

1 Syntax

$e ::= c \in \mathbb{Z}$	$e_e ::= \cdot +_1 e$	$s \in stat ::= skip$	$s_e ::= x :=_1 \cdot$
$ x \in Var$	$ \cdot +_2 \cdot$	$ s_1; s_2$	$ \cdot ;_1 s_2$
$ e_1 + e_2$	$ @_1(e_2)$	$ x := e$	$ if_1 s_1 s_2$
$ \lambda x. s$	$ @_2$	$ if (e > 0) s_1 s_2$	$ while_1 (e > 0) s$
$ e_1 (e_2)$	$ @_3$	$ while (e > 0) s$	$ while_2 (e > 0) s$
		$ return e$	$ return_1 \cdot$

2 Semantics

2.1 Expressions

$\frac{\text{RED-CONST}(c)}{H_e, \ell_e, \ell_c, c \Downarrow H_e, \ell_e, c}$	$\frac{\text{RED-VAR-LOCAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, \ell_c[x]} \quad x \in \text{dom}(H_e[\ell_c])$
$\frac{\text{RED-VAR-GLOBAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, E[x]} \quad x \in \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$	
$\frac{\text{RED-VAR-UNDEF}(x)}{H_e, \ell_e, \ell_c, x \Downarrow err} \quad x \notin \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$	
$\frac{\text{RED-ADD}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, \cdot +_1 e_2 \Downarrow r'}{H_e, \ell_e, \ell_c, e_1 + e_2 \Downarrow r'}$	$\frac{\text{RED-ADD-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad v_1, r, \cdot +_2 \cdot \Downarrow r'}{\ell_c, (H_e, \ell_e, v_1), \cdot +_1 e_2 \Downarrow r'}$
$\frac{\text{RED-ADD-2}}{v_1, (H_e, \ell_e, v_2), \cdot +_2 \cdot \Downarrow H_e, \ell_e, v_1 + v_2}$	$\frac{\text{RED-LAMBDA}(x, s)}{H_e, \ell_e, \ell_c, \lambda x. s \Downarrow H_e, \ell_e, (\ell_c, \lambda x. s)}$
$\frac{\text{RED-APP}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, @_1(e_2) \Downarrow r'}{H_e, \ell_e, \ell_c, e_1(e_2) \Downarrow r'}$	$\frac{\text{RED-APP-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad \ell'_c, x, s, r, @_2 \Downarrow r'}{\ell_c, (H_e, \ell_e, (\ell'_c, \lambda x. s)), @_1(e_2) \Downarrow r'}$
$\frac{\text{RED-APP-2}(s)}{\ell'_c = \text{fresh}(H_e) \quad C = H_e[\ell_c]}{H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c, s \Downarrow r \quad r, @_3 \Downarrow r'}{\ell_c, x, s, (H_e, \ell_e, v), @_2 \Downarrow r'}$	$\frac{\text{RED-APP-3-RET}}{ret(H_e, \ell_e, v), @_3 \Downarrow H_e, \ell_e, v}$
$\frac{\text{RED-APP-3-NO-RET}}{H_e, \ell_e, \ell_c, @_3 \Downarrow err}$	

2.2 Statements

$$\begin{array}{c}
\text{RED-SKIP} \\
\frac{}{H_e, \ell_e, \ell_c, \text{skip} \Downarrow H_e, \ell_e, \ell_c}
\end{array}
\quad
\begin{array}{c}
\text{RED-SEQ}(s_1, s_2) \\
\frac{H_e, \ell_e, \ell_c, s_1 \Downarrow r \quad r, \cdot;_1 s_2 \Downarrow r'}{H_e, \ell_e, \ell_c, s_1; s_2 \Downarrow r'}
\end{array}$$

$$\begin{array}{c}
\text{RED-SEQ-1}(s_2) \\
\frac{H_e, \ell_e, \ell_c, s_2 \Downarrow r}{H_e, \ell_e, \ell_c, \cdot;_1 s_2 \Downarrow r}
\end{array}
\quad
\begin{array}{c}
\text{RED-ASN}(x, e) \\
\frac{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, x :=_1 \cdot \Downarrow r'}{H_e, \ell_e, \ell_c, x := e \Downarrow r'}
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1}(x) \\
\frac{\ell'_e = \text{fresh}(H_e) \quad E = H_e[\ell_e]}{\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_e \leftarrow E[x \leftarrow v]], \ell'_e, \ell_c} \quad x \notin \text{dom}(H_e[\ell_c])
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1-LOCAL}(x) \\
\frac{\ell'_c = \text{fresh}(H_e) \quad C = H_e[\ell_c]}{\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c} \quad x \in \text{dom}(C)
\end{array}$$

$$\begin{array}{c}
\text{RED-IF}(e, s_1, s_2) \\
\frac{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{if}_1 s_1 s_2 \Downarrow r'}{H_e, \ell_e, \ell_c, \text{if}(e > 0) s_1 s_2 \Downarrow r'}
\end{array}
\quad
\begin{array}{c}
\text{RED-IF-1-POS}(s_1, s_2) \\
\frac{H_e, \ell_e, \ell_c, s_1 \Downarrow r}{\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r} \quad v > 0
\end{array}$$

$$\begin{array}{c}
\text{RED-IF-1-NEG}(s_1, s_2) \\
\frac{H_e, \ell_e, \ell_c, s_2 \Downarrow r}{\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r} \quad v \leq 0
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE}(e, s) \\
\frac{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{while}_1(e > 0) s \Downarrow r'}{H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r'}
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE-1-NEG}(e, s) \\
\frac{}{\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow H_e, \ell_e, \ell_c} \quad v \leq 0
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE-1-POS}(e, s) \\
\frac{H_e, \ell_e, \ell_c, s \Downarrow r \quad r, \text{while}_2(e > 0) s \Downarrow r'}{\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow r'} \quad v > 0
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE-2}(e, s) \\
\frac{H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r}{H_e, \ell_e, \ell_c, \text{while}_2(e > 0) s \Downarrow r}
\end{array}
\quad
\begin{array}{c}
\text{RED-RETURN}(e) \\
\frac{H_e, \ell_e, \ell_c, e \Downarrow r \quad r, \text{return}_1 \cdot \Downarrow r'}{H_e, \ell_e, \ell_c, \text{return } e \Downarrow r'}
\end{array}$$

$$\begin{array}{c}
\text{RED-RETURN-1} \\
\frac{}{(\ell_c, (H_e, \ell_e, v), \text{return}_1 \cdot \Downarrow \text{ret}(H_e, \ell_e, v))}
\end{array}$$

2.3 Aborting Rules

$$\begin{array}{c}
\frac{\text{RED-ERROR-EXPR}(e)}{\sigma, e \Downarrow \text{err}} \quad \mathbf{abort} \, \sigma \wedge \neg \mathbf{intercept}_e \, \sigma \qquad \frac{\text{RED-ERROR-STAT}(s)}{\sigma, s \Downarrow \text{err}} \quad \mathbf{abort} \, \sigma \\
\\
\frac{\sigma = C[\text{err}]}{\mathbf{abort} \, \sigma} \qquad \frac{}{\mathbf{intercept}_{@_3} \, \text{ret}(H_e, \ell_e, v)}
\end{array}$$