

Welcome to

CS 3516:
Computer Networks

Prof. Yanhua Li

Time: 9:00am –9:50am M, T, R, and F

Location: AK219

Fall 2018 A-term

Extra office hour on Monday 9/10

Office hours,

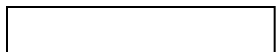
Email to cs3516-ta@cs.wpi.edu, and me at yli15@wpi.edu

Canvas discussion forum,

	Mondays	Tuesdays	Wednesdays	Thursdays	Fridays
9-9:50am	Lecture FL320	Lecture FL320		Lecture FL320	Lecture FL320
10-10:30am	Prof. Li, AK130	Prof. Li, AK130	TA: Menghai Sub-basement in Fuller		Prof. Li, AK130
10:30-11:30	Prof. Li, AK130		9:30-11:30AM		10-11AM
1-3pm	TA: Marissa Sub-basement in Fuller	TA: Sanket Sub-basement in Fuller	TA: Marissa Sub-basement in Fuller	TA:Sanket Sub-basement in Fuller	TA: Menghai Sub-basement in Fuller
		Project 1 Due			



Office hours for all questions, e.g., project/lab assignment related questions, like programming...



Office hours for lecture related questions, and general questions for labs and projects.

Chapter 2: outline

2.1 principles of network applications

- app architectures
- app requirements

2.2 Web and HTTP

2.5 DNS

Service Overview, Structure

Resolution process

Data Format

DNS: domain name system

people: many identifiers:

- SSN, name, passport #

Internet hosts, routers:

- IP address (32 bit) - used for addressing datagrams
- “name”, e.g., `www.yahoo.com` - used by humans

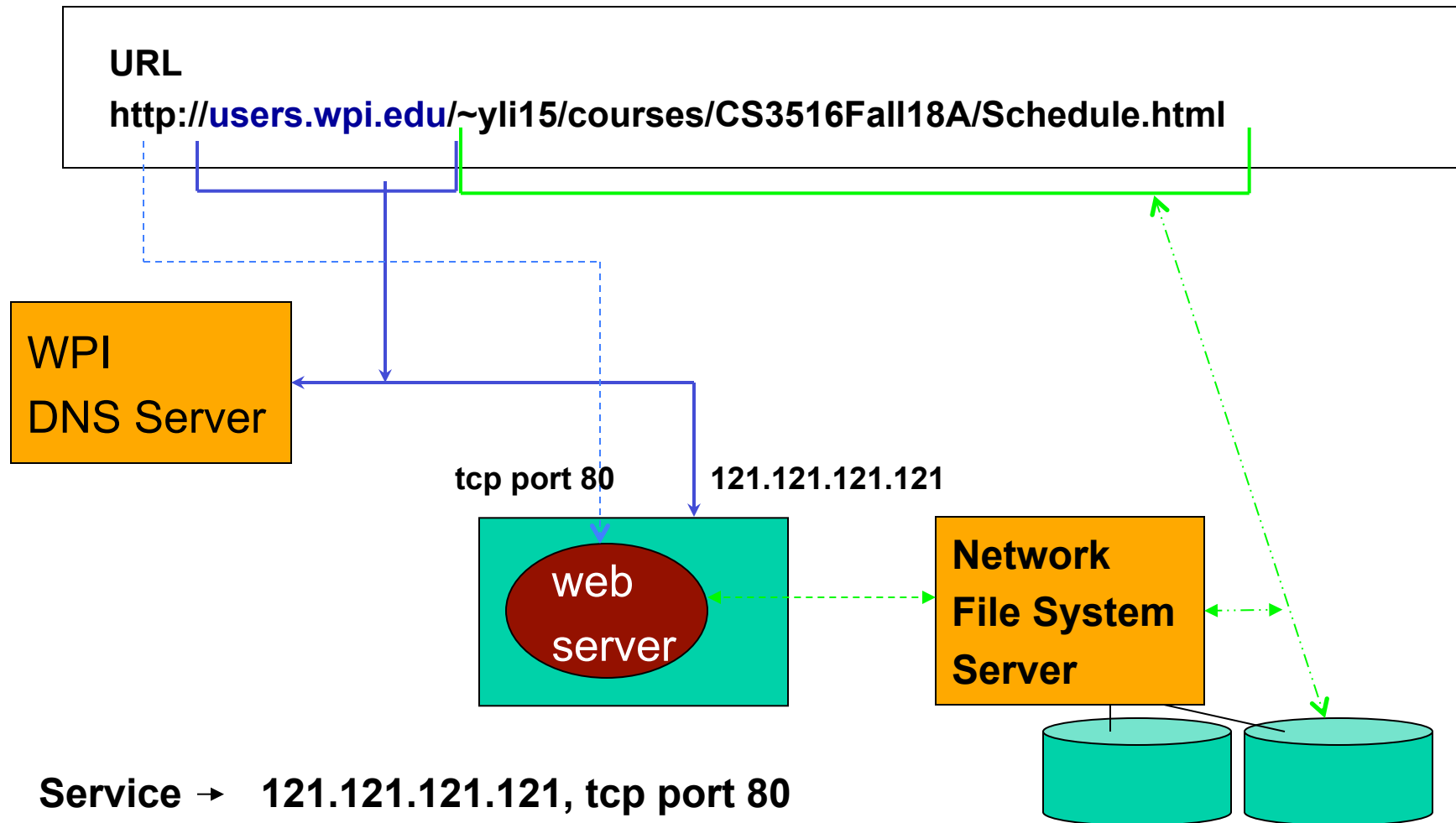
Q: how to map between IP address and name, and vice versa ?

Domain Name System:

- ❖ *distributed database*
implemented in hierarchy of many *name servers*
- ❖ *application-layer protocol:* hosts, name servers communicate to *resolve* names (address/name translation)
 - note: core Internet function, implemented as application-layer protocol
 - complexity at network's “edge”

Analogy: Marshalls -> Physical Address

Resolving Name, Locating Service/Object



Service → 121.121.121.121, tcp port 80

Object → ~yli15/courses/CS3516Fall18A/Schedule.html



DNS: services, structure

DNS services

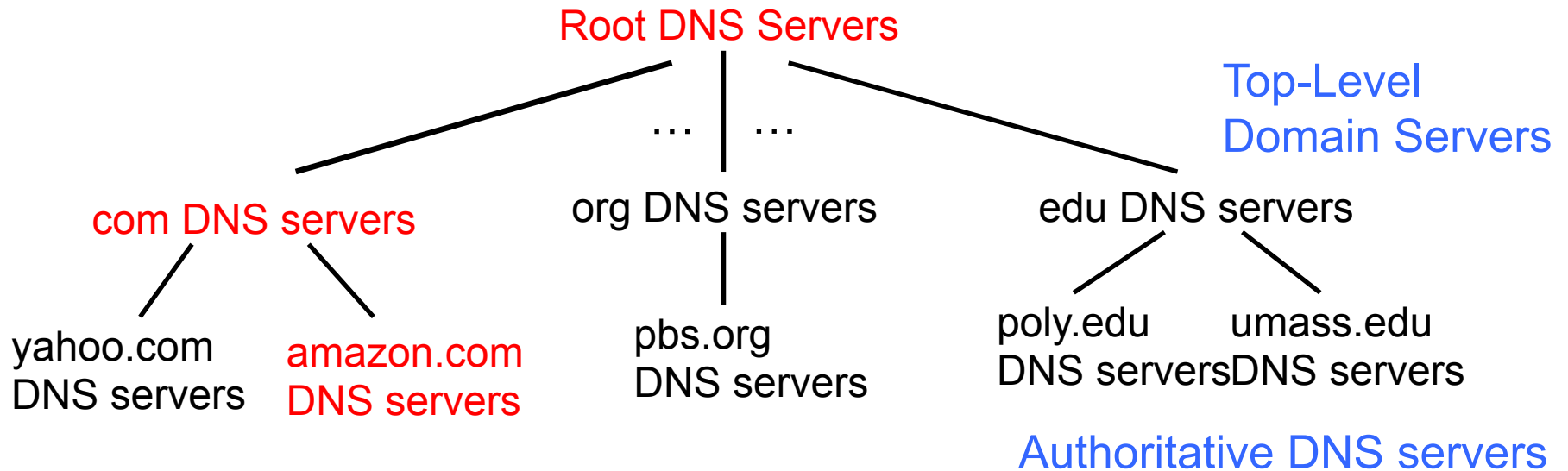
- ❖ hostname to IP address translation
- ❖ host aliasing
 - canonical, alias names
- ❖ mail server aliasing
- ❖ load distribution
 - replicated Web servers: many IP addresses correspond to one name

why not centralize DNS?

- ❖ single point of failure
- ❖ traffic volume
- ❖ distant centralized database
- ❖ maintenance

A: doesn't scale!

DNS: a distributed, hierarchical database



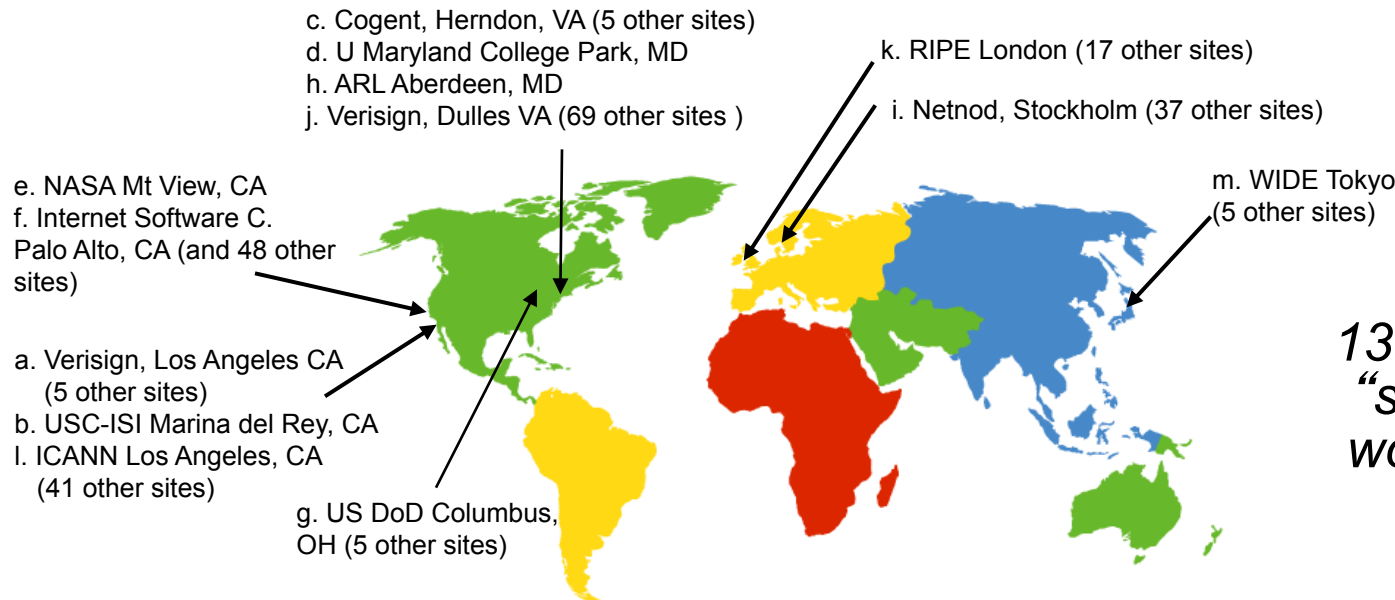
client wants IP for www.amazon.com; 1st approx:

- ❖ client queries root server to find com DNS server
- ❖ client queries .com DNS server to get amazon.com DNS server
- ❖ client queries amazon.com DNS server to get IP address for www.amazon.com

Analogy: Marshalls -> Physical Address

DNS: root name servers

- ❖ contacted by local name server that cannot resolve name
- ❖ root name server:
 - contacts authoritative DNS server
 - if name mapping not known, gets mapping
 - returns mapping to local name server



*13 root name
“servers”
worldwide*

TLD, authoritative servers

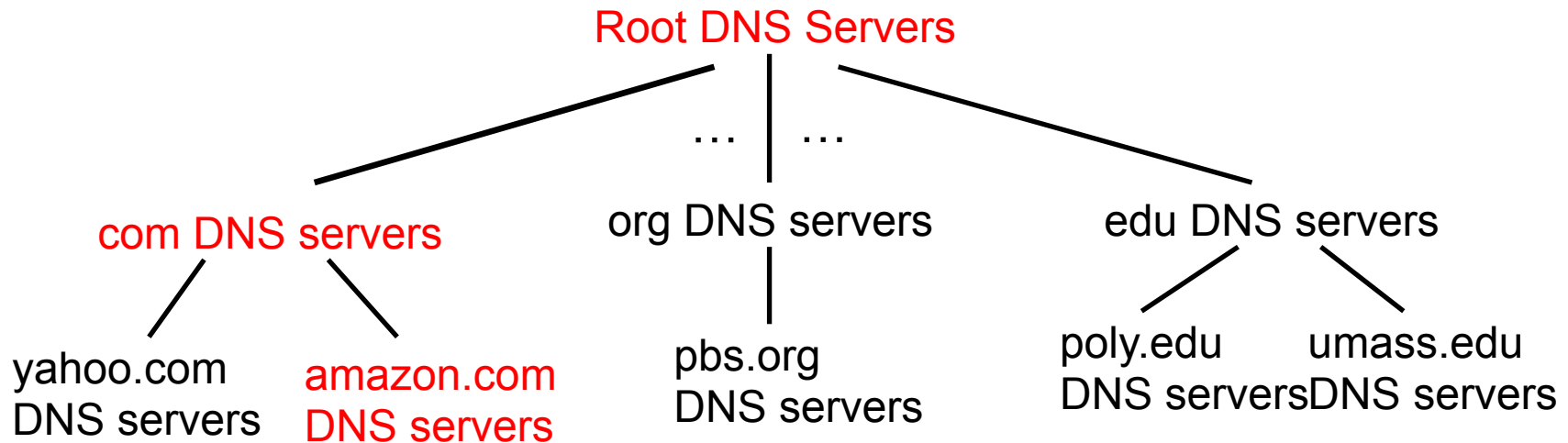
top-level domain (TLD) servers:

- responsible for com, org, net, edu, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, ca, jp
- **Network Solutions** maintains servers for .com TLD
- **Educause** for .edu TLD

authoritative DNS servers:

- organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- can be maintained by organization or service provider

DNS: a distributed, hierarchical database



client wants IP for www.amazon.com; 1st approx:

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- ❖ client queries amazon.com DNS server to get IP address for www.amazon.com

Analogy: Marshalls -> Physical Address



Local DNS name server

- ❖ does not strictly belong to hierarchy
- ❖ each ISP (residential ISP, company, university) has one
 - also called “default name server”
- ❖ when host makes DNS query, query is sent to its local DNS server
 - has **local cache** of recent name-to-address translation pairs (but may be out of date!)
 - **acts as proxy**, forwards query into hierarchy
- ❖ Difference btw Local DNS and Authoritative DNS server?
 - Given an organization, e.g., WPI, one for its internal users, one for external users

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Resolution process

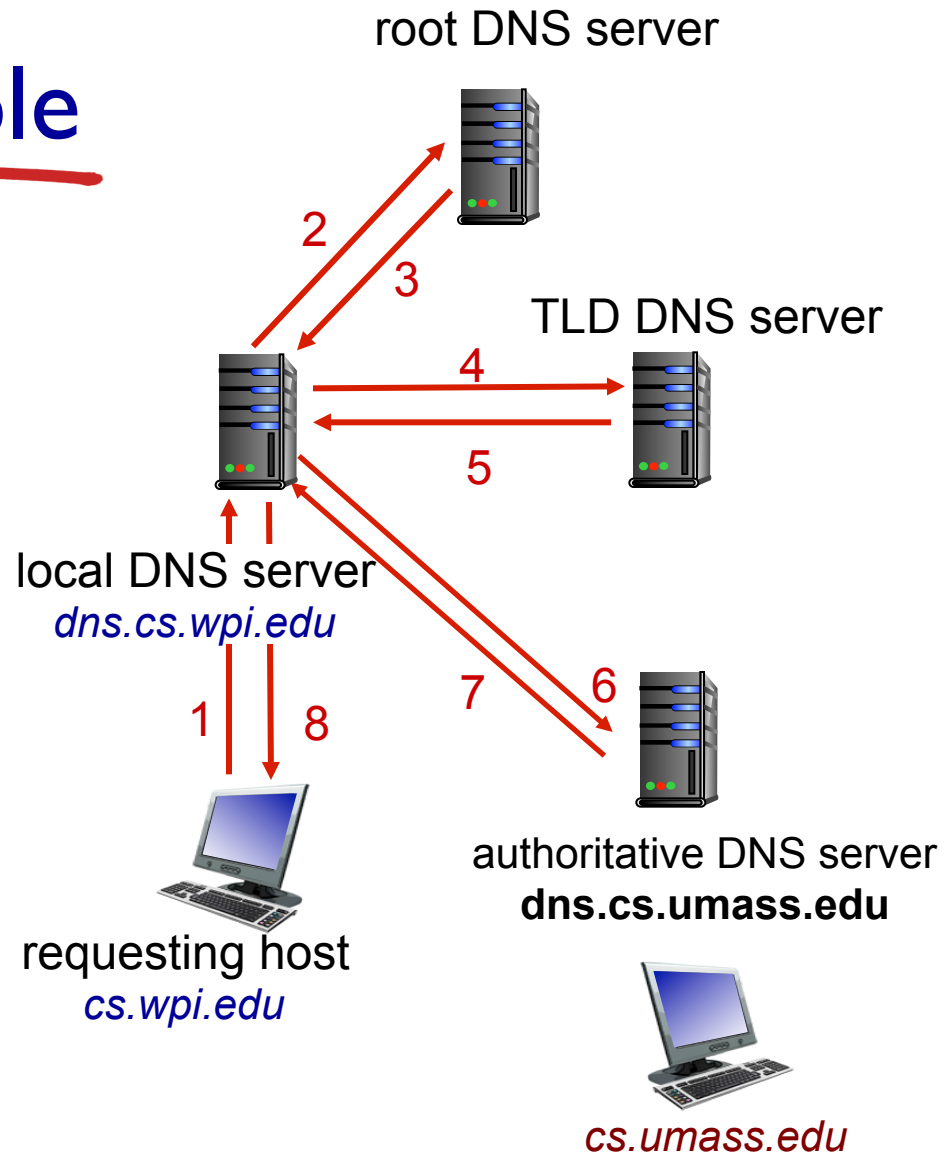
Data Format

DNS name resolution example

- ❖ host at *cs.wpi.edu* wants IP address for *cs.umass.edu*

iterated query:

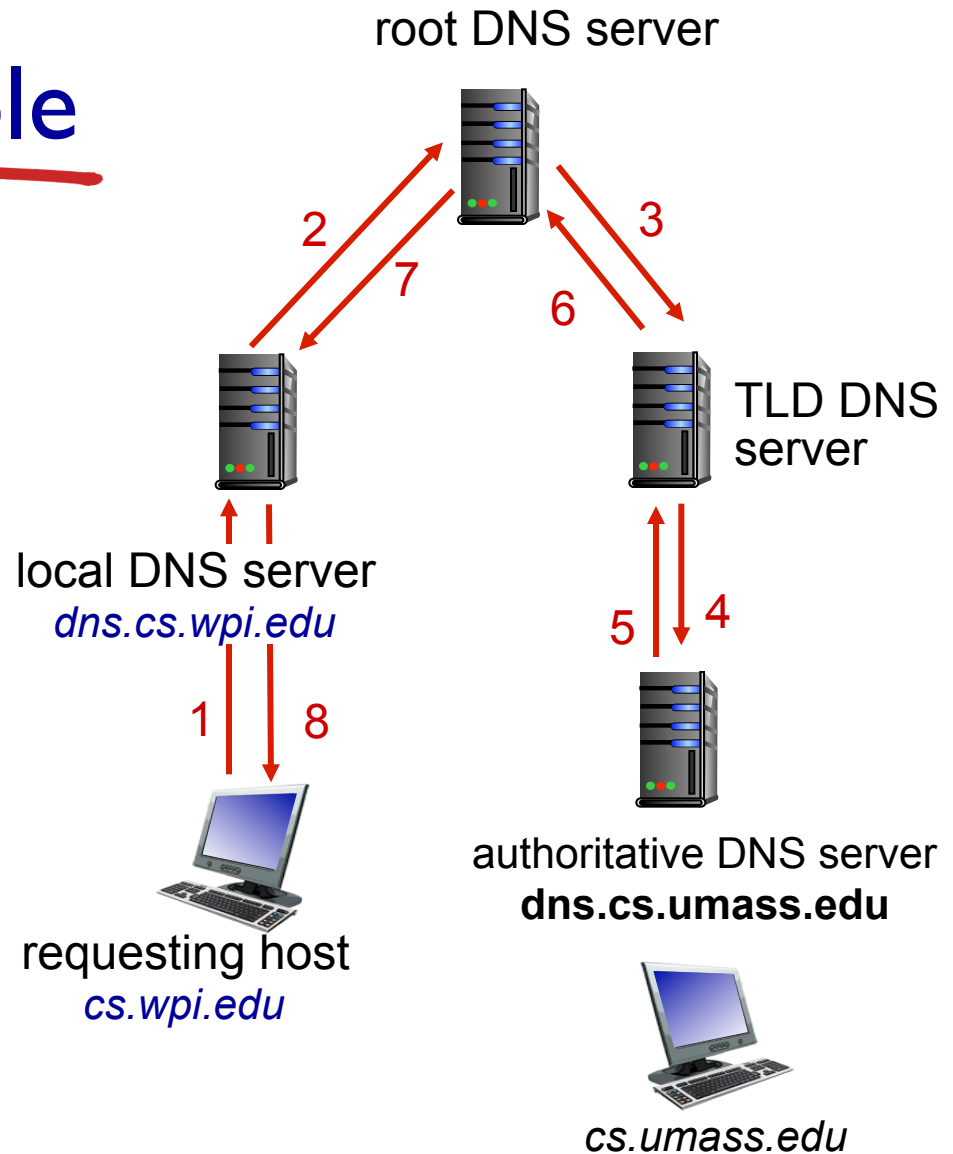
- ❖ contacted server replies with name of server to contact
- ❖ “I don’t know this name, but ask this server”



DNS name resolution example

recursive query:

- ❖ puts burden of name resolution on contacted name server
- ❖ **Cons:** heavy load at upper levels of hierarchy?



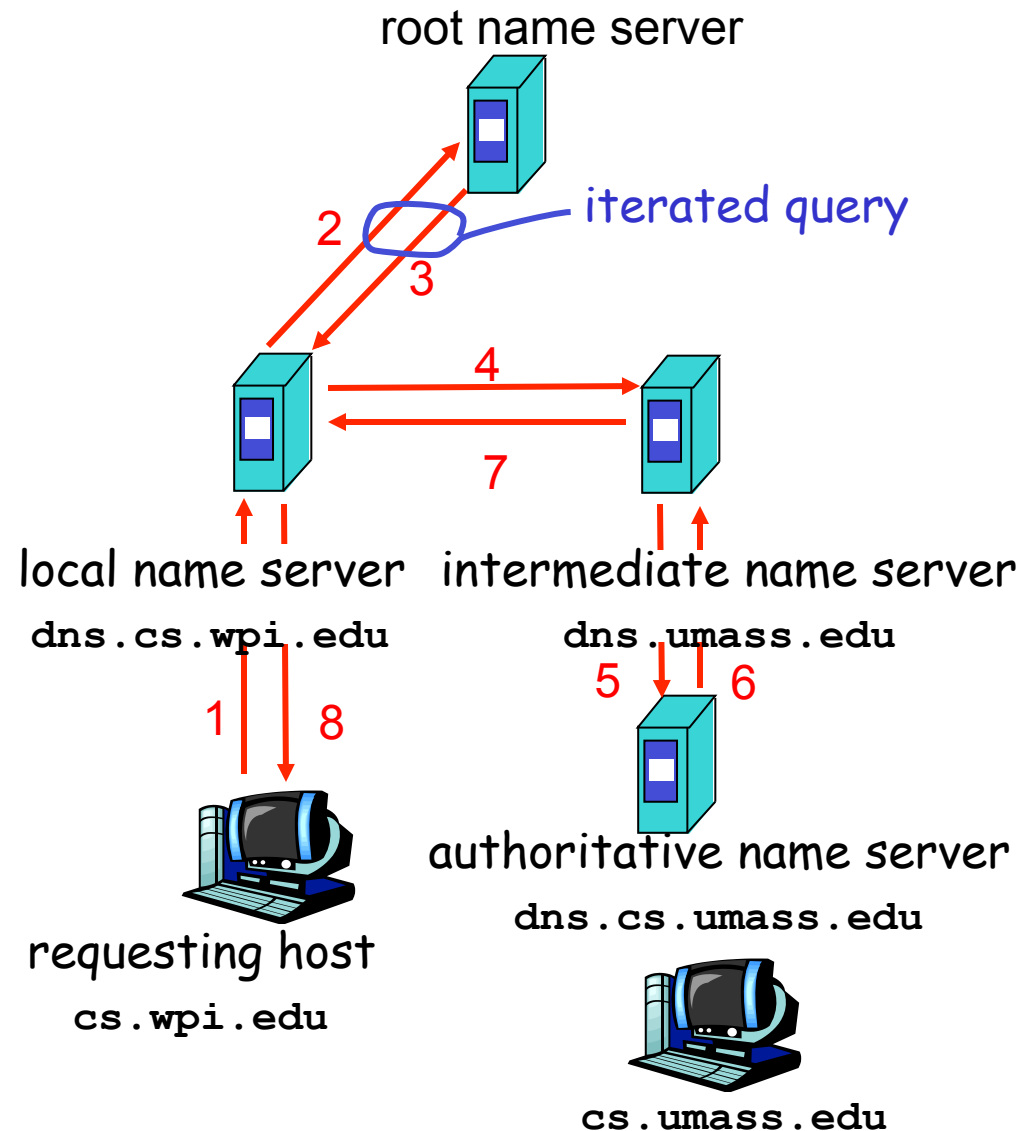
DNS queries

recursive query:

- ❖ puts burden of name resolution on contacted name server
- ❖ heavy load?

iterated query:

- ❖ contacted server replies with name of server to contact
- ❖ “I don’t know this name, but ask this server”



DNS: caching, updating records



- ❖ once (any) name server learns mapping, it *caches* mapping
 - **cache entries timeout** (disappear) after some time (TTL)
 - **TLD servers typically cached** in local name servers
 - thus root name servers not often visited
- ❖ cached entries may be *out-of-date* (best effort name-to-address translation!)
 - if name host changes IP address, it may not be known Internet-wide until all TTLs expire

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DNS records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

type=A

- **name** is hostname
- **value** is IP address

type=NS

- **name** is domain (e.g., foo.com)
- **value** is hostname of authoritative name server for this domain

type=CNAME

- **name** is alias name for some “canonical” (the real) name
- **www.ibm.com** is really **servereast.backup2.ibm.com**
- **value** is canonical name

type=MX

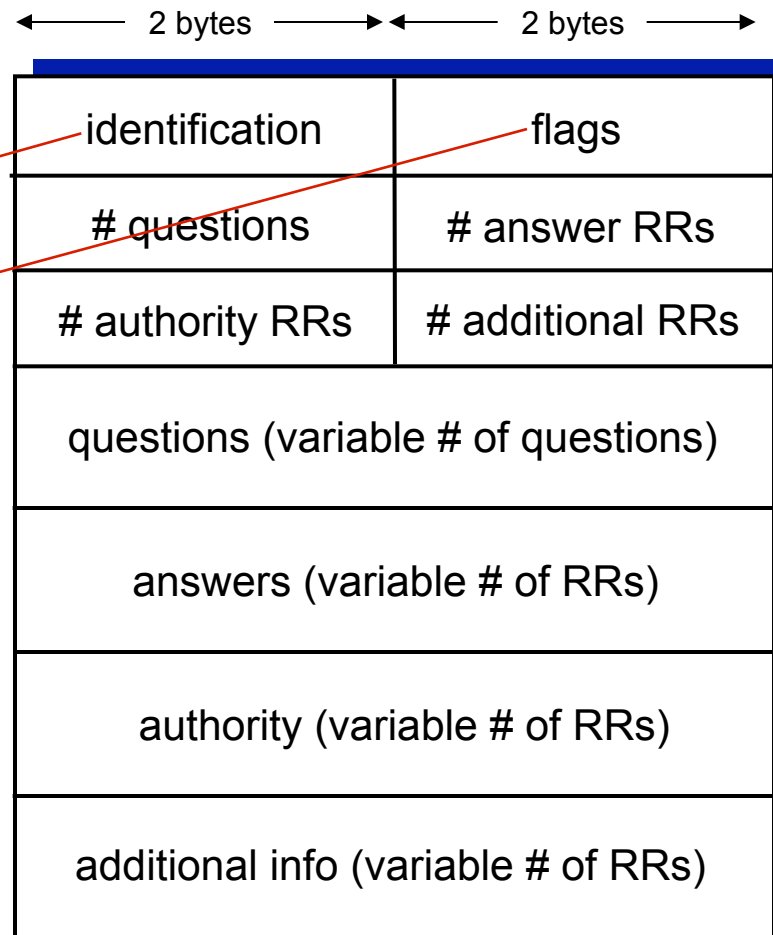
- **value** is name of mailserver associated with **name**

DNS protocol, messages

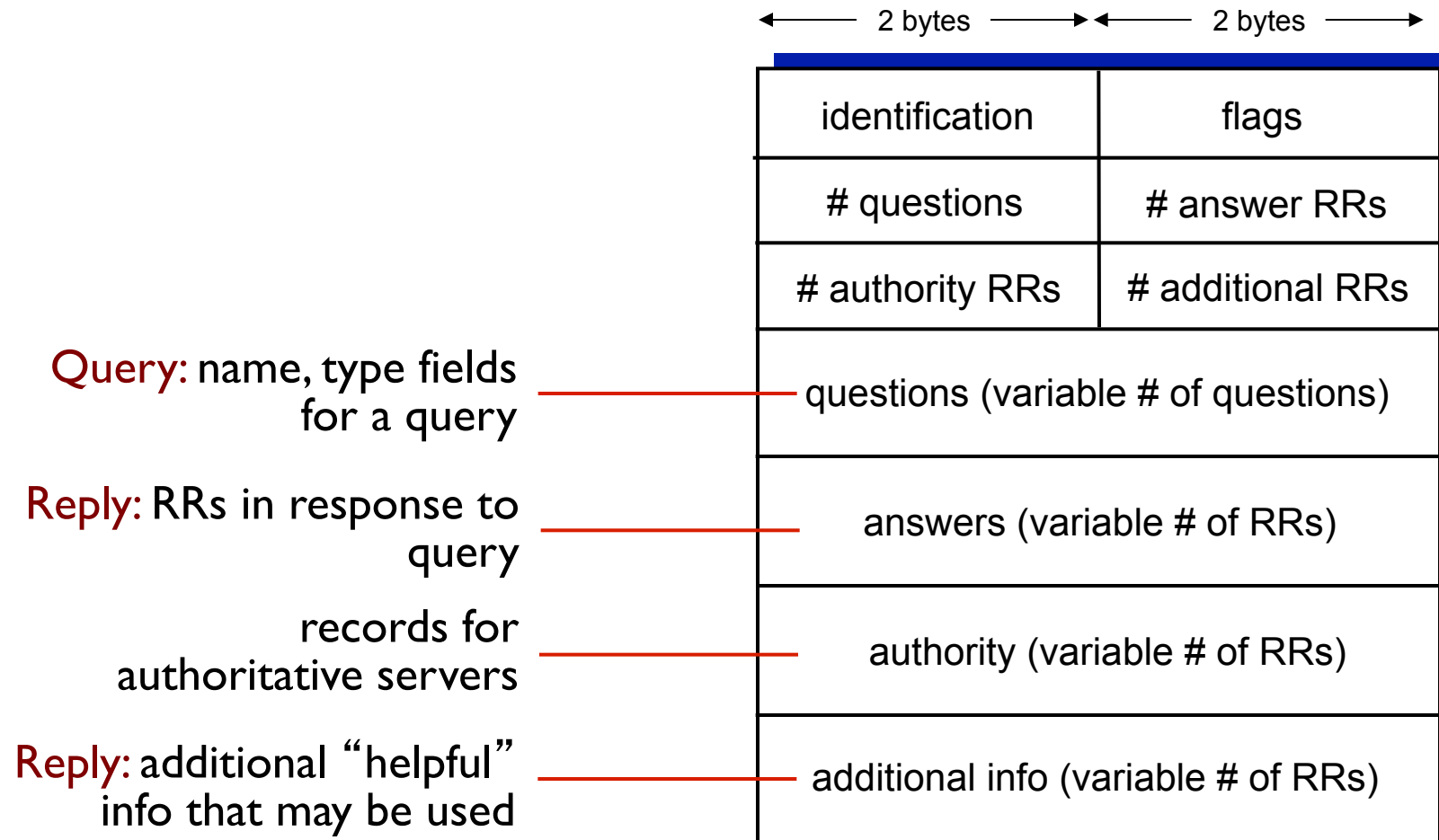
- ❖ *query* and *reply* messages, both with same *message format*

msg header

- ❖ **identification:** 16 bit # for query, reply to query uses same #
- ❖ **flags:**
 - query or reply
 - recursion desired (query)
 - recursion available (reply)
 - reply is authoritative (reply)
(DNS is an authoritative DNS to a queried name)



DNS protocol, messages



Inserting records into DNS



- ❖ example: new startup “Networkabc”
- ❖ register name networkabc.com at **DNS registrar** (e.g., Network Solutions) (and pay a fee for it.)
 - provide **names, IP addresses** of **authoritative name server** (primary and secondary)
 - registrar inserts two RRs into .com TLD server:
(networkabc.com, dns1.networkabc.com, NS)
(dns1.networkabc.com, 212.212.212.1, A)
- ❖ **Authoritative server**
 - create type A record for www.networkabc.com;
 - create type MX record for networkabc.com

Attacking DNS

DDoS attacks

- ❖ Bombard root servers with traffic
 - Not successful to date
 - Traffic Filtering
 - Local DNS servers cache IPs of TLD servers, allowing root server bypass
- ❖ Bombard TLD servers
 - Potentially more dangerous

Questions?

Quiz 4 and Lab 2

Quiz 4, 9/11, Tuesday

- Topic: DNS

Lab 2: DNS

Due 9/14 Friday at 23:59PM

Link:

<https://users.wpi.edu/~yli15/courses/CS3516Fall18A/Assignments.html>