Tutorial 01 - Basic C++, Basic OOP, Analysis

CS2040C Semester 2 2018/2019

Question 1

(Basic) List ADT

List Array ADT

Imagine *list* as a *chain* of beads

You can add beads to the front and back

(and in between)

We will elaborate more about "ADT" in Tut 02.

For now, understand the syntax of C++ classes.

List Array ADT - Methods

get(i)

Gets the *i*-th element from the front (0-indexed)

search(v)

Return the first index which contains v

List Array ADT - Methods

insert(i, v)

Insert element *v* at index *i*.

remove(i)

Remove the element at index i.

List Array ADT - Methods

printList()

Prints the list from front to back.

sortList()

Sort the list, in default ascending order.

Question 1a, 1b, 1c

```
class ListArray {
 private:
   int N;
   int A[10];
                         // question 1a
 public:
    ListArray(): N(0) {} // question 1b
   int get(int i) {
       return A[i]; // question 1c
```

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

Anything wrong with this line?

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

Anything wrong with this line?

Limited to 10 items!

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

What does this line mean?

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

What does this line mean?

Pass 0 to constructor of integer N. Effectively:

```
ListArray() {
    N = 0;
}
```

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

Anything potential issues with this line?

```
class ListArray {
  private:
    int N;
    int A[10];
  public:
    ListArray() : N(0) {}
    int get(int i) {
        return A[i];
```

Anything potential issues with this line?

No "safeguard"!

What if i > N - 1?
What if i is negative?

Question 1d, 1e, 1f

```
void insert(int i, int v) {
 if ((N == 10) || (i < 0) || (i > N)) // question 1d
    return;
 for (int j = i; j \le N-1; j++) // question 1e
   A[j+1] = A[j];
 A[i] = v;
 N++;
void remove(int i) {
 for (int j = i; j < N-1; j++)
                                       // question 1f
   A[j] = A[j+1];
 N--;
```

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
        A[j+1] = A[j];
    A[i] = v;
    N++;
}</pre>
```

What does this line mean?

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
        A[j+1] = A[j];
    A[i] = v;
    N++;
}</pre>
```

What does this line mean?

If ...
N is 10 or
i is negative or
i > N

Stop inserting.

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
        A[j+1] = A[j];
    A[i] = v;
    N++;
}</pre>
```

What does this line mean?

If ...
N is 10 or
i is negative or
i > N

Stop inserting.

Is there anything wrong with this?
How many possible values of

i are accepted?

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
        A[j+1] = A[j];
    A[i] = v;
    N++;
}</pre>
```

What does this line mean?

If ...
N is 10 or
i is negative or
i > N

Stop inserting.

i are accepted?

Is there anything wrong with this?
How many possible values of

N+1 possible insertion points

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
    A[j+1] = A[j];
  A[i] = v;
  N++;
}</pre>
```

Any potential issues with this line?

```
void insert(int i, int v) {
  if ((N == 10) || (i < 0) || (i > N))
    return;
  for (int j = i; j <= N-1; j++)
    A[j+1] = A[j];
  A[i] = v;
  N++;
}</pre>
```

Any potential issues with this line?

Values are being overwritten in the wrong order!

```
A[i+1] = A[i];
A[i+2] = A[i+1];
A[i+3] = A[i+2];
```

... etc

```
void remove(int i) {
  for (int j = i; j < N-1; j++)
     A[j] = A[j+1];
  N--;
}</pre>
```

Any potential issues with this line?

No. This is correct.

```
A[i] = A[i+1];
A[i+1] = A[i+2];
A[i+2] = A[i+3];
... etc;
```

Can also add safeguard i in [0..N-1]

```
void sortList() { // sort array A, question 1g
    // ...
}
```

Implement this routine using any sorting algorithm that you know!

```
void sortList() { // sort array A, question 1g
   for (int i = 0; i < N; i++) {
        for (int j = 1; j < N; j++) {
            if (A[j-1] > A[j])
                swap(A[j-1], A[j]);
```

Implement this routine using any sorting algorithm that you know!

Approach: manual implementation

```
void sortList() { // sort array A, question 1g
    sort(A, A+N);
}
```

Implement this routine using any sorting algorithm that you know!

Approach: Use std library

```
int main() {
    ListArray* LA = new ListArray();
    LA->insert(0, 5);
    LA->insert(0, 1);
    // ...
}
```

Can we just write ListArray LA; in this line?

```
int main() {
    ListArray* LA = new ListArray();
    LA->insert(0, 5);
    LA->insert(0, 1);
    // ...
}
```

Can we just write ListArray LA; in this line?

Yes, but we need to modify the way we call methods:

```
ListArray LA;
LA.insert(0, 5);
cout << LA.get(1) << endl;
```

What's the difference between the two methods?

```
ListArray LA;
LA.insert(0, 5);
cout << LA.get(1) << endl;</pre>
```

```
ListArray *LA = new ListArray();
LA->insert(0, 5);
cout << LA->get(1) << endl;</pre>
```

Question 3

Analysis/ Order of Growth

Complexity analysis

- Is a rough estimate of how execution time will grow with size of input. I.E. Order of growth
- Time complexity is commonly used as a metric for comparing the performance of different algorithms on the same task
- The expression ignores constants
- Just retain the largest power for each variable in the expression
- When order of growth is applied to measure memory consumption of algorithms, we call that space complexity

Real life application?

Single thread computer: ~108 operations/sec

Sorting **N** = 106 numbers

 $O(N^2)$

~ 2.5 hours

O(N log N)

~ 0.2 sec

Question 3

What is the bound to the following functions?

a.
$$F(n) = log(2^n) + \sqrt{n} + 100\ 000\ 000$$

b. $F(n) = n + \frac{1}{2}n + \frac{1}{3}n + \frac{1}{4}n + \dots + 1$

c.
$$G(n) = n + \frac{1}{2}n + \frac{1}{4}n + \frac{1}{8}n + \dots + 1$$

Question 3.a)

$$F(n) = log(2^{n}) + \sqrt{n} + 100\ 000\ 000$$

$$O(F(n)) = O(log(2^{n}) + \sqrt{n} + 100\ 000\ 000)$$

$$= O(n * log 2 + \sqrt{n})$$

$$= O(n + \sqrt{n})$$

$$= O(n)$$

Question 3.b)

$$F(n) = n + \frac{1}{2} n + \frac{1}{3} n + \frac{1}{4} n + \dots + 1$$

$$O(F(n)) = O(n + \frac{1}{2} n + \frac{1}{3} n + \frac{1}{4} n + \dots + 1)$$

$$= O(n (1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}))$$
How to deal with this?

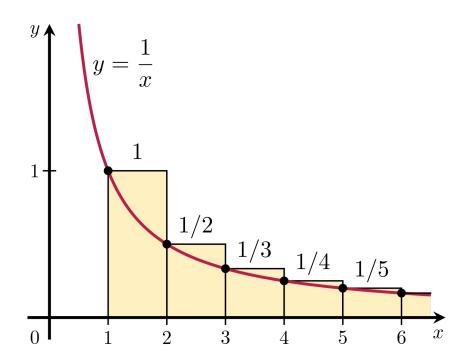
Harmonic Series

$$\sum_{n=1}^{\infty} \frac{1}{n} = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \cdots$$

$$y = \frac{1}{x}$$

Using integration:

$$\sum_{n=1}^k rac{1}{n} > \int_1^{k+1} rac{1}{x} \, dx = \ln(k+1).$$



Question 3.b)

From Harmonic Series

$$O(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}) = O(\log n)$$

$$O(F(n)) = O(n (1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}))$$

= $O(n \log n)$

Question 3.c)

$$G(n) = n + \frac{1}{2} n + \frac{1}{4} n + \frac{1}{8} n + \dots + 1$$

$$O(G(n)) = O(n + \frac{1}{2} n + \frac{1}{4} n + \frac{1}{8} n + \dots + 1)$$

$$= O(n (1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{n}))$$
How to deal with this?

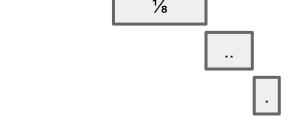
Convergent Geometric Series

1 1/2 1/4 1/4

How do we calculate the sum?

We can prove that the sum of the infinite geometric series exists if the **ratio** is a number between (but not including) -1 and 1, and r should not be equal to 0. The sum is given by the formula:

$$\sum_{k=0}^{\infty} ar^k = a + ar + ar^2 + ar^3 + \dots = \frac{a}{1-r}$$



Question 3.c)

$$O(1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{n}) = O(1 / (1 - \frac{1}{2}))$$

$$= O(2)$$

$$= O(1)$$

$$O(G(n)) = O(n (1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{n}))$$

$$= O(n)$$

Additional questions

Example (AY17/18 S1 Midterm Paper)

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = N; i >= 1; i--) {
    for (int j = 1; j \leftarrow N/i; j++) {
         counter++;
cout << counter << endl;</pre>
```

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < i; j++) {
        counter += j;
cout << counter << endl;</pre>
```

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < N; j++) {
        counter++;
        i++;
cout << counter << endl;</pre>
```

Questions