

**CS 3133 Foundations of Computer Science**  
**C term 2018**

**Solutions for Homework 5**

1. Exercise 17.b. on page 249.

**Solution:** (with the pumping lemma for context-free languages) Let us assume indirectly that the language  $L = \{a^i b^j c^i d^j \mid i, j \geq 0\}$  is context-free. Then by the pumping lemma for context-free languages there exists a number  $k$ , such that every string  $z \in L$  of length  $k$  or more can be decomposed into substrings  $u, v, w, x$  and  $y$  such that  $\text{length}(vwx) \leq k$ ,  $\text{length}(v) + \text{length}(x) > 0$  and  $uv^i wx^i y \in L$  for all  $i \geq 0$ .

Consider the string  $z = a^k b^k c^k d^k$ . Clearly  $z \in L$  and  $\text{length}(z) \geq k$ . Using the pumping lemma we decompose  $z$  into substrings  $u, v, w, x$  and  $y$ , where  $0 < \text{length}(vwx) \leq k$ . But in this case  $uv^2 wx^2 y$  cannot be in  $L$ , a contradiction. Indeed, consider the possibilities for  $v$  and  $x$ . If either of these contains more than one type symbol, then  $uv^2 wx^2 y$  is not in  $L$ . So  $v$  and  $x$  must be substrings of one of  $a^k, b^k, c^k$ , or  $d^k$  and if they contain different symbols, then they have to be two “consecutive” symbols, so  $a$  and  $b$ , or  $b$  and  $c$ , or  $c$  and  $d$ . But then again  $uv^2 wx^2 y$  is not in  $L$ , since either the number of  $a$ ’s is different from the number of  $c$ ’s, or the number of  $b$ ’s is different from the number of  $d$ ’s. Thus  $L$  is not context-free. (20 points)

2. Let  $M$  be the Turing machine defined by

| $\delta$ | B             | a             | b             | c             |
|----------|---------------|---------------|---------------|---------------|
| $q_0$    | $(q_0, B, R)$ | $(q_0, a, R)$ | $(q_0, b, R)$ | $(q_1, c, L)$ |
| $q_1$    | $(q_2, B, R)$ | $(q_1, b, L)$ | $(q_1, a, L)$ | -             |
| $q_2$    | -             | -             | -             | -             |

- (a) Trace the computation for the input string  $abcb$ .  
 (b) Trace the first six transitions of the computation for the input string  $abab$ .

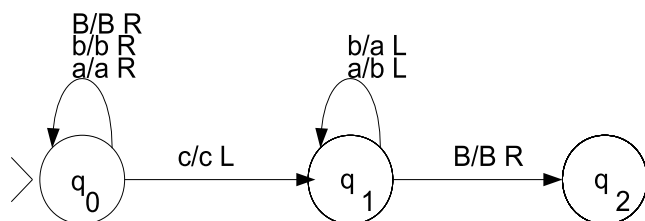
- (c) Give the state diagram of  $M$ .  
 (d) Describe the result of a computation in  $M$ .

**Solution:**

a,b.)

|                     |                     |
|---------------------|---------------------|
| $q_0 Babcab$        | $q_0 Babab$         |
| $\vdash q_0 abcab$  | $\vdash q_0 abab$   |
| $\vdash aq_0 bcab$  | $\vdash aq_0 bab$   |
| $\vdash abq_0 cab$  | $\vdash abq_0 ab$   |
| $\vdash aq_1 bcab$  | $\vdash abaq_0 b$   |
| $\vdash q_1 aacab$  | $\vdash ababq_0 B$  |
| $\vdash q_1 Bbacab$ | $\vdash ababBq_0 B$ |
| $\vdash q_2 bacab$  |                     |

c.)



d.) If there is a  $c$ , then swap the  $a$ 's and  $b$ 's before the first  $c$ . Otherwise, if there is no  $c$ , then go to the right infinitely. (20 points)

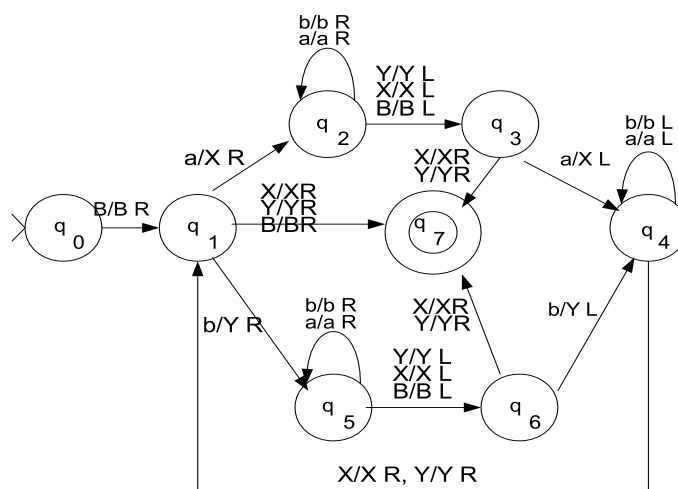
3. Construct a Turing machine with input alphabet  $\{a, b, c\}$  that accepts strings in which the first  $c$  is immediately preceded by the substring  $aaa$ . A string must contain a  $c$  to be accepted by the machine.

**Solution:**



dromes over  $\{a, b\}$ .

**Solution:**



(20 points)