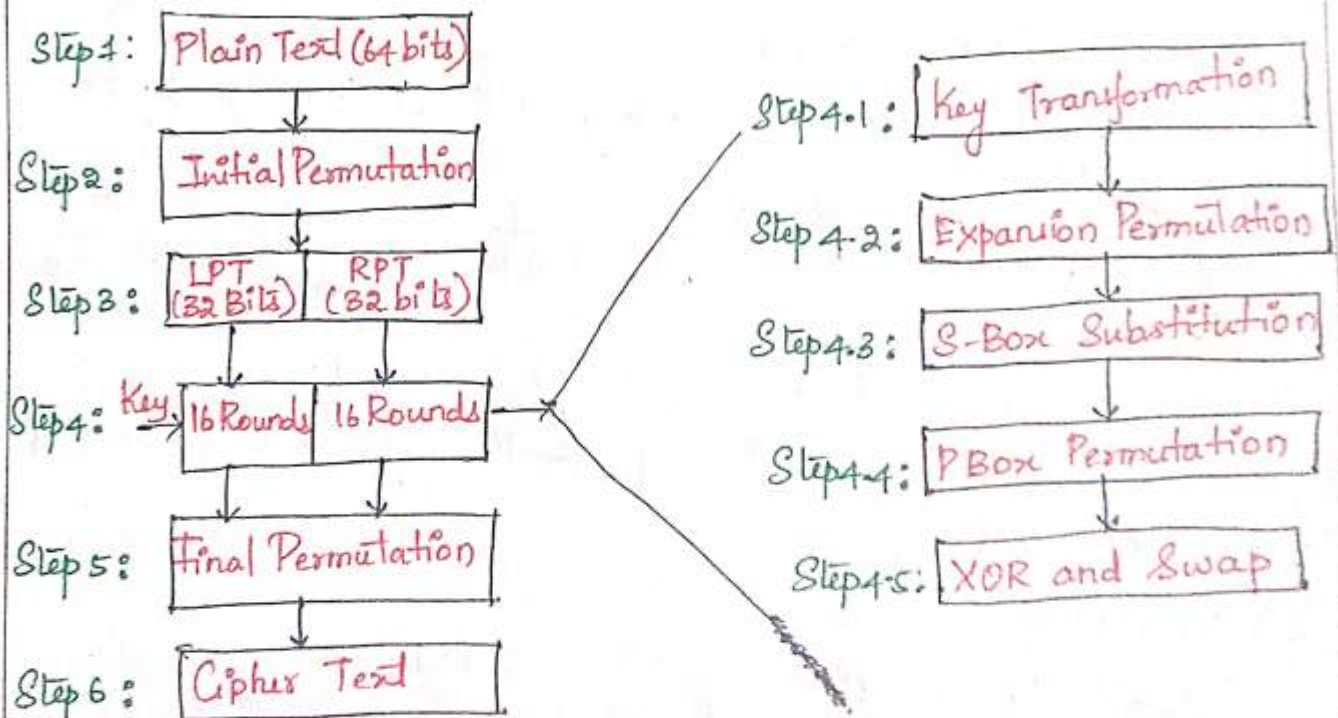


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DES (Data Encryption Standard) :

- * DES is a block cipher.
- * It encrypts data in blocks of size of 64 bits
- * 64 bits of plaintext goes as an input to DES, which produces 64 bits of Cipher text.
- * The same algorithm and key are used for encryption and decryption, with minor difference.
- * The key length is 56 bits.

Broad level steps in DES (or) Block diagram of DES.Step 2: Initial Permutation

- * There are pre-defined permutation rules for initial permutation. (Initial Permutation Table) Refer PPT.
- * For example, the 58th bit position in the plain text takes 1st position in new plain text, the 50th bit position in

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the plain-text ~~becomes~~ takes 2nd bit position in the new generated plain text.

* After initial permutation, a new 64 bit plain text is derived.

Step 3: LPT and RPT

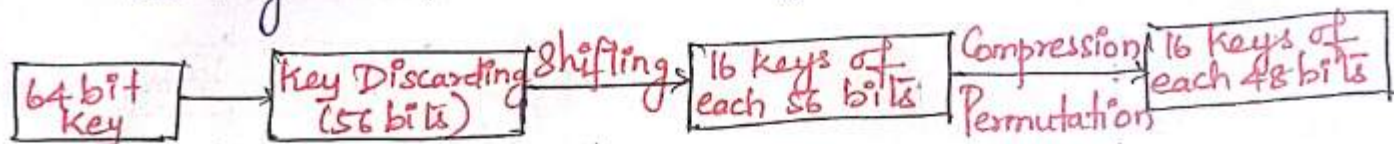
* The resulting 64-bit permuted text block is divided into two half blocks.

* Each half block consists of 32 bits.

* LPT (Left Plain Text \rightarrow 32 bits) and RPT (Right Plain Text \rightarrow 32 bits).

Step 4.1: Key Transformation

The key transformation undergoes the following steps:



* DES uses 16 rounds.

* For each round a unique key is used.

* Using 1 key a total of 16 keys are generated for each round.

Key Discarding:

* Initially a 64 bits key is selected.

* Every 8th bit of original key is discarded.

* ~~From~~ After discarding one key of 56 bits are generated.

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Generate 16 keys of 56 bits:

To generate 16 keys, it will follow the below rule. It will perform left circular shift according to this table.

Round	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
#key bits shifted	1	1	2	2	2	2	2	2	1	2	2	2	2	2	2	1

Example:

1111000^7 1010101^{14} 0011010^{21} 0000101^{28}
 0101010^{35} 1100110^{42} 0011110^{49} 1100101^{56}

According to the table, for the 1st key it will do 1 bit left circular shift.

$L_1: 1110001$ 0101010 0110100 0001011

$R_1: 1010101$ 1001100 0111011 1001010

⋮

 L_{16} R_{16}

$L_2 R_2$ is derived from $L_1 R_1$ by performing 1 bit left circular shift.
 $L_3 R_3$ is derived from $L_2 R_2$ by performing 2 bit left circular shift.

The same procedure is repeated for 16 times and 16 different keys of each 56 bits are generated.

Compression Permutation:

* Apply compression permutation using compression permutation table on $L_1 R_1, L_2 R_2, \dots, L_{16} R_{16}$ (Refer PPT)

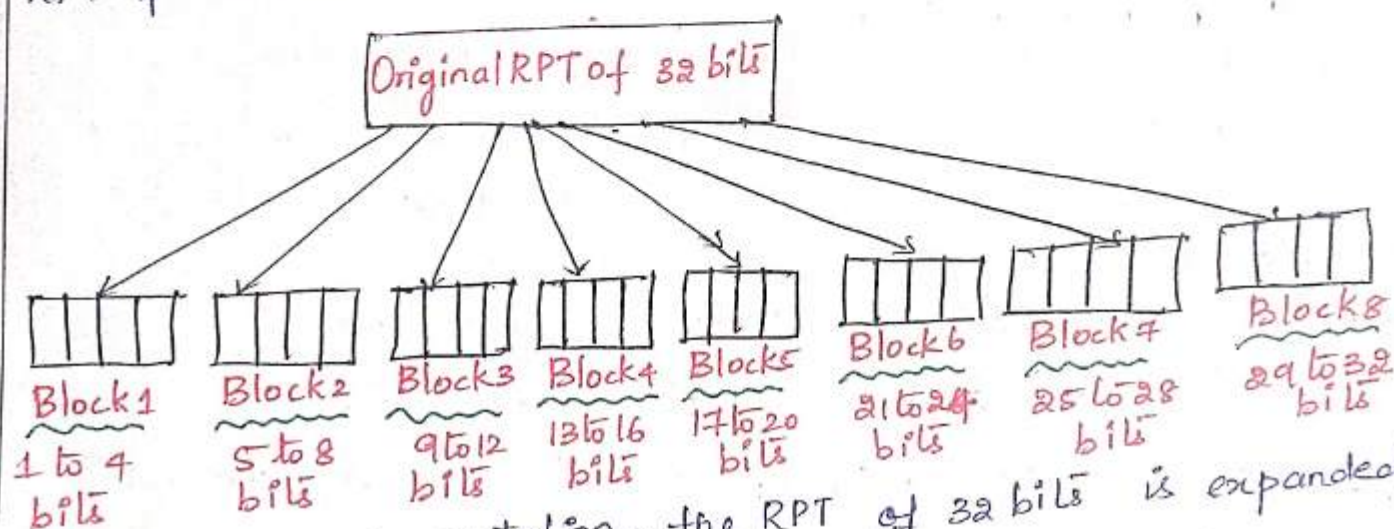
* The process is continued and 16 keys of 48 bits ($K_1, K_2, K_3, K_4, \dots, K_{16}$) are generated.

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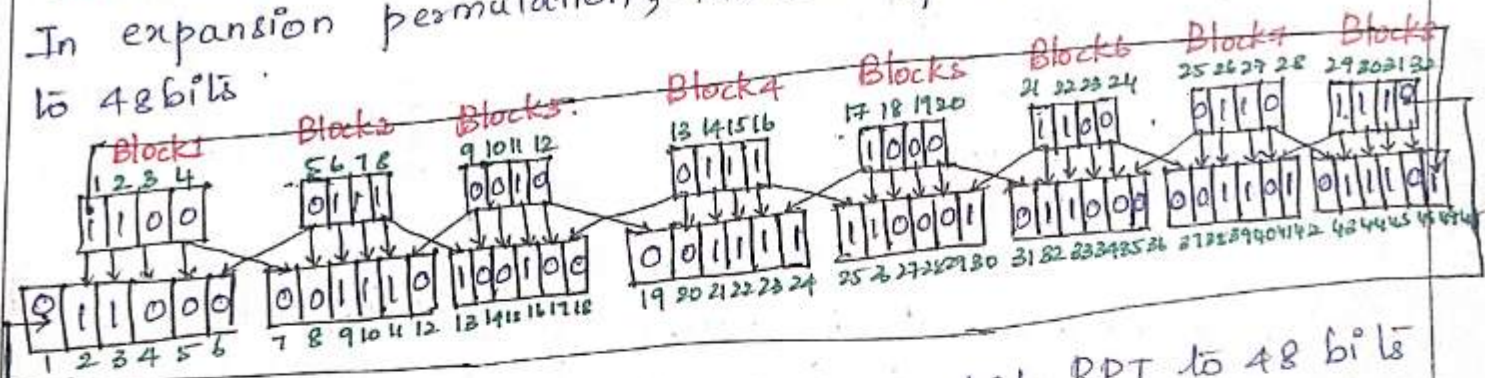
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Step 4.2 :- Expansion Permutation:

The expansion permutation is applied on the original RPT of 32 bits.



In expansion permutation, the RPT of 32 bits is expanded to 48 bits.

Example:

Using Expansion permutation expand 32bit RPT to 48 bits

1100 0111 0010 0111 1000 1100 0110 1110 → 32 bits

011000 001110 100100 001111 110001 011000 001101 011011

48 bits RPT

Way to S-Box Substitution:

- * The first key (K_1) of 48 bits is XORed with the 48 bit RPT generated in Expansion permutation.
- * A new 48 bits RPT is generated which will undergo



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S-Box substitution.

Key Transformation
(48 bits) K_1

Expansion Permutation
(48 bits) RPT

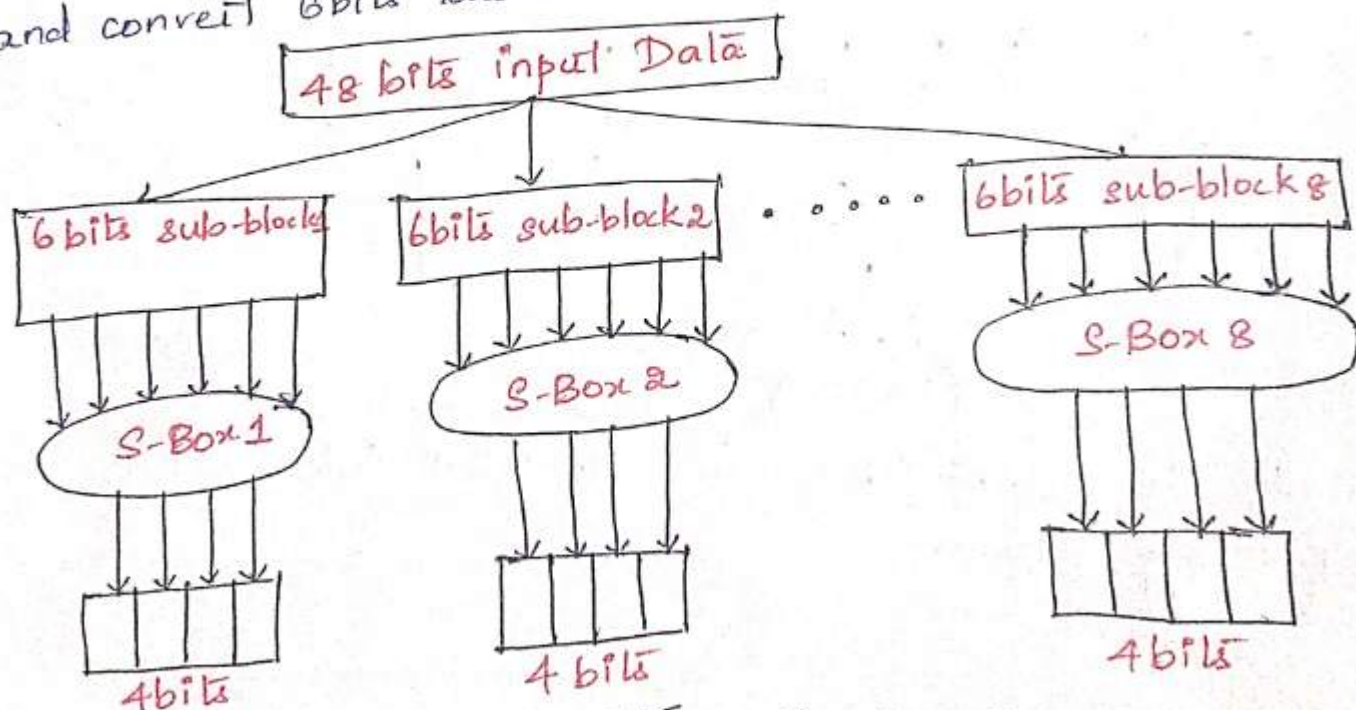
XOR.

48 bits RPT.
(The input for S-Box substitution)

S-Box Substitution:

* The 48 bits RPT is divided into 8 sub-blocks with 6 bits in each block.

* Each sub-block will refer to S-Box 1 to S-Box 8 and convert 6 bits block to 4 bits block.



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* The 1st and the last bit will refer the row value of respective S-Box, whereas the middle four bits will refer the column value of that S-Box.

* For S-Box 1 refer S-Box 1 table, for S-Box-2 refer S-Box 2 table, for S-Box 3 refer S-Box 3 table and so on.

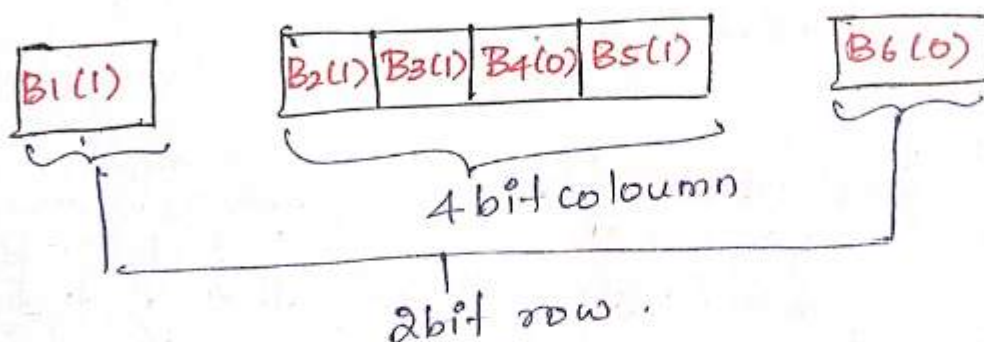
Example:

$\begin{array}{c} 111010 \\ \text{S-Box 1} \\ 1010 \\ 100111 \end{array}$
 $\begin{array}{c} 111110 \\ \text{S-Box 2} \\ 1111 \end{array}$
 $\begin{array}{c} 100101 \\ \text{S-Box 3} \\ 1101 \end{array}$
 $\begin{array}{c} 110000 \\ \text{S-Box 4} \\ 1111 \end{array}$
 $\begin{array}{c} 011101 \\ \text{S-Box 5} \\ 1000 \end{array}$
 $\begin{array}{c} 011111 \\ \text{S-Box 6} \\ 1000 \end{array}$
 $\begin{array}{c} 011100 \\ \text{S-Box 7} \\ 0110 \end{array}$

S-Box 8: 0111

S-Box 4: 111010

B ₁ (1)	B ₂ (1)	B ₃ (1)	B ₄ (0)	B ₅ (1)	B ₆ (0)
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(1010)

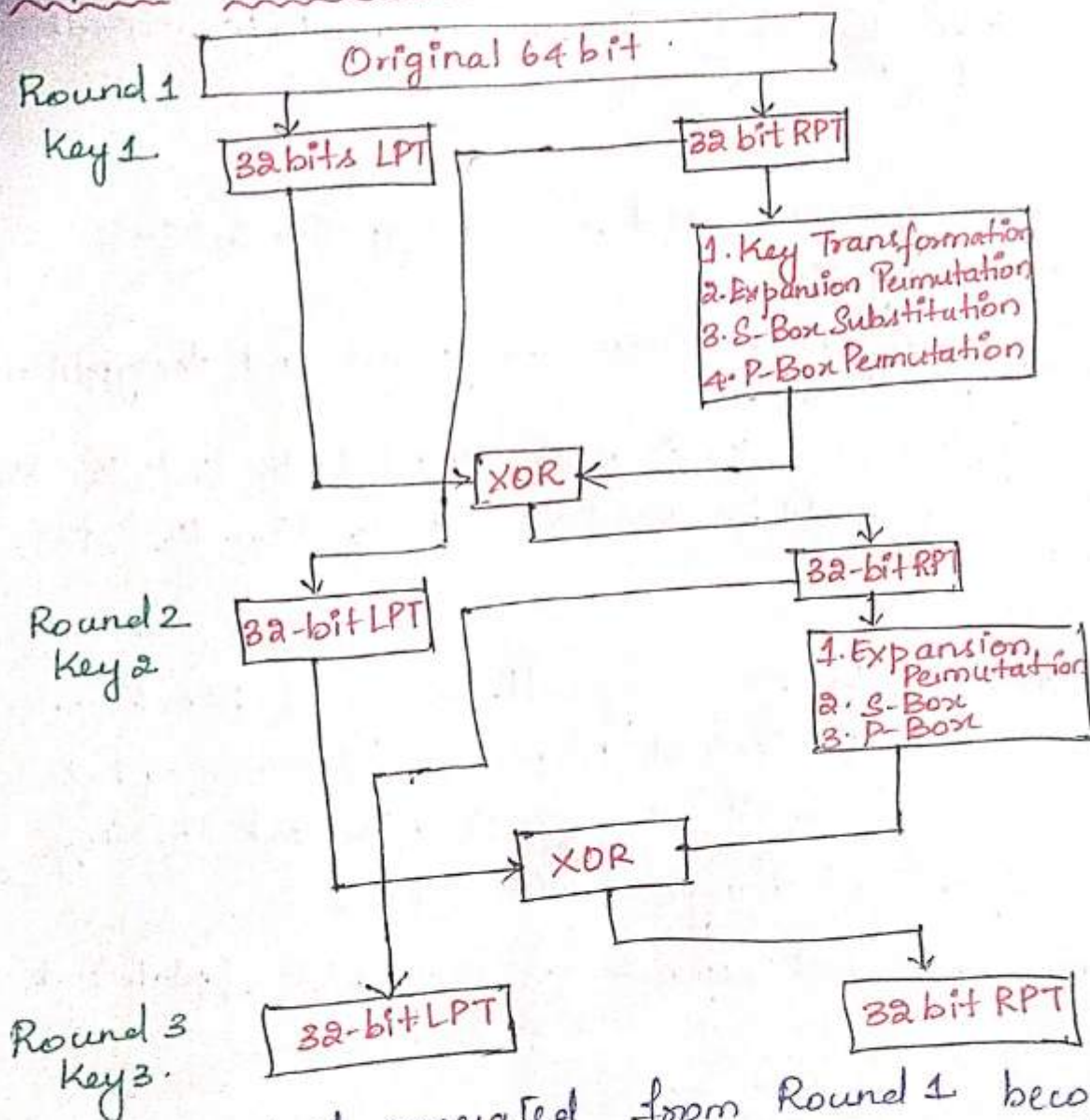
Finally the 48bit RPT is converted into 32 bits RPT using S-Box.

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Step 4.4 - P-Box Permutation:

The 32bits RPT generated from S-Box undergoes P-Box Permutation and new 32bits RPT is generated.

Step 4.5: XOR and swap

* The output generated from Round 1 becomes new RPT. and the previous 32bits RPT is new 32bits LPT.

* The process is repeated 16 times with 16 different keys and 64bit output is generated.

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Step 5 : Final Permutation :

* The 64 bits output generated in the previous step undergoes final permutation and new 64 bits is generated.

* The output generated after final permutation is the 64 bit Cipher Text

DES Decryption :

* The same algorithm used for encryption in DES also works for decryption.

* The only difference between encryption and decryption is reverse of key position.

* For encryption key (K) is ~~used~~ divided into $K_1, K_2, K_3 \dots K_{16}$ for decryption key should be used as $K_{16}, K_{15}, K_{14} \dots K_1$ for all 16 rounds.

DES Analysis :

DES has proved to be a very well-designed block cipher. DES satisfies both the desired properties of block cipher.

Avalanche effect :- A small change in plaintext (or) key results in significant change in the ciphertext.

Completeness :-

Each bit of cipher text depends on many bits of plaintext.