.
- An update only requires recomputing the values for the single block that was changed, also in $O(\sqrt{N})$ time.

```
cocess() {
  (int i = n - 1; i >= 0; i--) {
   int idx = i / SQ, r = min(n - 1, (i / SQ + 1) * SQ - 1LL);
   if (i + arr[i] > r) {
      jumps[i] = {i, 1};
} else {
      image for a second content of the second con
                                                                                                                                                                                                                jumps[i] = jumps[i + arr[i]];
jumps[i].second++;
                                                                                                                         }
12
                                                                                 void update(int idx, ll val) {
                                                                                                                         arr[idx] = val;
int blk_idx = idx / SQ;
int l = blk_idx * SQ, r = min(n - 1, (blk_idx + 1) * SQ - 1LL);
for (int i = r; i >= 1; i--) {
   if (i + arr[i] > r) {
      jumps[i] = {i, 1};
      }
}
14
15
19
                                                                                                                                                                                                                  jumps[i] = jumps[i + arr[i]];
```

```
jumps[i].second++;
22
23
                       }
               }
25
26
               pair < int , int > query(int 1) {
    pair < int , int > ans = {0,
    int idx = 1;
    while (idx < n) {</pre>
28
29
30
                               ans.first = jumps[idx].first;
ans.second += jumps[idx].seco
31
                                ans.second += jumps[idx].second;
idx = jumps[idx].first + arr[jumps[idx].first];
32
34
35
```

• Range Update: L R X, add X to heights of hills in range [L, R].

Jump Query: i k, find the final hill after making k jumps starting from hill i. A jump is to the nearest hill on the right that is strictly taller and within a distance of 100.

Implementation Details:

- Divide the array into blocks of size \sqrt{N} . Use a lazy array (e.g., lazy_add) for efficient range updates on full blocks.
- For each hill i, precompute: the next immediate jump (next_jump), the final hill reached if jumping only within the block (block_exit_node), and the number of jumps to get there (jumps_to_exit).
- This preprocessing is done efficiently in $O(\sqrt{N})$ per block by iterating backwards with a stack to find the next greater element.
- A jump query uses the precomputed data to skip across blocks and performs a small linear scan for inter-block jumps, leading to an overall $O(\sqrt{N})$ query time.

```
#include <bits/stdc++.h>
          using namespace std;
using l1 = long long;
2
3
          const int MAX_N = 100005, MAX_SQRT_N = 320, MAX_JUMP_DIST = 100;
6
          int n, q, block_size;
         11 heights[MAX_N], lazy_add[MAX_SQRT_N];
int next_jump[MAX_N], block_exit_node[MAX_N], jumps_to_exit[MAX_N];
9
          void rebuild_block(int block_idx) {
10
               int start_idx = block_idx * block_size;
12
               int end_idx = min((block_idx + 1) * block_size, n);
13
               if (lazy_add[block_idx] != 0) {
14
                    for (int i = start_idx; i < end_idx; ++i) heights[i] += lazy_add[block_idx];
lazy_add[block_idx] = 0;</pre>
15
16
17
19
               for (int i = end_idx - 1; i >= start_idx; --i) {
   while (!s.empty() && (s.top() - i > MAX_JUMP_DIST || heights[s.top()] <= heights[i
20
^{21}
                          ])) {
                         s.pop();
22
24
                    if (!s.empty()) {
25
                         rest_jump[i] = s.top();
block_exit_node[i] = block_exit_node[next_jump[i]];
26
27
                         jumps_to_exit[i] = jumps_to_exit[next_jump[i]] + 1;
28
30
                         next_jump[i] = i;
                          block_exit_node[i] = i;
31
                         jumps_to_exit[i] = 0;
32
33
                    s.push(i):
34
               }
         }
37
          void update_range(int 1, int r, 11 val) {
   int start_block = 1 / block_size;
   int end_block = r / block_size;
38
39
40
41
               if (start_block == end_block) {
42
                    for (int i = 1; i <= r; ++i) heights[i] += val;</pre>
43
44
                    rebuild_block(start_block);
```

```
45
                    return;
 46
 47
               for (int i = 1; i < (start_block + 1) * block_size; ++i) heights[i] += val;</pre>
                rebuild_block(start_block);
 49
 50
               for (int i = start_block + 1; i < end_block; ++i) lazy_add[i] += val;</pre>
 51
 52
               for (int i = end_block * block_size; i <= r; ++i) heights[i] += val;</pre>
 53
               rebuild_block(end_block);
 55
          }
 56
          int query_jumps(int start_idx, int k) {
   int current_hill = start_idx;
 57
 58
 59
 60
                    // Case 1: We don't have enough jumps to take the full precomputed path
 61
                    // We must take a single, precomputed step instead.
if (jumps_to_exit[current_hill] > 0 && k < jumps_to_exit[current_hill]) {</pre>
 62
 63
                         current_hill = next_jump[current_hill];
 64
 65
                         continue;
 66
                    }
                    // Case 2: We have enough jumps to use the precomputed path.
if (jumps_to_exit[current_hill] > 0) {
 69
 70
                         k -= jumps_to_exit[current_hill];
current_hill = block_exit_node[current_hill];
 71
 72
 73
 74
 75
                    // Case 3: We are now at a block's exit point and still have jumps left
                    // We must perform one jump by scanning manually. if (k > 0) {
 76
 77
                         bool jumped = false;
 78
                         11 current_true_height = heights[current_hill] + lazy_add[current_hill /
 79
                               block_size];
 80
                         for (int j = current_hill + 1; j < min(n, current_hill + 1 + MAX_JUMP_DIST); ++</pre>
 81
                               j) {
                              11 next_true_height = heights[j] + lazy_add[j / block_size];
 82
                              if (next_true_height > current_true_height) {
 83
                                   current_hill = j;
 85
                                   jumped = true;
 86
 87
                                   break;
 88
 89
                          if (!jumped) break; // No further jump possible, we are stuck.
 91
                    }
 92
               return current_hill + 1;
 93
 94
 95
           int main() {
                ios_base::sync_with_stdio(false), cin.tie(NULL);
 98
               cin >> n >> q;
block_size = static_cast<int>(sqrt(n));
if (block_size == 0) block_size = 1;
 99
100
101
102
                for (int i = 0; i < n; ++i) {</pre>
103
104
                    cin >> heights[i];
               }
105
106
               for (int i = (n - 1) / block_size; i >= 0; --i) {
107
                    rebuild_block(i);
108
               }
109
110
               while (q--) {
111
                    int type;
cin >> type;
if (type == 1) {
112
113
114
                         int i, k;
                         cin >> i >> k;
116
                         cout << query_jumps(i - 1, k) << " \n";
117
                    } else {
118
                         int 1, r;
119
                         11 x;
120
                         cin >> 1 >> r >> x;
update_range(1 - 1, r - 1, x);
121
122
123
                    }
124
               }
               return 0:
125
126
```