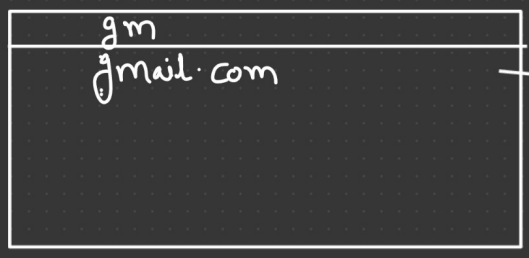
String Algorithms



→ Auto - completion.

Prateek

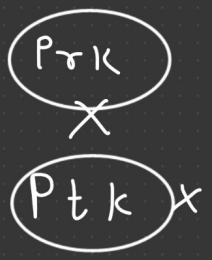
Prefix

P
Pr
Pra
Prat
Prate
Pratee
Prateek

Suffix

K
eK
eeK
teeK
ateeK
rateeK
PrateeK

Compare (prefix, Text list)
{
return list
}



Text = "My name is PrateeK" → string Matching

Pattern = "PrateeK"

1) Brute force approach →

2) Robin Karp

3) Kmp

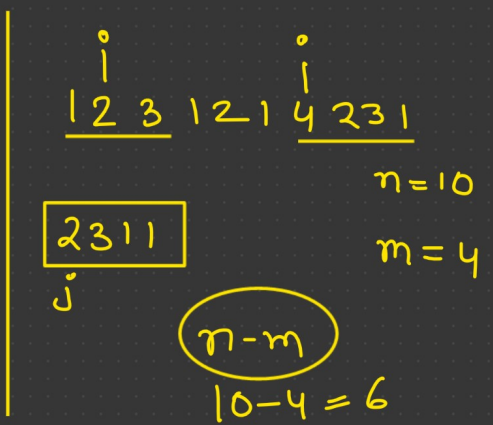
Brute force Approach :-

Text = "ABCAABCD" = n

Pattern = "AAB" = m

```
int BruteforceStringMatch(Char T[], char P[], int n, int m)
{
    for(int i=0; i<=n-m; i++)
    {
        int j=0;
        while( j<m && T[i+j]==P[j])
        {
            j++;
        }
        if(j==m)
            return i;
    }
    return -1;
}
```

$O(n \times m)$



2) Rabin Karp :- $O(n+m)$

Hashfunction (T(i, j) , p) → O(1) → may be

- (0,4) (1,5)
- p → m = 4
- Not possible

T = 1 2 3 1 2 1 1 3 , 121 → p, m=3

i, j

1×10+2 = 12×10+3 = 123

2×10+3 = 23×10+1 = 231

if (T % 10 == P % 10) { } → may be → O(1)

T = 123 % 10 = 3

P = 121 % 10 = 1

123 % 100 = 23×10+1 = 231

$O(n+m)$

T = 231 % 10 = 1

P = 121 % 10 = 1 → may be

231 % 100 = 31×10+2 = 312

12 % 10 = 2 × 10 = 20 + 1 = 21

123 % 100 = 23 × 10 = 230 + 1 = 231

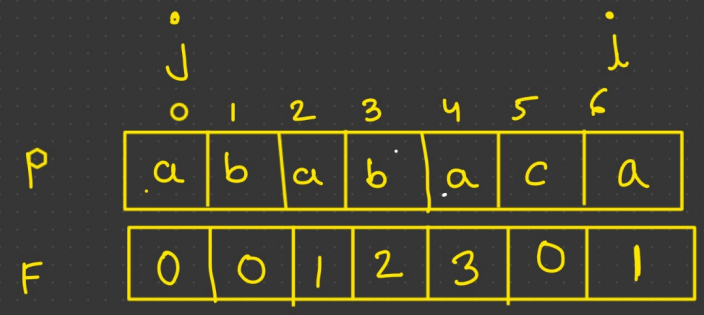
5641 % 1000 = 641×10 = 6410 + 1 = 6411 → O(1)

KMP Algorithm

Knuth Morris Pratt

Prefix Table

Ex:-




```
if(j==0)
{
  F[i]=0
    i++;
}
if( j>0)
  j=F[j-1]
```

→ ← a b a

a	a ✓
ab	ba

a	a ✓
ab	ba ✗
aba	aba ✓
abab	baba ✗

```
Void prefixTable(int P[], int m)
{
    int  i=1, j=0, F[0]=0;

                                  ⟶ O(m)
    while(i<m)
    {
        if(P[i]==P[j])
        {
            F[i]=j+1;
            i++;
            j++;
        }
        else if(j>0)
        {
            j=F[j-1];
        }
        else { F[i]=0;
               i++;
        }
```

```
int KMP( char T[], Char P[], int n, int m)
{
    int i=0, j=0;
    prefixTable( P, m); ⟶ O(m)

    while (i<n)  ⟶ O(n)
    {
        if ( T[i]== P[i])
        {
            i++;  j++;

            if (j== m)
                return i-j;
        }
        else if (j>0)
        {
            j = F[j-1];
        }
        else
            i++;
    }
    return -1;
}
```

O(n+m)