CSC360 Operating Systems

MapReduce / Dining Philosophers

Today!

- A2 part2: how to approach it
 - Ideas from MapReduce (many examples out there, will look at one)
- A2 part3: dining philosophers
 - Race conditions
 - Deadlock
 - Starvation
- Blockchain!
 - What does it have to do with concurrency?

MapReduce A2/P2

Splitting

Original Deer Bear River

Deer Bear River

Deer Car Bear

River Bear River

Deer Car Bear

River Bear River

Mapping

Deer: 1 Bear: 1

River: 1

Deer: 1 Car: 1 Bear: 1

River 1 Bear: 1 River: 1 Shuffling

Bear: 1 Bear: 1

Bear: 1

Car: 1

Deer: 1

Deer: 1

River: 1

River: 1 River: 1 Reducing

Bear: 3

Final Output

Car: 1

Bear: 3

Deer: 2

River: 3

Car: 1

Deer: 2

River: 3

What does this code do?

```
void** map(void** things, void* (*f)(void*), int length){
         void** results = malloc(sizeof(void*)*length);
 5
         for(int i = 0; i < length; i++){</pre>
             void* thing = things[i];
             void* result = (*f)(thing);
             results[i] = result;
 9
         }
10
11
12
         return results;
13
```

http://www.datastuff.tech/programming/how-to-code-mapreduce-in-c-from-scratch-using-threads-pt-1-map/

What about this code?

```
1 // returns 1 if the number is a prime, 0 otherwise
    void* naivePrime(void* number){
        int n = *((int*) number);
        int* res = malloc(sizeof(int));
        *res = 1;
        for(int i = 2; i < n; i++){
            if(n%i==0){
8
                *res=0;
                return res;
10
        return res;
13 }
14
     int main(int argc, char** argv){
16
      int N = 1000;
      int** numbers = malloc(sizeof(int*)*N);
      for(int i = 0; i < N; i++){
19
          int* n = malloc(sizeof(int));
20
          *n = i;
          numbers[i] = n;
24
      int** resulting_numbers = NULL;
      void** is_prime = map((void**) numbers, naivePrime, N);
       resulting_numbers = (int**) is_prime;
28 }
```

How can we use threads to speed this up?

- Partition the data to be processed into separate threads
 - How many threads?
- Experiment! For example, the naïve prime single versus multi-threaded

Time (s):	150000 elements	300000 elements
single-threaded	5.02	18.73
2-threads	3.76	13.78
4-threads	2.73	10.14
8-threads	2.43	8.70

"Embarrassingly Parallelizable!"

```
struct map_argument {
   void** things;
   void** results;
   void* (*f)(void*);
    int from;
    int to;
                                   void* chunk_map(void* argument){
   };
                                       struct map_argument* arg = (struct map_argument*) argument;
                                       for(int i = arg->from; i < arg->to; i++){
                                           arg->results[i] = (*(arg->f))(arg->things[i]);
                                       }
                                       return NULL;
```

```
\verb|concurrent_map((\verb|void**|) | numbers, twice, N, NTHREADS)|\\
```

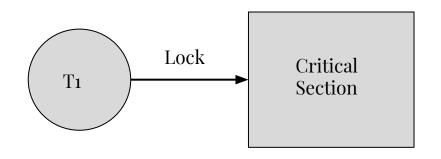
```
int nthreads){
 3
         void** results = malloc(sizeof(void*)*length);
 4
         struct map_argument arguments[nthreads];
 5
         pthread_t threads[nthreads];
 6
         int chunk_size = length/nthreads;
 8
         for(int j = 0 ; j < nthreads; j++){</pre>
9
             int from = j*chunk_size;
10
             struct map_argument* argument = &arguments[j];
12
             init_map_argument(argument, things, results, f, from, from+chunk_size);
             pthread_create(&threads[j], NULL, chunk_map, (void*) argument);
14
16
         for(int i = 0; i < length % nthreads; i++){</pre>
             int idx = chunk_size*nthreads + i;
18
             results[idx] = (*f)(things[idx]);
19
20
         for(int k = 0; k < nthreads; k++){</pre>
22
             pthread_join(threads[k], NULL);
24
         return results;
26
```

27 }

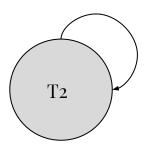
void** concurrent_map(void** things, void* (*f)(void*), int length,

Mutex

- Either lock or unlock

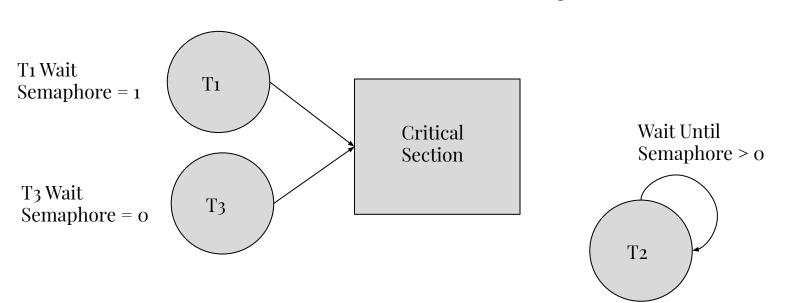


Wait Until Critical Section is Unlocked



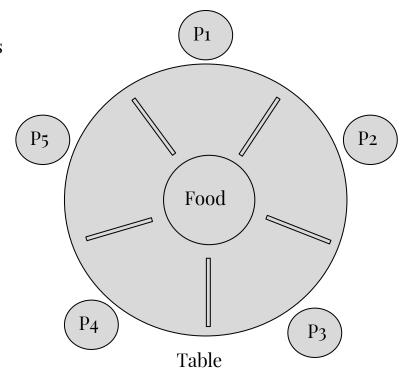
Semaphores

- Wait (-1) or Signal (+1)
- ex) Semaphore = 2
 - This means 2 Threads can be in Critical Section
 - T1 enters after Wait, Semaphore now equals 1
 - T₃ enters, Semaphore = 0, no room left
 - T2 waits until Semaphore > o



Dining Philosophers A2/P3

5 Philosophers5 Chopsticks



- A Philosophers needs 2 Chopsticks to eat.
- We need to implement this solution using Semaphores or Mutexes.
- A Philosopher must Lock a Left and a Right Chopstick down until done eating.
- Start: Put a lock on each chopstick

Dining Philosophers - Solution 1

- 1. Think until the left fork is available; when it is, pick it up
- 2. Think until the right fork is available; when it is, pick it up
- 3. With both forks, eat for a fixed amount of time
- 4. Put the right fork down
- 5. Put the left fork down
- 6. Repeat

Issues

Deadlock

- If each Philosopher takes the left chopstick, no one can take a right chopstick

Starvation

- 2 Philosophers are fast thinkers and fast eaters
- The other 3 Philosophers have no chance to eat

Solution 1 (Bad Solution) for a philosopher

```
Think();
Pick up left chopstick;
Pick up right chopstick;
Eat();
Put down right chopstick;
Put down left chopstick;
```

- Problem: POSSIBLE Deadlock
- How?
 - Each philosopher can pick up left fork before anyone picks up their right fork
 - Now everyone is waiting for right fork!

```
#include <stdio.h>
#include <pthread.h>
#include "dphil.h"
```

using pthreads...

```
typedef struct sticks {
 pthread mutex t *lock[MAXTHREADS];
 int phil count;
} Sticks;
void pickup(Phil struct *ps)
 Sticks *pp;
 int i;
 int phil count;
 pp = (Sticks *) ps->v;
 phil count = pp->phil count;
 pthread mutex lock(pp->lock[(ps->id+1)%phil count]); /* lock up right stick */
```

```
void putdown(Phil struct *ps)
 Sticks *pp;
  int i;
  int phil count;
 pp = (Sticks *) ps->v;
 phil count = pp->phil count;
  pthread mutex unlock(pp->lock[(ps->id+1)%phil count]); /* unlock right stick */
 pthread mutex unlock(pp->lock[ps->id]); /* unlock left stick */
void *initialize state(int phil count)
  Sticks *pp;
  int i;
 pp = (Sticks *) malloc(sizeof(Sticks));
 pp->phil count = phil count;
  for (i = 0; i < phil count; i++) {
   pp->lock[i] = (pthread mutex t *) malloc(sizeof(pthread mutex t));
  for (i = 0; i < phil count; i++) {
   pthread mutex init(pp->lock[i], NULL);
 return (void *) pp;
```

Solution 2: Global lock... what are the tradeoffs?

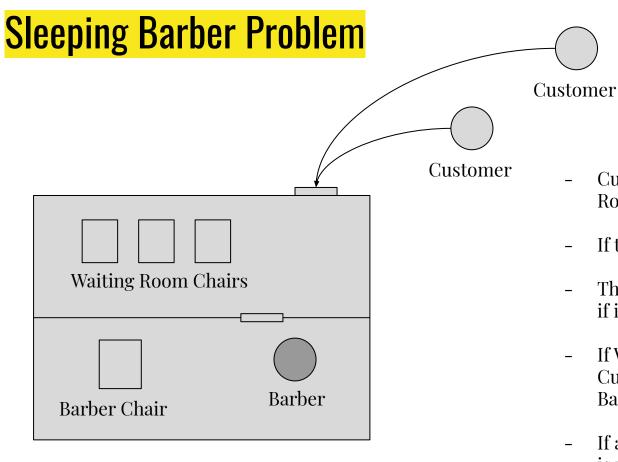
```
Think();
table.lock();
while (!both chopstick available)
  chopstickPutDown.wait();
Pick up left chopstick;
Pick up right chopstick;
table.unlock();
Eat();
Put down right chopstick;
Put down left chopstick;
chopstickPutDown.signal();
```

Solution 3: Reactive

```
Think();
Pick up left chopstick;
if(right chopstick available) {
    Pick up right chopstick;
} else {
    Put down left chopstick;
    continue; //Go back to Thinking
Eat();
```

Solution 4: Global ordering

```
Think();
Pick up "smaller" chopstick from left and right;
Pick up "bigger" chopstick from left and right;
Eat();
Put down "bigger" chopstick from left and right;
Put down "smaller" chopstick from left and right;
```



Barber Shop

- If Waiting Room not empty Customer sits in Barber Chair,

if it's empty goes to sleep.

Barber cuts hair.

Room.

- If a Customer comes in and no one is waiting, the Customer wakes up the Barber.

Customers must wait in Waiting

If the waiting room is full - leave.

The Barber checks Waiting Room -