

Exam #2 Solutions

1. **Iterative Optical Flow (30 pts):** An iterative method for computing optical flow updates $u(x, y), v(x, y)$ at each iteration according to

$$\begin{bmatrix} u(x, y) \\ v(x, y) \end{bmatrix}^{\text{new}} = \begin{bmatrix} \lambda I_x^2 + 4 & \lambda I_x I_y \\ \lambda I_x I_y & \lambda I_y^2 + 4 \end{bmatrix}^{-1} \begin{bmatrix} \sum_{n \in \text{neighbors}(x, y)} u^{\text{old}}(n) - \lambda I_x I_t \\ \sum_{n \in \text{neighbors}(x, y)} v^{\text{old}}(n) - \lambda I_y I_t \end{bmatrix}$$

Consider a local coordinate frame (x', y') where x' is aligned with the image gradient and y' is perpendicular to the image gradient. Likewise, $(u', v') = \left(\frac{dx'}{dt}, \frac{dy'}{dt}\right)$ are the image velocities in this frame. In this coordinate frame,

$$I_{x'} = \sqrt{I_x^2 + I_y^2} \text{ and } I_{y'} = 0$$

Show that the update equations in this coordinate frame

$$\begin{bmatrix} u'(x, y) \\ v'(x, y) \end{bmatrix}^{\text{new}} = \begin{bmatrix} \lambda I_{x'}^2 + 4 & \lambda I_{x'} I_{y'} \\ \lambda I_{x'} I_{y'} & \lambda I_{y'}^2 + 4 \end{bmatrix}^{-1} \begin{bmatrix} \sum_{n \in \text{neighbors}} u'^{\text{old}}(n) - \lambda I_{x'} I_t \\ \sum_{n \in \text{neighbors}} v'^{\text{old}}(n) - \lambda I_{y'} I_t \end{bmatrix}$$

reduce to

$$\begin{aligned} u'^{\text{new}}(x, y) &= \bar{u}'^{\text{old}} - \frac{I_{x'}^2 \bar{u}'^{\text{old}} + I_{x'} I_t}{I_{x'}^2 + \frac{4}{\lambda}} \\ v'^{\text{new}}(x, y) &= \bar{v}'^{\text{old}} \end{aligned}$$

Hint: This is not as hard as it looks.

$$\begin{aligned} \begin{bmatrix} u'(x, y) \\ v'(x, y) \end{bmatrix}^{\text{new}} &= \begin{bmatrix} \lambda I_{x'}^2 + 4 & \lambda I_{x'} I_{y'} \\ \lambda I_{x'} I_{y'} & \lambda I_{y'}^2 + 4 \end{bmatrix}^{-1} \begin{bmatrix} \sum_{n \in \text{neighbors}} u'^{\text{old}}(n) - \lambda I_{x'} I_t \\ \sum_{n \in \text{neighbors}} v'^{\text{old}}(n) - \lambda I_{y'} I_t \end{bmatrix} \\ &= \begin{bmatrix} \lambda I_{x'}^2 + 4 & 0 \\ 0 & 4 \end{bmatrix}^{-1} \begin{bmatrix} \sum_{n \in \text{neighbors}} u'^{\text{old}}(n) - \lambda I_{x'} I_t \\ \sum_{n \in \text{neighbors}} v'^{\text{old}}(n) \end{bmatrix} \\ &= \begin{bmatrix} \frac{1}{\lambda I_{x'}^2 + 4} & 0 \\ 0 & \frac{1}{4} \end{bmatrix} \begin{bmatrix} \sum_{n \in \text{neighbors}} u'^{\text{old}}(n) - \lambda I_{x'} I_t \\ \sum_{n \in \text{neighbors}} v'^{\text{old}}(n) \end{bmatrix} \\ u'^{\text{new}}(x, y) &= \frac{4\bar{u}'^{\text{old}} - \lambda I_{x'} I_t}{\lambda I_{x'}^2 + 4} = \frac{4\bar{u}'^{\text{old}} + \lambda I_{x'}^2 \bar{u}'^{\text{old}} - \lambda I_{x'}^2 \bar{u}'^{\text{old}} - \lambda I_{x'} I_t}{\lambda I_{x'}^2 + 4} \\ &= \bar{u}'^{\text{old}} \frac{4 + \lambda I_{x'}^2}{\lambda I_{x'}^2 + 4} - \lambda \frac{I_{x'}^2 \bar{u}'^{\text{old}} - I_{x'} I_t}{\lambda I_{x'}^2 + 4} = \bar{u}'^{\text{old}} - \frac{I_{x'}^2 \bar{u}'^{\text{old}} + I_{x'} I_t}{I_{x'}^2 + \frac{4}{\lambda}} \\ v'^{\text{new}}(x, y) &= \bar{v}'^{\text{old}} \end{aligned}$$

2. **Object Representation (40 pts):** We can represent an object by its boundary $(x(s), y(s)), 0 \leq s \leq S$ where S is the length of the object's boundary and s is distance along that boundary from some arbitrary starting point. We can combine x and y into a single complex function $z(s) = x(s) + jy(s)$. The Discrete Fourier Transform (DFT) of z is

$$Z(k) = \sum_{s=0}^{S-1} e^{-2\pi j \frac{ks}{S}} z(s), 0 \leq k \leq S-1$$

We can use the coefficients $Z(k)$ to represent the object boundary. The limit on s is $S-1$ because for a closed contour $z(s) = z(0)$. The Inverse Discrete Fourier Transform is

$$z(s) = \frac{1}{S} \sum_{k=0}^{S-1} e^{+2\pi j \frac{ks}{S}} Z(k), 0 \leq k \leq S-1$$

- a. (10 pts) Suppose that the object is translated by $(\Delta x, \Delta y)$, that is, $z'(s) = z(s) + \Delta x + j\Delta y$. How is z' 's DFT $Z'(k)$ related to $Z(k)$?

Only the DC $k=0$ component is affected. $Z'(\mathbf{0}) = Z(\mathbf{0}) + S(\Delta x + j\Delta y)$, $Z'(\mathbf{k}) = Z(\mathbf{k})$ for $\mathbf{k} > \mathbf{0}$

- b. (15 pts) Suppose that the object is scaled by integer constant c , that is, $z'(s) = cz(s)$. For simplicity, assume that $S' = S$. How is z' 's DFT $Z'(k)$ related to $Z(k)$?

$$Z'(\mathbf{k}) = \sum_{s=0}^{S'-1} e^{-2\pi j \frac{ks}{S'}} z'(s) = \sum_{s=0}^{S-1} e^{-2\pi j \frac{ks}{S}} cz(s) = c \sum_{s=0}^{S-1} e^{-2\pi j \frac{ks}{S}} z(s) = cZ(\mathbf{k})$$

- c. (10 pts) What object has $z(s) = [x_0 + R \cos \frac{2\pi s}{S}] + j[y_0 + R \sin \frac{2\pi s}{S}]$? Sketch it.

A circle with radius R centered at (x_0, y_0) .

- d. (5 pts) What is $Z(k)$ corresponding to $z(s)$ from Part c? Hint: Most coefficients are 0.

$$Z(\mathbf{0}) = S(x_0 + jy_0), Z(\mathbf{1}) = SR, Z(\mathbf{k}) = \mathbf{0} \text{ for } \mathbf{k} > \mathbf{1}$$

3. **Chain Code (30 pts):** An object boundary is represented by an 8-connected chain code C with values $c_i \in \{0, \dots, 7\}$. Each c_i represents a segment at $45^\circ = \pi/4$ radians, with $c_i = 0$ representing $(1, 0)$, $c_i = 1$ representing $(1, 1)$, $c_i = 2$ representing $(0, 1)$, etc. Note that segments have length 1 when c_i is even, and length $\sqrt{2}$ when c_i is odd. Computer Vision researcher Ima Robot claims that the chain code for a closed contour of length N must obey Eq (1):

$$\sum_{i=0}^{N-1} e^{-j\frac{\pi}{4}c_i} = 0$$

- a) (15 pts) Show that Eq (1) holds for a square

$$C = [0, 0, 2, 2, 4, 4, 6, 6]$$

and octagon

$$C = [0, 1, 2, 3, 4, 5, 6, 7]$$

Square: $\mathbf{1} + \mathbf{1} + j + j - \mathbf{1} - \mathbf{1} - j - j = \mathbf{0}$. **Octagon:** $\mathbf{1} + \frac{1+j}{\sqrt{2}} + j + \frac{-1+j}{\sqrt{2}} - \mathbf{1} - \frac{1+j}{\sqrt{2}} - j - \frac{-1+j}{\sqrt{2}} = \mathbf{0}$.

- b) (15 pts) Show that Eq (1) holds for any closed contour or provide a counterexample.

Counterexample: Triangle = [1, 4, 6].  $\frac{1+j}{\sqrt{2}} - \mathbf{1} - j \neq \mathbf{0}$

4. **Teamwork (1 pt):** On a scale of 1 to 5, with 1 being the lowest, 3 is neutral and 5 being the highest, rate how well your project team is working together.

No right / wrong answer.