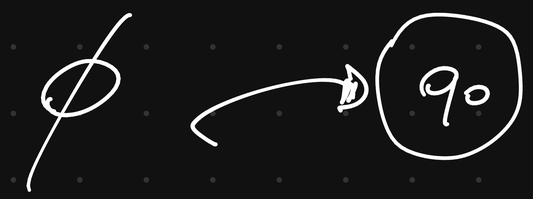
Passage de ER à ε-NFA

$a \in \Sigma$

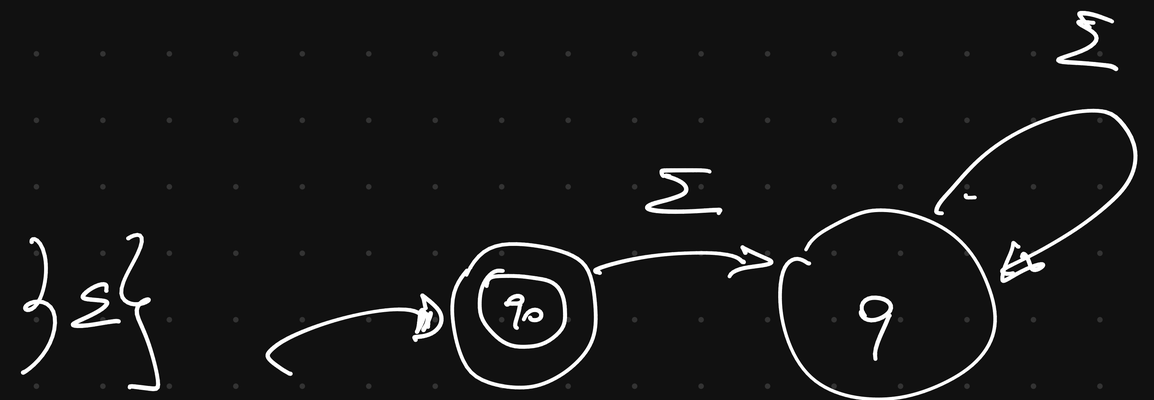


(ER)

$ER_1 \mid ER_2$

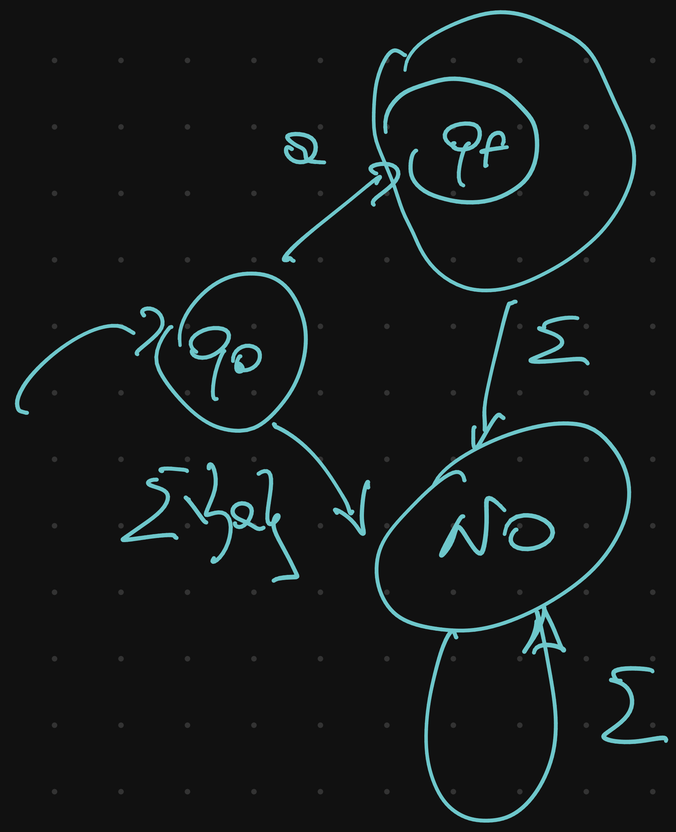
$ER_1 \, ER_2$

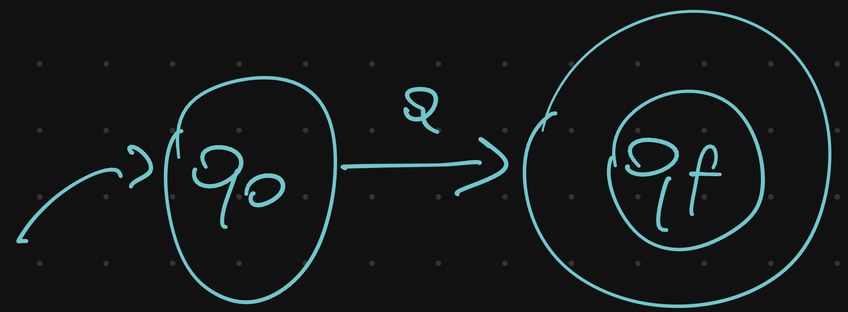
ER^*



$a \in \Sigma$

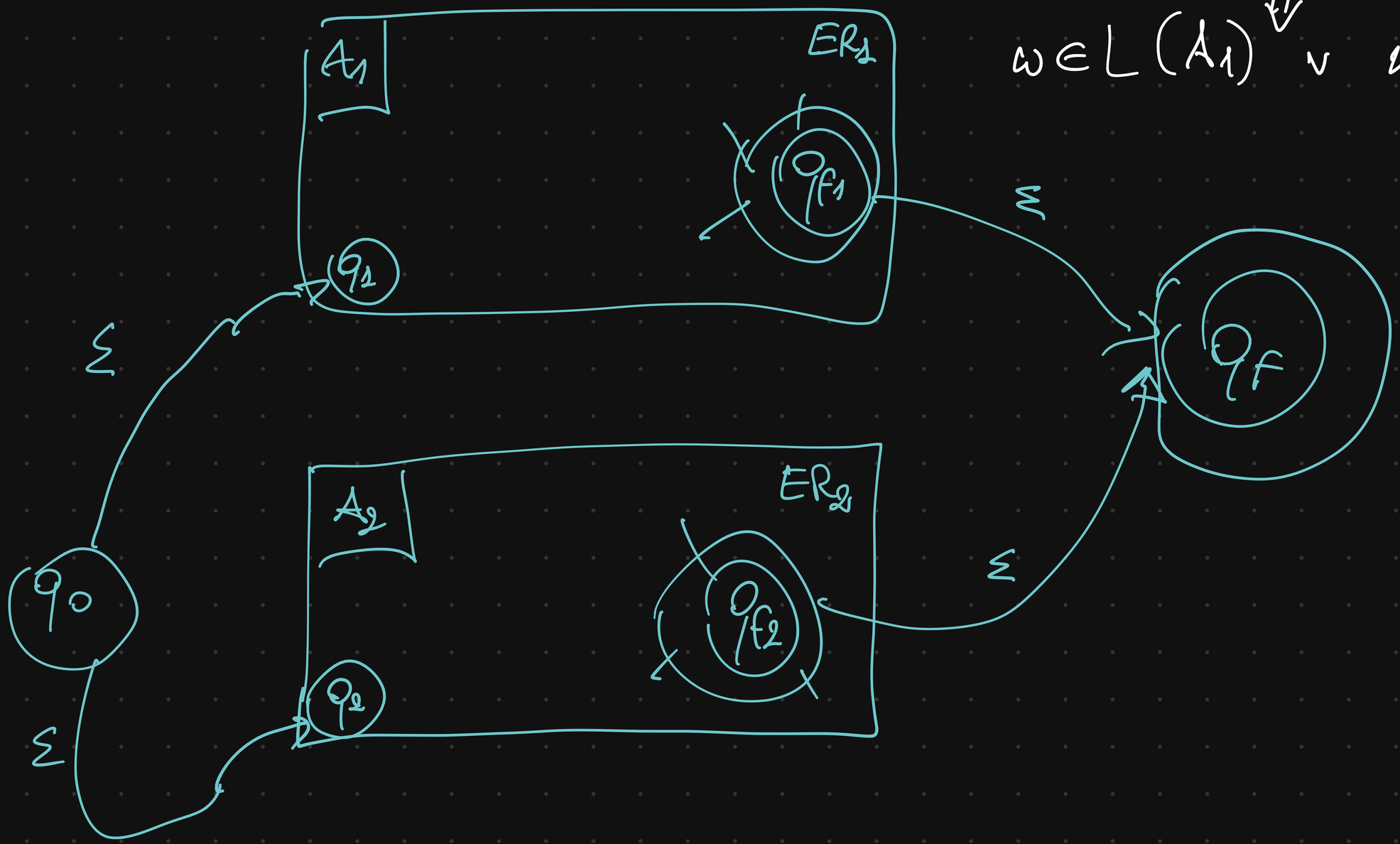
$\{a\}$



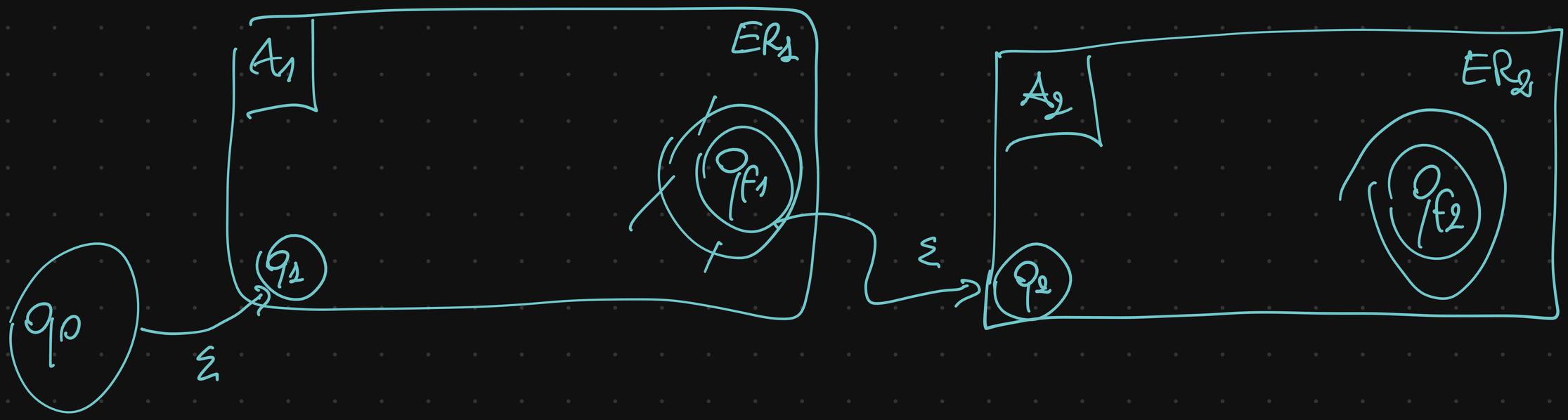


$ER_1 \mid ER_2$

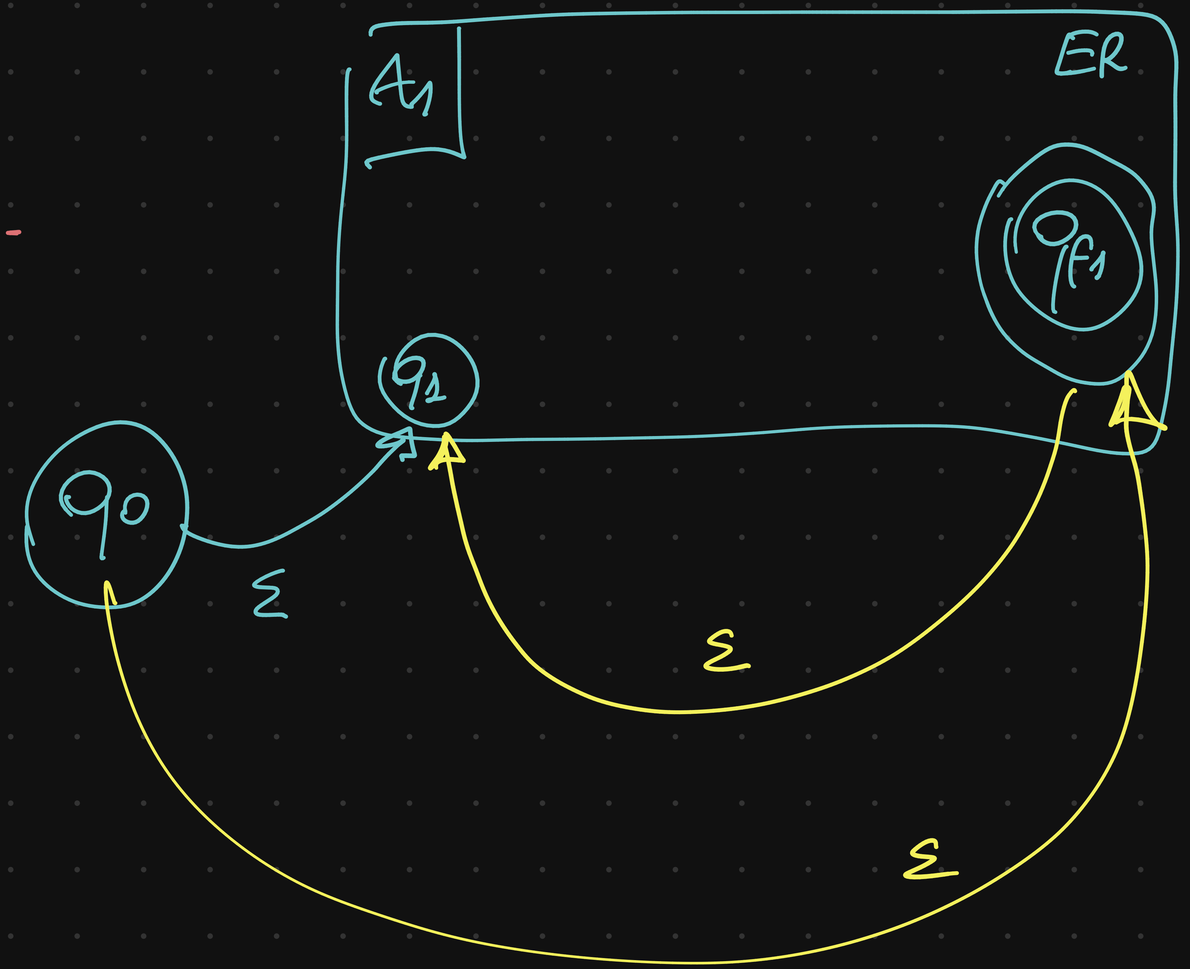
$w \in L(\text{Automa}) \Leftrightarrow w \in L(A_1) \vee w \in L(A_2)$

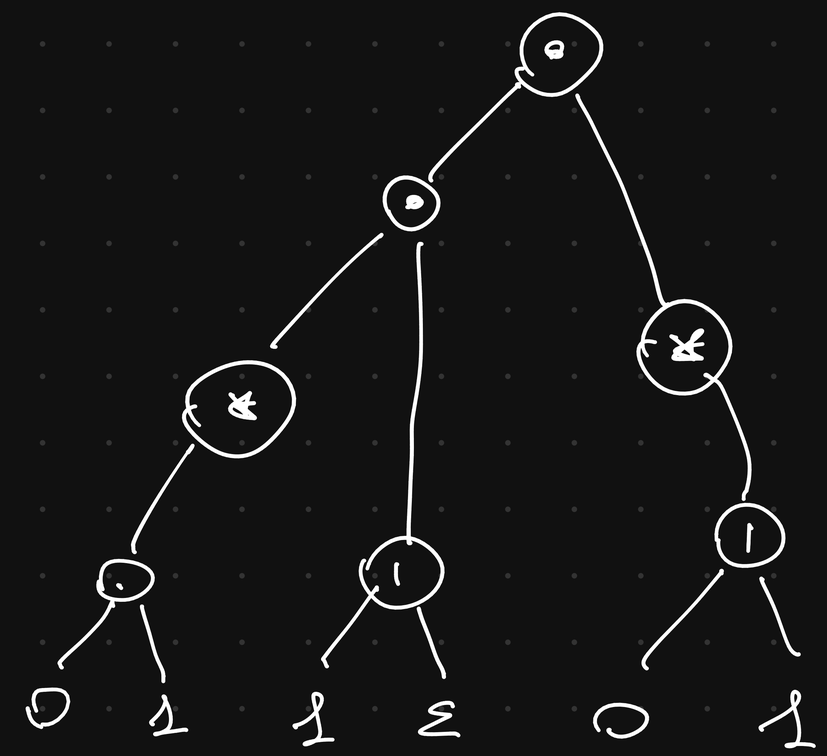


$ER_1 ER_2$



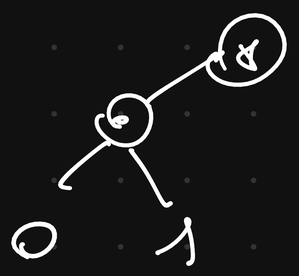
$ER_1^{\star}$

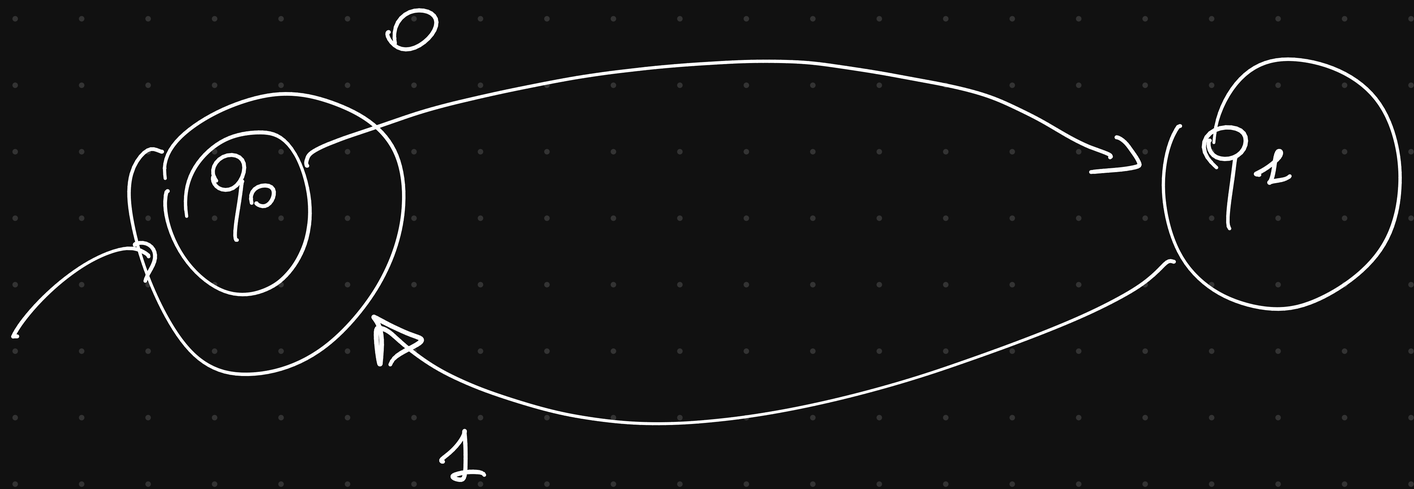




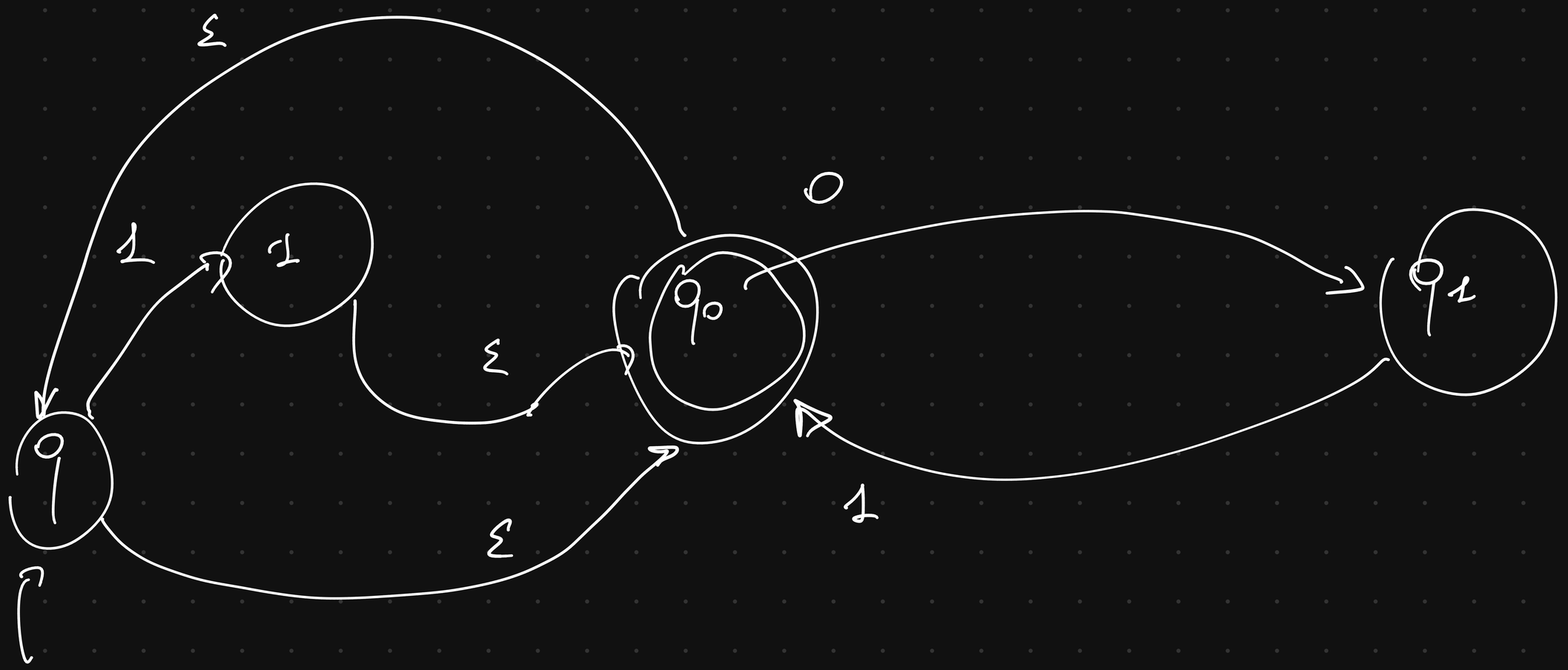
$$(01)^* (1|\varepsilon)(0|1)^*$$

Costruire un ε-NFA che riconosce il linguaggio associato a $(01)^*$

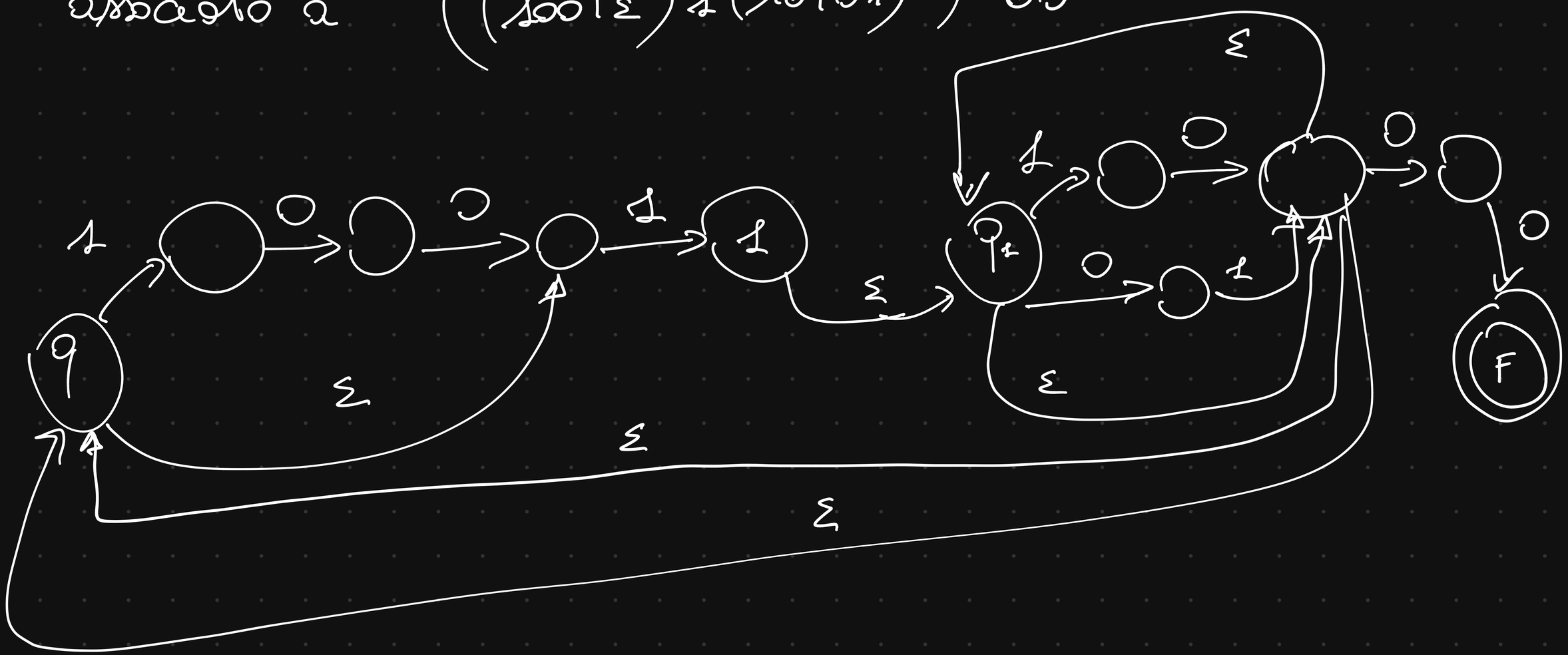




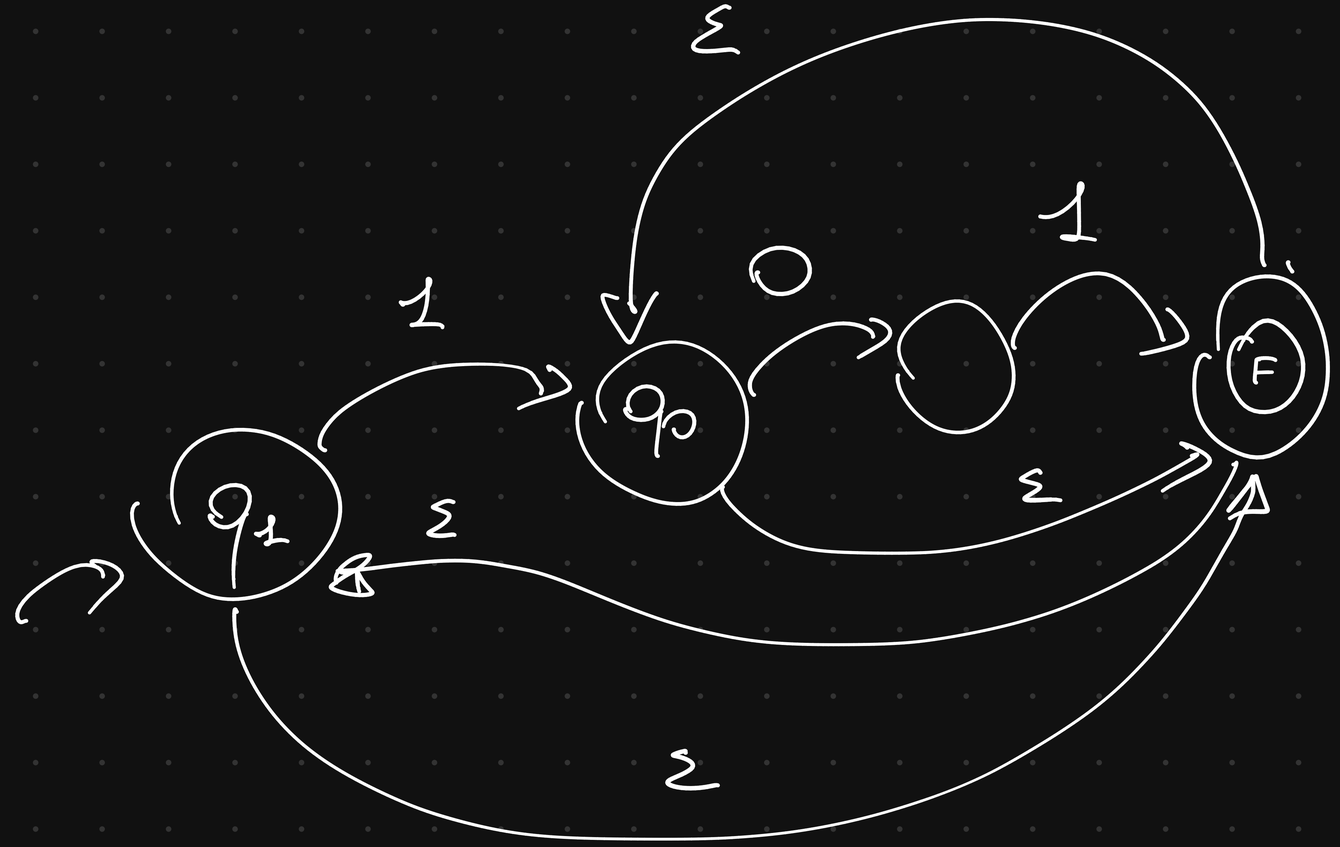
Costruire un ε-NFA che riconosce il linguaggio associato a $((1|\varepsilon)(01)^*)^*$



Costruire un ε-NFA che riconosce il linguaggio associato a $((1001|\varepsilon)1(1|01)^*)^*00$



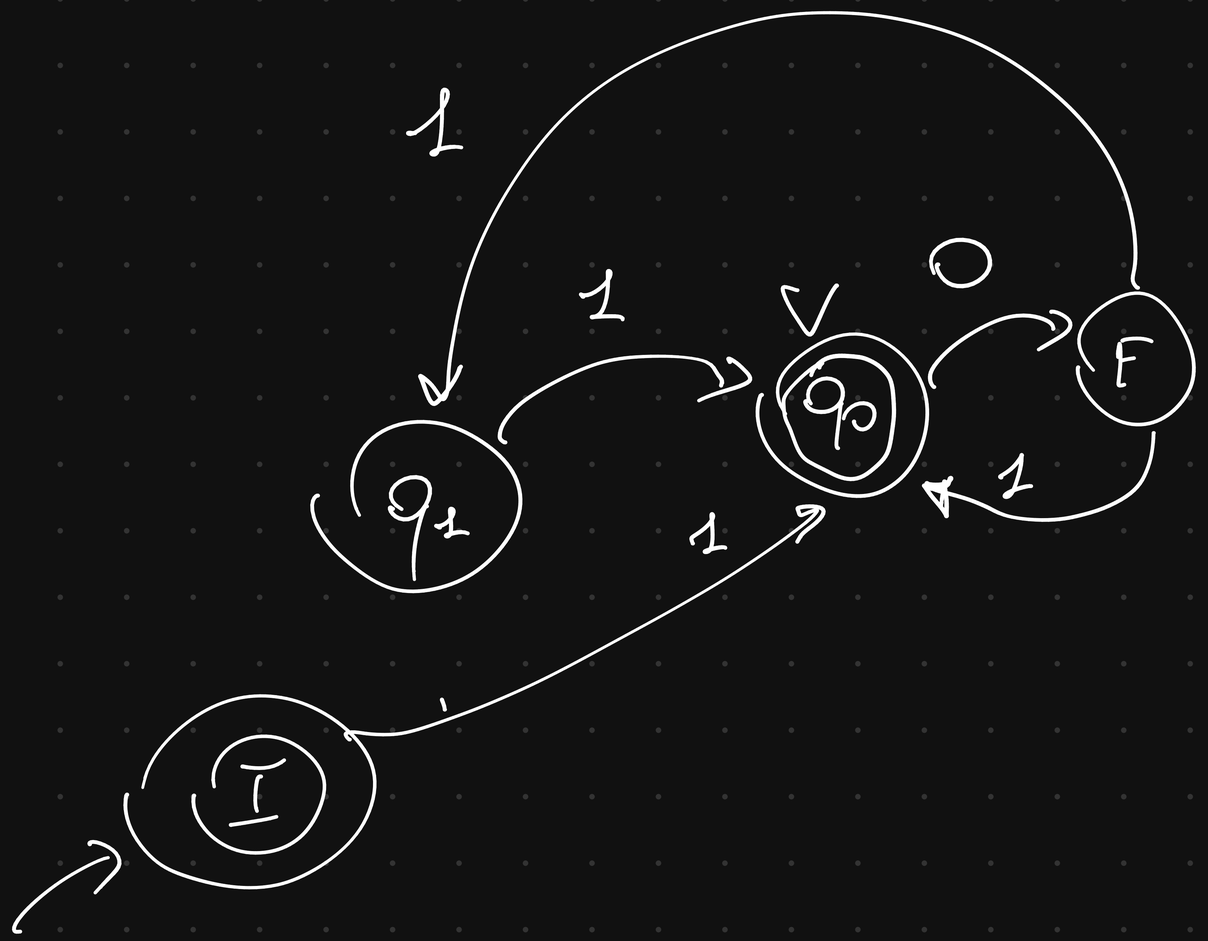
Costruire un DFA che riconosce il linguaggio associato all'espressione regolare $(1(01)^*)^*$

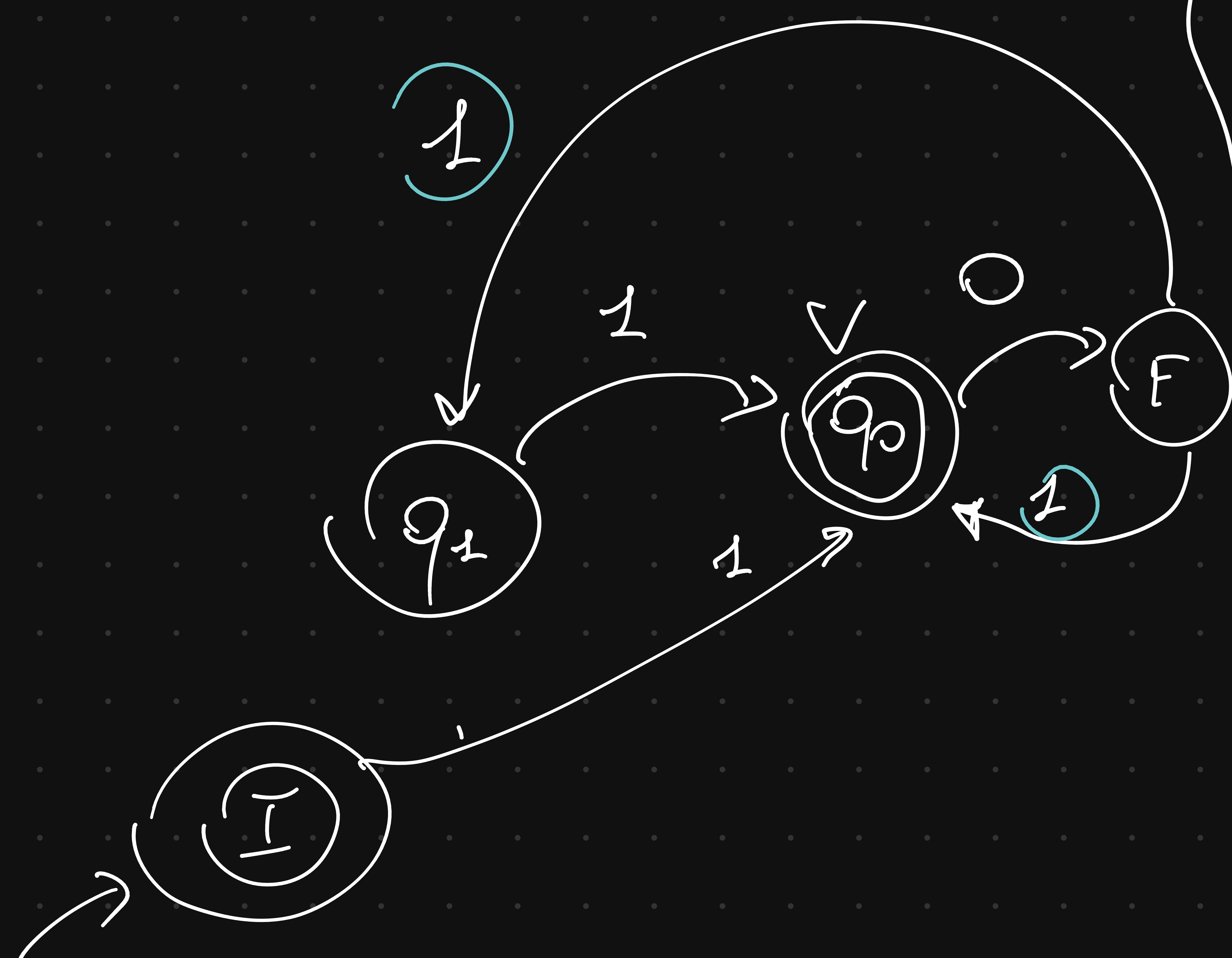


ε- NFA

NFA

$(1(01)^*)^*$




I
q1
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q0
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