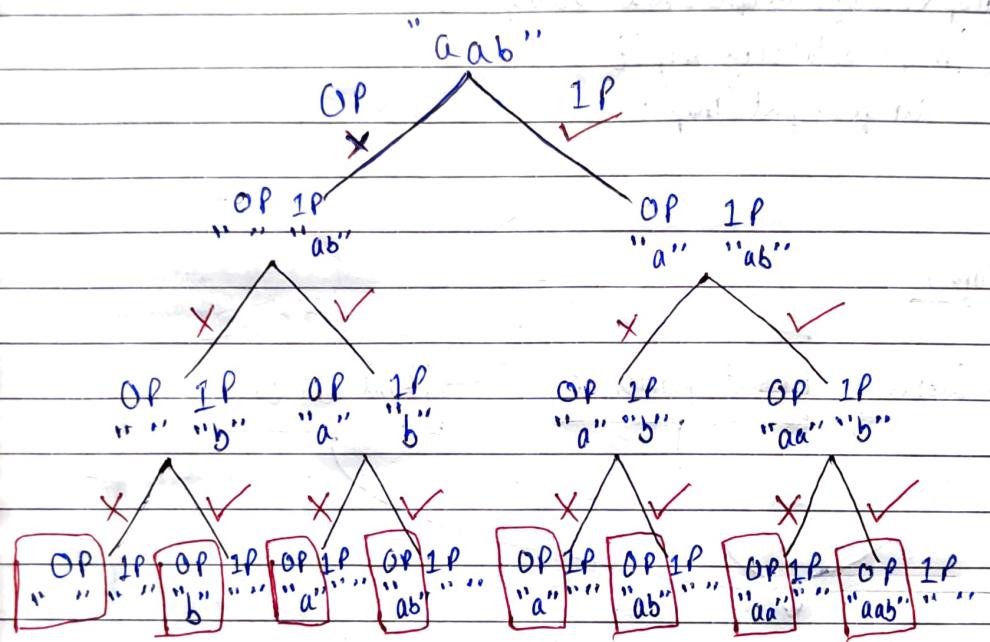


Print Unique Subsets & Variation

As we already know how to print subsets, but what does Print Unique subsets mean over here.

Let's see with an example:-



So from this we get subsets something like this,

```
" "
b
a
ab
a
ab
aa
aab
```

If you carefully look there are some repetition & we only want unique one. So, for that what we will do instead of printing OP. We will create one Hashmap. And fill that Hashmap with all subsets.

We now create one more Hashmap & fill that one only with unique subsets.

```

if (ip.length() == 0) {
    HashMap<String> map = new HashMap<>();
    return;
}

```

3

~~Hash~~

```

HashMap<String> unique = new HashMap<>();
if (!map.containsKey(map)) {
    unique.put(map);
}

```

}

Variations

Print all Subset

Print powerset

Print all Subsequence

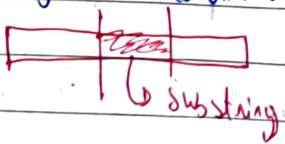
These 3 variations are not same. Print all subset is same as powerset. But Subsequence is a different thing.

Now, what is powerset?

→ We have given a set & we have to print its all subset
 $\{a, b\} \rightarrow \{a, b, "", a, b\}$
 $S \rightarrow PS$

Now, what is subsequence?

→ To understand this, we have to understand 1st substring.
 Substring is just a continuous part of string



~~1st~~ ~~2nd~~

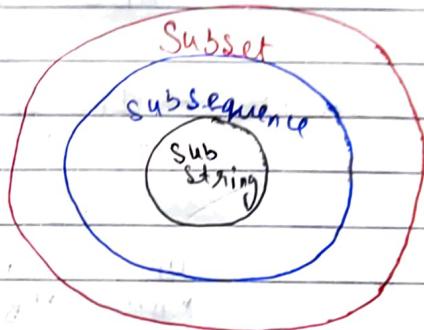
+ ↗ Is not a substring,
But it's a subsequence.

"abc" → ac is a subsequence
but ca is not a subsequence.

Venn's Diagram

All substrings are → Subsequence

All subsequences are → Subsets



Sum-up

→ If interviewer says to print Powerset, subset OR subsequence
we have to print subset only.

Now here's 2 conditions

→ Have Duplicate

↳ Not a unique set

→ Print exactly tree

→ Print lexicographically

→ Not have Duplicate

↳ A unique set

→ Print exactly tree

→ Print lexicographically

→ For making unique use HashSet

→ Forming lexicographical order. Put into an array, then sort it.

PERMUTATION WITH SPACES

Problem Statement:-

Input:- ABC

Output:- A-B-C

A-BC

AB-C

ABC

Recursive Tree \rightarrow IP OP

ABC " "

↓ A

OP IP

"A" "BC"

-B B

"A-B" "C" "AB" "C"

-C C

"A-B-C" "A-BC" "AB-C" "ABC" " "

→ Same as Output

A-B-C

A-BC

AB-C

ABC

Choice

- Include that letter with space
- Include that letter without space

int main() {

PSEUDO CODE

 input ip= " ";

 String op= " ";

 op = ip.chrAt(0).pushBack(ip[0]) // including A as it is

 ip.erase(ip.begin() + 0); // removing A for ABC = BC

 solve(ip, op)

```

Void solve (String ip, String op) {
    if (ip.length() == 0) {           // Base Condition
        System.out.print (op + " ");
        return;
    }
    String op1 = op;
    String op2 = op;
    op1.push_back ('-');           } // Including with space
    op1.push_back (ip[0]);          } // Including without space
    op2.push_back (ip[0]);          // Erasing from input to make smaller
    ip.erase (ip.begin() + 0);      // Erasing from input to make smaller
    solve (ip, op1);
    solve (ip, op2);
    return;
}

```

3 Let's Code it:-

```

public static void main (String [] args) {
    Scanner scn = new Scanner (System.in);
    String ip = scn.nextLine();
    String op = "";
    op += ip.charAt(0);
    ip = ip.substring (1);
    solve (ip, op);
}

```

```

3
public static void solve (String ip, String op) {
    if (ip.length() == 0) {
        System.out.println (op);
        return;
    }
}

```

```

String op1 = op;
String op2 = op;

```

op1 = "-" + ip.charAt(0);

op2 = ip.charAt(0);

ip = ip.substring(1);

solve(ip, op1);

solve(ip, op2);

}

PERMUTATION WITH CASE CHANGE

Problem statement:-

IP :- ab

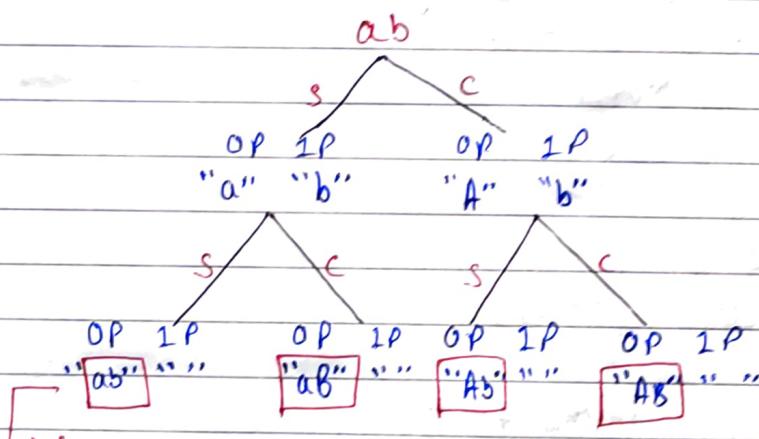
OP :- ab

aB

Ab

AB

Recursive Tree → OP IP



Same as output

ab

aB

Ab

AB

Void solve (String ip, String op) {

if (ip.length() == 0) {

print (op);

return;

}

String op1 = op;

String op2 = op;

op1.push_back(ip[0]);

op2.push_back(toupper(ip[0]));

ip.erase(ip.begin() + 0);

solve(ip, op1);

solve(ip, op2);

3

Let's Code it:-

```
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
    String ip = scn.nextLine();  
    String op = " ";  
    solve(ip, op);  
}
```

```
}
```

```
public static void solve(String ip, String op) {  
    if (ip.length() == 0) {  
        System.out.println(op);  
        return;  
    }
```

```
    String op1 = op + ip.charAt(0);
```

```
    String op2 = op;
```

```
    op1 += ip.charAt(0);
```

```
    op2 += Character.toUpperCase(ip.charAt(0));
```

```
    ip = ip.substring(1);
```

```
solve(ip, op1);
```

```
solve(ip, op2);
```

```
}
```

LETTER CASE PERMUTATION

Problem Statement:-

It's similar to previous question, but have some conditions in it.

→ Letter could be smaller or upper

→ And has digit as well

IP :- a1B2

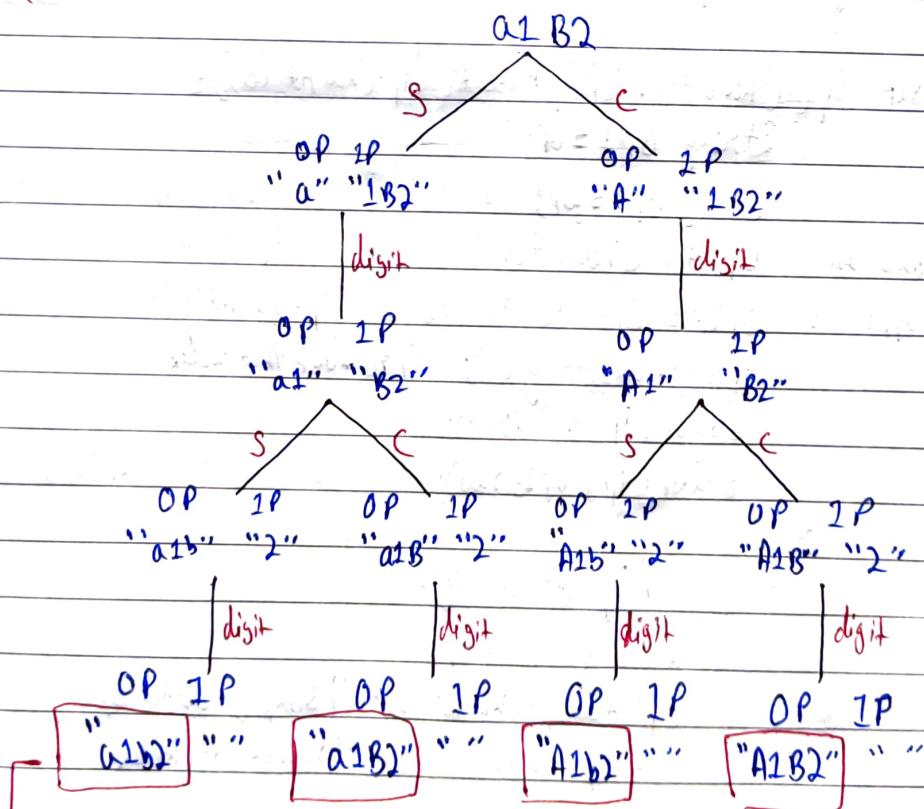
OP :- a1b2

a1B2

A1b2

A1B2

Recursive Tree → DP ~~OP IP~~



→ Some of Outputs:-
a1b2
a1B2
A1b2
A1B2

This is a LeetCode Problem,
LeetCode 178

```
public List<String> letterCasePermutation (String S) {
    String ip = S;
    String op2 = "";
    List<String> res = new ArrayList<>(); // result array
    solve(ip, op2, res); // func call
    return res;
}
```

```
3
public void solve (String ip, String op, List<String> res) {
    if (ip.length() == 0) { // Base Condition
        res.add(op); // add output to res
        return;
    }
```

```
// if its Alphabet
if (Character.isAlphabetic(ip.charAt(0))) {
    String op1 = op;
    String op2 = op;
    // convert to lowercase
    op1 += Character.toLowerCase(ip.charAt(0));
    // convert to uppercase
    op2 += Character.toUpperCase(ip.charAt(0));
    ip = ip.substring(1); // remove that index
    solve(ip, op1, res);
    solve(ip, op2, res);
}
```

```
3
else { // if it's a digit
    String op1 = op;
    op1 += ip.charAt(0); // simply add to op1
    ip = ip.substring(1); // remove its index
    solve(ip, op1, res);
}
```

3

LETTER CASE PERMUTATION

Problem Statement:-

It's similar to previous question, but have some conditions in it.

→ Letter could be smaller or upper

→ And has digit as well

IP :- a1B2

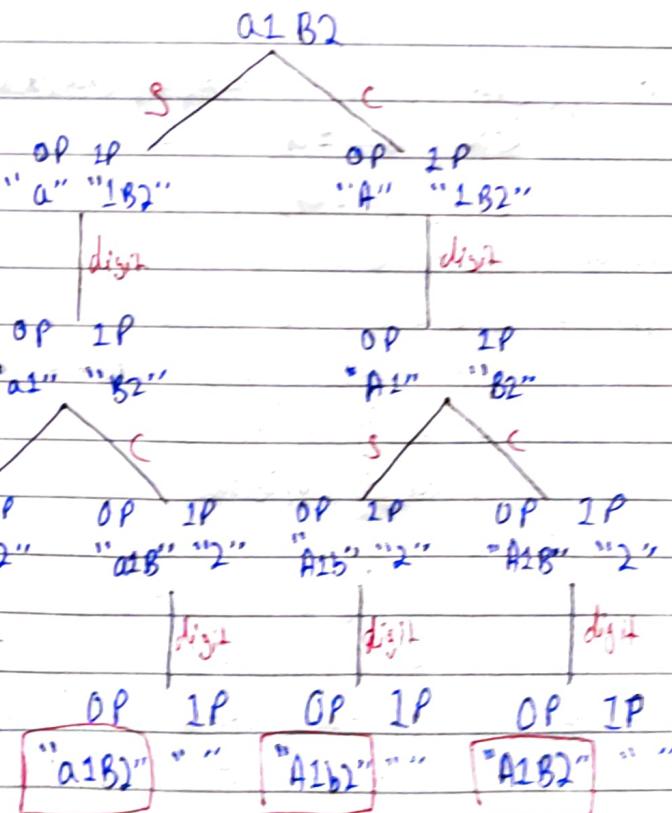
OP :- a1b2

a1B2

A1b2

A1B2

Recursive Tree → DP ~~OP IP~~



→ Some of outputs :- a1b2

a1B2

A1b2

A1B2

This is a LeetCode Problem,

Let's Code It :-

```
public List<String> letterCasePermutation (String S) {  
    String ip = S;  
    String op = "";  
    List<String> res = new ArrayList<>(); // result array  
    solve(ip, op, res); // func call  
    return res;
```

```
}  
public void solve (String ip, String op, List<String> res){  
    if (ip.length() == 0) { // Base Condition  
        res.add(op); // add output to res  
        return;
```

}

```
// if its Alphabet if (Character.isAlphabetic(ip.charAt(0))) {  
    String op1 = op;  
    String op2 = op;  
    // convert to lowercase op1 += Character.toLowerCase(ip.charAt(0));  
    // convert to uppercase op2 += Character.toUpperCase(ip.charAt(0));  
    ip = ip.substring(1); // remove that index  
    solve(ip, op1, res);  
    solve(ip, op2, res);
```

}

```
else { // if it's a digit
```

```
    String op1 = op;  
    op1 += ip.charAt(0); // simply add to op1  
    ip = ip.substring(1); // remove its index  
    solve(ip, op1, res);
```

}

3

GENERATE ALL BALANCED PARENTHESES

Problem Statement :-

I/P: $n=2 \rightarrow$ Balanced Parentheses,

↳ 2 open \boxed{CC}

2 close $\boxed{))}$

O/P:-

$\boxed{(C)} \quad \boxed{(C))}$

All possible balanced parentheses.

for $n=2$ boxes, there will be 4 :-

$\boxed{\square} \quad \boxed{C} \quad \boxed{D} \quad \boxed{D} \leftarrow$ Valid parentheses

$\boxed{D} \quad \boxed{C} \quad \boxed{D} \quad \boxed{C} \leftarrow$ Invalid parentheses

Let's take one more example, to understand it:-

$n=3$

For $n=3$, possible balanced parentheses will be 5:-

3 open

\boxed{CCC}

3 Close

$\boxed{)))}$

$\boxed{(CCC))})$
$\boxed{(C)C(C)}$
$\boxed{(C)(C)C}$
$\boxed{(C)(C))}$
$\boxed{(C)C(C)}$

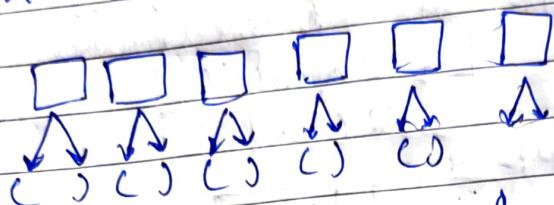
\Rightarrow In the output we have to return Array of String,
Where every string is balanced

Now, I identify how this is recursion.

for that we have to think about Choices + Decision

To understand choices & Decision let's take $n=3$

for $n \geq 3$, we have 6 boxes



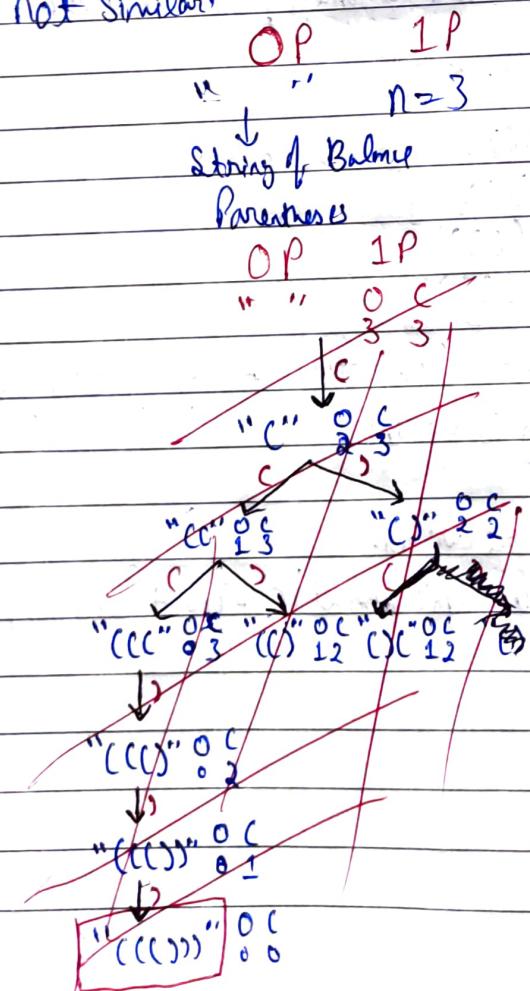
Now for every boxes we have 2 choices & we have to take decision accordingly.

So, our decision will give us balanced, stable decision according to the rule.

So, our decision will give us balanced.
But we have to be carefully while taking decision.

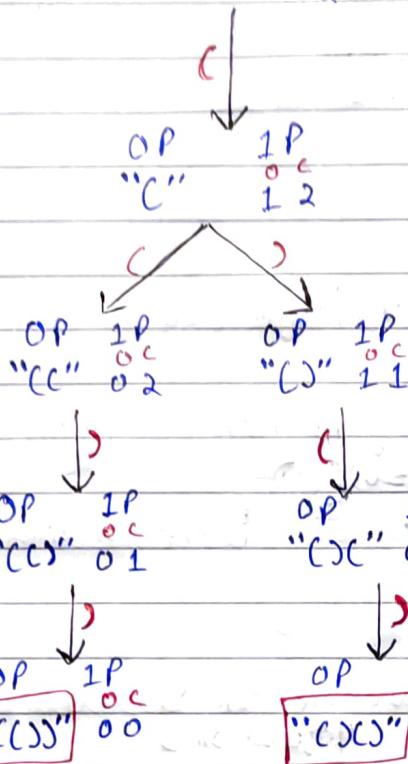
Let's Build Recursive Tree

To build it's recursive tree. If you see data types are not similar.



$n=2$

OP IP
" " 0 0
 2 2



Leaf Node $O=0$
 $C=0$

If you carefully look, sometimes we have 2 choices
sometimes 1 choice only.

What does it seem like we can only have closing bracket when we have open bracket. Like opening brackets we have all time choice but for closing we get sometime or not!
So if open & close brackets are equal we will have 2 choices & if close brackets are more than open we will have closing choice.

If like $()()$ → use close bracket

But $()$ so use open when they are balance.

Balance ↑↑ so that close become greater.

```
vector<string> bPar(int n) {
```

```
    vector<string> v;
```

```
    int open = n;
```

```
    int close = n;
```

```
    string op = "";
```

} Initialize

Solve(open, close, op, v) || fnc call

} return v;

```
void solve(int open, int close, string op, vector<string> v) {
```

```
    if (open == 0 & & close == 0) {
```

```
        v.push_back(op);
```

```
        return;
```

} Base condition

```
} if (open != 0) { // we have open still left
```

```
    string op1 = op; // equal or 0
```

```
    op1.push_back('(')
```

```
solve(open - 1, close, op1, v);
```

}

```
} if (close > open) { // we have close more than open
```

```
    string op2 = op;
```

```
    op2.push_back(')')
```

```
solve(open, close - 1, op2, v);
```

}

It's an LeetCode Problem,
Let's Code It:-

```

public List<String> generateParenthesis(int n) {
    List<String> res = new ArrayList<>();
    int open = n;
    int close = n;
    String op = "";
    solve(open, close, op, res);
    return res;
}

```

```

}
public void solve(int open, int close, String op, List<String> res) {
    if (open == 0 && close == 0) {
        res.add(op);
    }
}

```

```

}
if (open != 0) {
    String op1 = op;
    op1 += '(';
    solve(open - 1, close, op1, res);
}

```

```

}
if (close > open) {
    String op2 = op;
    op2 += ')';
    solve(open, close - 1, op2, res);
}

```

}