Swap And Maximize

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↔ difficulty	Easy
_≔ tags	Sorting
🖍 language	C++
solved on	@01/11/2024
⊘ link	https://www.geeksforgeeks.org/problems/swap-and-maximize5859/1

Intuition

The problem involves maximizing the sum of absolute differences between adjacent elements when arranging the elements of an array in a circular manner. The idea is to position the smallest and largest elements alternately to maximize the differences between adjacent elements.

Approach

- 1. Sort the Array: Start by sorting the array in ascending order.
- 2. **Rearrange Elements**: Create a new array temp by placing elements alternately from the start and end of the sorted array. This ensures that the difference between adjacent elements is maximized.
- 3. **Sum of Absolute Differences**: Iterate through the rearranged array temp and calculate the sum of absolute differences between consecutive elements, treating the array as circular by wrapping the last element to the first.
- 4. Return the Result: Return the calculated sum.

Complexity

Time Complexity:

- Sorting the array takes O(n log n).
- Rearranging the elements and calculating the sum take O(n).
- ullet Overall, the time complexity is $O(n \ log \ n)$, where n is the length of the input array.

Space Complexity:

• The space complexity is O(n) due to the use of the temp array for rearranging elements.

Code

```
class Solution {
  public:
    long long maxSum(vector<int>& arr) {
       long long answer = 0;
       sort(arr.begin(), arr.end());
}
```

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```
int n = arr.size();
    vector<int> temp;

int i, j;
    for(i = 0, j = n - 1; i < j; i++, j--) {
          temp.push_back(arr[i]);
          temp.push_back(arr[j]);
    }
    if(i == j) temp.push_back(arr[i]);

    for(int i = 0; i < n; i++) {
         answer += abs(temp[i] - temp[(i + 1) % n]);
    }
    return answer;
}</pre>
```

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