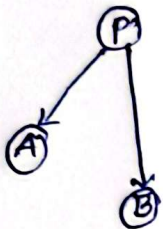


1(a)



1(b) This is parent

3  
2  
1  
0

This is a child

This is another child

This is another program

This is parent again.

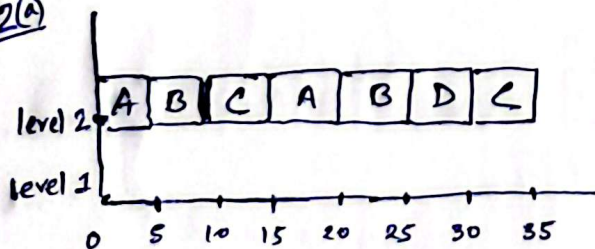
1(c) \* shell figures out the path of the executable based on users command.

\* calls fork to create a child process to run the command.

\* calls some variant of exec to run the cmd.

\* waits for the command to complete by calling wait().

2(a)



B<sub>2</sub>: A, B, C, A, B, D, C

B<sub>1</sub>: A, B, C

2(b) Avg. time =  $\frac{20 + 25 + (30 - 14) + 35}{4}$   
= 24 ms

2(c) Yes, because I/O bound processes have a small CPU burst and hence a higher f.

3(a)

def uni-shuttle-thread:

lock(&m)

shuttle-arrived = true

int n = min(waiting-count, K)

for(int i=1; i<=n; i+):

signal(&cv-shuttle-arrived)

wait(&cv-student-boarded, &n)

waiting-count -= n

shuttle-arrived = false

depart()

broadcast-signal(&cv-queue-open)

unlock(&m)

3(b) Problem: Both threads can grab the lock.

Thread 1

while(\*lock);

Thread 1 Descheduled

Thread 2

while(\*lock);

\*lock = true;

Thread 2 Desched.

\*lock = true

Soln

TestAndSet function which works as an atomic instruction for checking availability of lock and acquiring it.

3(c)

So that threads can have separate execution state and run independently.

2(d) No, processes with smaller CPU burst will get scheduled more number of times.

2(e) Yes, a CPU bound process with a low priority can starve if higher priority I/O bound processes keep arriving continuously.