Confuzer: System Call Fuzzer for Understanding Secure Container Mechanism

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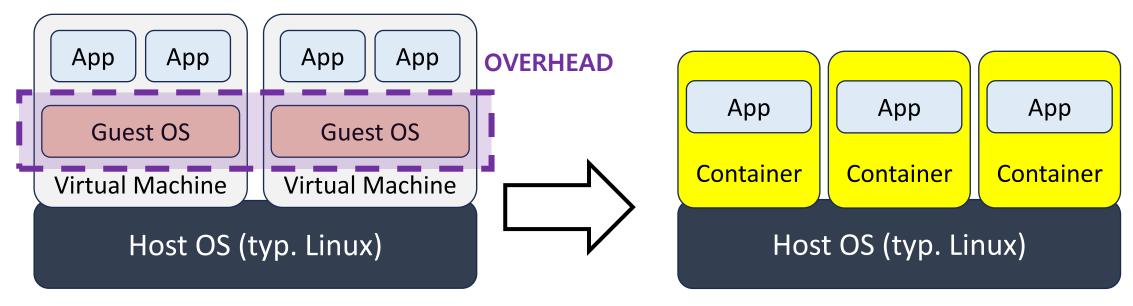
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Container

In the past, cloud environments have used VM technology, but nowadays, mainly use Container technology

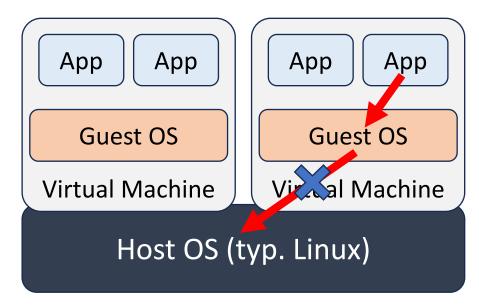


VMs are Heavy because they contain guest OS

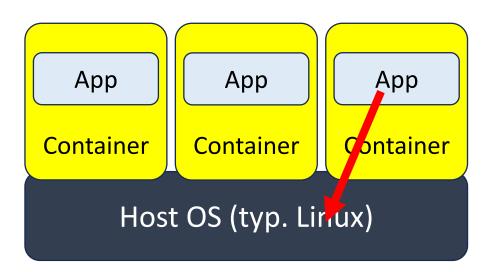
Containers only share a single Host OS

Container Security Problem

Because Container shares the host OS, it is possible to access the host kernel and affect the entire system through Privilege Escalation by invoking dangerous system calls

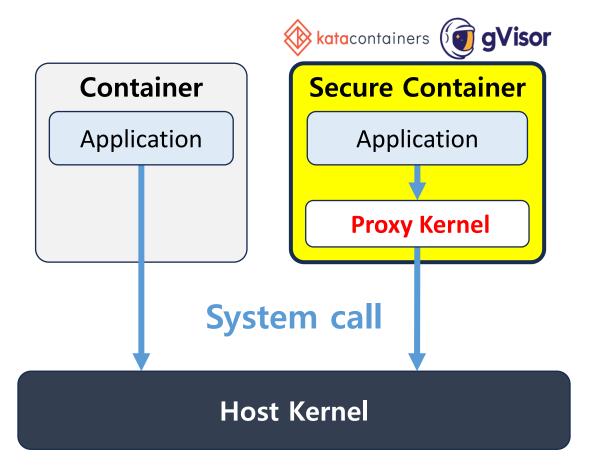


VM's application can't access to Host OS



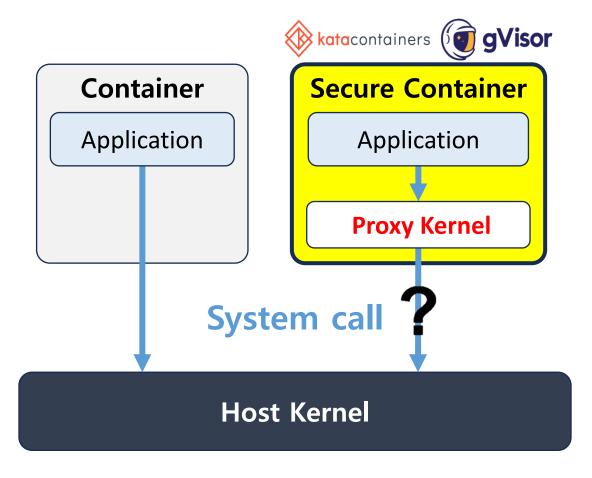
Container's application can access to Host OS

What is Secure Container?



- Secure containers add a proxy kernel to prevent the container from directly accessing the host kernel
- Currently, this approach has led to various secure containers, such as gVisor and Kata Containers
- The Proxy Kernel is unable to determine which system calls are being invoked

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- Secure containers add a proxy kernel to prevent the container from directly accessing the host kernel
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Goals

Goal 1. Fuzzing on A// System Call

Utilize *Confuzer* to fuzz all system calls for all possible arguments

Goal 2. Expand *Observability* for Secure Container Mechanism

Analyze output system calls and *understand the secure* container's mechanism

Can help find Contradictions by comparing the obtained system call data to the expected design of the Proxy Kernel Logic

Accuracy

Improve the Logic of the Proxy Kernel Layer from the Developer's Perspective

Clarity

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Efficiency

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Clarity Accuracy

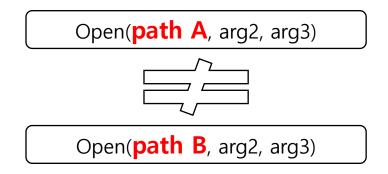
Data

Obtained System Call Data can be Used as Primary Data to Quantify the Size of the Container Attack Surface

Of the many system calls observed in the Host Kernel, it should be possible to

Extract Only Certain Test System Calls





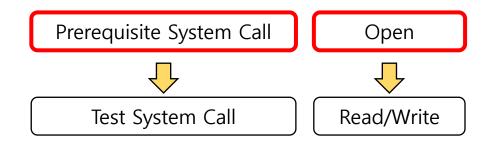
Ex)

Some arguments affect observed system calls, so **Tests Should be Conducted with All Possible Arguments**

Open system call produces *Very Different Results*when inserting a path inside a container than when inserting a path to storage shared with the host

Since there are *System Calls that Require Preparation* rather than system calls that can be executed alone, the *Preparation Process Must be Resolved*

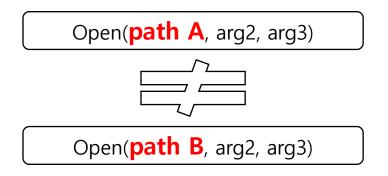
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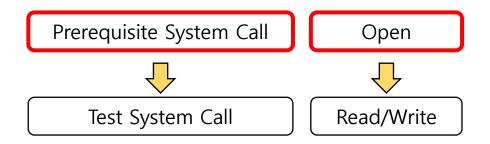
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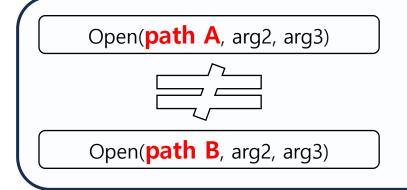
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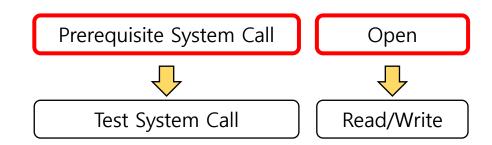
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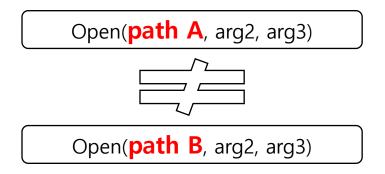
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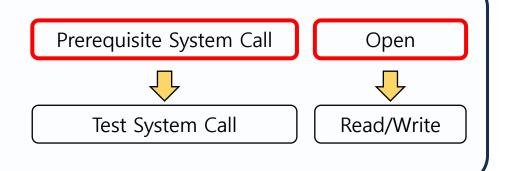
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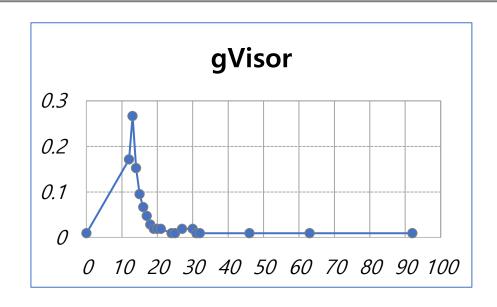
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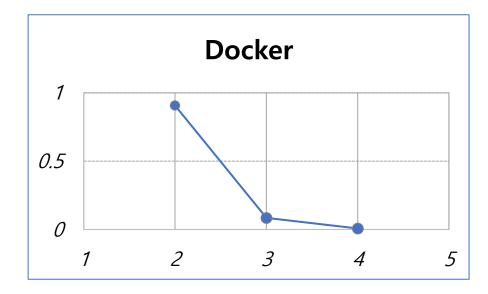
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Current Results and Findings



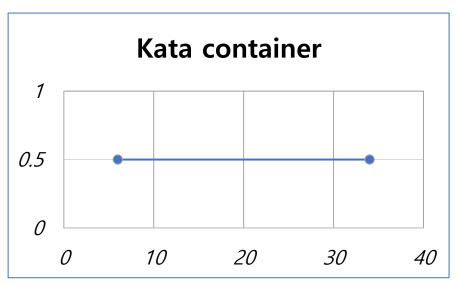


X-axis

The number of system calls generated by the runtime as a response to single system call

Y-axis

The proportion of system calls



The End

