Course Load in a Nutshell

Term Schedule

Dates are a guideline and are subject to change.

Assignment/Exam	Weight	Due Date	
Tutorials	10%	1% each	
Assignment 1	8%	January 31	
Term Test 1	12%	February 3	
Assignment 2	15%	February 28	
Term Test 2	12%	March 2	
Assignment 3	18%	April 3	
Final Exam	25%	To be scheduled	

Text: readings and forum questions (will be part of assignments!) Silberschatz, Galvin, and Gagne 7th edition or later...

2020-01-05

Tutorials

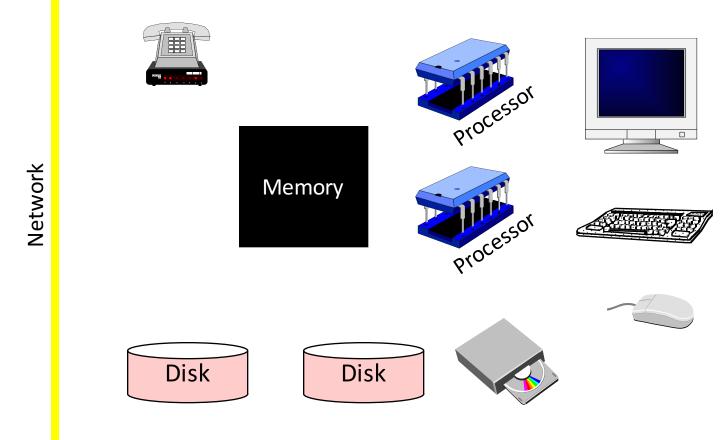
YES!!!! THEY ARE WORTH GOING TO!!! 1% of your mark each, good practice for tests/exam

T01	ELL 061	2020-01-13	2020-04-03	W	10:30-11:20
T02	ECS 108	2020-01-13	2020-04-03	W	13:30-14:20
T03	DSB C116	2020-01-13	2020-04-03	R	14:30-15:20

In the Beginning ...

- There was hardware
 - processor
 - storage
 - card reader...
- And not much else
 - no operating system
 - no libraries
 - no compilers

Hardware



MAGICAL Abstractions

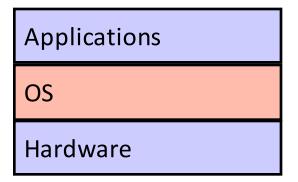
Hardware

- disks
- memory
- processors
- network
- monitor
- keyboard
- mouse

Operating system

- files
- programs
- threads of control
- communication
- windows, graphics
- input
- locator

The traditional picture



- "The OS is everything you don't need to write in order to run your application"
 - In some ways, it is:
 - all operations on I/O devices require OS calls (syscalls)
 - In other ways, it isn't:
 - you use the CPU/memory without OS calls
 - it intervenes without having been explicitly called

The OS and hardware

- An OS mediates programs' access to hardware resources (sharing and protection)
 - computation (CPU)
 - volatile storage (memory) and persistent storage (disk, etc.)
 - network communications (TCP/IP stacks, Ethernet cards, etc.)
 - input/output devices (keyboard, display, sound card, etc.)
- The OS abstracts hardware into logical resources and well-defined interfaces to those resources (ease of use)
 - processes (CPU, memory)
 - files (disk)
 - programs (sequences of instructions)
 - sockets (network)

Privileged instructions

- some instructions are restricted to the OS
 - known as privileged instructions
- e.g., only the OS can:
 - directly access I/O devices (disks, network cards)
 - why?
 - manipulate memory state management
 - page table pointers, TLB loads, etc.
 - why?
 - manipulate special 'mode bits'
 - interrupt priority level
 - why?

OS protection

- So how does the processor know if a privileged instruction should be executed?
 - the architecture must support at least two modes of operation: kernel mode and user mode
 - VAX, x86 support 4 protection modes
 - mode is set by status bit in a protected processor register
 - user programs execute in user mode
 - OS executes in kernel (privileged) mode (OS == kernel)
- Privileged instructions can only be executed in kernel (privileged) mode

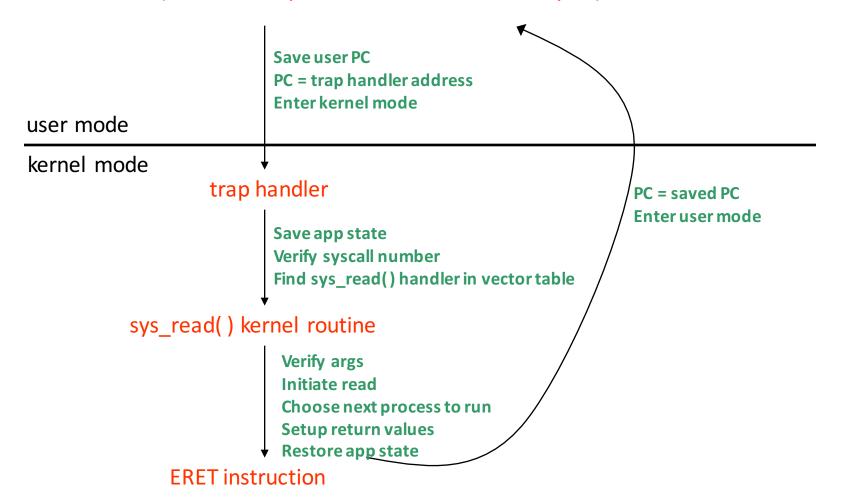
Crossing protection boundaries

- So how do user programs do something privileged?
 - e.g., how can you write to a disk if you can't execute an I/O instructions?
- User programs must call an OS procedure that is, get the OS to do it for them
 - OS defines a set of system calls
 - User-mode program executes system call instruction
- Syscall instruction
 - Like a <u>protected</u> procedure call

- The syscall instruction atomically:
 - Saves the current PC
 - Sets the execution mode to privileged
 - Sets the PC to a handler address
- With that, it's a lot like a local procedure call
 - Caller puts arguments in a place callee expects (registers or stack)
 - One of the args is a syscall number, indicating which OS function to invoke
 - Callee (OS) saves caller's state (registers, other control state) so it can use the CPU
 - OS function code runs
 - OS must verify caller's arguments (e.g., pointers)
 - OS returns using a special instruction
 - Automatically sets PC to return address and sets execution mode to user

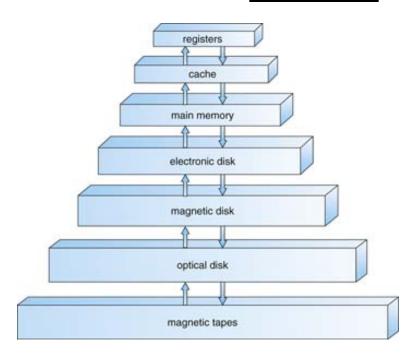
A kernel crossing illustrated

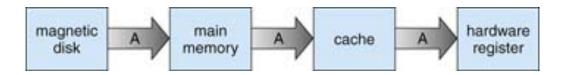
Firefox: read(int fileDescriptor, void *buffer, int numBytes)



Memory?







operating system

job 1

job 2

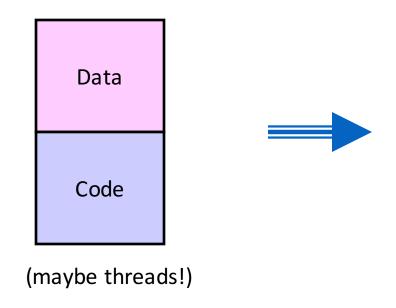
job 3

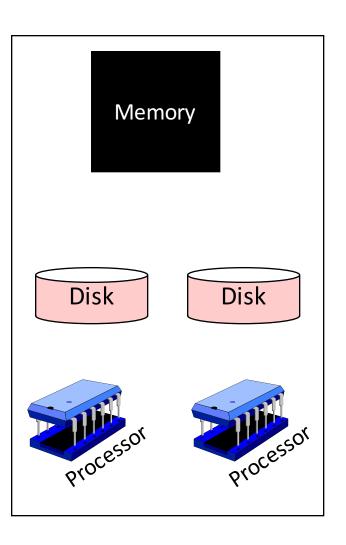
job 4

512M

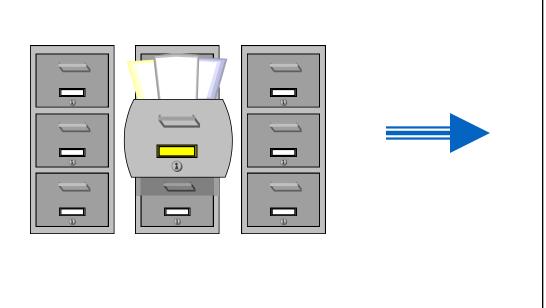
Level	1	2	3	4
Name	registers	cache	main memory	disk storage
Typical size	< 1 KB	< 16 MB	< 64 GB	> 100 GB
Implementation technology	custom memory wi multiple ports,		-COMCED DRAM	magnetic disk
Access time (ns)	0.25- 0.5	0.5- 25	80- 250	5,000.000
Bandwidth (MB/se	c2)0,009-100,000	5000- 10,000	1000- 5000	20- 150
Managed by	compiler	hardware	operating sys	coperating sys
Backed by	cache	main memory	disk	CD or tape

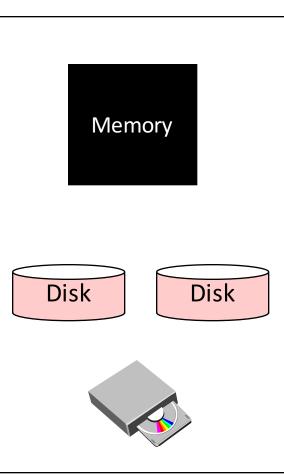
Programs (Assignment 2)





Files (Assignment 3)

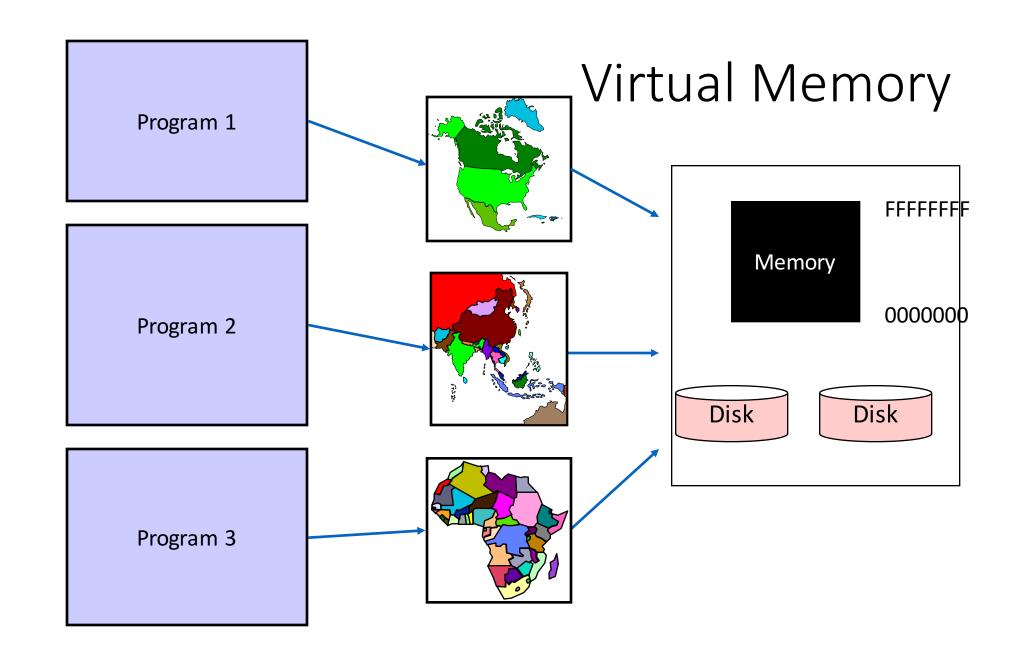




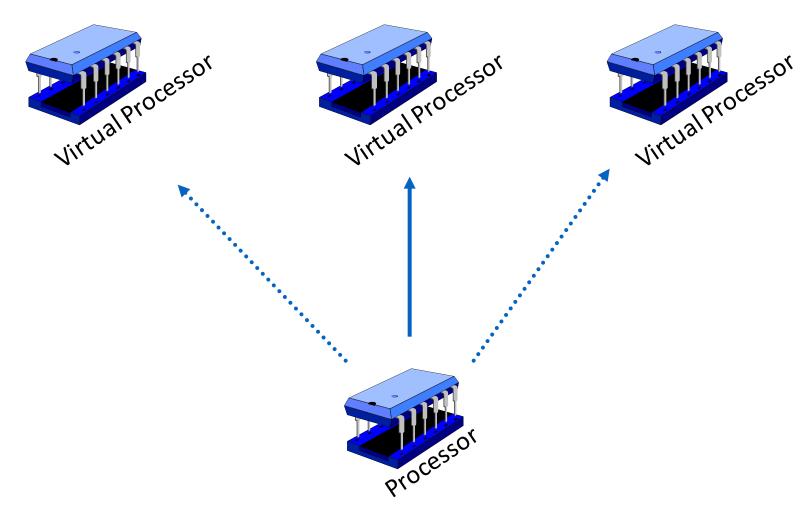
There are a lot of Issues...

- Naming!
- Allocating space on disk (permanent storage)
 - organized for fast access
 - minimize waste
- Shuffling data between disk and memory (high-speed temporary storage)
- Coping with crashes

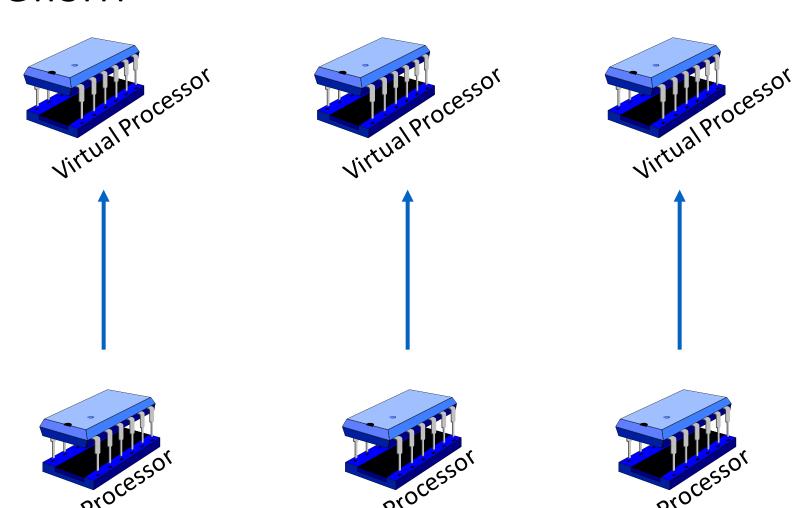
Memory Sharing (2) Program 1 Program 2 Memory Program 3



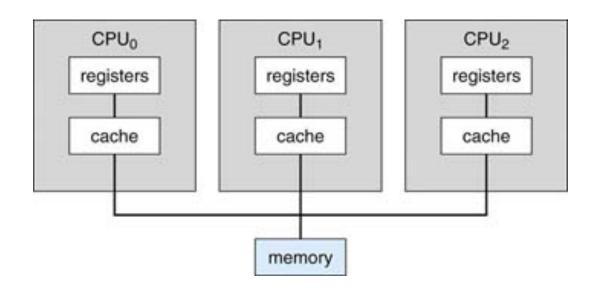
Concurrency (more in Assignment 2)

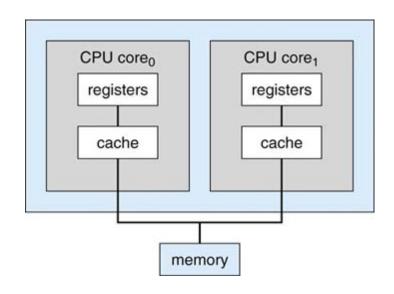


Parallelism



Processors and cores! (Assignment 2)





TO DO!!!!

- Be sure you're on CSC 360 CourseSpaces and upload your "picture"
 - You should have received email by now!
- Please keep up with the programming!
- Forum (reading/questions) is posted on the web now
 - Due as part of Assignment 1
 - https://www.g2.com/categories/operating-system
- Assignment 0 (parts 1 and 2) posted on CousreSpaces now
 - Tutorial question next week!
 - Upload a super duper d-list!!!