

31.12.24

Operating System

Virtual Memory = 2^{48} Bytes

$$1KB = 2^{10} B$$

$$\approx 2^9 \cdot 2^{40} \text{ Bytes}$$

$$1MB = 2^{20} B$$

$$= 256 \text{ TB}$$

$$1GB = 2^{30} B$$

$$1TB = 2^{40} B$$

a) Offset size = \log_2 (page size in Bytes)

$$= \log_2 (4KB)$$

$$= \log_2 (4 \times 1024)$$

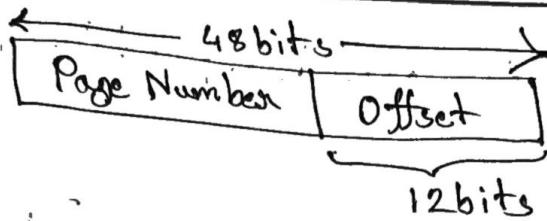
$$= 12 \text{ bits}$$

b) No. of virtual pages = $\frac{\text{Virtual Mem. Size}}{\text{Page size}}$

$$= \frac{2^{48} \text{ Bytes}}{2^{12} \text{ Bytes}}$$

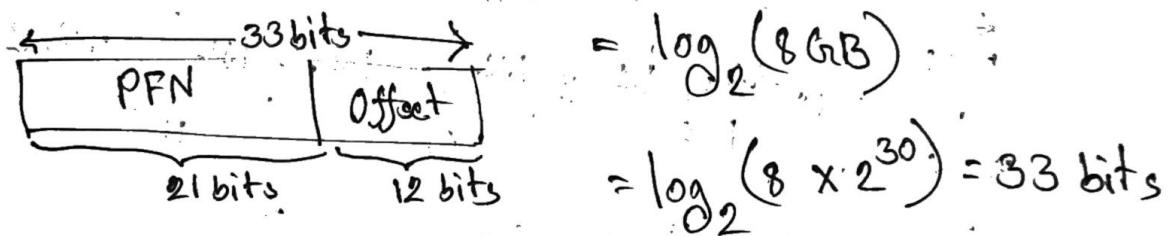
$$= 2^{36} \text{ Bytes}$$

c)



$$\text{No. of bits for VPN} = 48 - 12 = 36 \text{ bits}$$

d) Physical address = $\log_2(\text{Physical Mem. Size})$



$$\therefore \text{Number of bits for PFN} = 33 - 12 = 21$$

e) Number of PTEs = No. of Virtual Pages

$$= \frac{\text{Virtual Memory Size}}{\text{Page Size}}$$
$$= \frac{2^{48}}{2^{12}} = 2^{36} \text{ Bytes}$$

$$\therefore \text{Page table size} = \text{No. of PTEs} \times \text{PTE size}$$

$$= 2^{36} \times 4 \text{ bytes}$$

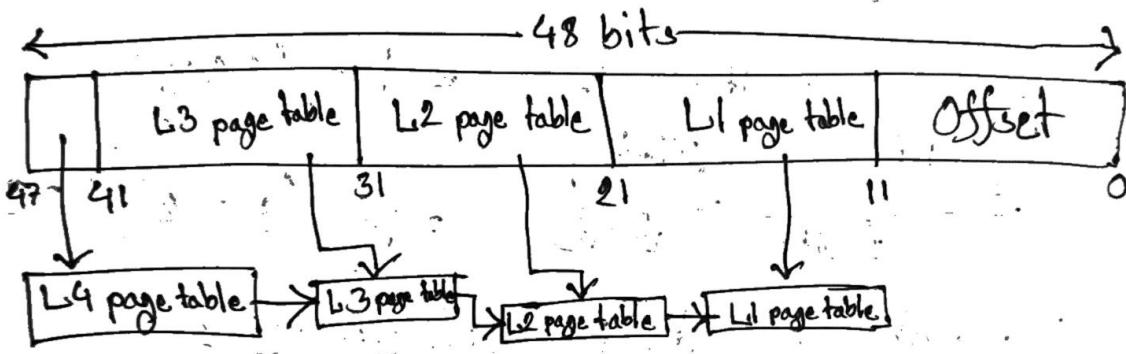
$$= 2^{30} \times 2^6 \times 2^2 \text{ bytes}$$

$$= 1 \text{ MB} \times 256 = 256 \text{ MB} \cancel{\text{GB}}$$

c) No. of bits for each level page table = $\log_2 \left(\frac{\text{Page Size}}{\text{PTE size}} \right)$

$$= \log_2 \left(\frac{4 \times 1024 \text{ bytes}}{4 \text{ bytes}} \right)$$

$$= 10 \text{ bits}$$



4.c) Using data from ①,

Maximum required pages for L2 page table,

$$= \frac{\text{No. of PTEs}}{\text{No. of PTEs in each page}}$$

$$= \frac{2^{36}}{2^{10}} = 2^{26} = 2^6 \cdot 2^{20}$$

~~in 64MB pages~~

No of PTEs in each page = $\frac{2^{12}}{4} = 2^{10}$
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Maximum required pages for L2 page table = $\frac{2^{26}}{2^{10}} = 2^{16}$
= ~~4K~~ 4096 pages

Similarly, for L3 = $\frac{2^{16}}{2^{10}} = 2^6$ ~~and 4 pages~~ pages

$$\therefore L4 = \left\lceil \frac{2^6}{2^{10}} \right\rceil = 1$$

∴ Total pages required = $2^{26} + 2^{16} + 2^6 + 1$
= $2^{48} + 1$

Minimum page requirement for the multi-level page
tables = Number of levels = 4