



Semester : VI

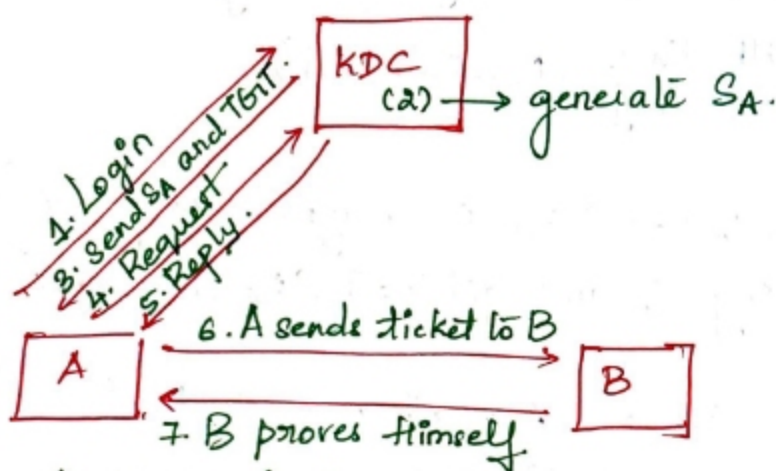
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### KERBEROS:

It performs the following tasks:

- \* Kerberos login.
- \* Kerberos ticket



#### Step 1: Kerberos login.

KDC → Key Distribution Center.

SA → Session Key of A.

TGT → Ticket Granting Ticket.

$K_A$  → Symmetric key of A.

$K_B$  → Symmetric key of B.

$K_{AB}$  → Session key for both A and B.

#### Step 1: Kerberos Login.

A  $\xrightarrow{\text{login}}$  KDC

$h[\text{password}] \Rightarrow K_A$  is derived from password.  
 $K_A$  is the hash value of the password.

$K_A$  is known by A and KDC.



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Step 2: KDC generates  $S_A$ .

After A logs in KDC generates session key for A.

Step 3: Send  $S_A$  and TGT.

$S_A = E("A", S_A, K_A) \rightarrow S_A$  is encrypted using  $K_A$  (Sender).

$\Rightarrow$  A will decrypt with  $K_A$  (Receiver end).  
A will get  $S_A$ .

$TGT = E("A", TGT, S_A) \rightarrow TGT$  is encrypted using  $S_A$  (Sender).

$\Rightarrow$  A will decrypt using  $S_A$  and get TGT (Receiver end).

Step 4: Request

A sends request to TGT  $\Rightarrow$  A wants to communicate with B.

REQUEST = (TGT, authenticator).

authenticator =  $E(\text{timestamp}, S_A)$ .

$= E(TGT, E(\text{timestamp}, S_A))$

KDC will decrypt using  $S_A$  and will receive timestamp and TGT.

Step 5: Reply  $\rightarrow$  KDC issues a ticket to A and B.

Reply =  $E("B", K_{AB}, \text{Ticket to B}, S_A)$ .

Ticket to B =  $E("A", K_{AB}, K_B)$ .

Reply =  $E("B", K_{AB}, E("A", K_{AB}, K_B), S_A)$ .

A will decrypt with  $S_A$ , he will get  $K_{AB}$  and Ticket to B.

A cannot decrypt Ticket to B, because he don't have  $K_B$ .





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Step 6: A sends ticket to B

$A \rightarrow B$ . (Ticket to B, authenticator).

$(E("A", K_{AB}, K_B), E(\text{timestamp}, K_{AB}))$

1. B will decrypt using  $K_B$  and get  $K_{AB}$ .
2. B will decrypt using  $K_{AB}$  and get timestamp.

Step 7: B proves himself.

Now both got  $K_{AB}$ .

$E(\text{timestamp} + 1, K_{AB})$ .

B will send the timestamp by encrypting using  $K_{AB}$ .

A will decrypt using  $K_{AB}$ .

B proved that he got the same session key as A has got. Now they are ready to communicate.