#### **COMP1001**

#### **Computer Systems**

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- Von Neumann architecture
- What is the CPU?
- How CPU works?
  - Arithmetic Logic Unit (ALU)

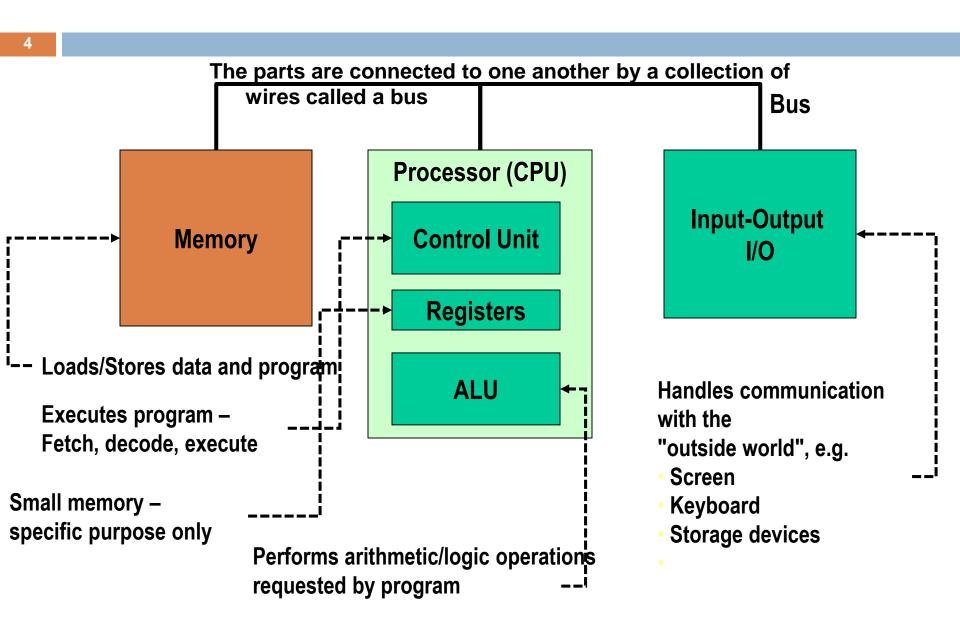
Outline

- Control Unit (CU)
- Bus
- Memory
- Registers
- Clock
- Memory Hierarchy
- Secondary Memory

#### The Von Neumann Architecture

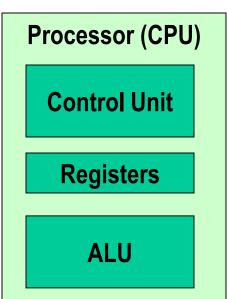
- All computers more or less based on the same basic design, the Von Neumann architecture
- It is a Model for designing and building computers, based on the following three characteristics:
  - 1. The computer consists of four main sub-systems:
    - Memory
    - ALU (Arithmetic/Logic Unit)
    - Control Unit
    - Input/Output System (I/O)
  - 2. Program is stored in memory during execution
  - 3. Program instructions are executed sequentially
- ✓ The architecture is named after the mathematician, John Von Neumann
- ✓ A variation of this architecture is the **Harvard** architecture which separates data and instructions into two pathways

#### The Von Neumann Architecture



# Central Processing Unit (CPU)

- CPU is responsible for fetching program instructions, decoding each instruction that is fetched, and executing the indicated sequence of operations on the correct data
- □ The key parts of the CPU are
  - 1. Arithmetic Logic Unit (ALU)
  - 2. Control Unit (CU)
  - 3. Registers
  - 4. Clock

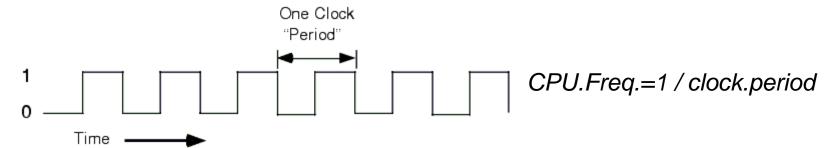


# Central Processing Unit (CPU)

- Carries out the program's instructions!
- Operates on data it finds in the computer's memory
  - Includes all binary circuits that carry out arithmetic & logic operationsreduced to a single Integrated Circuit
- CPU has four key parts that we will examine: Control Unit, Arithmetic & Logic Unit, Registers, Clock
- CPUs support a set of very simple instructions that typically fall into the following categories:
  - Data movement (load, store, copy...)
  - Arithmetic/logical (add, subtract, compare..)
  - Program control (branch, jump...)
- Very primitive commands (operations) executed by the CPU
- These commands are implemented as electronic binary circuits which can transform the 0s and 1s.

#### CPU Clock (1)

- Every computer contains an internal clock that regulates the rate at which instructions are executed and synchronizes all the various computer components
- All CPU and bus operations are synchronized to the clock
- Clock speeds are expressed in megahertz (MHz) or gigahertz ((GHz)
  - □ the beginning of each cycle is when the clock signal goes from "0" to "1"
  - e.g. CPU frequency 2 GHz -> clock cycle 0.5 ns
  - In the computer, all timings are measured in terms of clock cycles, e.g., an addition needs 2 cycles



# CPU Clock (2)

- □ The faster the clock, the more instructions the CPU can execute per second
- But, to think that clock and performance is the same thing is the most common misconception about processors
- CPU frequency is not necessarily indicative of the execution speed; e.g. complex operations, cache misses etc
  - FLOPS: floating operations per second, a useful measure for (super)computers dedicated to extensive computations
- A typical modern PC now has either four or five different clocks, running at different (but related) speeds, e.g., system (memory) bus, L2 cache bus, PCI bus
- □ The entire system is tied to the speed of the system clock.
  - Normally, the processor spends a significant amount of time waiting on data and signals from much slower devices, e.g., memory

- In a computer, a register is the fastest memory
- Registers are fast stand-alone storage locations that hold data temporarily
- Multiple registers are needed to facilitate the operation of the CPU
- The main registers are:
  - The Instruction Register (IR)
  - The Program Counter (PC) or Instruction Pointer (IP)
  - Data registers

# **ALU** (1)

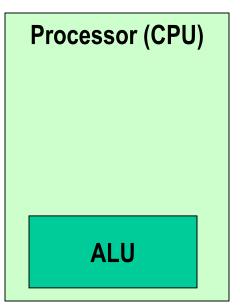
The arithmetic logic unit (ALU) carries out the **logic operations** (such as comparisons) and **arithmetic operations** (add, shift, multiply) required during the program execution

#### The ALU must know

- 1. Which operation to perform
- Where are the input data
- 3. Where to store the output data
- All the above are provided by the CU

#### The ALU performs

- 1. Integer arithmetic operations
  - Add, subtract, increment, decrement
- 2. Bitwise logical operations
  - AND, OR, XOR, NOT, Arithmetic shift, logical shift, rotate



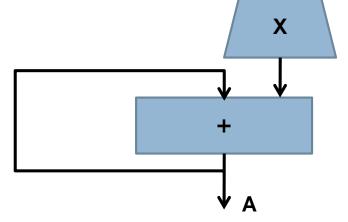
- ALUs often handle the multiplication of two integers, since the result is also an integer
- ALUs typically do not perform division operations, since the result may be a fraction, or a "floating point" number
- The floating-point unit (FPU) performs operations on floating point numbers
- Modern CPUs have separate unit to perform multiplication faster
- Processors include a coprocessor hardware unit which is used to perform more complex mathematical operations such as arcsine, cosine, floatingpoint division, etc

# ALU (3)

- The multiply-accumulate operation is a common step that computes the product of two numbers and adds that product to an accumulator
- It speeds up many computations that involve the accumulation of products, e.g.,
   Matrix-matrix Multiplication
  - The hardware unit that performs the operation is known as a Multiplier
     Accumulator unit (MAC, or MAC unit)

    B C
  - the operation is called MAC or MAC operation
  - The MAC operation performs

$$A = A + B \times C$$



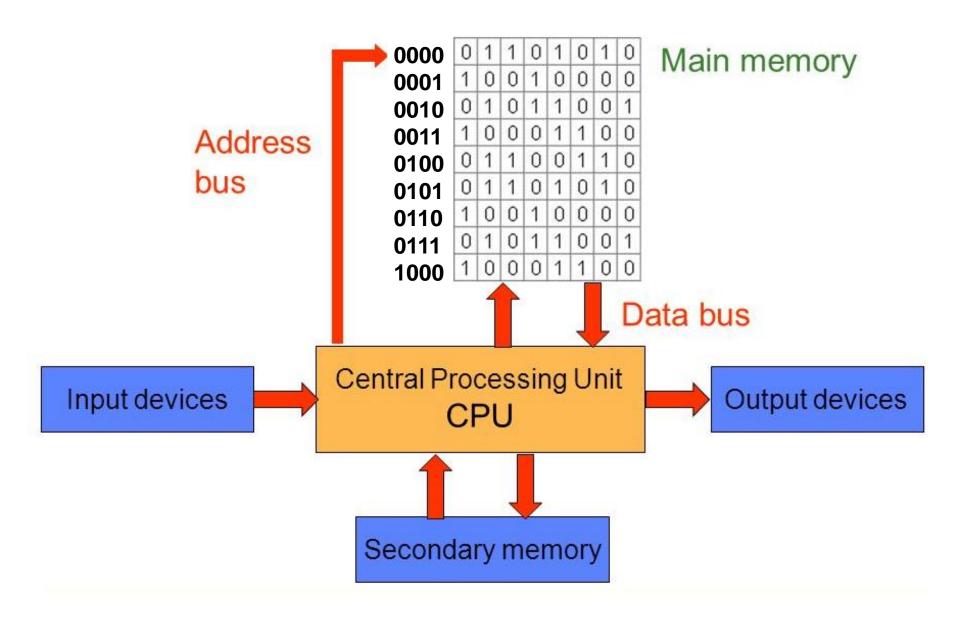
- Arithmetic shift: when shifting to the right, the leftmost bit (the vacant MSB) is filled with the value of the previous MSB (sign)
  - Ideal for signed two's complement binary numbers
- Logical shift: The vacant bits are filled with zero
  - the logical and arithmetic left-shifts are exactly the same
  - Ideal for unsigned binary numbers
- □ Circular shift or bit rotation: In this operation, the bits are "rotated" as if the left and right ends of the register were joined. The value that is shifted in on the right during a left-shift is whatever value was shifted out on the left, and vice versa
  - frequently used in digital cryptography

# Control Unit (CU) (1)

- A Control Unit is the unit that handles the central work of the computer
- There are two registers in the control unit
  - The instruction register (IR) contains the instruction that is being executed
  - The **program counter (PC)** contains the address of the next instruction to be executed
- The CU is responsible of executing the right instruction and organizes the other function units appropriately
- In every clock cycle the CU:
  - Loads from memory the next instruction to be executed PC register contains that address
  - The instruction is stored into IR and is decoded
  - The CU sends the appropriate signals to the ALU, memory, I/O devices in order to execute the instruction
  - The CU increments PC to show next instruction

# Control Unit (CU) (2)

- Program is stored in memory
  - machine language instructions are in binary format
- The task of the control unit is to execute programs by repeatedly:
  - Fetch from memory the next instruction to be executed
  - Decode it, that is, determine what is to be done
  - Execute it by issuing the appropriate signals to the ALU, memory, and I/O subsystems
  - Continues until the program terminates (HALT instruction)



#### Main Memory

- Computer memory consists of a linear array of addressable storage cells that are similar to registers
- □ Both program and data are stored into memory
- Load/store operations are performed on both instructions and data
- Consists of many memory cells (storage units) of a fixed size
  - Each cell has an address associated with it
- All accesses to memory are to a specified address
- The time it takes to fetch/store a word is the same for all words

# Main Memory - Address Space

- To access a word in memory requires an identifier, e.g., int temp. Although programmers use a name to identify a word (or a collection of words), at the hardware level each word is identified by an address
- The total number of uniquely identifiable locations in memory is called the **address space**. For example, a memory with 1 kilobyte (2<sup>10</sup>) and a word size of 1 byte has an address space that ranges from 0 to 1023
- If memory contains N words, then log<sub>2</sub>N bits are needed to address all words in memory
- If the memory address space needs N bits, then there are 2<sup>N</sup> words in memory

# Main Memory - Address Space (2)

Mem. Addr.	Data
00	0
01	1
10	2
11	3

Mem. Addr.	Data
000	0
001	1
010	2
011	3
100	4
101	5
110	6
111	7

- If memory contains 4 words, then log<sub>2</sub>4 =2 bits are needed to address all words in memory
- If the memory address space needs 2 bits, then there are 2<sup>2</sup> words in memory

#### Think Pair Share

1. A computer has 32 MB (megabytes) of memory. How many bits are needed to address any single byte in memory?

The memory address space is 32 MB, or  $2^{25}$  bytes ( $2^5 \times 2^{20}$ ). This means that we need  $\log_2 2^{25}$ , or 25 bits, to address each byte

2. If main memory is of 64Mbyte and every word is of 2 bytes how many bits do we need to address any single word in memory?

The memory address space is 64 MB, which means  $2^{26}$  bytes. However, each word is two ( $2^{1}$ ) bytes, which means that we have  $2^{25}$  words. Note that (Mem.size=number.words x word.size)

This means that we need  $\log_2 2^{25}$ , or 25 bits, to address each word

#### Question1 in previous slide

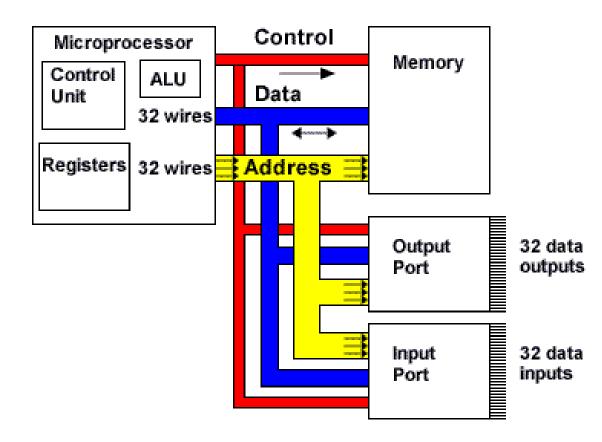
Mem. Addr.	Data
000	0
001	1
010	2
011	3
100	4
101	5
110	6
111	7

#### Question2 in previous slide

Mem.Addr.	Do	ıta
00	0	1
01	2	3
10	4	5
11	6	7

#### How Main Memory and CPU are connected?

- Control Bus sends appropriate signal whether store or load
- Address Bus sends the memory address
- Data Bus sends the data



#### Buses (1)

Bus: a group of wires that transfer data from one part to another (data, address, control)

#### ✓ Data bus:

- bi-directional (read/write)
- 8, 16, 32, 64-bit wide (same as 'word size')

#### ✓ Address bus:

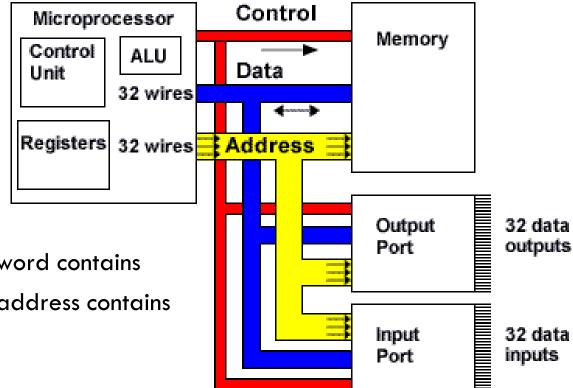
- specifies memory location in RAM/ROM/interface device to be accessed;
   monodirectional
- address space: 16-bit wide  $-> 2^{16}$  words=  $64x2^{10} = 64KB$
- 32-bit wide ->  $2^{32} = 4GB$  this is why in 32-bit PCs we cannot use more than 4Gbyte of RAM

#### ✓ Control bus:

carries commands from the CPU and returns status signals from the devices

# Buses (2)

- How many wires needed for the?
  - A. Data bus?
  - B. Address bus?
  - C. Control Bus?



- A. As many bits as the memory word contains
- B. As many bits as the memory address contains
- C. 1 bit is enough

#### Think Pair Share

Q: If main memory is of 64Kbyte and every word is of 8 bytes how many wires do we need for the address bus?

# Assembly basic instructions

Depending on the target CPU, different assembly instructions exist

Instr:	<u>Meaning:</u>

Store OxA2B, R1 Store R1 in A2B<sub>16</sub> memory location

ADD R1, R2 Add R1 and R2 and store the result in R1

SUB R1, R2, R3 R1 = R2 - R3

MUL R1, R2 R1 = R1 \* R2

DIV R1, R2 R1 = R1 / R2

INC R4 R4=R4+1

HALT Stops program execution

CMP R1, 10 If R1=10, then set the register EQ=1, else EQ=0

JMP EQ S1 Load next instruction from memory location S1, if EQ==1

 Assemblers translate instructions that are comprehensible to humans into the machine language that is comprehensible to computers

#### Machine Language or Machine Code (1)

- A <u>program</u> consists of a sequence of instructions (in binary)
- EACH instruction specifies both:
  - The operation to perform
  - The address of the data
- Instructions are stored and processed in machine language--also called microcode
- Like everything else (e.g. like ASCII characters) machine language consists solely of <u>bit patterns</u>

# Machine Language or Machine Code (2)

- □ A machine language instruction consists of:
  - Operation code, specifying which operation to perform
  - Address field(s), specifying the memory addresses of the values on which the operation works

Opcode (8 bits)	Address 1 (16 bits)	Address 2 (16 bits)
00001001	000000001100011	000000001100100

Instructions are given to the processor in the form of a program ... so it knows what circuits to use, in what order; and from where the data should be read or to where it should be stored

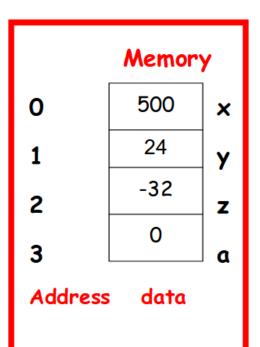
#### Machine Language or Machine Code (3)

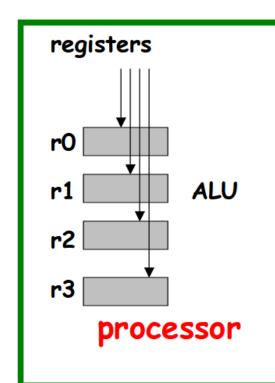
- Machine languages consist entirely of binary numbers and are almost impossible for humans to read and write
- Assembly languages have the same structure and set of commands as machine languages, but they enable a programmer to use names instead of numbers
- Each type of CPU has its own machine language and assembly language
  - an assembly language program written for one type of CPU won't run on another
- In the early days of programming, all programs were written in assembly language
- Now, most programs are written in a high-level language such as Java, Python, C/C++
- Programmers still use assembly language when speed is essential

#### High level language VS Assembly language

High Level Language (HLL)

$$a = x + y - z$$





Assembly Language (AL):

load x into r1

load y into r2

load z into r0

$$r3 \leftarrow r1 + r2$$

store r0 into a

- In HLL we write a=5 and we assume that it is stored somewhere – we don't care
- But the computer does not work like that – it has to store every variable in an exact memory location
- In assembly we must specify the Hardware registers as well as the memory locations

#### High level language VS Assembly language

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- Easily understandable
- Portable run on just any computer
- Debugging of the code is easy
- High level languages like java, C++, etc. have one to many relationship with assembly, i.e., one statement of java expands into many assembly language commands
- High level language is always converted into assembly language
- The high level language programmer doesn't need to know the HW details

- Hard to understand
- Runs only on the target CPU only
- Debugging is very hard
  - One to one or one to a few relationship
  - The assembly language programmer must know about the hardware such as registers, etc.
- Assembly language can control the machine code better
- Assembly is much faster

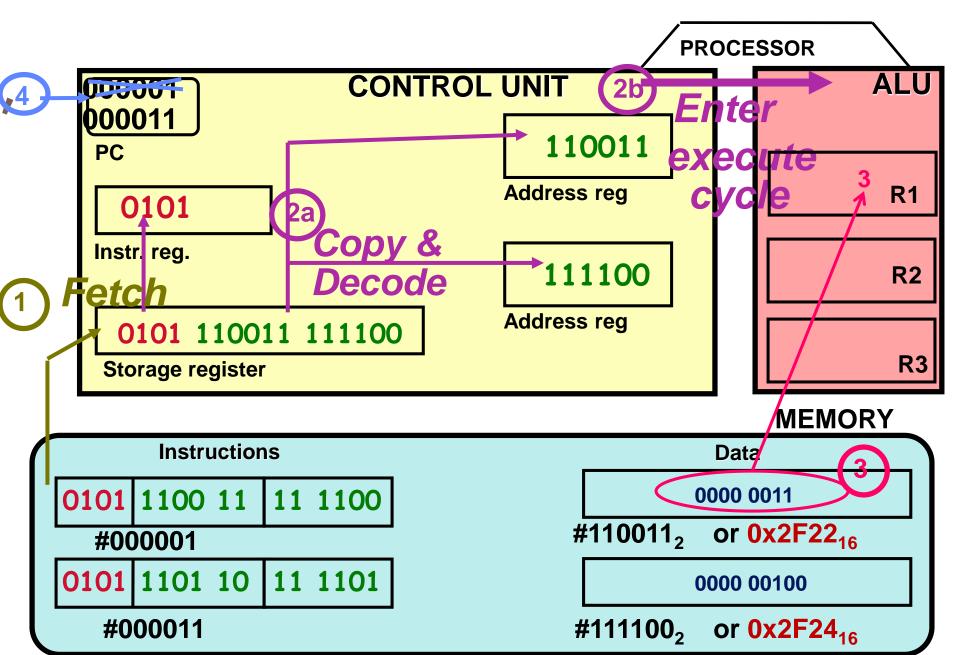
#### **Program Execution**

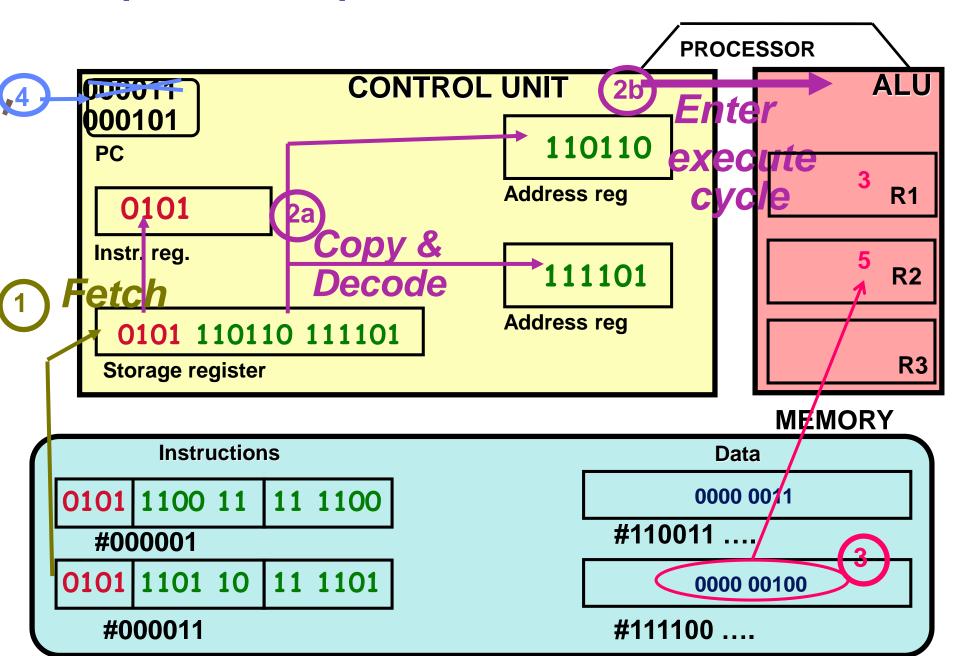
1. PC is set to the address where the first program instruction is stored in memory

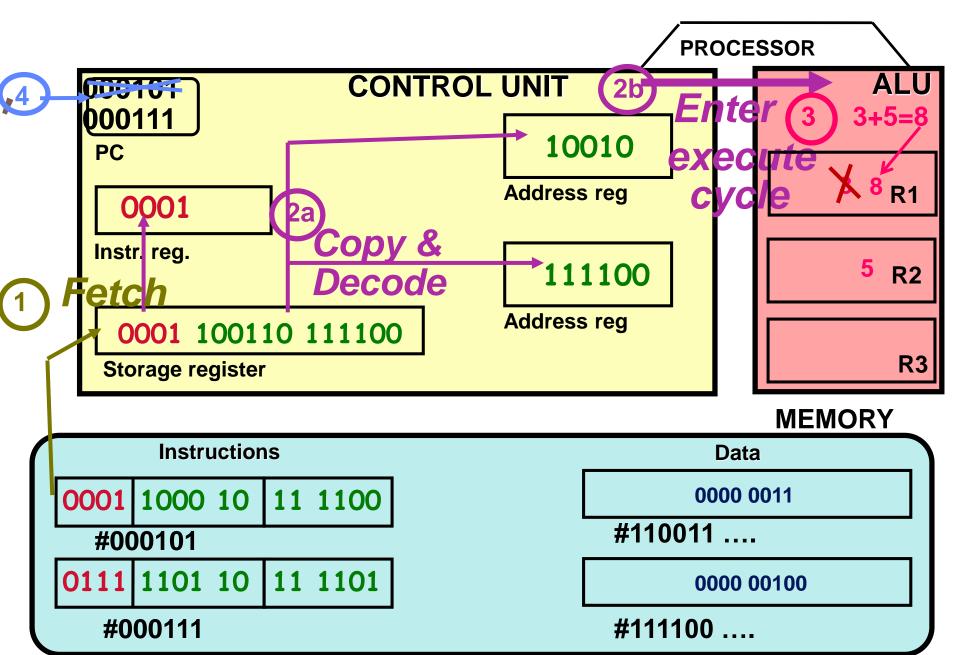
#### 2. Repeat until HALT instruction or fatal error

- 2a. Fetch instruction
  - Fetches instruction (from memory) at address given by PC; copies it into storage register
- 2b. Decode instruction
  - Copies op code into IR and operands into address registers
  - Interprets instruction
  - ALU is invoked
- 2c. Execute instruction
  - Execution cycles vary, depending on the op code (instruction), e.g., Load copies data from memory to ALU register, ADD adds values inside the ALU
- 2d. Increment PC by one, now contains the memory address of the *next* instruction that will be fetched

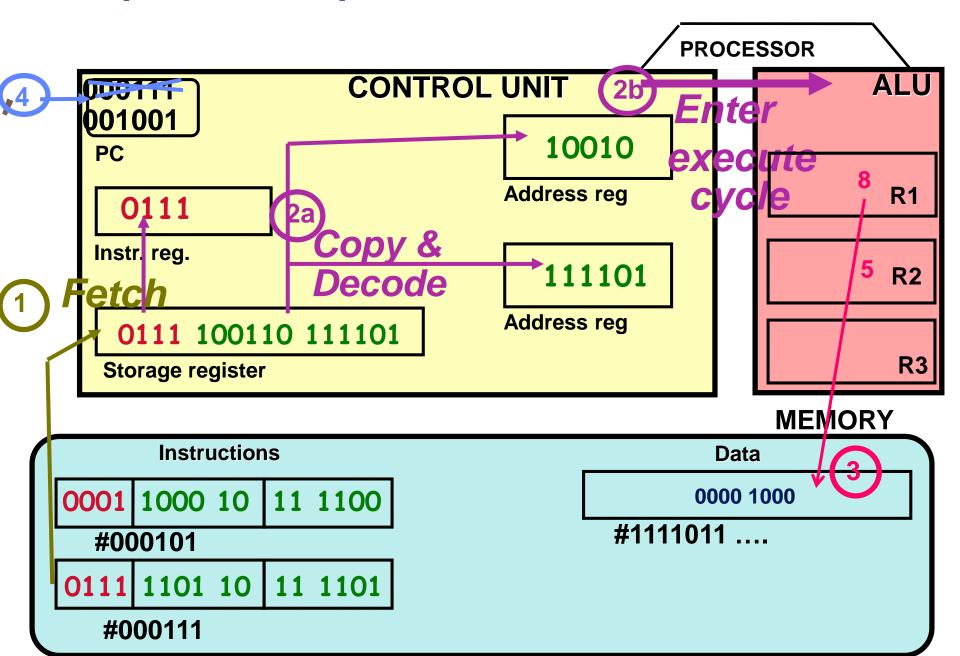
Load R1, 0x2F22
Load R2, 0x2F24
ADD R1, R2
Store 0x2F24, R1

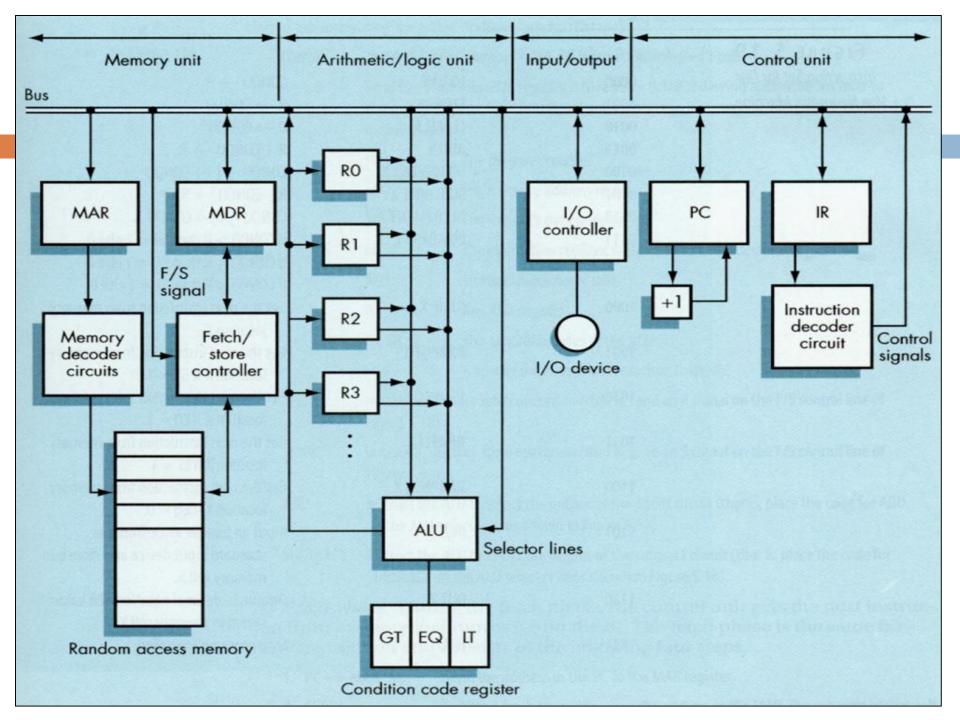






# A simplified example



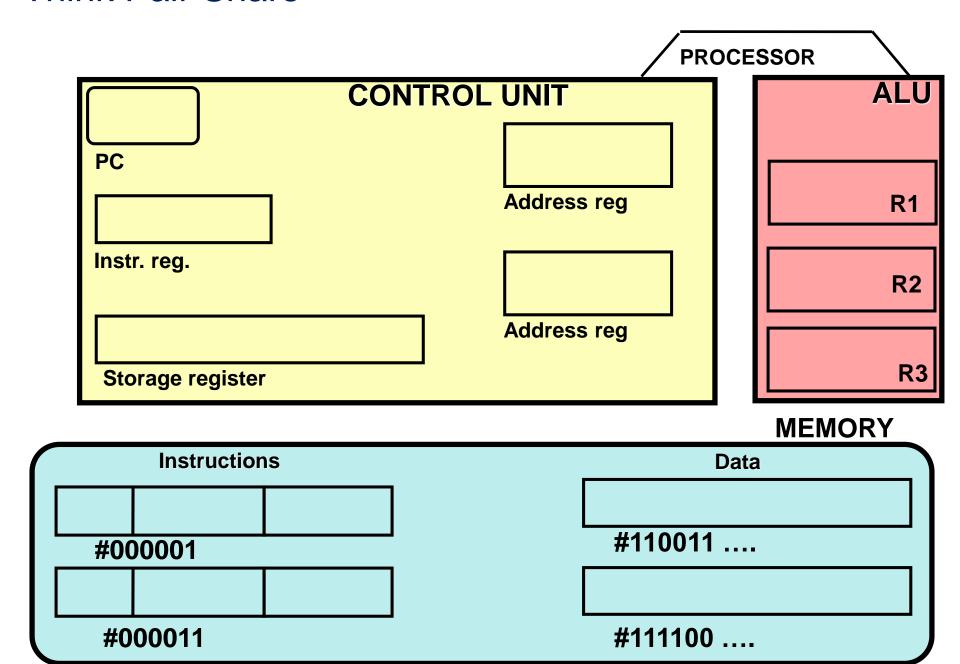


#### Think Pair Share

Given the following high level code, first translate it into assembly pseudo code and second explain the CPU steps as before. Consider that a=0, x=2, y=3, z=4.

$$a = x + y - z;$$

#### Think Pair Share



#### Memory Types (1)

- Random Access Memory (RAM) Alternatively referred to as main memory
  - Static RAM (SRAM) fast but more expensive
  - Dynamic RAM (DRAM) slow but cheap
  - Synchronous Dynamic RAM (SDRAM)
  - □ DDR SDRAM Double Data Rate SDRAM (transferring data on both the rising and failing edges of the clock signal)
- Read Only Memory (ROM)
  - Programmable read-only memory (PROM)
  - Erasable programmable read-only memory (EPROM)
  - □ Electrically erasable programmable read-only memory (EEPROM)
- RAM loses any information it is holding when the power is turned off
- □ ROM is meant for permanent storage, while RAM is for temporary storage

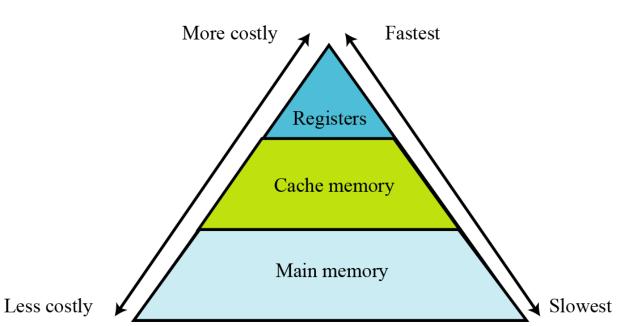
# Memory Types (2)

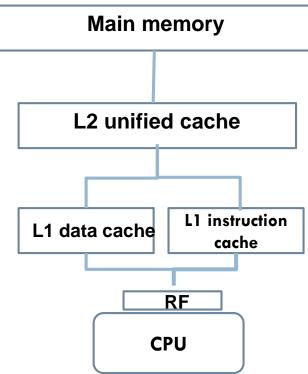
- ROM is used primarily in the start up process of a computer, whereas a RAM chip is used in the normal operations of a computer once the operating system has been loaded
- A good example of ROM is the computer BIOS, a PROM chip that stores the programming needed to begin the initial computer start up process
- Writing data to a ROM chip is a much slower process than writing it to a RAM chip
- A RAM chip can store multiple gigabytes (GB) of data, ranging from 1
   GB to 256 GB per chip.
- A ROM chip stores several megabytes (MB) of data, typically 4 MB or 8
   MB per chip

## Memory Hierarchy

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- users want unlimited fast memory
- fast memory is expensive, slow memory is cheap
- cache: small, fast memory near CPU
- main memory, disk : large, slow memory





# Cache Memory

- Cache is a high-speed static random access memory (SRAM) that a CPU can access more quickly than it can access regular random access memory (RAM)
- This memory is typically integrated directly into the CPU chip
- Cache memory is faster than main memory, but slower than the CPU and its registers
- Cache memory, which is normally small in size, is placed between the CPU and main memory
- The purpose of cache memory is to store program instructions and data that are used repeatedly - The computer processor can access this information quickly from the cache rather than having to get it from computer's main memory
- Fast access to these instructions increases the overall speed of the program

## Secondary memory (1)

- Secondary memory is where programs and data are kept on a longterm basis
  - Common secondary storage devices are the hard disk and optical disks
- The hard disk has enormous storage capacity compared to main memory
- The hard disk is used for long-term storage of programs and data
- Data and programs on the hard disk are organized into files
  - A file is a collection of data on the disk that has a name

#### Secondary memory (1)

- Running programs are always located in main memory
- When creating a new file and type something it is stored into main memory; When you "save" your document, the characters are copied to a file on the hard disk
- A permanent copy will also be in secondary memory on the hard disk

Main Memory Hard Disc

Fast Slow

Expensive Cheap

Low Capacity High Capacity

Question: Do you think that data transfer from the network is slower or faster than from main memory?

**Answer:** Data transfers from the network are much slower than from main memory

# Questions?

# Further Reading

 Chapter 14 in in 'Computer Organization and architecture' available at

http://home.ustc.edu.cn/~leedsong/reference\_books\_tools/Computer%20Organization%20and%20Architecture%2010th%20-%20William%20Stallings.pdf