
A Compiler for the FASTO Language

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1 MULTIPLICATION, DIVISION, BOOLEAN OPERATORS AND LITERALS

Implementing multiplication and division was a simple matter, when having the already implemented code for addition and subtraction to look at. They served as a great way of getting to know the fasto compiler, and how things operate.

Negation is implemented using the Mips.SUB instruction, where we pass the original argument to the operator, and subtracts this from zero.

So the instruction looks like:

```
1 Mips.SUB(place, "0", t1),
```

where t1 register that holds the argument x for $\sim x$, and place is the register in which we place the result.

Not was more complicated than the previous ones, given that this requires more than one instruction to execute. However, the pattern learned here, proved to be useful for implementing and and or as well.

```
1 [ Mips.LI (place, "0")  
  , Mips.BNE (b, "0", falseLabel)  
3   , Mips.LI (place, "1")  
  , Mips.LABEL falseLabel ]  
5
```

Place is the register in which we want to store our result. We start by putting 0 into place, we then check if our argument b, actually is 0. Since Mips.BNE branches if its arguments are not equal, we will jump to falseLabel if and only if our argument b is 1, thus ending with a 0 in the place register, given we never execute Mips.LI(place, "1").

And and Or are almost identical in their implementations. consider the expression $e1 \ \&\& \ e2$. For our code generation we then compile each of the expressions $e1$, $e2$, as conditionals, using compileCond, that lets us perform a jump based on the result of the first conditional. So for and, if $e1$ is false, we immediately jump to the end of the code. While for or, if $e1$ is true, we do not need to evaluate $e2$ since we know the entire or expression is already true. We hereby make use of short-circuiting.

1.1 TESTING

Since most of the tests for these features are simple, we will generalize some of them, and talk about the more interesting ones in particular.

The general idea behind the tests was to perform some operation with the feature to test in particular, and then assert the value with a comparator, like:

$$x \times y == z.$$

The assertions should always be true, so that we can write the result of all the tests with ands. Consider a test with 3 cases, result1, result2 and result3, we would then write the result of the entire test like:

```
2 write(result1 && result2 && result3),
```

which in turn should be true.

1.1.1 NEGATE

For negate we tested that $\sim 0 == 0$, and that $\sim \sim x == x$.

1.1.2 AND AND OR

For these functions, we made sure to test the entire truth table. The test and.fo:

```
1 fun bool main() =  
2   let r1 = (false && true == false) in  
3   let r2 = (true && false == false) in  
4   let r3 = (false && false == false) in  
5   let r4 = (true && true == true) in  
6  
7  
8  
9   write(r1 && r2 && r3 && r4)
```

The rest of the tests can be found in the appendix.

HAR INGEN IDE OM BOOL LITS

2 FILTER AND SCAN

2.1 FILTER

For filter we expect a function with a return type of bool, and some type of array. We make sure this is the case in our typechecker, raise errors if we get a non-array argument or a non-bool function. If everything checks up, we among other things pass the type of the array to the code generator.

Given that filter takes some function and an array as input, and runs this function over the array, we notice that filter and map are similar in ways, so we based the code for filter on the code for map. With the difference being, that we do not want to store some computed value, instead we want to store the original value of the element we might be iterating on, only if the input function returns true.

We achieved this by making changes to the load part, so that we now store the actual value (instead of the computed one), in a register.

Since our input and output arrays will be of different sizes, we created a new counter *j_reg*, which gets incremented only when the input function on an element from the array, returns true.

Between the load and save sequence of map, we made a *check_fun* instruction set;

```
2 val check_fun = [ Mips.BEQ (res_reg, "0", loop_beg)
                   , Mips.ADDI (j_reg, j_reg, "1") ]
```

where *res_reg* is the register in which we place the result of the function call on the input function. We branch to the loop beginning, if the result is false.

2.1.1 TESTING

We used several tests for filter, making sure we use different types, as well as regular and lambda functions.

A test with a regular function, on integers.

```
2 fun bool great(int n) = 5 < n
4 fun int writeInt(int n) = write(n)
6 fun [int] main() =
8   let a = filter(great, {6, 9, 8, 2}) in
    map(writeInt, a)
```

Similar test to the above, but with a lambda function, on integers.

```
1 fun int writeInt(int n) = write(n)
3 fun [int] main() =
5   let a = filter(fn bool (int x) => 5 < x, {6, 9, 8, 2, ~5}) in
    map(writeInt, a)
```

A test with a lambda function on bools.

```
1 fun bool writeBool(bool s) = write(s)
3 fun [bool] main() =
5   let a = {true, false, false, true} in
    let b = filter(fn bool (bool x) => not x, a) in
    map(writeBool, b)
```

All three tests perform as expected.

3 λ -EXPRESSIONS IN SOACS

4 COPY PROPAGATION AND CONSTANT FOLDING

We have implemented *copy propagation and constant foldning* in our compiler, though it does not handle shadowing. Furthermore have all the features from **REF TIL TASK 1**.

Below is the test-case for Times explained. This will indicate that both the cases work and it is able to actually optimize by propagating and fold correspondingly.

When we want to fold expressions, we can do so, by predicting what the result is going to be of a given expression and return this instead.

We thought through for each expression, what base-cases there could be and what we then should return.

Below is shown our code for the *Times*-expression.

```

1 | Times (e1, e2, pos) =>
2 | let val e1' = copyConstPropFoldExp vtable e1
   |   val e2' = copyConstPropFoldExp vtable e2
4 | in case (e1', e2') of
   |   (Constant (IntVal x, _), Constant (IntVal y, _)) =>
6 |     Constant (IntVal (x*y), pos)
   |   (Constant (IntVal 0, _), _) =>
8 |     e1'
   |   (_, Constant (IntVal 0, _)) =>
10 |    e2'
   |   (Constant (IntVal 1, _), _) =>
12 |    e2'
   |   (_, Constant (IntVal 1, _)) =>
14 |    e1'
   |   _ =>
16 |    Times (e1', e2', pos)
   | end

```

We start by optimizing the two subexpressions. Then we check if any of the two expressions evaluates to 0 or 1. If one of them is 0, we return 0. Also, if one of them is 1 then it will be the other expression that is returned, no matter what it states. If none of the above applies, then we want to compute the expression.

Below is the code for testing *Times*.

```
1 fun int main() =  
  let a = 5 in  
3   let b = a in  
    let c = b in  
5     write(b + c * 0)
```

If our propagation is correct, all the variables should evaluate to 5, due to *a* having that constant assigned.

In the *write*-statement we have some expressions which also can be optimized. Since we multiply by 0, that whole expression will evaluate to 0. Then we plus *b* with 0 and end up with *b* as the result, and we end up writing *b*.

In the end, our program will look like this when optimized using copy propagation and constant folding.

```
1 fun int main() =  
  let a = 5 in  
3   let b = 5 in  
    let c = 5 in  
5     write(5)
```

Lexer.lex

```
1 40 | "fn"          => Parser.FN pos
42 | "not"        => Parser.NOT pos
3
66 | "true"|"false" { case Bool.fromString (getLexeme lexbuf) of
5 67 |               NONE  => lexerError lexbuf "Bad bool"
68 |               | SOME b => Parser.BOOLEAN (b, getPos lexbuf) }
7
87 | '*'          { Parser.TIMES (getPos lexbuf) }
9 88 | '/'          { Parser.DIV (getPos lexbuf) }
89 | '~'          { Parser.NEG (getPos lexbuf) }
11 91 | "=>"         { Parser.ARROW (getPos lexbuf) }
94 | "||"          { Parser.OR (getPos lexbuf) }
13 95 | "&&"          { Parser.AND (getPos lexbuf) }
```

Parser.grm

```
1 10 token <(int*int)> IF THEN ELSE LET IN INT BOOL CHAR EOF
11 token <string*(int*int)> ID STRINGLIT
3 12 token <int*(int*int)> NUM
13 token <char*(int*int)> CHARLIT
5 14 token <bool*(int*int)> BOOLEAN
15 token <(int*int)> PLUS MINUS DEQ EQ LTH NEG NOT ARROW
7 16 token <(int*int)> TIMES DIV AND NOT OR LPAR RPAR LBRACKET RBRACKET LCURLY RCURLY
17 token <(int*int)> COMMA
9 18 token <(int*int)> FUN FN IOTA REPLICATE MAP REDUCE FILTER SCAN READ WRITE
19 token <(int*int)> OP
11
21 nonassoc ifprec letprec
13 22 left DEQ LTH
23 left AND
15 24 left OR
25 nonassoc NOT
17 26 left PLUS MINUS
27 left TIMES DIV
19 28 nonassoc NEG
```


TypeChecker.sml

```

255 | In.Times (e1, e2, pos)
256   => let val (_, e1_dec, e2_dec) = checkBinOp ftab vtab (pos, Int, e1, e2)
257       in (Int,
4 258         Out.Times(e1_dec, e2_dec, pos))
259       end
6 260
261 | In.Divide (e1, e2, pos)
8 262   => let val (_, e1_dec, e2_dec) = checkBinOp ftab vtab (pos, Int, e1, e2)
263       in (Int,
10 264         Out.Divide(e1_dec, e2_dec, pos))
265       end
12 266
267 | In.Negate (e1, pos)
14 268   => let val (t1, e1') = checkExp ftab vtab e1
269       in (Int, Out.Negate(e1', pos))
16 270       end
271
18 272 | In.Not (e1, pos)
273   => let val (t1, e1') = checkExp ftab vtab e1
20 274       in (Bool, Out.Not(e1', pos))
275       end
22 276
277 | In.Or (e1, e2, pos)
278   => let val (t1, e1') = checkExp ftab vtab e1
279       val (t2, e2') = checkExp ftab vtab e2
26 280       in
281         if (t1 = Bool andalso t2 = Bool) then
28 282           (Bool, Out.Or(e1', e2', pos))
283         else
30 284           raise Error("Type Error: Non-boolean arguments given to ||", pos)
285       end
32 286
287 | In.And (e1, e2, pos)
34 288   => let val (t1, e1') = checkExp ftab vtab e1
289       val (t2, e2') = checkExp ftab vtab e2
36 290       in
291         if (t1 = Bool andalso t2 = Bool) then
38 292           (Bool, Out.And(e1', e2', pos))
293         else
40 294           raise Error("Type Error: Non-boolean arguments given to &&", pos)
295       end
42 296
298 | In.Scan (f, n_exp, arr_exp, _, pos) 256
44 299   => let val (n_type, n_dec) = checkExp ftab vtab n_exp
300       val (arr_type, arr_dec) = checkExp ftab vtab arr_exp
46 301       val elem_type =
302         case arr_type of
48 303         Array t => t
304         | other => raise Error ("Scan: argument is not an array", pos)
50 305       val (f', f_arg_type) =
306         case checkFunArg (f, vtab, ftab, pos) of
52 307         (f', res, [a1, a2]) =>

```

```

308         if a1 = a2 andalso a2 = res
54 309         then (f', res)
310         else raise Error
56 311             ("Scan: incompatible function type of "
312              ^ In.ppFunArg 0 f ^ ": " ^ showFunType ([a1, a2], res),
pos)
58 313         | (_, res, args) =>
314             raise Error ("Scan: incompatible function type of "
60 315              ^ In.ppFunArg 0 f ^ ": " ^ showFunType (args, res),
pos)
316         fun err (s, t) =
62 317             Error ("Scan: unexpected " ^ s ^ " type " ^ ppType t ^
318                  ", expected " ^ ppType f_arg_type, pos)
64 319     in if elem_type = f_arg_type
320         then if elem_type = n_type
66 321             then (Array elem_type,
322                  Out.Scan (f', n_dec, arr_dec, elem_type, pos))
68 323             else raise (err ("neutral element", n_type))
324         else raise err ("array element", elem_type)
70 325     end

72 329 | In.Filter (f, arr_exp, _, pos) 260
330     => let val (arr_type, arr_exp_dec) = checkExp ftab vtab arr_exp
74 331
332         val elem_type =
76 333         case arr_type of
334             Array t => t
78 335         | other => raise Error ("Filter: argument is not an array", pos)
336
337         val (f', f_res_type, f_arg_type) =
338         case checkFunArg (f, vtab, ftab, pos) of
82 339             (f', Bool, [a1]) => (f', Bool, a1)
340         | (_, res, args) =>
84 341             raise Error ("Filter: incompatible function type of "
342              ^ In.ppFunArg 0 f ^ ": " ^ showFunType (args, res),
pos)
86 343
344         in (arr_type,
88 345             Out.Filter (f', arr_exp_dec, elem_type, pos))
346         end

90 362 | checkFunArg (In.Lambda (ret_type, params, body, funpos) 276
92 363     , vtab, ftab, pos) =
364     let val (Out.FunDec (name, _, params, body', pos)) = checkFunWithVtable (In.
FunDec ("Lambda", ret_type, params, body, pos), vtab, ftab, funpos)
94 365     val arg_types = map (fn (Param (_, ty)) => ty) params
366     in
96 367     (Out.Lambda(ret_type, params, body', funpos), ret_type, arg_types)
368     end

```

CodeGen.sml

```

1 614 | Constant (BoolVal b, pos) =>
615     if(b) then
3 616         [Mips.LI(place, "1")]
617     else
5 618         [Mips.LI(place, "0")]

7 624 | Times (e1, e2, pos) =>
625     let val t1 = newName "times_L"
9 626         val t2 = newName "times_R"
627         val code1 = compileExp e1 vtable t1
11 628         val code2 = compileExp e2 vtable t2
629     in code1 @ code2 @ [Mips.MUL (place, t1, t2)]
630     end
13 631 | Divide (e1, e2, pos) =>
632     let val t1 = newName "div_L"
15 633         val t2 = newName "div_R"
634         val code1 = compileExp e1 vtable t1
17 635         val code2 = compileExp e2 vtable t2
636     in code1 @ code2 @ [Mips.DIV (place, t1, t2)]
637     end
21 638
639 | Negate (e1, pos) =>
23 640     let val t1 = newName "negateVal"
641         val code1 = compileExp e1 vtable t1
25 642     in code1 @ [ Mips.SUB(place, "0", t1) ]
643     end
27
645 | Not (b_exp, pos) =>
29 646     let val b = "boolean"
647         val code1 = compileExp b_exp vtable b
31 648         val falseLabel = newName "false"
649     in code1 @
33 650         [ Mips.LI (place, "0")
651           , Mips.BNE (b, "0", falseLabel)
35 652           , Mips.LI (place, "1")
653           , Mips.LABEL falseLabel ]
37 654     end
655
39 656 | Or (e1, e2, pos) =>
657     let val trueLabel = newName "trueLabel"
41 658         val falseLabel = newName "falseLabel"
659         val endLabel = newName "endLabel"
43 660         val code1 = compileCond e1 vtable trueLabel falseLabel
661         val code2 = compileCond e2 vtable trueLabel endLabel
45 662     in [Mips.LI(place, "0")] @
663         code1 @
47 664         [Mips.LABEL falseLabel] @
665         code2 @
49 666         [Mips.LABEL trueLabel, Mips.LI(place, "1"), Mips.LABEL endLabel]
667     end
51 668
669 | And (e1, e2, pos) =>

```

```

53 670      let val trueLabel = newName "trueLabel"
671          val falseLabel = newName "falseLabel"
55 672          val endLabel = newName "endLabel"
673          val code1 = compileCond e1 vtable trueLabel falseLabel
57 674          val code2 = compileCond e2 vtable endLabel falseLabel
675      in [Mips.ADD(place, "0", "1")] @
59 676          code1 @
677          [Mips.LABEL trueLabel] @
61 678          code2 @
679          [Mips.LABEL falseLabel, Mips.ADD(place, "0", "0"), Mips.LABEL endLabel]
63 680      end

65 688      (* Scan(f, e, [a1, a2, ..., an]) = [e, f(e, a1), f(f(e, a1), a2), ...] *)
689      | Scan (farg, acc_exp, arr_exp, elem_type, pos) =>
67 690          let val in_size_reg = newName "in_size_reg" (* size of input array *)
691              val out_size_reg = newName "out_size_reg" (* size of output array *)
69 692              val acc_reg = newName "acc_reg" (* last computed value for output *)
693              val i_reg = newName "i_reg" (* Iterator register *)
71 694              val addr_reg = newName "addr_reg" (* address of element in output array
*)
695              val arr_reg = newName "arr_reg"
73 696              val elem_reg = newName "elem_reg"
697
75 698              val arr_code = compileExp arr_exp vtable arr_reg
699              val acc_code = compileExp acc_exp vtable acc_reg
77 700
701              val get_size = [ Mips.LW (in_size_reg, arr_reg, "0") (* Loads array-size
into in_size_reg *)
79 702                          , Mips.ADDI(out_size_reg, in_size_reg, "1") (* Puts size
into out_size_reg *)
703                          ]
81 704
705              (* Initiate registers.
83 706              Put address of place into addr_reg, so we return proper addresses.
707              Increment in_size_reg by 1 to determine out_size_reg, since output
array is
85 708              1 element longer. *)
709              val init_regs = [ Mips.ADDI (addr_reg, place, "4") (* point to the next
word *)
87 710                          , Mips.SW (acc_reg, addr_reg, "0")
711                          , Mips.ADDI (addr_reg, addr_reg, "4") (* point to next
word*)
89 712                          , Mips.MOVE(i_reg, "0") (* initialize iterator*)
713                          , Mips.ADDI (arr_reg, arr_reg, "4") (* Look at first
element in input array*)
91 714                          ]
715
93 716
717              val loop_beg = newName "loop_beg"
95 718              val loop_end = newName "loop_end"
719              val tmp_reg = newName "tmp_reg"
97 720
721              (* while i_reg < in_size_reg*)
99 722              val loop_head =

```

```

723         [ Mips.LABEL(loop_beg)
101 724         , Mips.SUB(tmp_reg, i_reg, out_size_reg) (* make statement*)
725         , Mips.BGEZ(tmp_reg, loop_end) ] (* if statement is equal to zero,
jump to end *)
103 726         (* loads next value in input array into tmp_reg *)
727         val load_value =
105 728         case getElemSize elem_type of
729             One => [ Mips.LB (elem_reg, arr_reg, "0")
107 730                     , Mips.ADDI (arr_reg, arr_reg, "1") ]
731             | Four => [ Mips.LW (elem_reg, arr_reg, "0")
109 732                        , Mips.ADDI (arr_reg, arr_reg, "4") ]
733
111 734         val apply_code =
735             applyFunArg(farg, [acc_reg, elem_reg], vtable, acc_reg, pos)
113 736
737         (* save the current accumulated value to a register for later
computations,
115 738         and to memory for later output *)
739         val save_value =
117 740         case getElemSize elem_type of
741             One => [ Mips.SB (acc_reg, addr_reg, "0")
119 742                     , Mips.ADDI (addr_reg, addr_reg, "1") ]
743             | Four => [ Mips.SW (acc_reg, addr_reg, "0")
121 744                        , Mips.ADDI (addr_reg, addr_reg, "4") ]
745
123 746
747
125 748         (* increments i_reg *)
749         val loop_foot =
127 750         [ Mips.ADDI(i_reg, i_reg, "1")
751           , Mips.J loop_beg
129 752           , Mips.LABEL loop_end ]
753
131 754     in [Mips.LABEL "Det_her_er_starten"]
755         @ arr_code
133 756         @[Mips.LABEL "array_kode"]
757         @ acc_code
135 758         @[Mips.LABEL "akku_kode"]
759         @ get_size
137 760         @[Mips.LABEL "Stoerrelsen_på_dyret"]
761         @ dynalloc (out_size_reg, place, elem_type)
139 762         @[Mips.LABEL "init_regz"]
763         @ init_regs
141 764         @[Mips.LABEL "starten_af_loopet"]
765         @ loop_head
143 766         @[Mips.LABEL "midten_af_loopet"]
767         @ load_value
145 768         @[Mips.LABEL "mere_midte"]
769         @ apply_code
147 770         @[Mips.LABEL "gem_den_akku"]
771         @ save_value
149 772         @[Mips.LABEL "foden_af_loopet"]
773         @ loop_foot
151 774         @ [Mips.LABEL "Det_her_er_slutningen_SCAN"]

```

```

775         end
153
783         (* filter(f(), acc, [a1,a2]) = [f(acc, a1), f(acc, a2)] *)
155 784     | Filter (farg, arr_exp, elem_type, pos) =>
785         let val size_reg = newName "size_reg" (* size of input/output array *)
157 786             val arr_reg = newName "arr_reg" (* address of input array *)
787             val elem_reg = newName "elem_reg" (* address of single element *)
159 788             val res_reg = newName "res_reg"
789             val val_reg = newName "val_reg"
161 790             val arr_code = compileExp arr_exp vtable arr_reg
791
163 792             val get_size = [ Mips.LW (size_reg, arr_reg, "0") ] (* *)
793
165 794             val addr_reg = newName "addr_reg" (* address of element in new array *)
795             val i_reg = newName "i_reg"
167 796             val j_reg = newName "j_reg"
797             val init_regs = [ Mips.ADDI (addr_reg, place, "4") (*point to the next
word, so we don't overwrite the size*)
169 798                             , Mips.MOVE (i_reg, "0")
799                             , Mips.LI (j_reg, "0")
171 800                             , Mips.ADDI (elem_reg, arr_reg, "4") ]
801
173 802             val loop_beg = newName "loop_beg"
803             val loop_end = newName "loop_end"
175 804             val tmp_reg = newName "tmp_reg"
805             val loop_header = [ Mips.LABEL (loop_beg)
177 806                             , Mips.SUB (tmp_reg, i_reg, size_reg)
807                             , Mips.BGEZ (tmp_reg, loop_end)
179 808                             , Mips.ADDI (i_reg, i_reg, "1") ]
809
181 810             (* map is 'arr[i] = f(old_arr[i])'. *)
811             val loop_load =
183 812                 case getElemSize elem_type of
813                     One => Mips.LB(val_reg, elem_reg, "0")
185 814                     :: applyFunArg(farg, [val_reg], vtable, res_reg, pos)
815                     @ [ Mips.ADDI(elem_reg, elem_reg, "1") ]
187 816                 | Four => Mips.LW(val_reg, elem_reg, "0")
817                     :: applyFunArg(farg, [val_reg], vtable, res_reg, pos)
189 818                     @ [ Mips.ADDI(elem_reg, elem_reg, "4") ]
819
191 820             val check_fun = [ Mips.BEQ (res_reg, "0", loop_beg)
821                             , Mips.ADDI (j_reg, j_reg, "1") ]
193 822
823             val loop_save =
195 824                 case getElemSize elem_type of
825                     One => [ Mips.SB (val_reg, addr_reg, "0") ]
197 826                 | Four => [ Mips.SW (val_reg, addr_reg, "0") ]
827
199 828             val loop_footer =
829                 [ Mips.ADDI (addr_reg, addr_reg,
201 830                     makeConst (elemSizeToInt (getElemSize elem_type)))
831                 , Mips.J loop_beg
832                 , Mips.LABEL loop_end
203 833                 , Mips.SW (j_reg, place, "0") ]

```

```

205 834         ]
835         in [Mips.LABEL "array_kode"]
207 836         @ arr_code (* make arr_reg point to input array*)
837         @[Mips.LABEL "measuring"]
209 838         @ get_size (* gets the size of the input array *)
839         @[Mips.LABEL "allocating_da_mem"]
211 840         @ dynalloc (size_reg, place, elem_type) (* return place, which is an
address where the ouput array is going to be *)
841         @[Mips.LABEL "initializing"]
213 842         @ init_regs
843         @[Mips.LABEL "loopin"]
215 844         @ loop_header
845         @ loop_load
217 846         @ check_fun
847         @ [Mips.LABEL "saving"]
219 848         @ loop_save
849         @ loop_footer
221 850     end

223 868     | applyFunArg (Lambda(ret_type, params, body', funpos), args, vtable, place,
pos) : Mips.Prog =
868     let
225 869         fun bindVars ([], [], vtable) = vtable
870         | bindVars([], args, vtable) = raise Error("stop det pjat", pos)
227 871         | bindVars(params, [], vtable) = raise Error("stop det pjat stadigvæk",
pos)
872         | bindVars(Param (name, paramtype)::params, arg::args, vtable) = SymTab.
bind name arg (bindVars(params, args, vtable))
229 873         val newVtable = bindVars(params, args, vtable)
874         val code1 = compileExp body' newVtable place
231 875     in
876         code1
233 877     end

```

Interpreter.sml

```

1 412 | evalExp ( Times(e1, e2, pos), vtab, ftab ) =
413     let val res1  = evalExp(e1, vtab, ftab)
3 414       val res2  = evalExp(e2, vtab, ftab)
415     in evalBinopNum(op *, res1, res2, pos)
5 416     end
417
7 418 | evalExp ( Divide(e1, e2, pos), vtab, ftab ) =
419     let val res1  = evalExp(e1, vtab, ftab)
9 420       val res2  = evalExp(e2, vtab, ftab)
421     in evalBinopNum(op div, res1, res2, pos)
11 422     end
423
13 424 | evalExp ( Negate(e1, pos), vtab, ftab ) =
425     let val res1 = evalExp(e1, vtab, ftab)
15 426       val res2 = evalExp(Constant(IntVal 0, pos), vtab, ftab)
427     in evalBinopNum(op +, res1, res2, pos)
17 428     end
429
19 430 | evalExp ( Not(e1, pos), vtab, ftab ) =
431     let val res1 = evalExp(e1, vtab, ftab)
21 432     in case res1 of
433         BoolVal true => BoolVal false
23 434       | BoolVal false => BoolVal true
435       | _ => raise Error("Input is not of type bool", pos)
25 436     end
437
27 438 | evalExp ( And(e1, e2, pos), vtab, ftab ) =
439     let val res1 = evalExp(e1, vtab, ftab)
29 440       val res2 = evalExp(e2, vtab, ftab)
441     in case res1 of
31 442         BoolVal false => res1
443       | BoolVal true => (case res2 of
33 444           BoolVal _ => res2
445           | _ => raise Error("Second argument is not of type
bool", pos))
35 446       | _ => raise Error("First argument is not of type bool", pos)
447
37 448     end
449
39 450 | evalExp ( Or(e1, e2, pos), vtab, ftab ) =
451     let val res1 = evalExp(e1, vtab, ftab)
41 452       val res2 = evalExp(e2, vtab, ftab)
453     in case res1 of
43 454         BoolVal true => res1
455       | BoolVal false => (case res2 of
45 456           BoolVal _ => res2
457           | _ => raise Error("Second argument is not of type
bool", pos))
47 458       | _ => raise Error("First argument is not of type bool", pos)
459
49 460     end

```



```
51 517 | evalFunArg (Lambda(ret_type, params, exp, pos), vtab, ftab, callpos) =  
518 |   let  
53 519 |     val fundec = FunDec ("Lambda", ret_type, params, exp, pos)  
520 |   in  
55 521 |     (fn aargs => callFunWithVtable(fundec, aargs, vtab, ftab, callpos), ret_type)  
522 |   end
```

CopyConstPropFold.sml

```

16 fun copyConstPropFoldExp vtable e =
2 17   case e of
18     Constant x => Constant x
4 19   | StringLit x => StringLit x
20   | ArrayLit (es, t, pos) =>
6 21     ArrayLit (map (copyConstPropFoldExp vtable) es, t, pos)
22   | Var (name, pos) =>
8 23     (* TODO TASK 4: This case currently does nothing.
24
10 25     You must perform a lookup in the symbol table and if you find
26     a Propagatee, return either a new Var or Constant node. *)
12 27     let
28       val value = SymTab.lookup name vtable
14 29     in
30       case value of
16 31         SOME (VarProp v)  => copyConstPropFoldExp vtable (Var(v, pos))
32       | SOME (ConstProp v) => copyConstPropFoldExp vtable (Constant(v, pos))
18 33       | NONE              => Var(name, pos)
34     end
20
102 let val e' = copyConstPropFoldExp vtable e
22 103   in
104     let
24 105       val vtable' =
106         case e' of
26 107           (Var (vname, p))    => (SymTab.bind name (VarProp(vname)) vtable)
108         | (Constant (vval, p)) => (SymTab.bind name (ConstProp(vval)) vtable)
28 109         | _ => vtable
110     in
30 111       Let (Dec (name, e', decpos), copyConstPropFoldExp vtable' body, pos)
112     end
32 113   end

```