#### **COMP1001**

## **Computer Systems**

Dr. Vasilios Kelefouras

Email: v.kelefouras@plymouth.ac.uk

Website:

https://www.plymouth.ac.uk/staff/vasilios -kelefouras

School of Computing (University of Plymouth)

#### Outline

- Positional Numbering Systems
- Converting Positional Numbering Systems
- Basic Binary Arithmetic Operations
- Signed Integer Representation
- Floating Point Representation
- Character Codes

## Basics (1)

- The bit is the most basic unit of information in a computer
  - Switching activity 0 or 1
- A Byte is a group of 8 bits
  - A byte is the smallest possible addressable unit of computer storage
  - The term, "addressable," means that a particular byte can be retrieved according to its location in memory
- A word is a contiguous group of bytes, e.g., an integer uses 4 bytes
- Word sizes of 4 or 8 bytes are most common

## Basics (2)

```
Kilo- (K) = 1 thousand = 10^3 and 2^{10}

Mega- (M) = 1 million = 10^6 and 2^{20}

Giga- (G) = 1 billion = 10^9 and 2^{30}

Tera- (T) = 1 trillion = 10^{12} and 2^{40}

Peta- (P) = 1 quadrillion = 10^{15} and 2^{50}

Exa- (E) = 1 quintillion = 10^{18} and 2^{60}

Zetta- (Z) = 1 sextillion = 10^{21} and 2^{70}

Yotta- (Y) = 1 septillion = 10^{24} and 2^{80}
```

Normally, powers of 2 are used for measuring capacity

Milli- (m) = 1 thousandth = 
$$10^{-3}$$
  
Micro- ( $\mu$ ) = 1 millionth =  $10^{-6}$   
Nano- (n) = 1 billionth =  $10^{-9}$   
Pico- ( $p$ ) = 1 trillionth =  $10^{-12}$ 

## Basics (3)

- □ Hertz = clock cycles per second (frequency)
  - $\square$  1MHz = 1,000,000Hz
  - Processor speeds are measured in MHz or GHz
- $\square$  Byte = a unit of storage
  - $\square$  1KB =  $2^{10}$  = 1024 Bytes
  - $\square$  1MB =  $2^{20}$  = 1,048,576 Bytes
  - $\square$  1GB =  $2^{30}$  = 1,099,511,627,776 Bytes
- Main memory (RAM) is measured in GB
- $\square$  Disk storage is measured in TB (2<sup>40</sup>)

#### Think Pair Share activity

□ How many milliseconds (ms) are in 1 second?

How many nanoseconds (ns) are in 1 millisecond?

- How many kilobytes (KB) are in 1 gigabyte (GB)?
- How many bytes are in 20 megabytes?

#### **POSITIONAL NUMBERING SYSTEMS**

- Positional numbering systems are systems in which the placement of a digit in connection to its intrinsic value determines its actual meaning in a numeral string
- The organization of any computer depends considerably on how it represents numbers, characters, and control information
  - There are several positional numbering systems such as Decimal, Binary,
     Octal, Hexadecimal etc
- The positioning system is provided as a subscript, e.g.,  $14_{10}$ ,  $10101_2$ ,  $82_{16}$
- Our decimal system is the base-10 system. It uses powers of 10 for each position in a number
- The binary system is also called the base-2 system
- The hexadecimal system is the base-16 system
- The Mayan and other Mesoamerican cultures used a number system based in a base-20 system

#### **Decimal System**

- Decimal system: Our well known and used system.
  - It uses 10 different digits: 0,1,2,3,4,5,6,7,8,9
  - Our decimal system is the base-10 system. It uses powers of 10 for each position in a number
  - For example, the decimal number 947 in powers of 10 is 947 =

$$=9\times100 + 4\times10 + 7\times1 =$$
  
 $=9\times10^2 + 4\times10^1 + 7\times10^0$ 

$$70216=7x10000+0x1000+2x100+1x10+6x1=$$

$$=7x10^{4}+0x10^{3}+2x10^{2}+1x10^{1}+6x10^{0}$$

The decimal number 3812.46 in powers of 10 is  $(3x10^3 + 8x10^2 + 1x10^1 + 2x10^0 + 4x10^{-1} + 6x10^{-2})$ 

```
10
```

- A binary number is a number expressed in the base-2 numeral system or binary numeral system, which uses only two symbols: typically 0 (zero) and 1 (one)
- The base is 2
- 2 different digits are used: 0,1
- For example,  $101_2 = 1x2^2 + 0x2^1 + 1x2^0$ = 1x4 + 0x2 + 1x1=  $5_{10}$
- The binary number 11001 in powers of 2 is:  $1x2^4 + 1x2^3 + 0x2^2 + 0x2^1 + 1x2^0 = 16 + 8 + 0 + 0 + 1 = 25_{10}$
- $1011.101_{2} =$   $= 1x2^{3} + 0x2^{2} + 1x2^{1} + 1x2^{0} + 1x2^{-1} + 0x2^{-2} + 1x2^{-3} =$  = 1x8 + 0x4 + 1x2 + 1x1 + 1x0.5 + 0x0.25 + 1x0.125  $= 11.625_{10}$

## **Binary System (2)**

2 <sup>n</sup> representation	<b>2</b> <sup>10</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>2</sup>	<b>2</b> <sup>1</sup>	<b>2</b> <sup>0</sup>	2-1	2-2	<b>2</b> -3
number	1024	16	8	4	2	1	0.5	0.25	0.125

#### Convert the following binary number 1101.101 to decimal

**1101.101** 
$$_{2}$$
=1x2 $^{3}$  + 1x2 $^{2}$  + 0x2 $^{1}$  + 1x2 $^{0}$  + 1x2 $^{-1}$  + 0x2 $^{-2}$  + 1x2 $^{-3}$  = 8 + 4 + 0 + 1 + 0.5 + 0 + 0.125 = 13.625 $_{10}$ 

## Octal system

- The base is 8
- 8 different digits are used only: 0,1,2,3,4,5,6,7
- For example:  $436_8 = 4x8^2 + 3x8^1 + 6x8^0$ = 4x64 + 3x8 + 6x1=  $286_{10}$

Convert the following octal number 205.24<sub>8,</sub>to decimal:

$$205.24_8 = 2x8^2 + 0x8^1 + 5x8^0 + 2x8^{-1} + 4x8^{-2}$$
$$= 2x64 + 0 + 5 + 2x0.125 + 4x0.015625$$
$$= 133.3125_{10}$$

## Hexadecimal system

- ☐ The base is 16
- 16 different digits are used: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F (we do not use numbers with 2 digits like 10,11,12,...), but

A instead of 10, B instead of 11, C instead of 12, etc)

Example: 
$$3B1_{16} = 3x16^2 + 11x16^1 + 1x16^0$$
  
=  $3x256 + 11x16 + 1 =$   
=  $768 + 176 + 1 =$   
=  $945_{10}$ 

Convert the following hexadecimal number 20C.2<sub>16</sub> to decimal

**20C.2**<sub>16</sub>= 
$$2x16^2 + 0x16^1 + 12x16^0 + 2x16^{-1} = 2x256 + 0 + 12x1 + 2x0.0625 = 512 + 12 + 0.125 = 524.12510$$

## Positional Numbering Systems - General case

- □ Base: r
- Uses r different digits: 0,1,2,3,..r-1

For example if  $234.03_5 = ?_{10}$  then n=3, m=2 and r=5

The left most digit (An-1) is called Most Significant Bit-(MSB) while the right most (A-m) Least Significant Bit-(LSB)

## Converting Positional Numbering Systems

#### From Decimal to Binary

The easiest method of converting integers from decimal to some other base uses division

#### **Procedure:**

- a. Divide the decimal by 2
- b. Write down the quotient and the remainder
- c. Divide quotient by 2
- d. Write down the quotient and the remainder
- e. Repeat the process (c)-(d) until the quotient becomes zero
- f. Write down the binary number from bottom (MSB) to top (LSB)

#### From Decimal to Binary, an example

```
83_{10} = ?_2
83 \div 2 = 41 remainder 1
                                 (LSB)
41 \div 2 = 20 remainder 1
20 \div 2 = 10 remainder 0
                                        If the decimal number is lower than 1,
                                        e.g., 0.25, or contains a fractional part
10 \div 2 = 5 remainder
                                        the procedure applied is different
 5 \div 2 = 2 remainder
 2 \div 2 = 1 remainder
  1 \div 2 = 0 remainder 1
                                 '(MSB)
83_{10} = 1010011_2
```

Our result, is the remainders in reverse order (reading from bottom to top)

#### **Convert From Decimal To Another Base**

- We follow the same procedure as in the previous slide (from decimal to binary) but instead of using 2-base we use the r-base.
- Convert the decimal 524<sub>10</sub> to hexadecimal

524:16= 32 remainder 12 (
$$12_{10} = C_{16}$$
)  
32:16=2 remainder 0  
2:16=0 remainder 2  
Thus,  $524_{10} = 20C_{16}$ 

#### Convert from binary to octal

#### **Procedure:**

 $101011_2 = 53_8$ 

- 1. To convert a binary number of octal, we group all the 1's and 0's in the binary number in sets of three, starting from the far right
- 2. Start from the right to make your groups
- 3. Add zeros to the left of the last digit if you don't have enough digits to

make a set of three

4. Write down the decimal representation of every group

$$10011011_{2} = ;_{8}$$

$$10011011_{2} = ;_{8}$$

$$10 011 011$$

$$= 101 011$$

$$= 5 3$$

$$10011011_{2} = ;_{8}$$

$$10011 011_{2} = ;_{8}$$

$$10011 011_{2} = ;_{8}$$

	Octal		
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

## Convert from Octal to binary

 Converting from octal to binary is as easy as converting from binary to octal. Simply look up each octal digit to obtain the equivalent group of three binary digits

```
317.2_8 = \frac{1}{2}
= 011 \ 001 \ 111.010
= 11001111.01_2
```

	Binary		Octal
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

#### Convert Binary to Hexadecimal

#### **Procedure:**

- 1. Cut your string of binary numbers **into groups of four,** starting from the right
- 2. Add extra zeros to the front of the first number if it is not four digits
- 3. Convert one 4-digit group at a time. To convert between Binary and Hexadecimal, you simply replace each Hex digit with its 4-bit binary equivalent and vice versa

$$10001101011_{2} = ;_{16}$$

$$= 100 0110 1011$$

$$= 0100 0110 1011 0110_{2} = 0 + 4 + 0 + 0 = 4_{16}$$

$$= 0100 0110 1011 0110_{2} = 0 + 4 + 2 + 0 = 6_{16}$$

$$= 4 6 B 1011_{2} = 8 + 0 + 2 + 1 = 11 = B_{16}$$

$$10001101011_{2} = 46B_{16}$$

Hexadecimal	Binary
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001
Α	1010
В	1011
С	1100
D	1101
Е	1110
F	1111

#### Convert Hexadecimal to Binary

- Likewise from Octal to Binary
- Simply look up each hexadecimal digit to obtain the equivalent group of four binary digits

That's why we use hexadecimal in computer systems! Humans can still understand it, and computers can calculate Hex faster than decimal values!

Hexadecimal	Binary	Decimal
0 ←	0000	0
1 👉	0001	1
2 👉	0010	2
3 🛑	0011	3
4 👉	0100	4
5 ←	0101	5
6 ←	<b>0110</b>	6
7 👉	0111	7
8 ←	1000	8
9 ←	1001	9
Α 👉	1010	10
В	1011	11
c <del>(</del>	1100	12
D -	1101	13
E -	1110	14
F ←	1111	15

#### Basic arithmetic operations

- The basic arithmetic operations are applied to all the previous numerical systems. There are:
  - Addition
  - Subtraction
  - Multiplication
  - Division
- For the reminder of this lecture we will focus on the binary system

#### Binary Addition (1)

#### Binary addition is like decimal addition:

- 1. Put the numbers in a vertical column, aligning the decimal points
- Add each column of digits, starting on the right and working left. If the sum of a column is more than ten, "carry" digits to the next column on the left.

```
457 +

148

605

(011 Carry)
```

- In the example above we add 8+7 and write 5 instead of 15. We propagate 10 (which is the base) to the left and we write the remainder
- For every propagation, we add 1 carry to the next addition (left)
- The same holds when adding different numerical system numbers too

#### Binary Addition (2)

#### **Binary addition:**

Note that: 0+0=0 and carry 0 1+0=0+1=1 and carry 0 1+1=0 and carry 1, as  $1+1=10_2=2_{10}$ 1+1+1=1 and carry 1, as  $1+1+1=11_2=3_{10}$ 

#### **Binary Multiplication**

(0111221100) Carry

As in the decimal system:

```
110111 X
         1101
    110111
               Note that 1+1+1+1=0 and carry
   000000
               2, as 1+1+1+1=100_2=4_{10}
  110111
 110111
1011001011
```

## Signed integer representation

#### **Introduction**

- In practice we have to use negative binary numbers too. We need to define signed binary numbers
- There are three ways in which signed binary integers may be expressed:
  - 1. Signed magnitude
  - 2. One's complement
  - 3. Two's complement

## Signed Magnitude Representation (1)

- Allocate the high-order (leftmost) bit to indicate the sign of a number
  - The high-order bit is the leftmost bit. It is also called the most significant bit
  - 0 is used to indicate a positive number; 1 indicates a negative number
- □ The remaining bits contain the value of the number
- Note that we also pay attention to the number of bits used to represent signed binary numbers
  - $\blacksquare$  i.e. if using 4 bit numbers, then we use  $0001_2$  rather than  $1_2$
- In an 8-bit word, signed magnitude representation places the absolute value of the number in the 7 bits to the right of the sign bit

#### For example:

+3 is: 00000011

- 3 is: 10000011

## Signed Magnitude Representation (2)

The "binary addition algorithm" does NOT work with sign-magnitude

0 0 1 
$$I_2 = 3_{10}$$
  
1 1 0  $0_2 = -4_{10}$   
0 0 1 1  
1 + 1 0 0  
1 1 1 1 this is wrong

# Signed Magnitude: intuitive for humans, difficult for computers

- Signed magnitude representation is easy for people to understand, but it requires complicated computer hardware
- Also it allows two different representations for zero: positive zero and negative zero
- As such, computer systems employ complement systems for signed number representation

# Signed Integer Representation Complement Systems

- In binary systems, these are:
  - One's Complement. To represent negative values, invert all the bits in the binary representation of the number (swapping 0s for 1s and vice versa)
    - 1 becomes 0 and 0 becomes 1
    - To represent positive numbers no change is applied

For example, using 8-bit one's complement representation

+ 3 is: 0000011

- 3 is: 11111100

More examples

X=11011100, 1C(X)=00100011

X=1011, 1C(X)=?

- One's complement still has the disadvantage of having two different representations for zero: positive zero and negative zero
- In addition positive and negative integers need to be processed separately
- Two's complement solves this problem
- Two's complement
  - One's Complement add 1

# Signed Integer Representation Two's Complement

#### Two's complement 2C(X)

- You represent positive numbers, just like the unsigned numbers
- To represent negative values, start with the corresponding positive number, invert all the bits. Then add 1
- For example, using 8-bit two's complement representation:

- -3 in 8-bit Two's Complement Representation is 11111101
- ✓ Negative numbers must always start with '1'
- ✓ Both positive and negative numbers must have the same number of bits

# Signed Integer Representation Two's Complement – Example 1

```
Example: X=01001 (+9<sub>10</sub>), n=5 bits -> Y=2C(X)= 10111 = -9<sub>10</sub>

Check: X+Y=

01001

+10111
```

The carry is discarded as the result must be 5 bits

100000=00000

- We can always check whether the two's complement result is correct by adding the two numbers. The result has to be zero. Note that the result of the addition must be of the n bits, where n is the number of bits of the inputs
- □ Find the negative binary number (two's complement) of the following positive number with 7bits  $0101101 (45_{10})$

Answer 1010011

# Signed Integer Representation Two's Complement – Example 2

#### Find the negative binary number $(-12_{10})$

- $\square$  Write down the positive  $12_{10}$  =  $01100_2$
- Decide the number of the bits. We can use 5 or more. Let say 8 bits
- Find the two's complement as follows

$$12_{10} = 00001100_2$$
,  $-12_{10} = 111110011 + 1 = 11110100$ 

-12<sub>10</sub> =11110100 (negative number)

## Wrap up

Given a positive binary number, we find its negative binary number by following the procedure:

- We decide the number of bits of the positive number. At least one'0' has to appear on its left.
- 2. We find its two's complement
- 3. If the MSB (the left most) is not '1' then we made a mistake...

#### Subtraction (with two's complement)

- $\square$  We know that A-B=A+(-B)
- So, instead of applying a subtraction we can make an addition with the opposite number, i.e., the two's complement. *The procedure follows:* 
  - 1. Find -B, i.e., its two's complement
  - 2. Add A with B's two complement
  - 3. The result has as many bits as the inputs

```
Make the subtraction Z=12-5 (use 5 digits) 12_{10}=01100_2, \, 5_{10}=00101_2, \, -5_{10}=11011_2 Z=01100 + 11011=100111, but we need 5 bits thus Z=00111_2=7_{10}
```

Z = 00111

Make the subtraction 9-12 (use 5 digits)

```
9_{10}=01001<sub>2</sub>, 12_{10}=01100<sub>2</sub>, -12_{10}=10100<sub>2</sub>
Z=01001 + 10100=11101 and thus Z=11101 (-3<sub>10</sub>)
Z=11101
```

#### Unsigned and Signed Integer Representation

- Both signed and unsigned numbers are useful
  - For example, memory addresses are always unsigned
- Using the same number of bits, unsigned integers can express twice as many "positive" values as signed numbers.
  - For example, the range of values that can be represented in 4-bits is:
    - $0000_2$  to  $1111_2$  (or 0 to 15) as unsigned
    - $0111_2$  to  $1111_2$  (or +7 to -7) as signed magnitude
    - $0111_2$  to  $1000_2$  (or +7 to -8) as two's complement

#### Example #1

#### What decimal value does the 8-bit binary number 10011110 have if

- a) it is interpreted as an unsigned number?
- b) it is on a computer using signed-magnitude representation?
- c) it is on a computer using one's complement representation?
- d) it is on a computer using two's complement representation?

#### Answer:

- a.  $100111110_2 = 1x2^7 + 1x2^4 + 1x2^3 + 1x2^2 + 1x2^1 = 128 + 16 + 8 + 4 + 2 = 158_{10}$
- b.  $100111110_2 = 1$  (negative)  $00111110_2 = -1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 = -30_{10}$
- Find the positive of  $100111110_2$ ; invert the bits of  $100111110_2$  and therefore  $01100001_2$ .  $01100001_2 = 1 \times 2^6 + 1 \times 2^5 + 1 \times 2^0 = 64 + 32 + 1 = 97_{10}$ . Since the original number was negative, the number is  $-97_{10}$
- Find the positive of  $100111110_2$ ; invert the bits of  $100111110_2$  and add 1; therefore  $01100010_2$ .  $01100010_2 = 1 \times 2^6 + 1 \times 2^5 + 1 \times 2^1 = 64 + 32 + 2 = 98_{10}$ . Since the original number was negative, the number is  $-98_{10}$

#### Example #2

#### What decimal value does the 8-bit binary number 00010001 have if

- a) it is interpreted as an unsigned number?
- b) it is on a computer using signed-magnitude representation?
- c) it is on a computer using one's complement representation?
- d) it is on a computer using two's complement representation?

#### **Answer:**

a. 
$$00010001_2 = 1 \times 2^4 + 1 \times 2^0 = 16 + 1 = 17_{10}$$

b. 
$$00010001_2 = 0$$
 (positive)  $0010001_2 = 1 \times 2^4 + 1 \times 2^0 = 16 + 1 = 17_{10}$ 

c. 
$$00010001_2 = 1 \times 2^4 + 1 \times 2^0 = 16 + 1 = 17_{10}$$

d. 
$$00010001_2 = 1 \times 2^4 + 1 \times 2^0 = 16 + 1 = 17_{10}$$

#### Example #3

# Perform the following binary subtraction using two's complement representation: Z=8 - 6

#### Answer:

- Instead of subtraction we do addition: 8 6 = 8 + (-6)
- 2. Choose the number of bits to represent 8 and 6 decimal numbers to binary. At least one zero has to appear on the left. Thus, we need 5 bits or more

$$8_{10} = 01000_2$$
  
 $6_{10} = 00110_2$ 

- 3. Find  $6_{10}$  (two's complement)  $6_{10} = 11010_2$
- Make the addition (The result has 5 bits)

## Signed Integer Representation Multiplication and Division by 2 (1)

- Binary multiplication and division by 2 is very easy. (as easy as it is to multiply and divide by 10 in the decimal system)
- Simply use an arithmetic shift operation

$$1_2 = 1_{10}$$
 $10_2 = 2_{10}$ 
 $100_2 = 4_{10}$ 
 $1000_2 = 8_{10}$ 
 $10000_2 = 16_{10}$ 

- A left arithmetic shift inserts a 0 in for the rightmost bit and shifts everything else left one bit; in effect, it multiplies by 2
- A right arithmetic shift shifts everything one bit to the right, but copies the sign bit; it divides by 2

## Signed Integer Representation Multiplication and Division by 2 (2)

- Multiplication by 2
  - Shift left by one place
  - E.g. to calculate 2 \* 7
- Division by 2
  - Shift right by one place
  - $\blacksquare$  E.g., to calculate 14/2

Binary Decimal 0000 0111 +7 0000 1110 +14

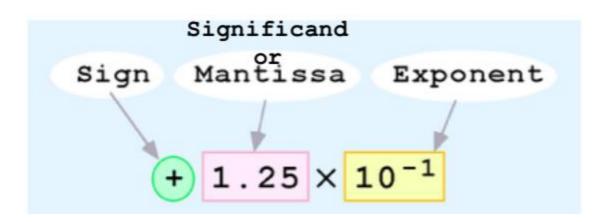
- To multiply by 4, we perform a left shift twice
- To divide by 4, we perform a right shift twice
- Using arithmetic shifting, perform the following:
  - a) double the value 00010101<sub>2</sub>
  - b) quadruple the value 00110111<sub>2</sub>
  - divide the value  $11001010_2$  in half

#### Floating-Point Representation (1)

- To represent real numbers with fractional values, floating-point representation is used
- □ Floating-point numbers are often expressed in scientific notation
  - □ For example:  $0.125 = 1.25 \times 10^{-1}$
- Remember that when a number is multiplied by its base, e.g., 10, then we add a zero or we move the ',' by one position to the right
  - $\square$  235x10 = 2350
  - □ 1.345×10=13.45
  - $110_2 \times 2 = 1100_2 (6 \times 2 = 12_{10})$
  - $\square$  101.11<sub>2</sub>×2=1011.1 (5.75×2=11.5<sub>10</sub>)

#### Floating-Point Representation (2)

- Computers use a form of scientific notation for floating-point representation
  - Single Precision floating point format 32-bit
  - Double Precision floating point format 64-bit
- Numbers written in scientific notation have three components:



#### Single precision Floating-Point format (1)

A binary number is represented in FP format as follows:

- 1. We write the number using only a single non-zero digit before the radix point : e.g.,  $1011010010001=1.011010010001 \times 2^{12}$  $1101.10111=1.101101111 \times 2^{3}$
- 1101.10111 = 1.101101111 x  $2^3$  r.

  2. Then we transform the number to the following format using 32 bits  $N = (-1)^{S} (1+F)(2^{E-127})$

Sign-S	Exponent-E	Mantissa (Fraction) - F
1-bit	8 - bits	23 - bits

**S: Sign,** 0/1 for positives/negatives, respectively

E: Exponent. E-127=exp, where exp is the corresponding exponent

F: Significant or Mantissa. We write the fractional part in 23 bits

E=127+exp in order to avoid using negative numbers. exp=[-127,128] and therefore E=[0,255]-255 needs 8 bits

## Single precision Floating-Point format (2)

Convert the positive number N=1011010010001 in Floating point format

```
Step 1: 1011010010001 = 1.011010010001 \times 2^{12}
```

Step 2: 
$$N = (-1)^{S} (1+F)(2^{E-127})$$

S = 0 (positive number)

$$E - 127 = 12$$
, and thus  $E = 139_{10}$  and  $E = 10001011_2$ 

Therefore N in FP format is:

0	10001011	0110100100010000000000
---	----------	------------------------

#### Single precision Floating-Point format (3)

Suppose that the 32-bit floating-point representation pattern is the following. Find the binary number

1	10010001	100011100010000000000
---	----------	-----------------------

S is 1 and thus the number is negative

E is  $10010001 = 145_{10}$ , and thus the exponent is exp=E-127=145-127=18

 $N = (-1)^{S} (1+F)(2^{E-127})$ 

N = -110001110001000000

#### Floating-Point Representation (1)

- No matter how many bits we use in a FP representation, the model is finite
  - The real number system is, of course, infinite, so our models can give nothing more than an approximation of a real value
- At some point, every model breaks down, introducing errors into our calculations
  - By using a greater number of bits in our model, we can reduce these errors, but we can never totally eliminate them

# Why is 0.1+0.2 not equal to 0.3 in most programming languages?

- computers use a binary floating point format that cannot accurately represent a number like 0.1<sub>10</sub>
- $\square$  0.1<sub>10</sub> is already rounded to the nearest number in that format
- 0.1<sub>10</sub> doesn't exist in the FP representation
- 0.1<sub>10</sub> is already rounded to the nearest number in that format, which results in a small rounding error
- This means that  $0.1_{10}$  is converted to a binary number that's just very close to  $0.1_{10}$
- The error is tiny since 0.1<sub>10</sub> is
   0.100000000000000055511151231257827
- The constants  $0.2_{10}$  and  $0.3_{10}$  are also approximations to their true values
- $\square$  So,  $0.1_{10} + 0.2_{10} == 0.30000000000000044408920985006_{10}$

#### **Character Codes**

- So far, we have learnt how to represent numbers. How about text?
- To represent text characters, we use character codes
  - Essentially, we assign a number for each character we want to represent
- As computers have evolved, character codes have evolved. Larger computer memories and storage devices permit richer character codes
- Some of the character codes are
  - 1. BCD
  - 2. ASCII (American Standard Code for Information Interchange) (7 bits)
  - 3. Extended ASCII (8-bits)
  - 4. Unicode
  - 5. and others
- A binary number of n bits gives 2<sup>n</sup> different codes
  - □ For n=2 there are  $2^2$  =4 different codes, i.e., bit combinations {00, 01, 10, 11}

#### Binary Coded Decimal (BCD) code

- when numbers, letters or words are represented by a specific group of symbols, it is said that the number, letter or word is being encoded. The group of symbols is called as a code
- Binary Coded Decimal (BCD) code
  - In this code each decimal digit is represented by a 4-bit binary number
  - BCD is a way to express each of the decimal digits with a binary code
  - In the BCD, with four bits we can represent sixteen numbers (0000 to 1111)

$$256_{10} = 0010 \ 0101 \ 0110_{BCD}$$

And vise versa

$$0011\ 1000\ 1001_{BCD} = 389_{10}$$

#### **ASCII Code**

- The most widely accepted code is called the American Standard Code for Information Interchange (ASCII).
- The ASCII code associates an integer value for each symbol in the character set, such as letters, digits, punctuation marks, special characters, and control characters
- The ASCII table has 128 characters, with values from 0 through 127.
  Thus, 7 bits are sufficient to represent a character in ASCII

#### **ASCII Code**

```
Dec Hx Oct Char
                                      Dec Hx Oct Html Chr
                                                            Dec Hx Oct Html Chr Dec Hx Oct Html Chr
                                       32 20 040 @#32; Space
                                                            64 40 100 @ 0
    0 000 NUL (null)
                                                                                96 60 140 4#96:
                                       33 21 041 6#33; !
    1 001 SOH (start of heading)
                                                            65 41 101 A A
                                                                                97 61 141 @#97;
    2 002 STX (start of text)
                                       34 22 042 @#34; "
                                                               42 102 &#66: B
                                                                                98 62 142 6#98;
                                                                                99 63 143 4#99;
    3 003 ETX (end of text)
                                       35 23 043 # #
                                                               43 103 C C
             (end of transmission)
                                       36 24 044 $ $
                                                               44 104 D D
                                                                              100 64 144 d d
    4 004 EOT
    5 005 ENQ
             (enquiry)
                                       37 25 045 @#37; %
                                                               45 105 E E
                                                                              101 65 145 @#101; e
                                                               46 106 @#70; F
                                                                              102 66 146 f f
    6 006 ACK (acknowledge)
                                       38 26 046 @#38; @
    7 007 BEL
              (bell)
                                       39 27 047 @#39; '
                                                            71 47 107 &#71: 🕃
                                                                              103 67 147 @#103; g
                                                               48 110 H H
                                                                              104 68 150 @#104; h
                                       40 28 050 @#40; (
    8 010 BS
              (backspace)
                                       41 29 051 6#41; )
              (horizontal tab)
                                                               49 111 @#73; I
                                                                              105 69 151 i i
    9 011 TAB
                                       42 2A 052 6#42; *
                                                               4A 112 6#74; J
                                                                              106 6A 152 @#106; j
    A 012 LF
              (NL line feed, new line)
10
                                       43 2B 053 &#43: +
                                                               4B 113 K K
11
    B 013 VT
              (vertical tab)
                                                                              |107 6B 153 k k
                                                               4C 114 L L
12
    C 014 FF
              (NP form feed, new page)
                                       44 2C 054 @#44;
                                                                              |108 6C 154 l <mark>1</mark>
    D 015 CR
              (carriage return)
                                       45 2D 055 - -
                                                               4D 115 %#77; M
                                                                              109 6D 155 m ™
                                       46 2E 056 .
                                                               4E 116 N N
                                                                              110 6E 156 n n
    E 016 SO
              (shift out)
   F 017 SI
              (shift in)
                                       47 2F 057 / /
                                                               4F 117 &#79: 0
                                                                              111 6F 157 @#111; o
                                       48 30 060 4#48; 0
                                                               50 120 P P
                                                                              112 70 160 @#112; p
16 10 020 DLE
              (data link escape)
17 11 021 DC1
             (device control 1)
                                       49 31 061 4#49; 1
                                                            81 51 121 Q 🔾
                                                                              113 71 161 q q
18 12 022 DC2 (device control 2)
                                       50 32 062 4#50; 2
                                                            82 52 122 @#82; R
                                                                              114 72 162 @#114; r
                                                            83 53 123 4#83; $
                                                                              115 73 163 @#115; 3
19 13 023 DC3 (device control 3)
                                       51 33 063 & #51; 3
20 14 024 DC4 (device control 4)
                                       52 34 064 4#52; 4
                                                            84 54 124 T T
                                                                              116 74 164 @#116; t
                                                                              117 75 165 u u
21 15 025 NAK (negative acknowledge)
                                       53 35 065 4#53; 5
                                                               55 125 U U
                                       54 36 066 @#54; 6
                                                               56 126 V V
                                                                              118 76 166 v V
22 16 026 SYN (synchronous idle)
23 17 027 ETB
             (end of trans. block)
                                       55 37 067 4#55; 7
                                                               57 127 W ₩
                                                                              119 77 167 w ₩
                                                               58 130 X X
                                                                              120 78 170 @#120; X
24 18 030 CAN (cancel)
                                       56 38 070 4#56; 8
                                                               59 131 4#89; Y
                                                                              121 79 171 @#121; Y
25 19 031 EM
              (end of medium)
                                       57 39 071 4#57; 9
                                       58 3A 072 @#58; :
                                                               5A 132 @#90; Z
                                                                              122 7A 172 @#122; Z
26 1A 032 SUB
              (substitute)
                                                                              |123 7B 173 @#123: {
                                       59 3B 073 &#59; ;
                                                            91 5B 133 [ 「
27 1B 033 ESC
              (escape)
                                                                              124 7C 174 @#124;
28 1C 034 FS
                                       60 30 074 < <
                                                               5C 134 @#92; \
              (file separator)
                                       61 3D 075 = =
                                                               5D 135 ] 1
                                                                              125 7D 175 @#125; )
29 1D 035 GS
              (group separator)
                                       62 3E 076 >>
                                                            94 5E 136 ^ ^
                                                                              126 7E 176 ~ ~
30 1E 036 RS
              (record separator)
                                       63 3F 077 4#63; ?
                                                            95 5F 137 6#95;
                                                                              127 7F 177  DEL
31 1F 037 US
              (unit separator)
```

Source: www.LookupTables.com

#### **Extended ASCII Characters**

- ASCII was designed in the 1960s for teleprinters and telegraphy, and some computing
- The number of printable characters was deliberately kept small, to keep teleprinters and line printers inexpensive
- When computers and peripherals standardized on eight-bit bytes, it became obvious that computers and software could handle text that uses 256-character sets at almost no additional cost in programming, and no additional cost for storage
- An eight-bit character set (using one byte per character) encodes 256
   characters, so it can include ASCII plus 128 more characters
- The extra characters represent characters from foreign languages and special symbols for drawing pictures

A set of codes that extends the basic ASCII set. The extended ASCII character set uses 8 bits, which gives it an additional 128 characters

128	Ç	144	É	160	á	176		192	L	208	Ш	224	Œ	240	=
129	ü	145	æ	161	í	177	******	193	Т	209	₹	225	ß	241	±
130	é	146	Æ	162	ó	178		194	т	210	π	226	$\Gamma_{\scriptscriptstyle{\parallel}}$	242	≥
131	â	147	ô	163	ú	179		195	H	211	Ш	227	π	243	≤
132	ä	148	ö	164	ñ	180	4	196	- (	212	Ŀ	228	Σ	244	ſ
133	à	149	ò	165	Ñ	181	4	197	+	213	F	229	σ	245	J
134	å	150	û	166	•	182	-	198	F	214	П	230	μ	246	÷
135	ç	151	ù	167	۰	183	П	199	\ <b> </b>	215	#	231	τ	247	æ
136	ê	152	ÿ	168	ż	184	1	200	Ŀ	216	#	232	Φ	248	۰
137	ë	153	Ö	169		185	4	201	F	217	J	233	Θ	249	
138	è	154	Ü	170	(4)	186		202	<u> IL</u>	218	Г	234	Ω	250	
139	ï	155	¢	171	1/2	187	a	203	īF	219		235	δ	251	V
140	î	156	£	172	1∕4	188	1	204	ŀ	220		236	00	252	ъ
141	ì	157	¥	173	i	189	Ш	205	=	221		237	ф	253	2
142	Ä	158	R.	174	«	190	4	206	#	222		238	ε	254	
143	Å	159	f	175	»	191	٦	207	±	223		239	$\wedge$	255	

Source: www.LookupTables.com

#### UNICODE

- Many of today's systems embrace Unicode that can encode the characters of every language in the world
  - The Java programming language, and some operating systems now use Unicode as their default character code
    - UTF-8 (8-bits: essentially the extended ASCII Table)
    - UTF-16 (16 bits: Most spoken languages in the world, widely used)
    - UTF-32 (32 bits: includes past languages, space inefficient)

# Any questions?



## Further Reading

 Chapter 9 and chapter 10 in 'Computer Organization and architecture' available at

http://home.ustc.edu.cn/~louwenqi/reference\_books tools/Computer%20Organization%20and%20Archite cture%2010th%20-%20William%20Stallings.pdf