

# Maximum XOR For Each Query

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🌐 Platform	LeetCode
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# Serial	1829
≡ tags	Bit Manipulation
🗨 language	C++
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🔗 link	<a href="https://leetcode.com/problems/maximum-xor-for-each-query/description/">https://leetcode.com/problems/maximum-xor-for-each-query/description/</a>
☑ Completion	✓

## Intuition

The problem revolves around calculating the maximum XOR value for each prefix of the given list of integers, where the XOR is computed with a mask that keeps the bits within a specified bit range (`maximumBit`). The result for each position is the maximum XOR value that can be obtained using the given numbers with respect to the bit limit.

- The approach focuses on maintaining a running XOR (`XOR`) of the elements in the array.
- The mask is used to limit the size of the XOR result to fit within the given `maximumBit`.
- The key insight is that for each element, the XOR operation flips bits, and by taking the complement (`~XOR`), the goal is to maximize the XOR value.

## Approach

- Mask Calculation:** First, calculate a mask based on the given `maximumBit`. The mask ensures that we only consider the lower `maximumBit` bits when calculating the XOR result.
- Iterate Over the Array:**
  - Maintain a running XOR of the array's elements.
  - For each element in the array, the result for that element is `(~XOR) & mask`. The result is stored in the corresponding position, starting from the end of the array.
- Why Reverse the Iteration:** The result is filled in reverse order because, for each position `i`, the maximum XOR up to that position needs to be calculated. Since the last element's result depends on all previous ones, iterating backward ensures we respect the prefix XOR.
- Return the Answer:** After populating the results in reverse, return the list.

## Complexity

### Time Complexity:

- The time complexity is  $O(n)$ , where `n` is the size of the input array `nums`. The loop runs once over the entire array, performing constant-time operations like XOR and bitwise operations.

### Space Complexity:

- The space complexity is  $O(n)$  due to the `answer` vector storing the results. No additional significant space is used beyond this.

## Code

```
class Solution {
public:
    vector<int> getMaximumXor(vector<int>& nums, int maximumBit) {
        vector<int> answer(nums.size());
        int mask = (1 << maximumBit) - 1;

        int XOR = 0;

        for (int i = 0; i < nums.size(); i++) {
            XOR ^= nums[i];
            answer[nums.size() - 1 - i] = (~XOR) & mask;
        }

        return answer;
    }
};
```