Chapter 3

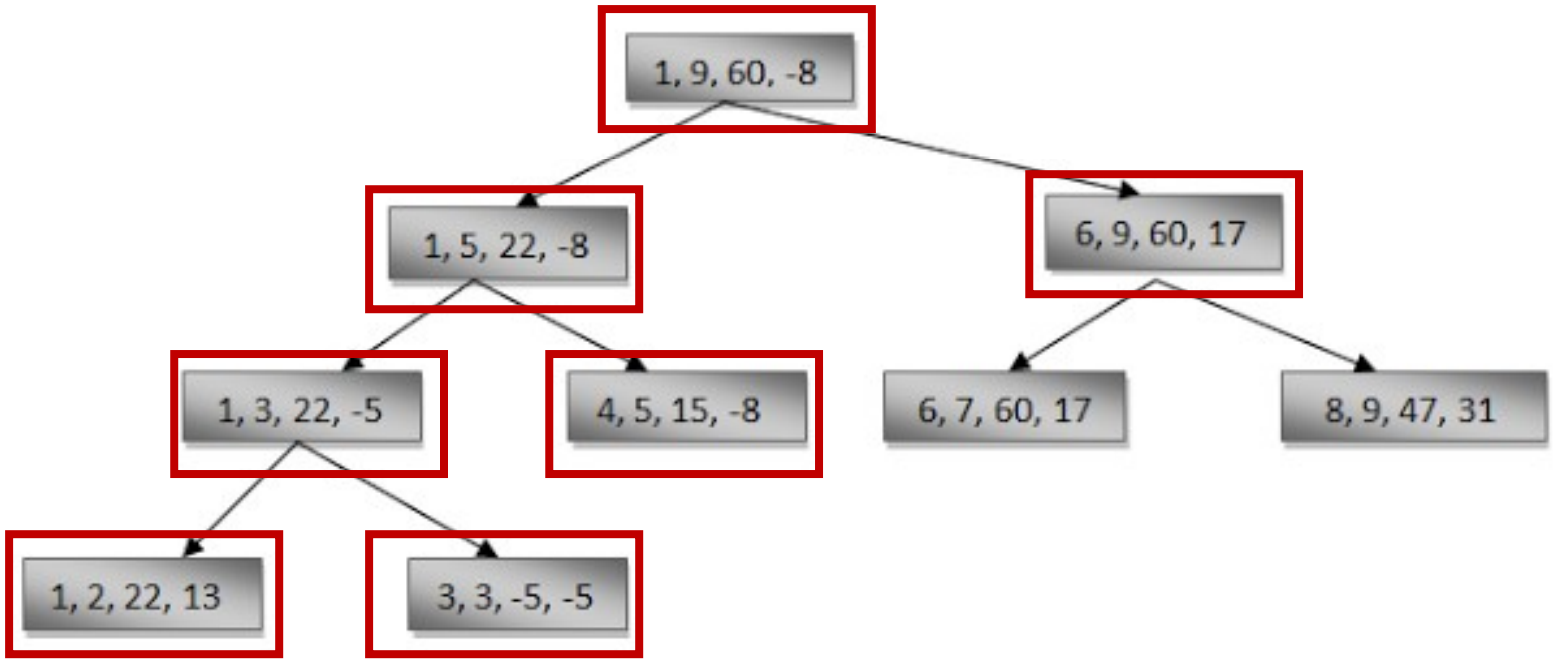
Divide-and-Conquer

MaxMin Algorithm

```
MaxMin(i, j, max, min)
{
    if (i=j) then max := min := a[i]; //Small(P)
    else if (i=j-1) then // Another case of Small(P)
        {
            if (a[i] < a[j]) then max := a[j]; min := a[i];
            else max := a[i]; min := a[j];
        }
    else
    {
        // if P is not small, divide P into sub-problems. Find where to split the set.
        mid := ( i + j )/2;     // Solve the sub-problems.
        MaxMin( i, mid, max, min );
        MaxMin( mid+1, j, max1, min1 );
        // Combine the solutions.
        if (max < max1) then max := max1;
        if (min > min1) then min := min1;
    }
}
```

A simple example

- Finding the maximum and minimum of a set **S** of **n** numbers



Time complexity

$T(n)=2T(n/2)+2T(n/2)+2 \quad \text{if } n>=2$

$=1 \quad \text{if } n=2$

$=0 \quad \text{if } n=1$

- Calculation of T(n):

Assume $n = 2^k$,

$T(n) = 2T(n/2)+2$

$= 2(2T(n/4)+2)+2$

$= 4T(n/4)+4+2$

:

$=2^{k-1}T(2)+2^k-2$

$=3n/2-2$

Merge Sort

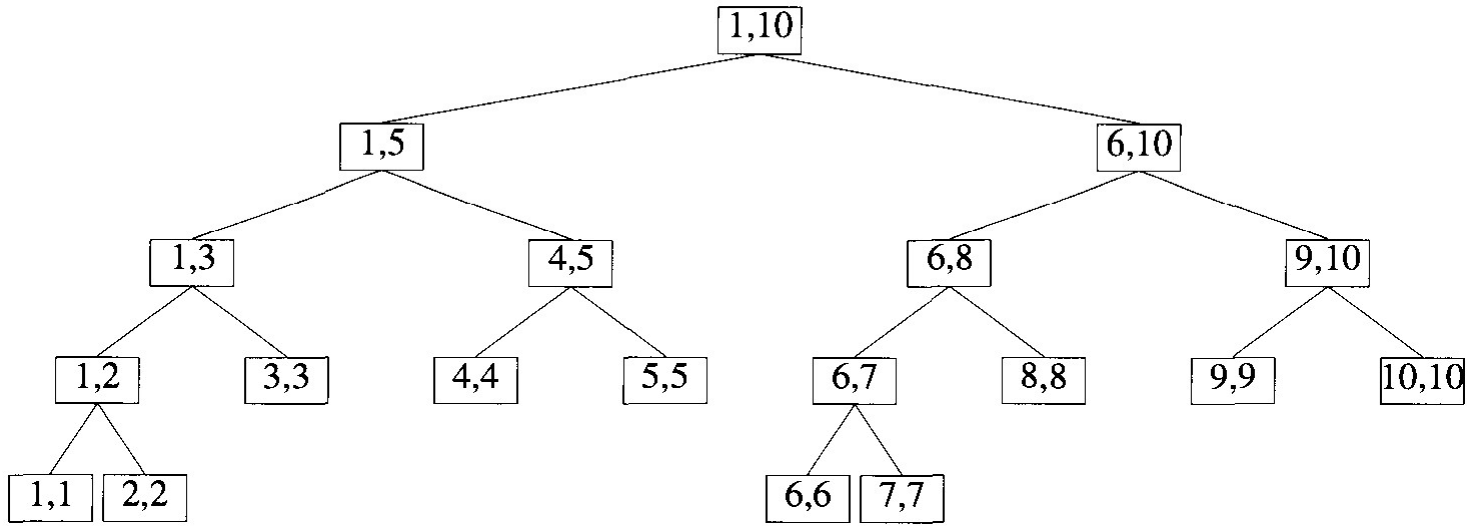
```
1   Algorithm MergeSort(low, high)
2   // a[low : high] is a global array to be sorted.
3   // Small(P) is true if there is only one element
4   // to sort. In this case the list is already sorted.
5   {
6       if (low < high) then  // If there are more than one element
7       {
8           // Divide P into subproblems.
9               // Find where to split the set.
10                  mid := ⌊(low + high)/2⌋;
11          // Solve the subproblems.
12              MergeSort(low, mid);
13              MergeSort(mid + 1, high);
14          // Combine the solutions.
15              Merge(low, mid, high);
16      }
17  }
```

Algorithm: Merge Element

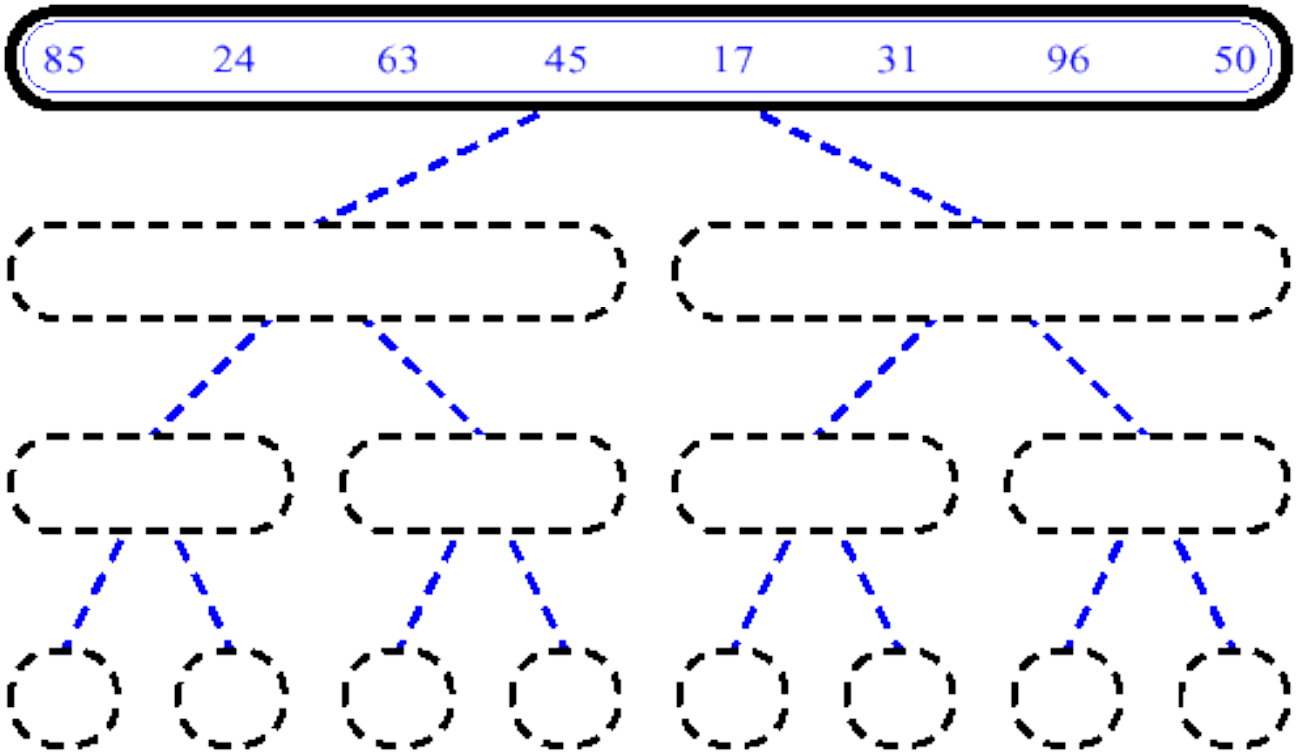
```
1   Algorithm Merge(low, mid, high)
2   // a[low : high] is a global array containing two sorted
3   // subsets in a[low : mid] and in a[mid + 1 : high]. The goal
4   // is to merge these two sets into a single set residing
5   // in a[low : high]. b[ ] is an auxiliary global array.
6   {
7       h := low; i := low; j := mid + 1;
8       while ((h ≤ mid) and (j ≤ high)) do
9       {
10          if (a[h] ≤ a[j]) then
11          {
12              b[i] := a[h];h := h + 1;
13          }
14          else
15          {
16              b[i] := a[j]; j := j + 1;
17          }
18          i := i + 1;
19      }
20      if (h > mid) then
21          for k := j to high do
22          {
23              b[i] := a[k]; i := i + 1;
24          }
25      else
26          for k := h to mid do
27          {
28              b[i] := a[k]; i := i + 1;
29          }
30      for k := low to high do a[k] := b[k];
31  }
```

Tree of Calls of Merge Sort

$$(310 \mid 285 \mid 179 \mid 652, 351 \mid 423, 861, 254, 450, 520)$$



Merge Sort (Example) Cont..



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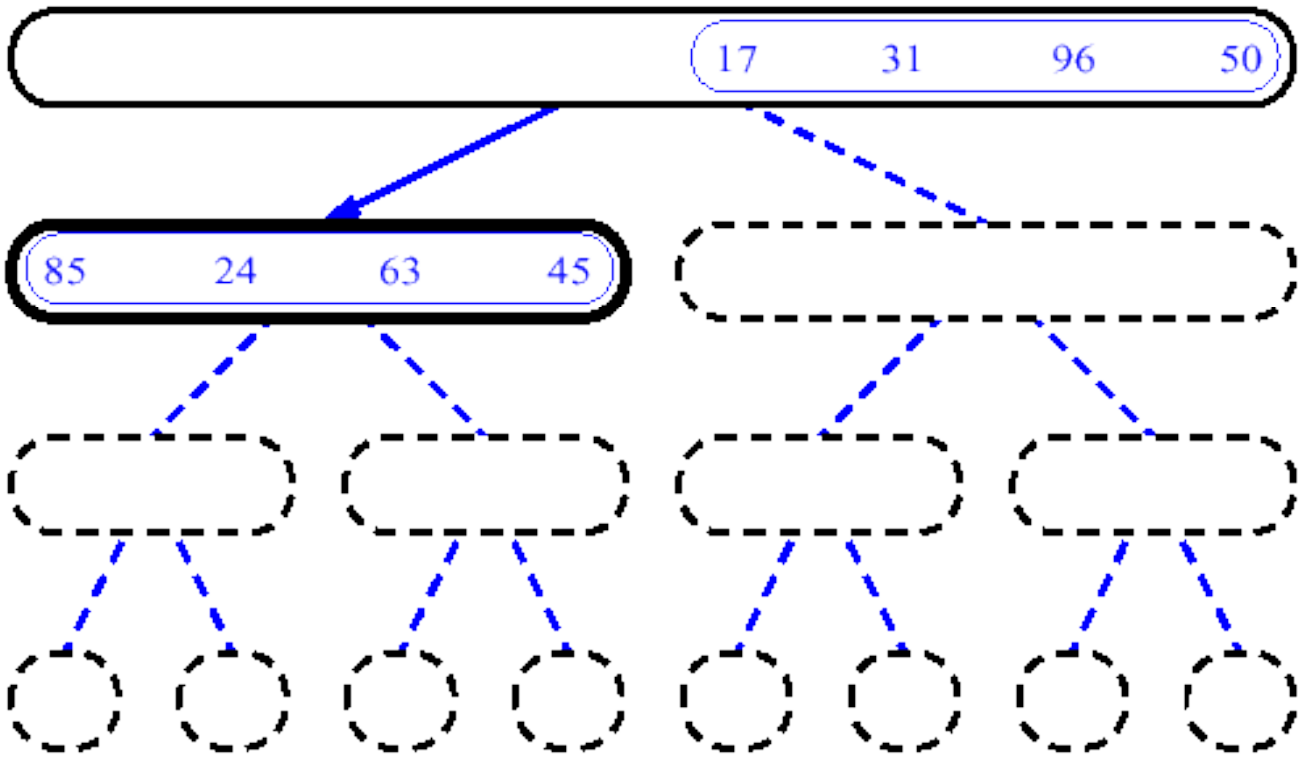

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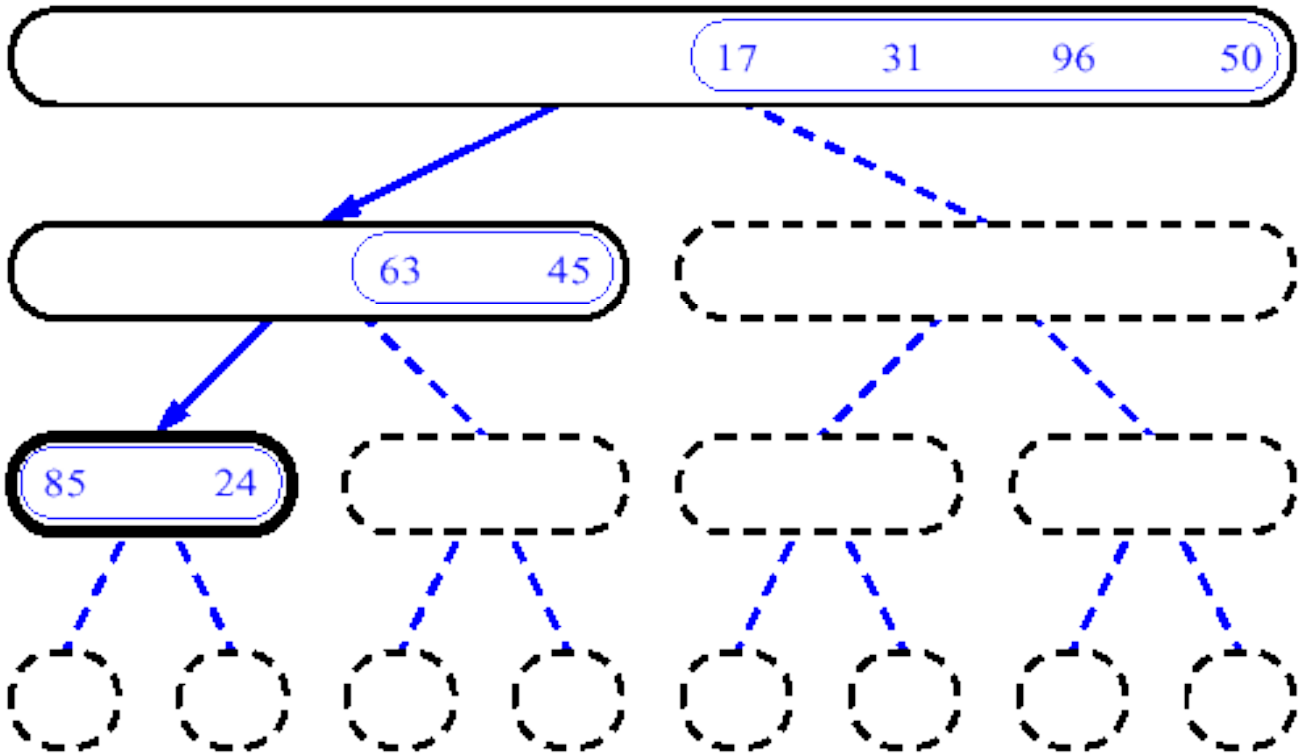

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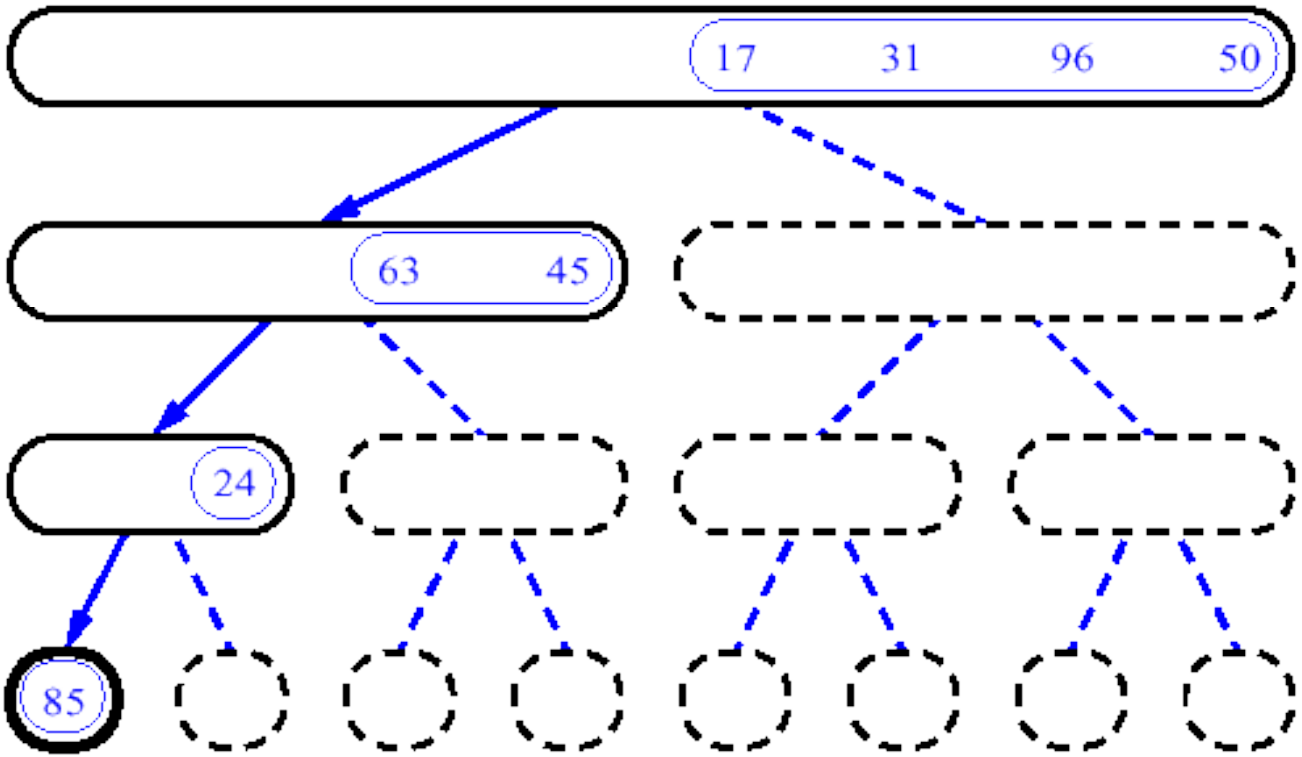

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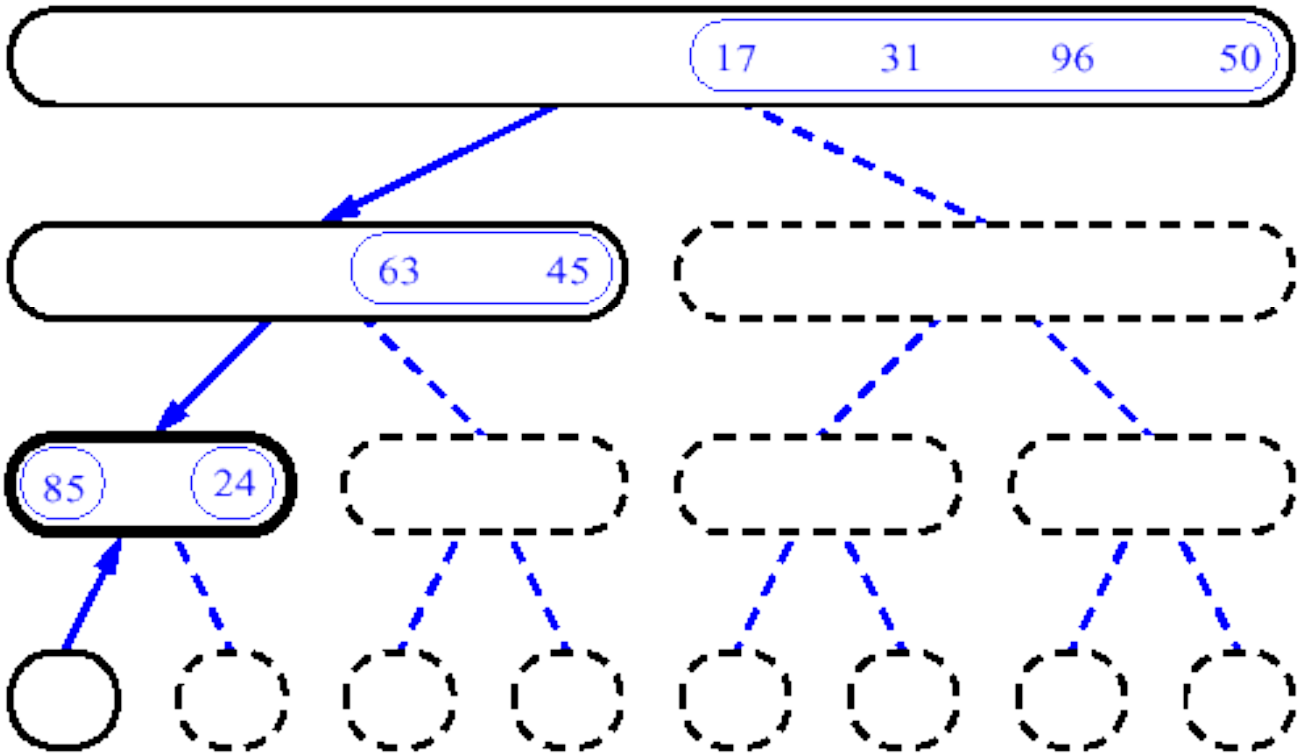

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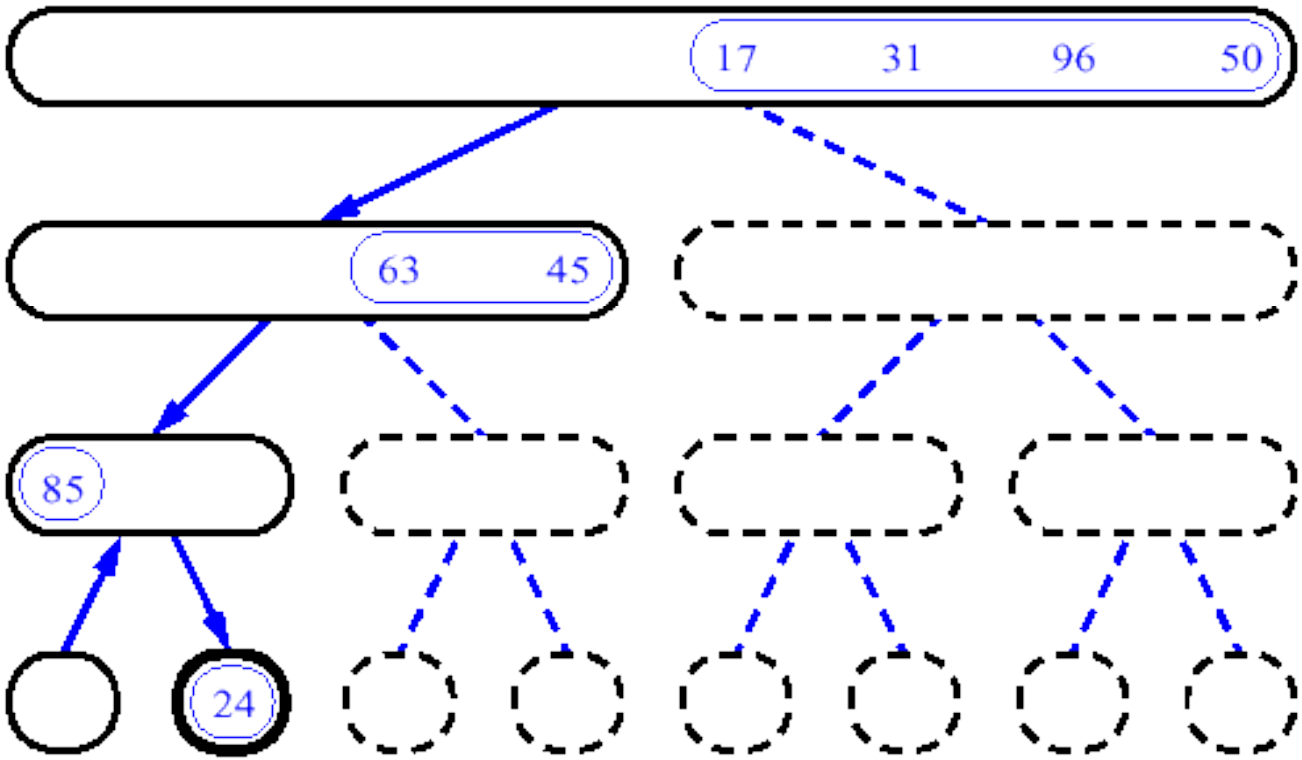

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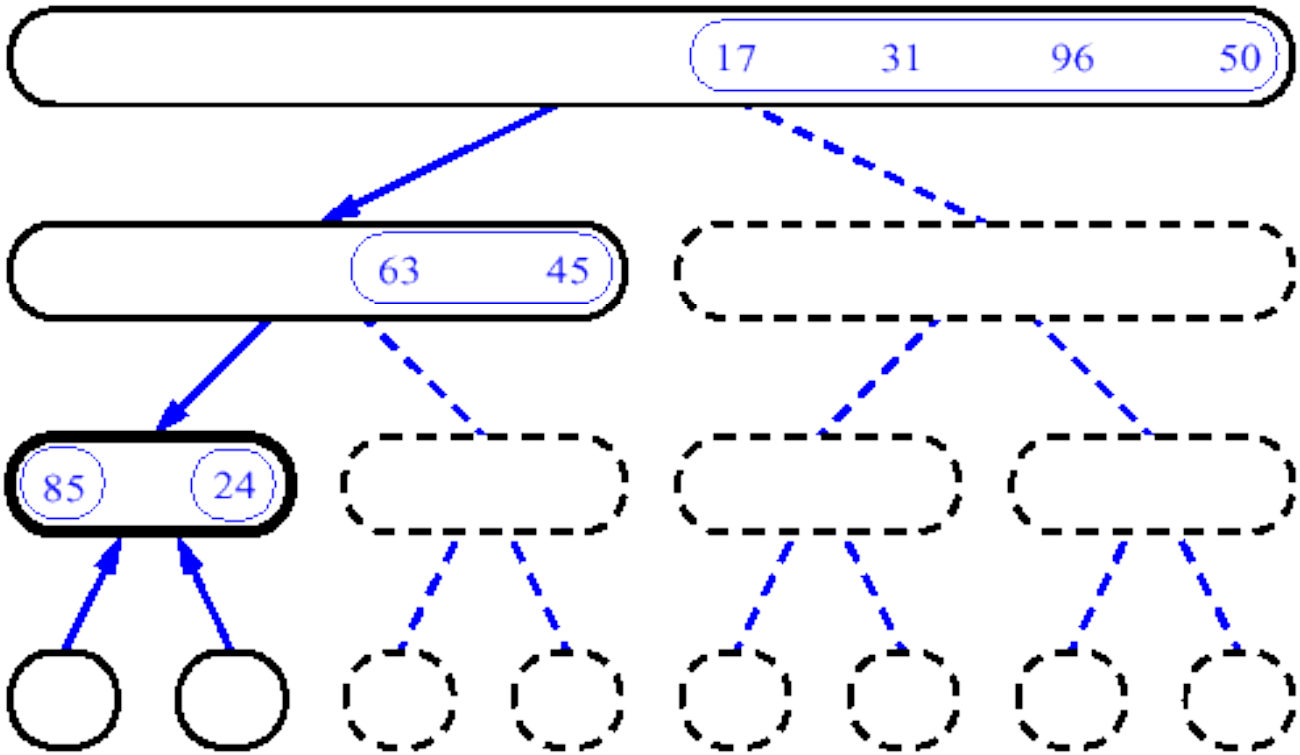

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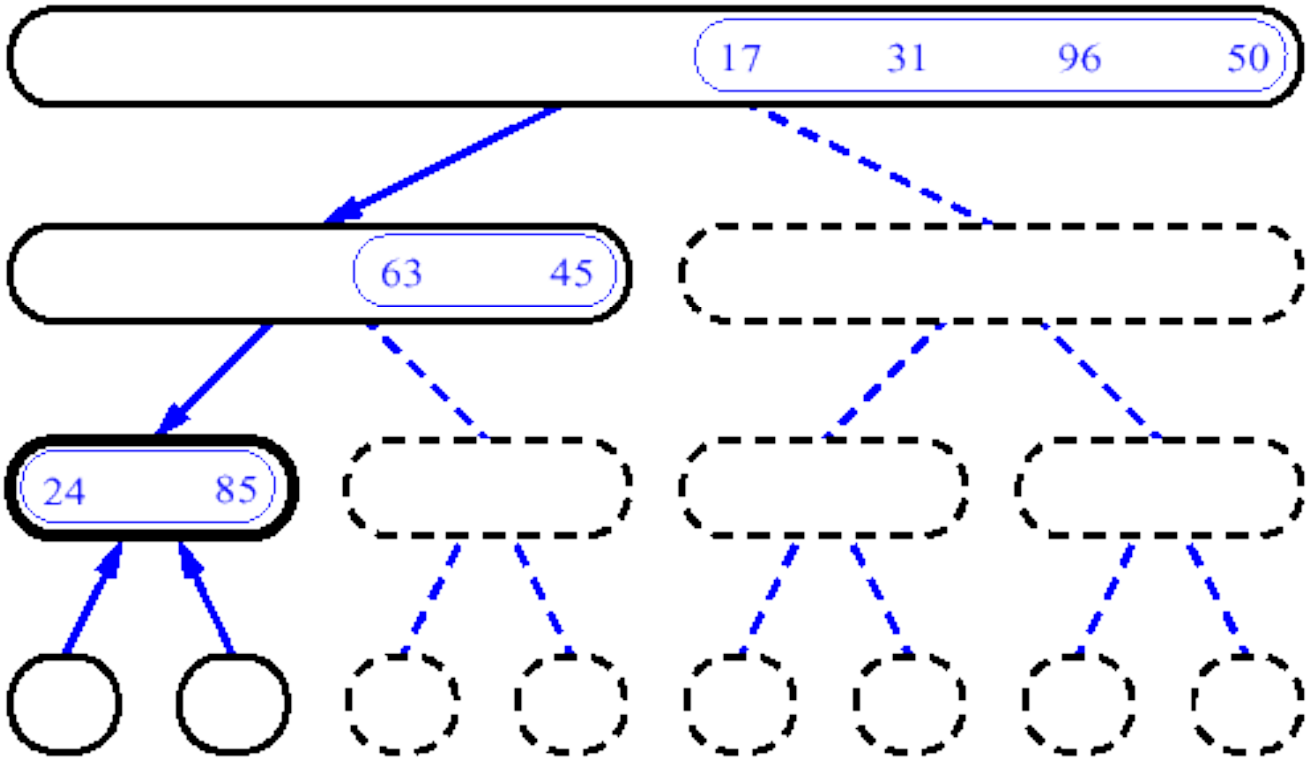

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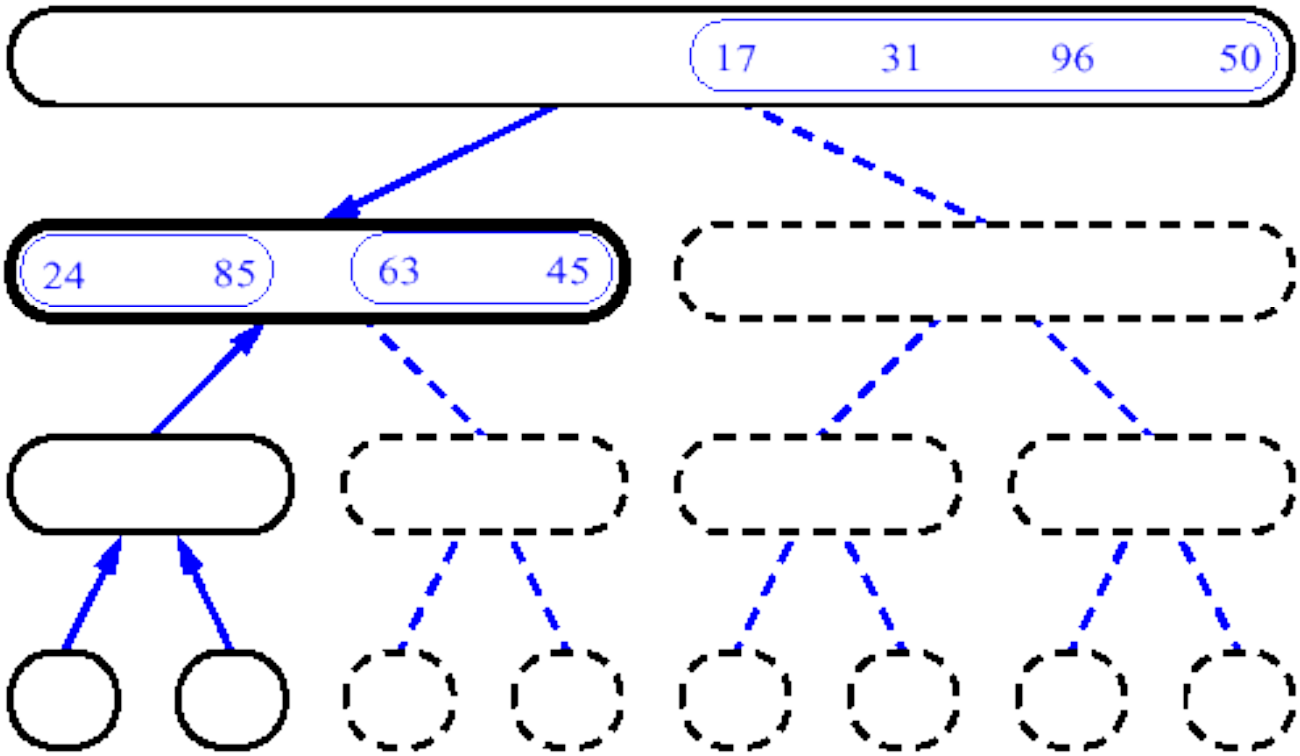

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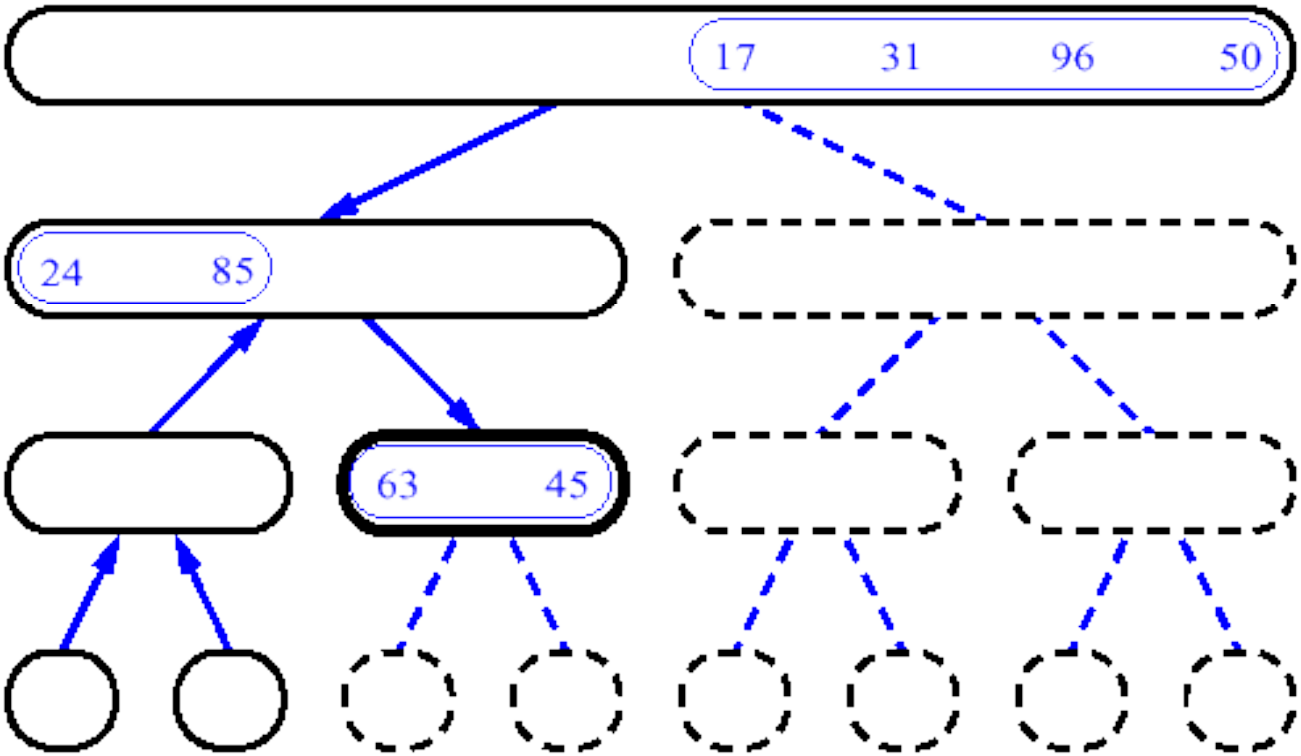

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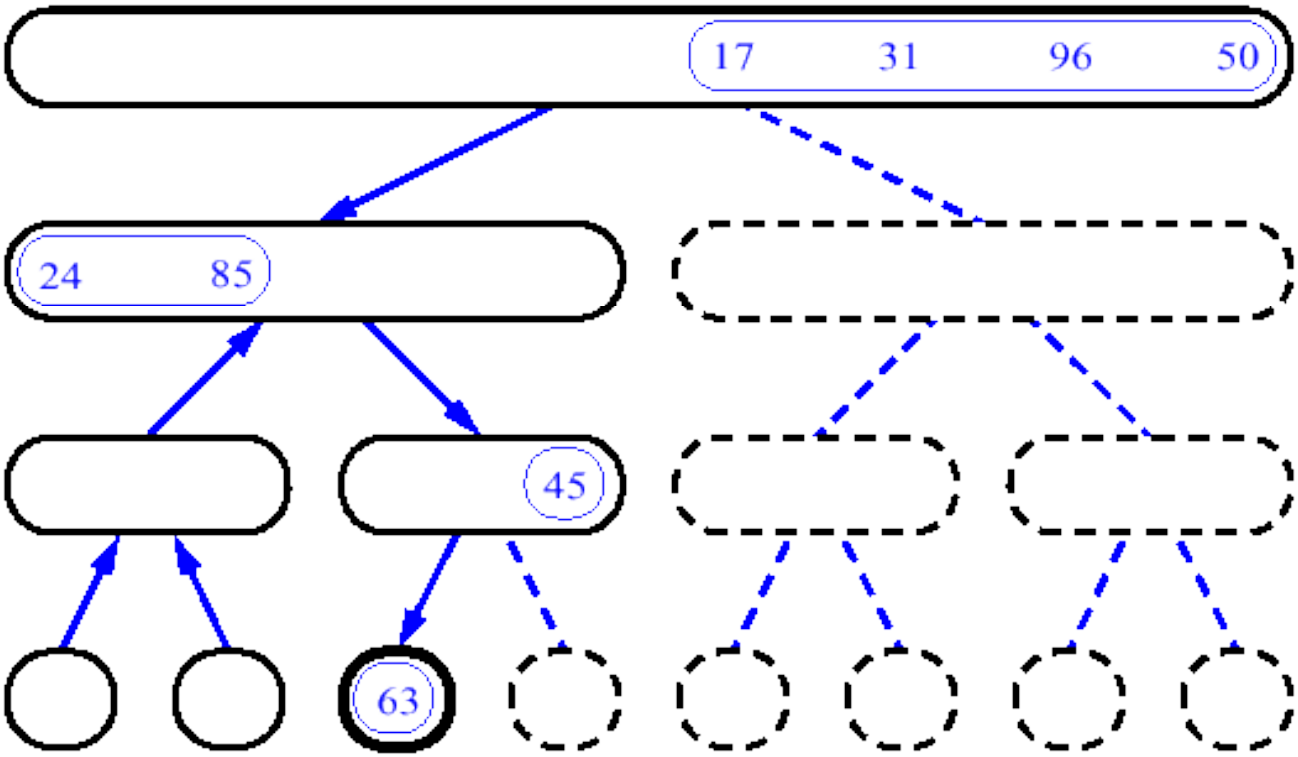

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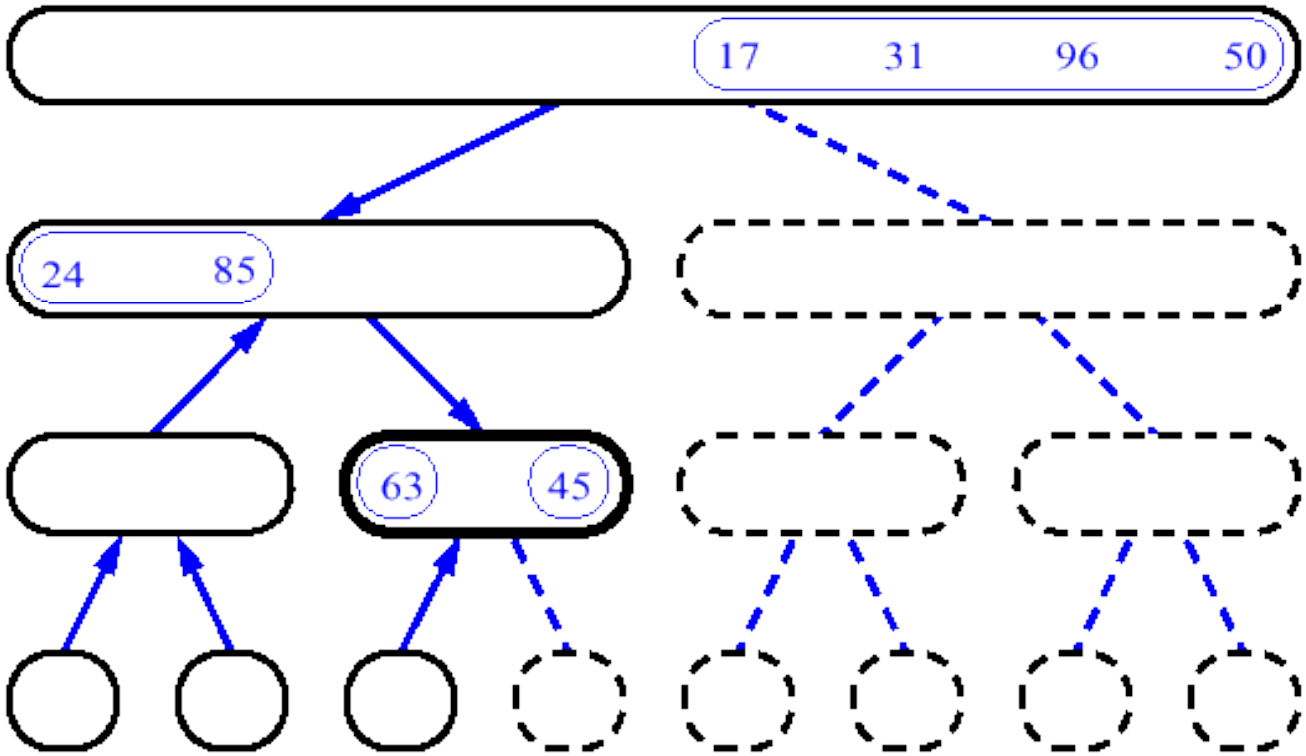

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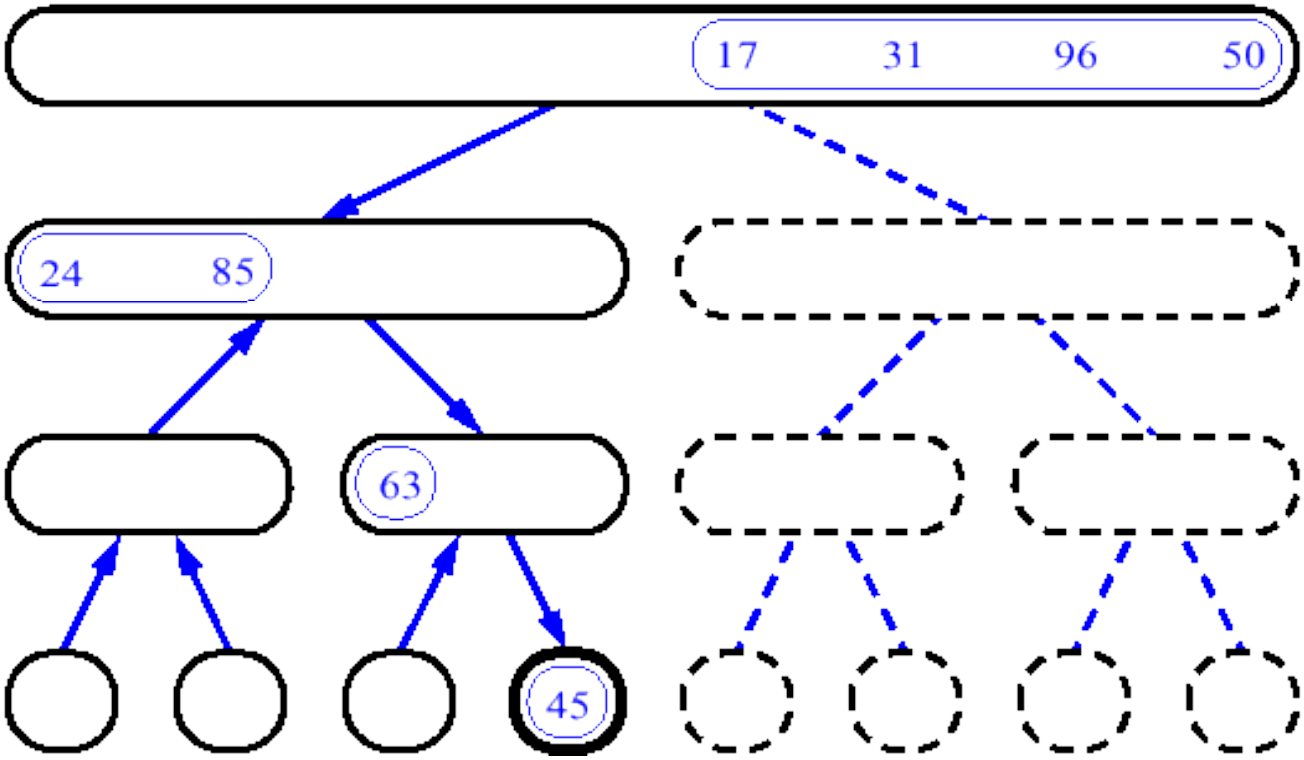

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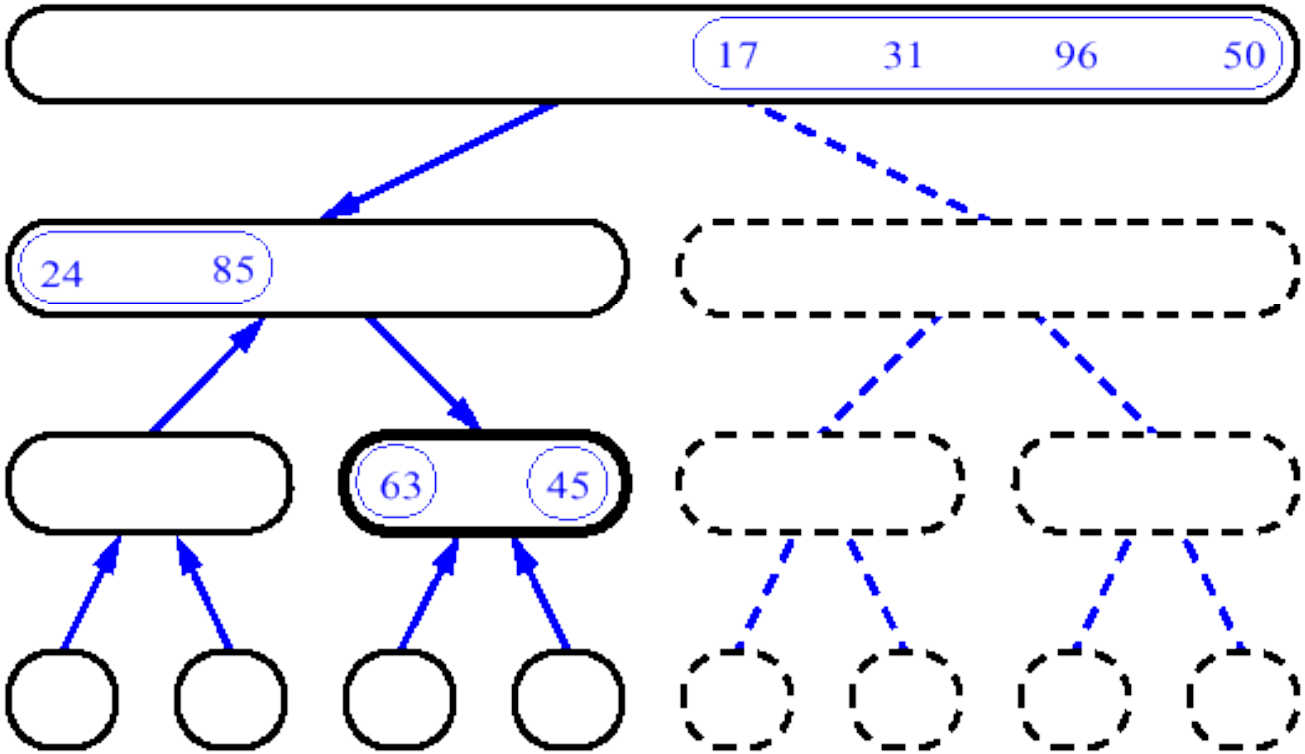

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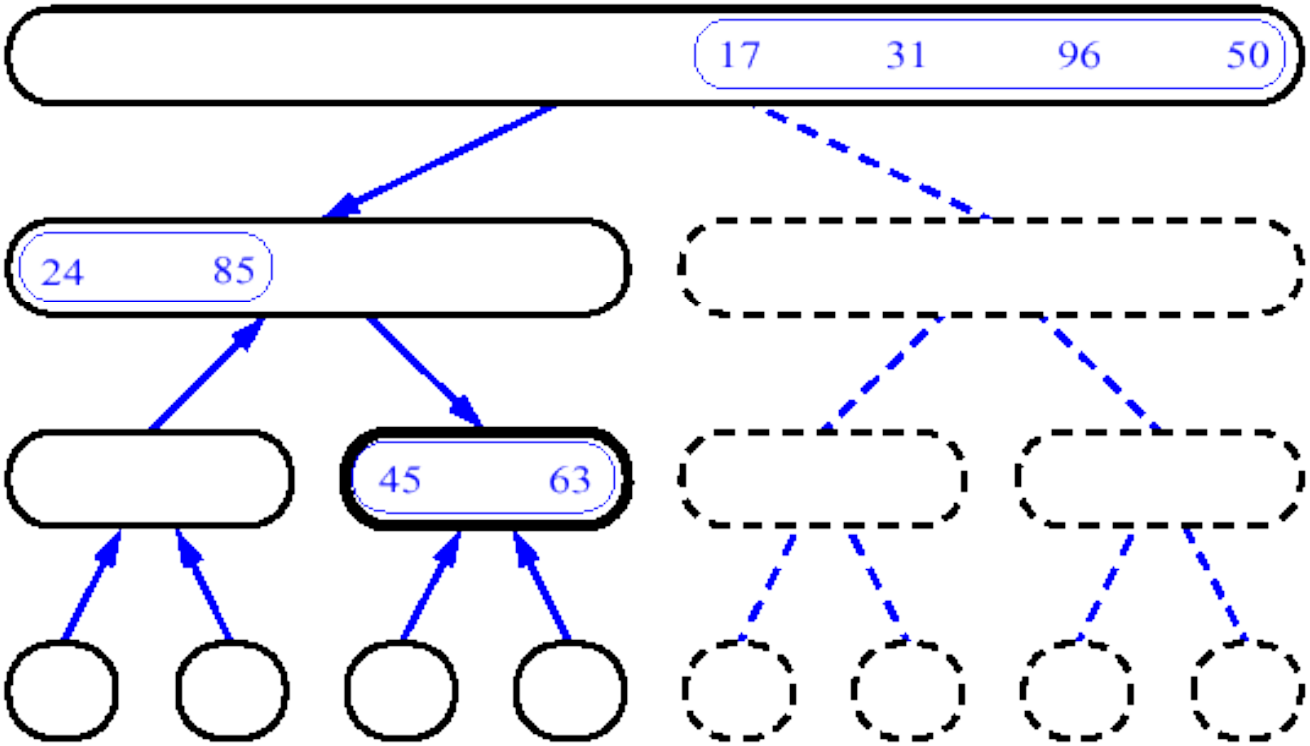

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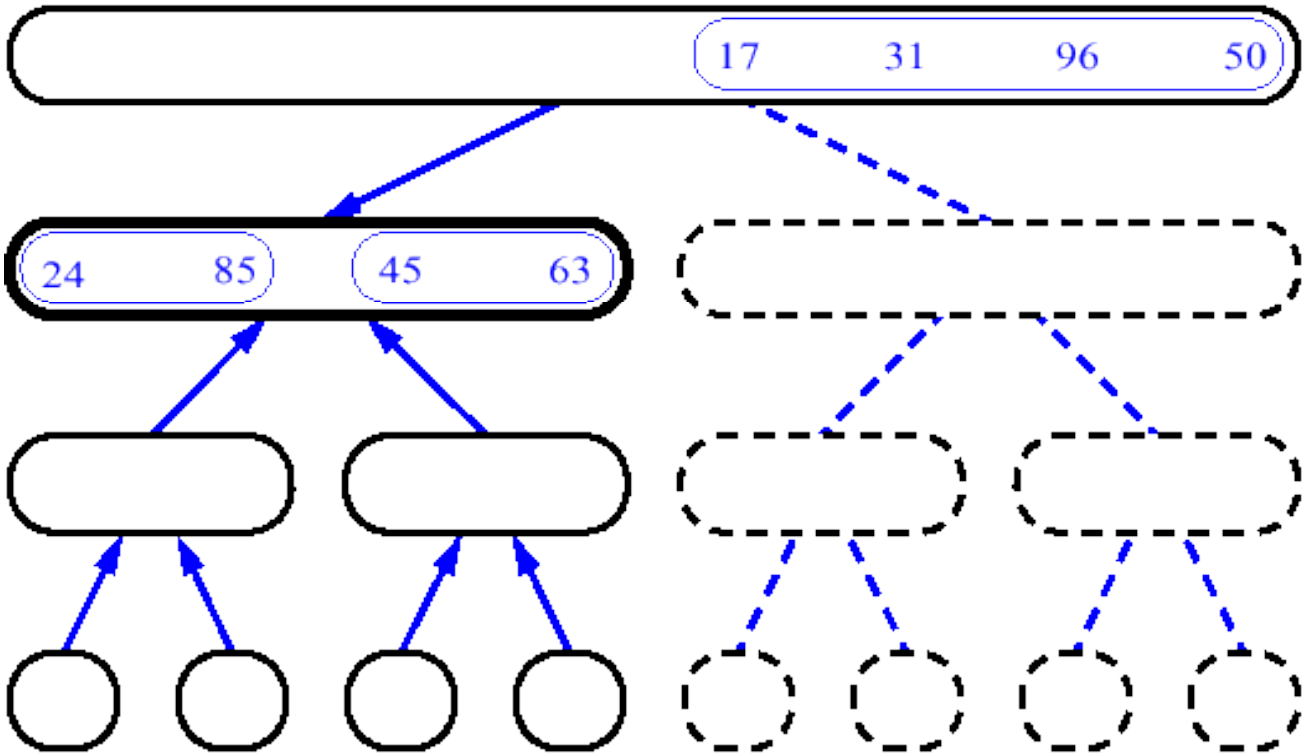

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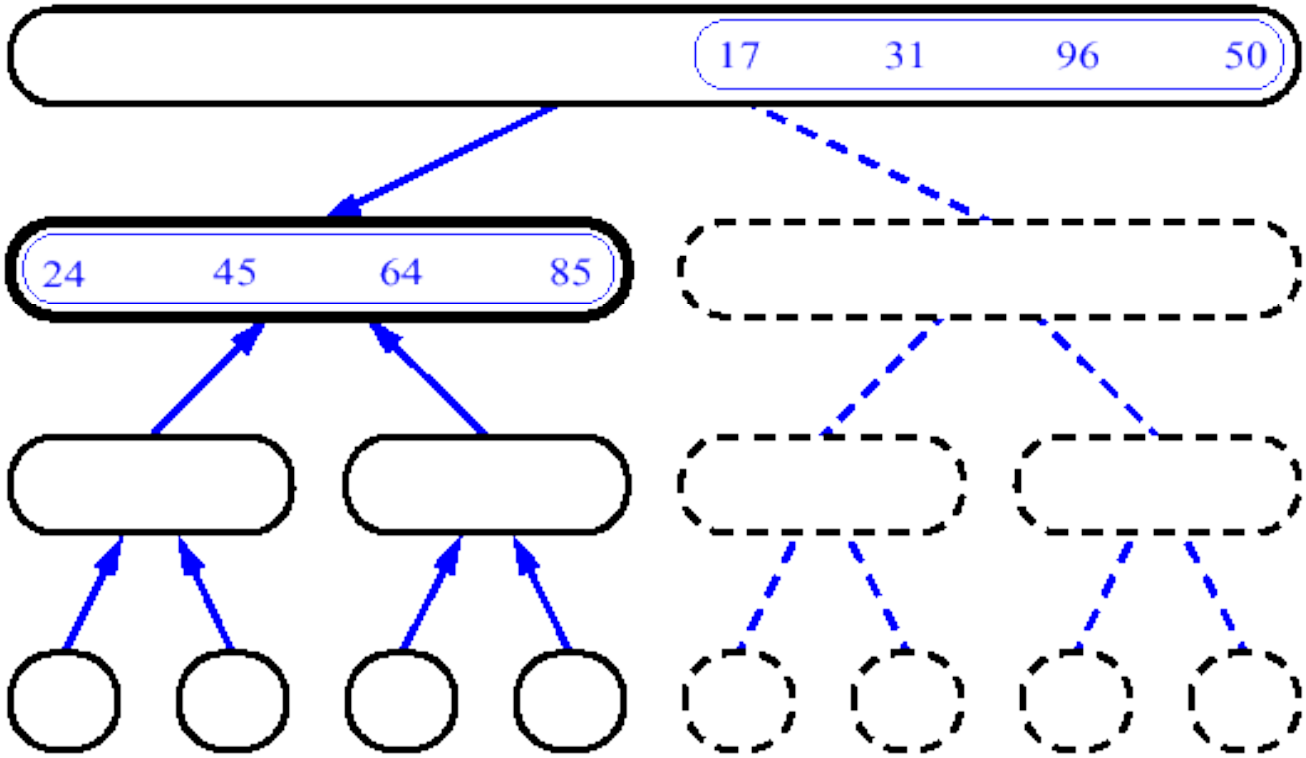

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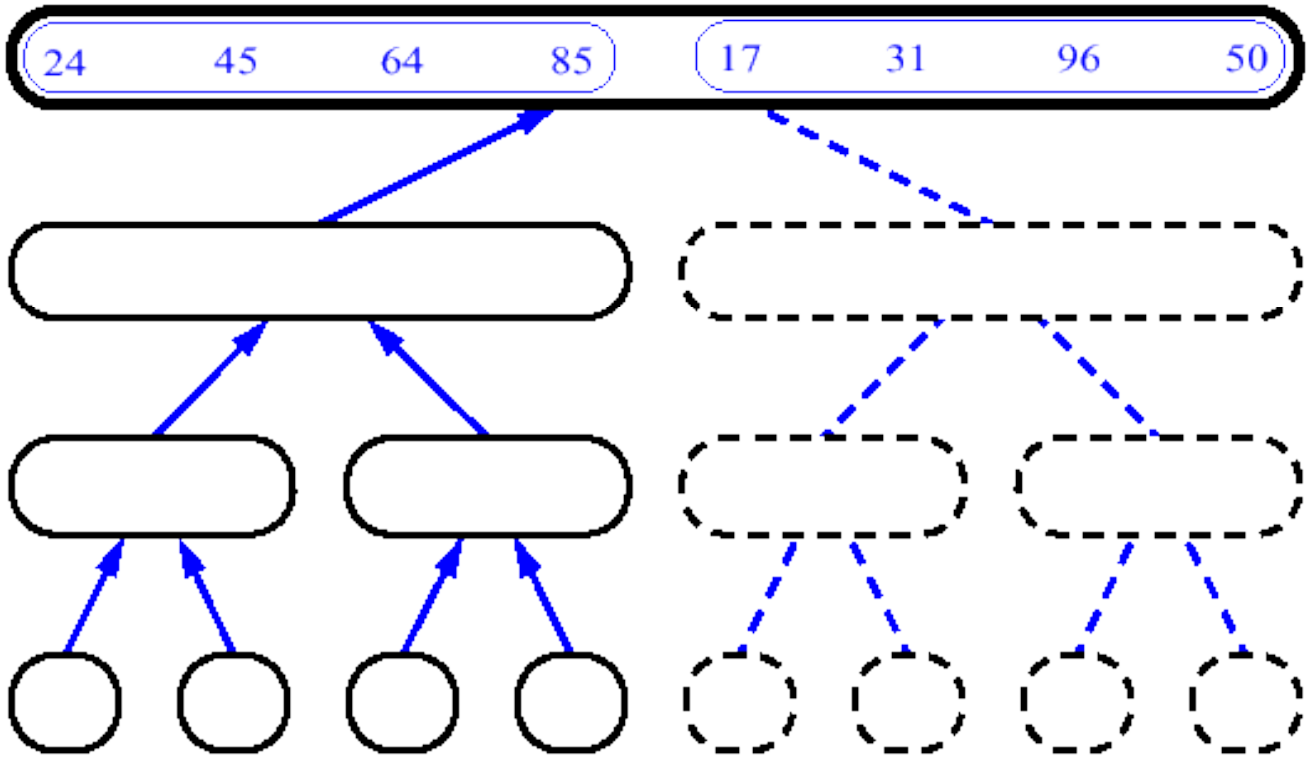

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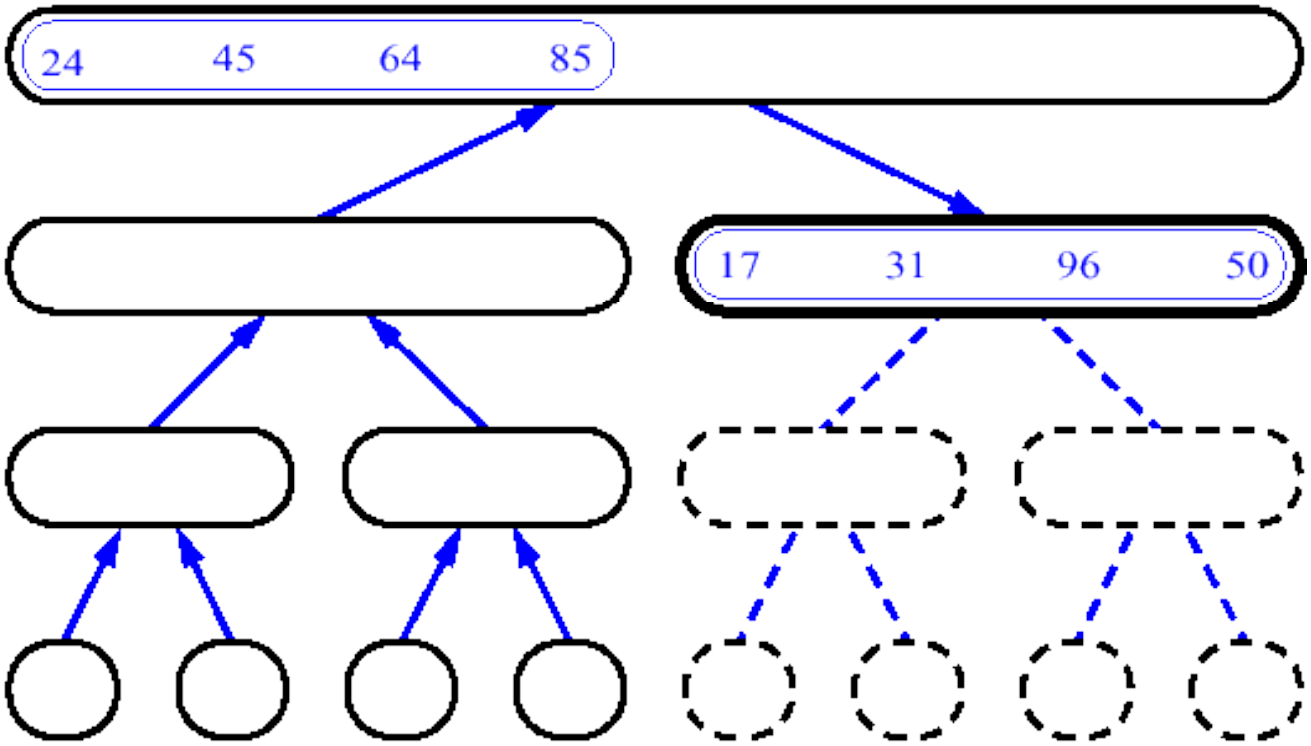

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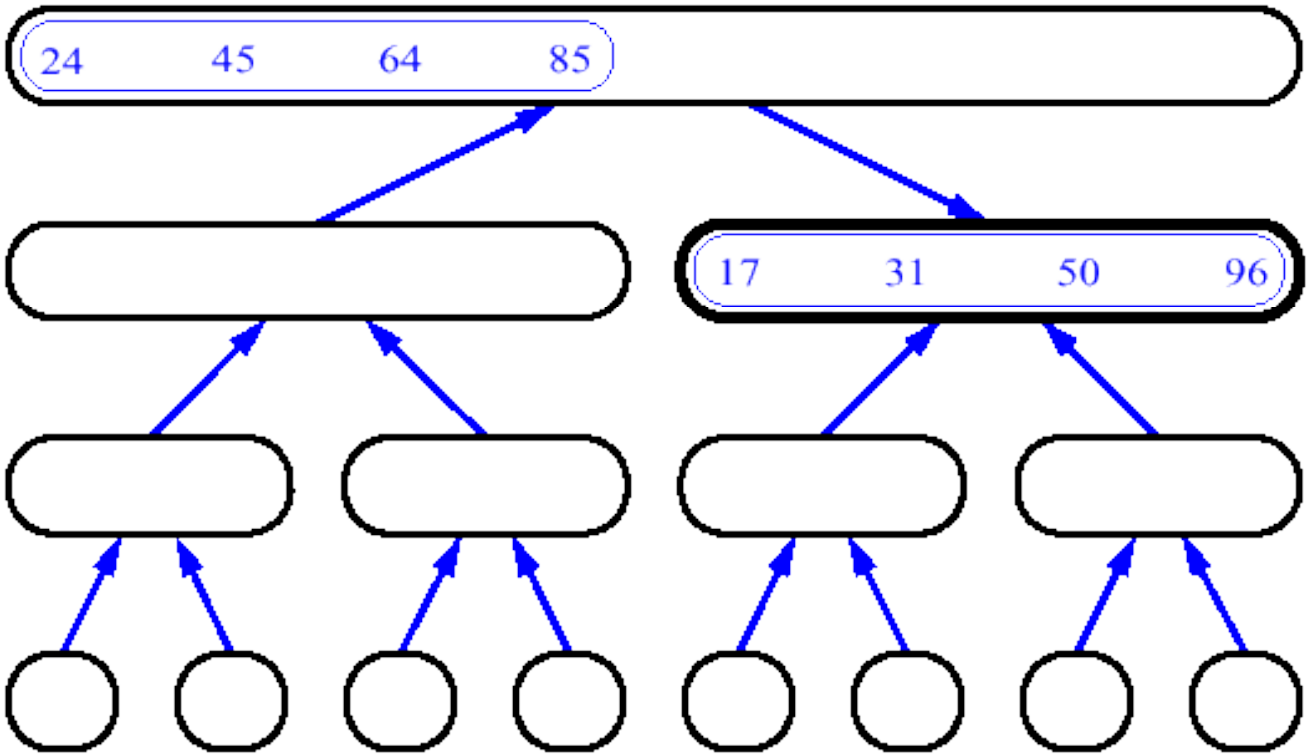

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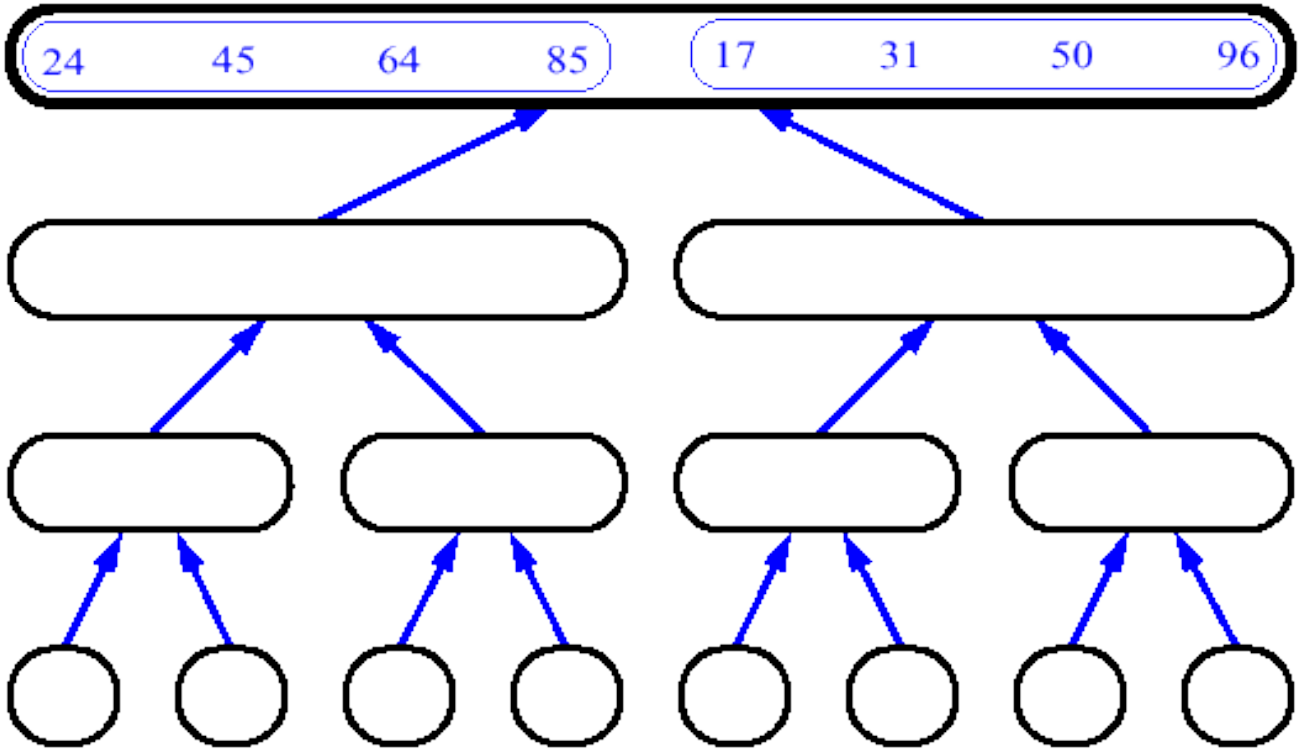

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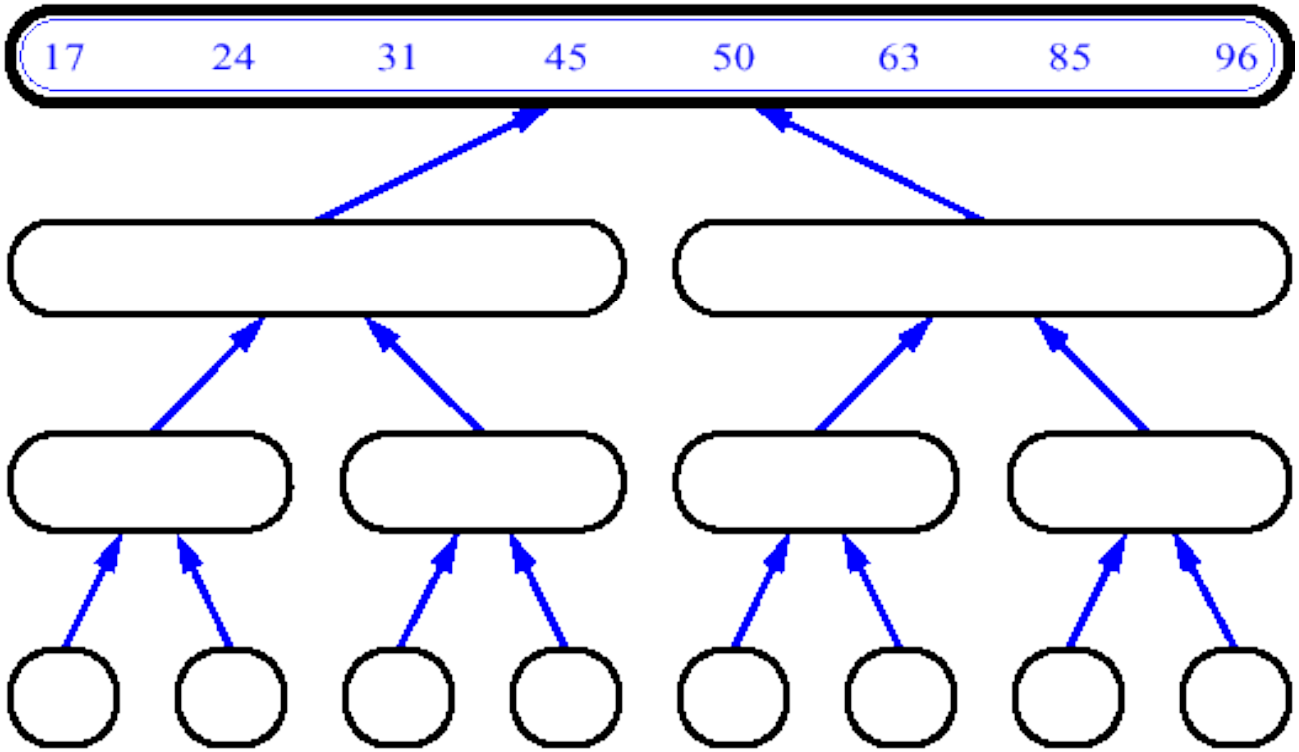

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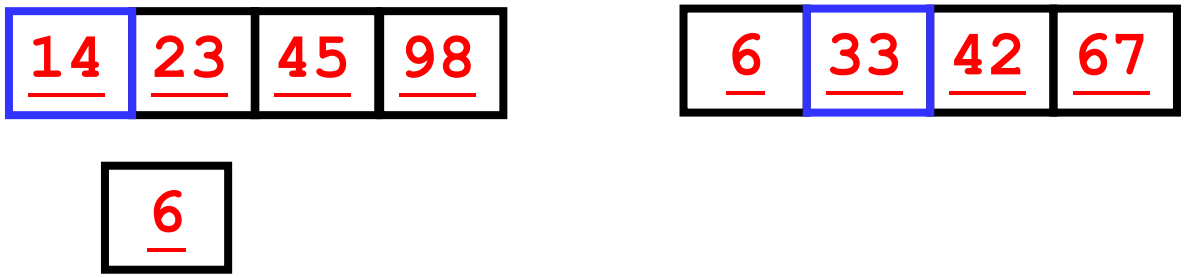

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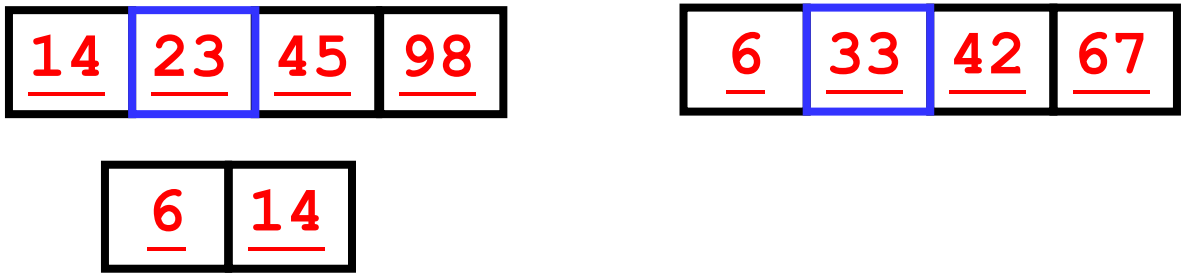

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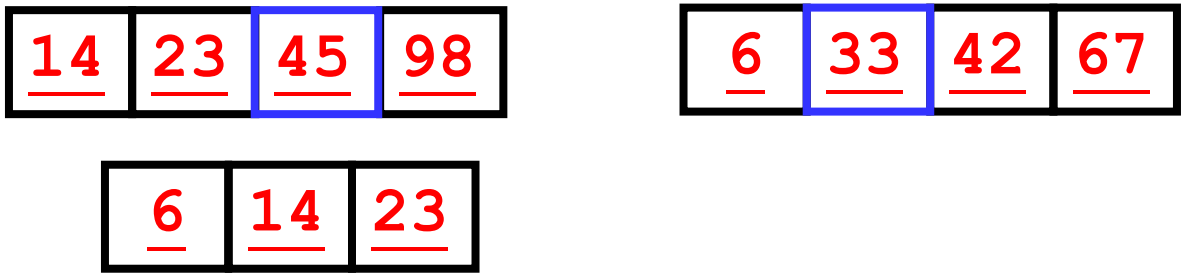

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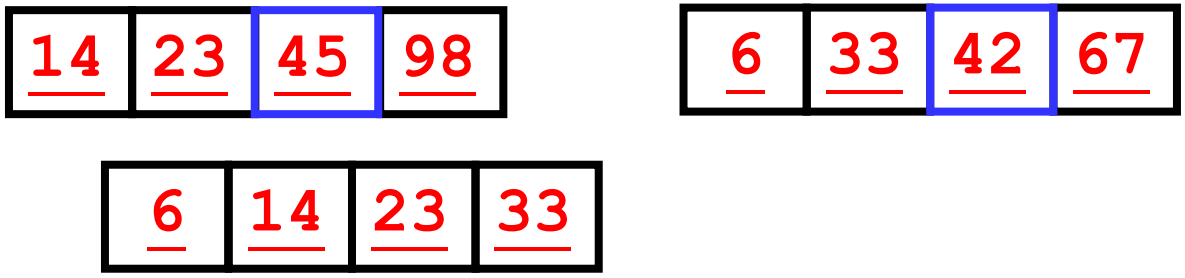

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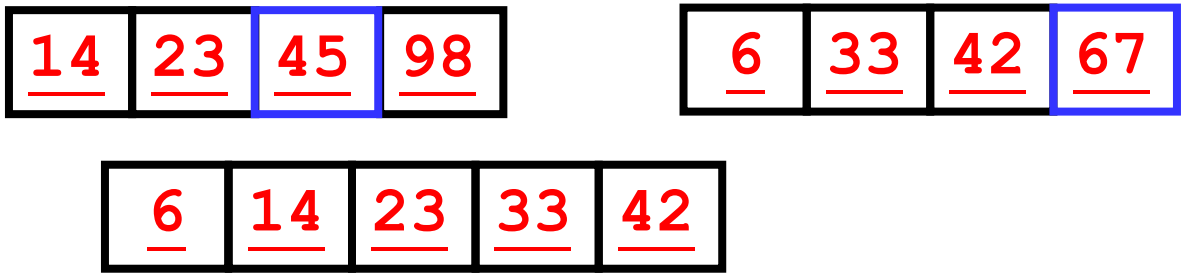

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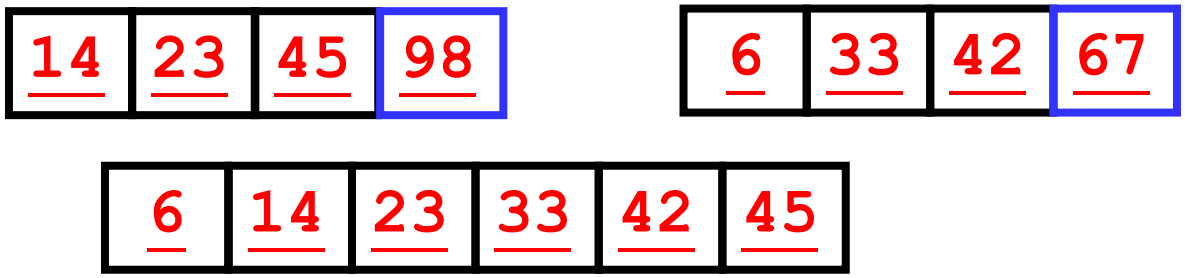

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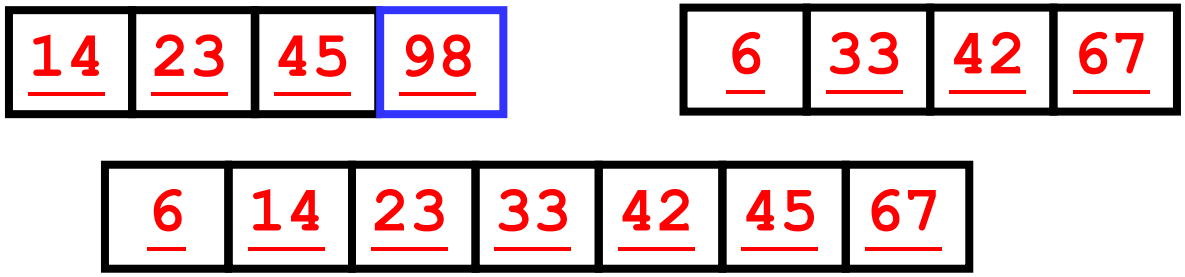

Merge