



# **Lecture - 07**

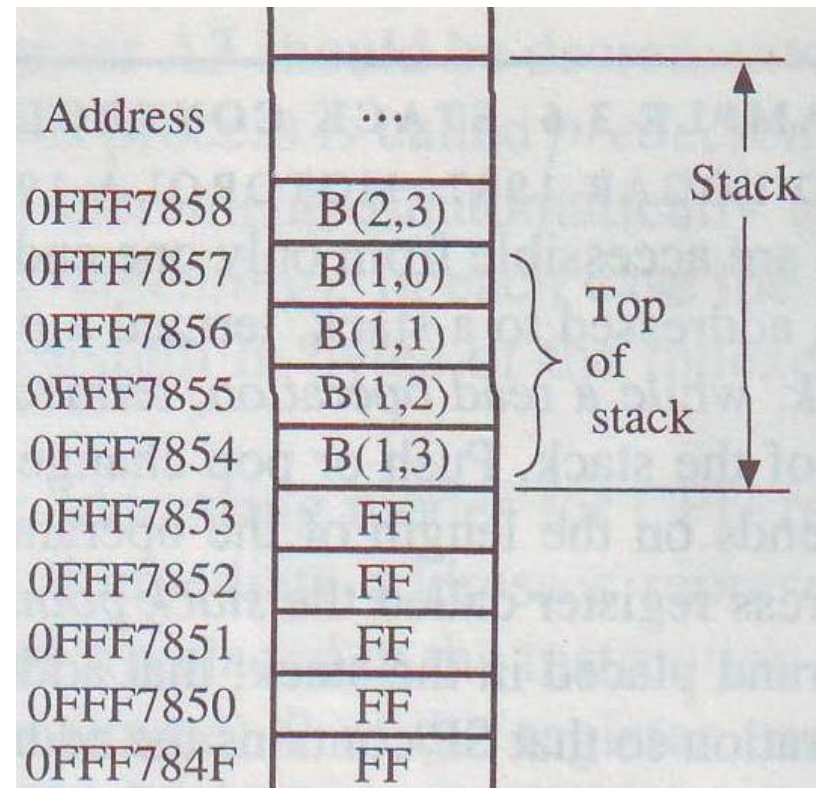
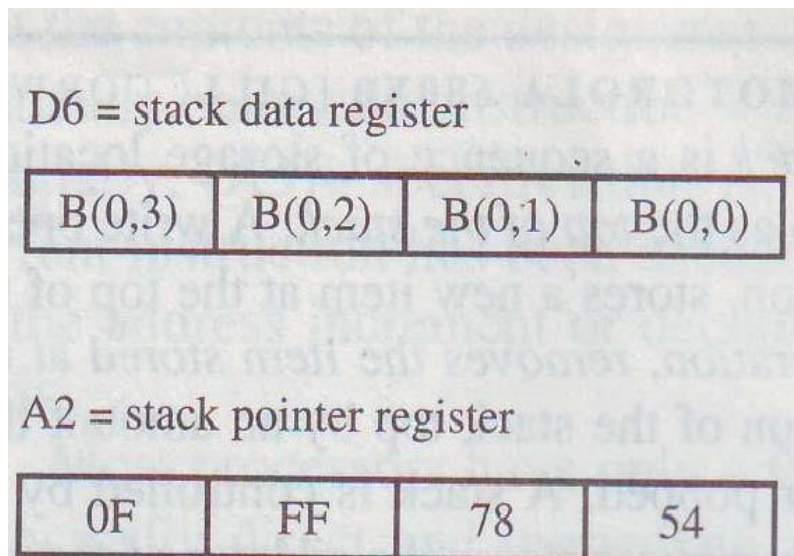
# Stack Control

- A **stack** is a sequence of storage locations that are accessible from only one end referred to as the **top of the stack**.
- Two operations: **Push** and **Pop**.
- Push or pop changes the position of the stack top by an amount that depends on the length of the operand pushed or popped.
- A stack is controlled by an address register called the **stack pointer SP**. It contains the address of the new stack top.

# Stack Control in Motorola 680X0

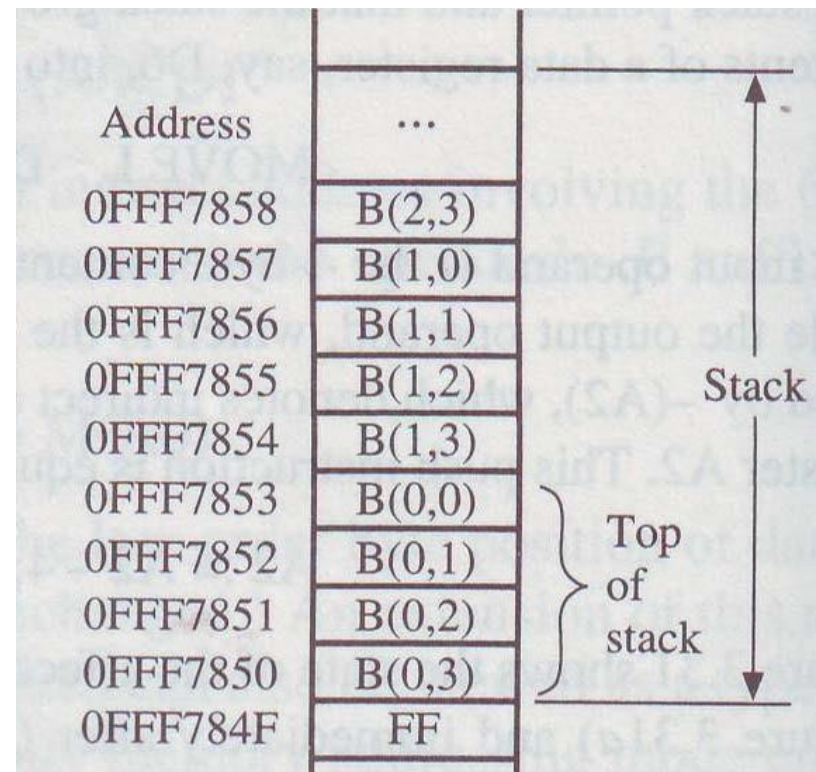
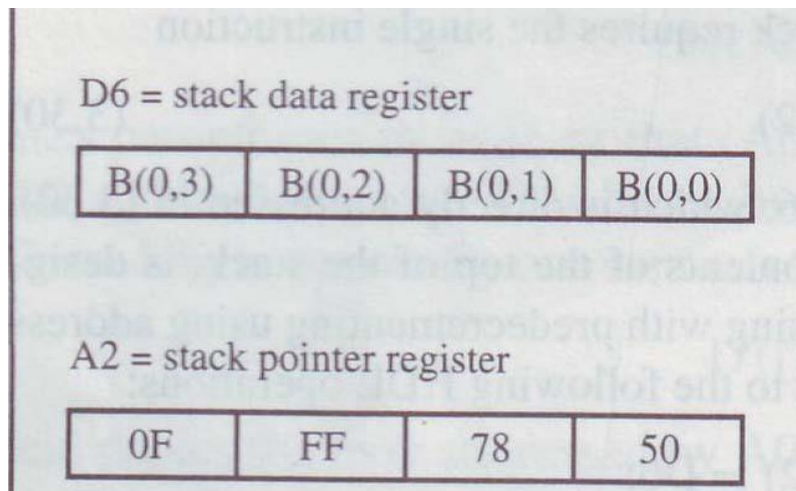
- Motorola 680X0 has no explicit hardware for stack support. But its various addressing modes make it easy to treat any contiguous region of its external memory M as a stack.
- Suppose the address register A2 of the 680X0 is designated as stack pointer and the stack grows toward the low addresses of M.
- To push the content of D6 (data register) into the stack requires a single instruction `MOVE.L D6, -(A2)` which is equivalent to  $A2 := A2 - 4$  and  $M(A2) := D6$ .
- The pop instruction is `MOVE.L (A2)+, D6` which is equivalent to  $D6 := M(A2)$  and  $A2 := A2 + 4$

# Stack Control in Motorola 680X0



**Prior the execution of `MOVE.L D6, -(A2)`**

# Stack Control in Motorola 680X0



**After the execution of MOVE.L D6, -(A2)**

# Number of Addresses

- **Three-address instruction:**

ADD Z, X, Y

$Z := X + Y.$

- **Two-address instruction:**

ADD X, Y

$X := X + Y.$

- **One-address instruction:**

ADD X

$AC := AC + X.$

# Number of Addresses

**Implement  $X := A \times B + C \times C$ , where A, B, C and X are stored in the memory.**

Instruction	Comment
LOAD A	$AC := A$
MULTIPLY B	$AC := AC * B$
STORE T	$M(T) := AC$
LOAD C	$AC := C$
MULTIPLY C	$AC := AC * C$
ADD T	$AC := AC + T$
STORE X	$M(X) := AC$

**One address machine**

Instruction	Comment
MOVE T, A	$T := A$
MULTIPLY T, B	$T := T * B$
MOVE X, C	$X := C$
MULTIPLY X, C	$X := X * C$
ADD X, T	$X := X + T$

**Two address machine**

Instruction	Comment
MULTIPLY T, A, B	$T := A * B$
MULTIPLY X, C, C	$X := C * C$
ADD X, X, T	$X := X + T$

**Three address machine**

# Accumulator Architectures

- **Instruction set:**

add A, sub A, mult A, div A, . . .

load A, store A

- **Example:  $A*B - (A+B*C)$**

load B

mul C

add A

store D

load A

mul B

sub D







# Accumulator Architecture : Pros and Cons

## ■ Pros

- Very low hardware requirements
- Easy to design and understand

## ■ Cons

- Accumulator becomes the bottleneck
- Little ability for parallelism or pipelining
- High memory traffic

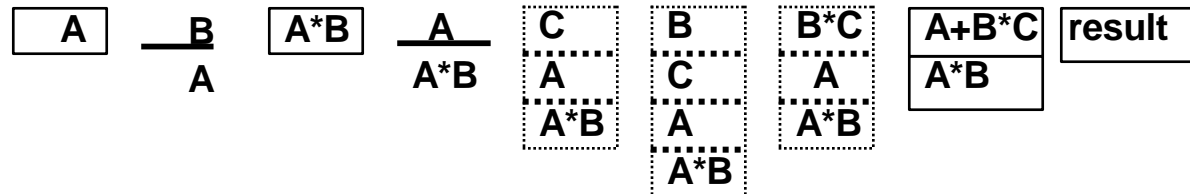
# Zero Address Machine (Stack Based Architecture)

- Addresses are eliminated by storing operands in a push-down stack.
- All the operands are required to be in the top locations in the stack.
- Stack pointer automatically keeps track of the stack top.
- Push and pop operations are needed to transfer data to and from the stack.
- For example,  $X+Y$  is invoked by ADD that causes the top two operands, which should be  $X$  and  $Y$ , to be removed from the stack and added. The resulting sum  $X+Y$  is then placed at the top of the stack.

# Stack Architectures

**Example:  $A*B - (A+C*B)$**

push A  
push B  
mul  
push A  
push C  
push B  
mul  
add  
sub





# Stack Architecture: Pros and Cons

## ■ Pros

- implicit operand addressing. Maintained by top of the stack.
- Low hardware requirements.

## ■ Cons

- Stack becomes the bottleneck.
- Little ability for parallelism or pipelining.
- Data is not always at the top of stack when need, so additional instructions like POP and SWAP are needed.



# The Requirements for Instruction Set

- It should be **complete**. We should be able to construct a machine language program to evaluate any function that is computable using a reasonable amount of memory space.
- It should be **efficient**. The frequently required functions can be performed rapidly using relatively few instructions.
- It should be **regular**. The instruction set should contain expected opcodes and addressing modes.
- The instruction should be **compatible** with those of existing machine.

# Classification of Instruction

- **Data transfer instructions** – Copy information.  
Example: *load, store, move* instruction.
- **Arithmetic instructions** – Perform operations on numerical data. Example: *add, sub, mul, div* etc.
- **Logical instructions** - Boolean and non-numerical operations. Example: **AND, NOT** etc.
- **Program control instructions** – It change the sequence in which program are executed. Example: Branch instruction.
- **Input-Output (IO) instructions** – It cause information to be transferred to or from external IO devices.  
Example: PRINT LINE.