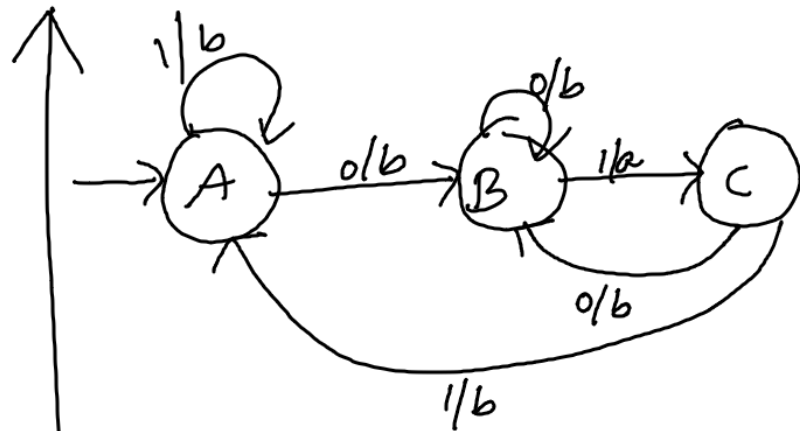


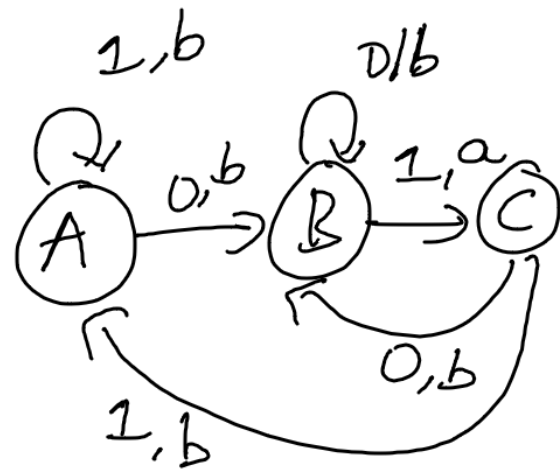
state	0	1	output
→ A	B	A	<u>b</u>
B	B	C	<u>b</u>
C	B	A	a

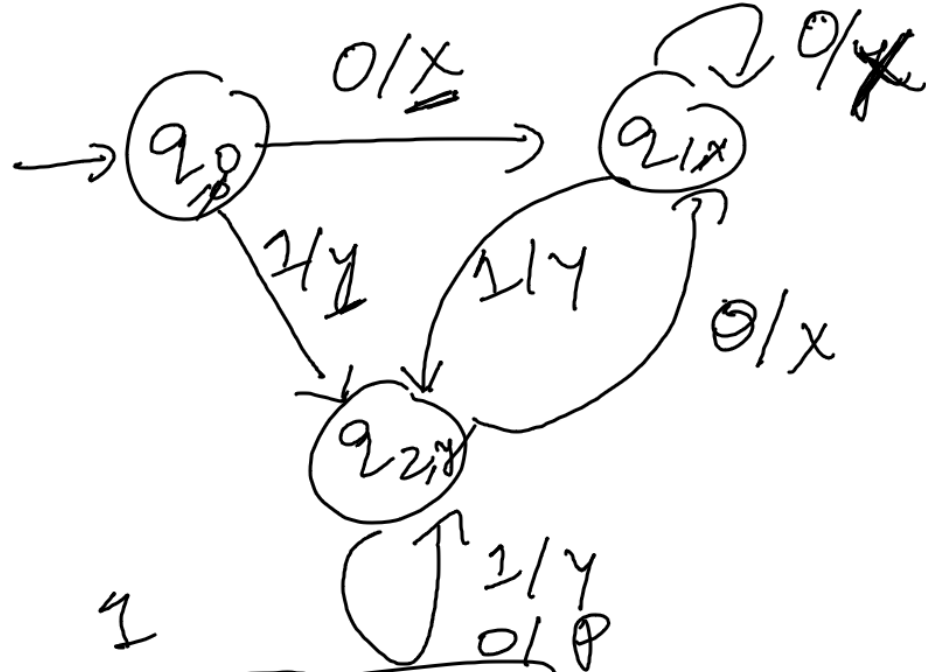
Moorey machine



state	0	1
→ A	B, <u>b</u>	A, b
B	B, <u>b</u>	C, a
C	B, <u>b</u>	A, b

Meelay machine



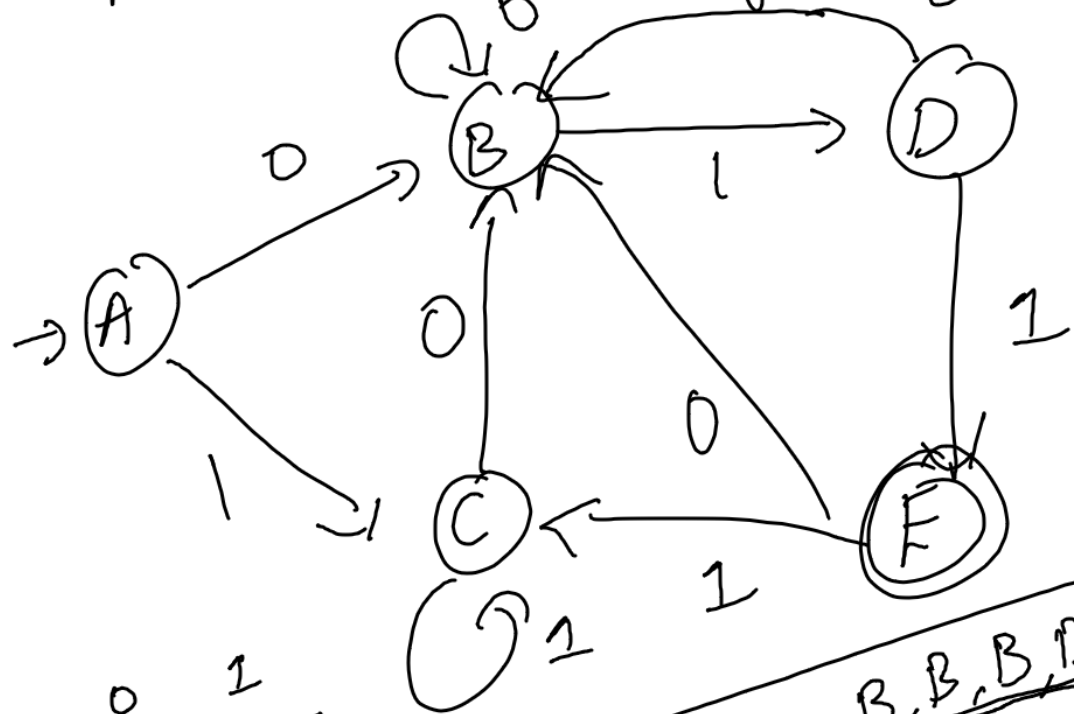


	0	1	
q <sub>0</sub>	q <sub>1</sub>	q <sub>2</sub>	∅
q <sub>1</sub>	q <sub>1</sub>	q <sub>2</sub>	X
q <sub>2</sub>	q <sub>1</sub>	q <sub>2</sub>	Y

Moore Mach

	<del>0</del>	1	O/P
q <sub>0</sub>	q <sub>1</sub> , X	q <sub>2</sub> , Y	∅
q <sub>1</sub>	q <sub>1</sub> , <del>X</del>	q <sub>2</sub> , Y	X
q <sub>2</sub>	q <sub>1</sub> , X	q <sub>2</sub> , Y	Y

# Minimization of DFA



	0	1
A	B	C
B	B	D
C	B	E
D	B	E
E	B	C

$A \rightarrow \underline{B, B, B, B}$   
 $A \rightarrow \underline{C, D, C, E}$

$A, C \rightarrow B$   
 $A, C \rightarrow C$   
 $A, B, C \rightarrow B, B, B$   
 $A, B, C \rightarrow C, D, C$

0-Equivalent

$\{A, B, C, D\}, \{E\}$

1-Equivalent

$\sim \{A, B, C, D\}, \{E\}$

$\approx \boxed{\{A, B, C\} \{D\} \{E\}}$

2-Equivalent

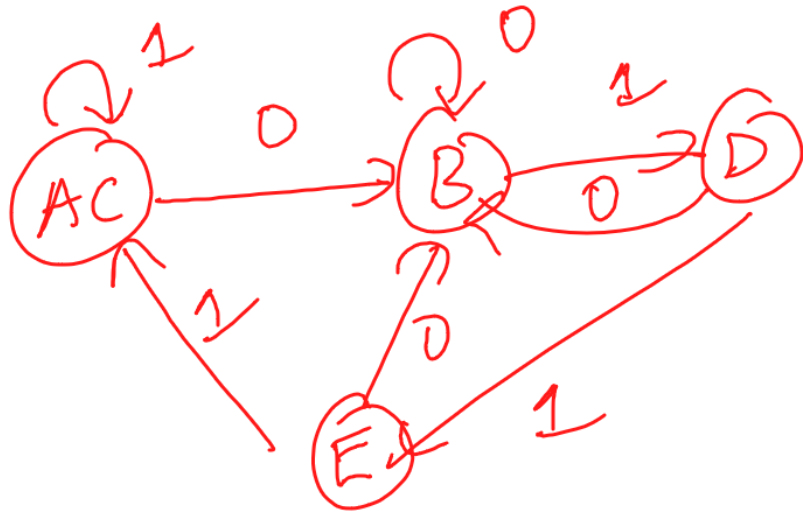
$\sim \{A, B, C\} \{D\} \{E\}$

$\approx \{A, C\} \underline{\{B\}} \underline{\{D\}} \{E\}$

3-Equivalent

$= \underline{\{A, C\}} \underline{\{B\}} \underline{\{D\}} \{E\}$

$\{AC, B, D, E\}$



	0	1
A	B	F
B	G	C
C	A	C
D	C	G
E	H	F
F	C	G
G	G	E
H	G	C

0-Equivalent  
 $\{A, B, D, E, F, G, H\}, \{C\}$  ✓

1-Equivalent  
 $\{A, B, E, G, H\}, \{D, F\}, \{C\}$  ✓  
 $\{A, E\}, \{B, H\}, \{G\}, \{D, F\}, \{C\}$  ✓

2-Equivalent  
 ✓ ~  $\{A, E\}, \{B, H\}, \{G\}, \{D, F\}, \{C\}$   
 ✓ ~  $\{A, E\}, \{B, H\}, \{G\}, \{D, F\}, \{C\}$

AE, BH, G, DF, C

$\{A, B, D, E, F, G, H\} \setminus \{C\}$

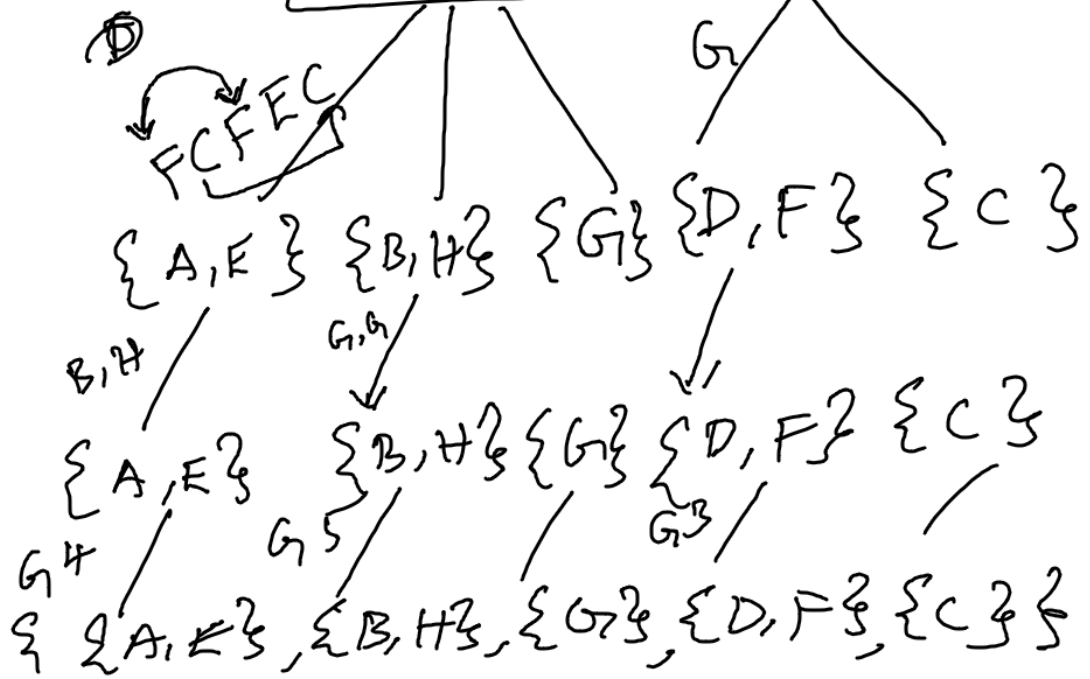
$\{A, B, D, E, F, G, H\}$

0  $\{B, G, C, H, C, G, G\}$

1  $\{\}$

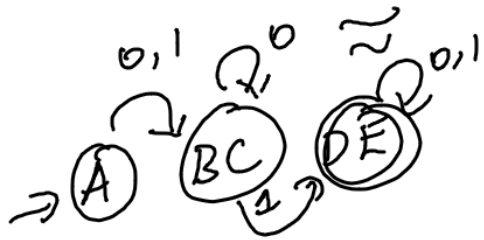
$\{A, B, E, G, H\}$   $\{D, F\}$   $\{C\}$

$A, B, E, G, H \xrightarrow{1}$



	0	1
A	B	C
B	C	D
C	B	E
D	D	D
E	E	E

	0	1
A	BC	BC
BC	BC	DE
DE	DE	DE



0-Equivalent  
 $\{A, B, C\}, \{D, E\}$

1-Equivalent  
 $\{A, B, C\}, \{D, E\}$

$\sim \frac{\{A, B, C\}, \{D, E\}}{\{A\}, \{B, C\}, \{D, E\}}$

2-Equivalent  
 $\sim \{A\}, \{B, C\}, \{D, E\}$   
 $\approx \{A\}, \{B, C\}, \{D, E\}$   
 $\{A, BC, DE\}$

## Pumping Lemma:-

If  $A$  is Regular Language, then  
 $A$  has a pumping length  $p$   
such that any string  $s$  where  
 $|s| \geq p$  may be divided into  
3 parts  $s = xyz$

$$(1) \ x y^i z \in A \text{ for } i \geq 0$$

$$(2) \ |y| > 0$$

$$(3) \ |xy| \leq p$$

$$(2) \ |y| > 0$$

$$(3) \ |xy| \leq 4$$

$$\textcircled{Q} \ A = \{a^n b^n\}$$

Let  $p$  be pumping length

$$A = a^p b^p$$

$$\text{Let } p = 4$$

$$A = a^4 b^4$$

$$\mathcal{S} = \underbrace{\{aaaa\}}_x \underbrace{bbbb}_y z$$

$$x = aa, y = aa$$

$$z = bbbb$$



let  $i = 2$

$x y^2 z$

$x y y z$   
 $\downarrow \quad \downarrow$   
aaaaa bbbb

$\therefore A$  is not Regular

⑥ Apply pumping lemma for  
language  $L = \{a^n / n \text{ is prime}\}$   
and prove it's not regular.

$$L = a^n$$

let  $P$  be pumping length

$$L = a^P$$

let  $P = 5$

$$L = a^5$$

$$L = a$$

$S = \boxed{a} \boxed{aa} \boxed{aaa}$   
 $\quad x \quad y \quad z$

$x = a$   
 $y = aa$   
 $z = aa$

$x y^i z$        $i=2$

$x y y z$

aaaaaa

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$x y y y z$

$i=3$

aaaaaaa

$\therefore L$  is not Regular.

$S = \underline{aaaaa}$

$x = aaa, y = a, z = a$

$x y^i z$        $i=2$

$x y y z$

aaaaa